

# Monic Morphisms

A morphism  $\alpha : x \rightarrow y$  is *monic* iff in the situation

$$z \begin{array}{c} \xrightarrow{\beta_1} \\ \xrightarrow{\beta_2} \end{array} x \xrightarrow{\alpha} y ,$$

$$\alpha\beta_1 = \alpha\beta_2 \rightarrow \beta_1 = \beta_2,$$

*i.e. it can be cancelled on the left.*

*Remark.* This definition is [dual](#) to that of [epic morphisms](#).

**Theorem 1.** *In a concrete category, every injective morphism is monic.*

*Proof.* Let  $\alpha : X \rightarrow Y$  be an injective morphism in a concrete category  $\mathcal{C}$ . Let  $Z \in \mathcal{C}$  and  $\beta_1, \beta_2 : Z \rightarrow X$  such that  $\alpha\beta_1 = \alpha\beta_2$ . For each  $z \in Z$ ,

$$\alpha(\beta_1(z)) = \alpha\beta_1(z) = \alpha\beta_2(z) = \alpha(\beta_2(z)),$$

meaning  $\beta_1(z) = \beta_2(z)$  by the injectivity of  $\alpha$ . Since this applies for all  $z \in Z$ , this gives  $\beta_1 = \beta_2$ .

**Theorem 2.** *In the [category](#) of [sets](#), [topologies](#), groups, and rings, a morphism is injective iff it is monic.*