

ADVANCE COMPUTER NETWORK LAB. (CSPC-352)



Submitted to: Dr. Kunwar Pal

Submitted By: BALJOT SINGH 18103026

GAGANDEEP SINGH 18103031

GAURAV BHARDWAJ 18103034

JAGNOOR SINGH 18103045

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Group No: 9

Group: G2

Performance Comparison of TCP Variants using ns3.

Theory:

TCP stands for **Transmission Control Protocol**. It is a transport layer protocol that facilitates the transmission of packets from source to destination. It is a connection-oriented protocol that means it establishes the connection prior to the communication that occurs between the computing devices in a network.

Features of TCP protocol

- **Transport Layer Protocol**
- **Reliable**
- **Order of the data is maintained**
- **Connection-oriented**
- **Full duplex**
- **Stream-oriented**

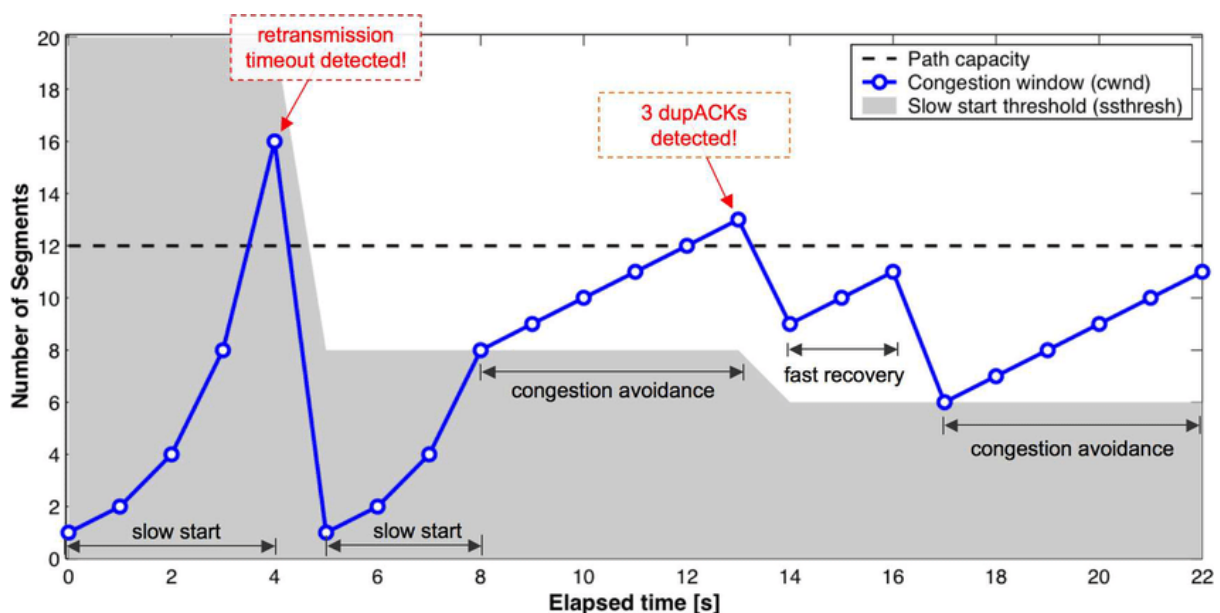
TCP Congestion Policy:

- TCP uses a congestion window and a congestion policy that avoid congestion
- **Slow Start, Exponential Increase:**
 - This algorithm is based on the idea that the size of the congestion window (cwnd) starts with one maximum segment size (MSS). The MSS is determined during connection establishment by using an option of the same name. The size of the window increases one MSS each time an acknowledgment is received. As the name implies, the window starts slowly, but grows exponentially.
 - Slow start cannot continue indefinitely. There must be a threshold to stop this phase. The sender keeps track of a variable named ssthresh (slow-start threshold). When the size of window in bytes reaches this threshold, slow start stops and the next phase starts. In most implementations the value of ssthresh is 65,535 bytes.
- **Congestion Avoidance, Additive Increase:**
 - If we start with the slow-start algorithm, the size of the congestion window increases exponentially. To avoid congestion before it happens, one must slow down this exponential growth. TCP defines another algorithm called congestion avoidance, which undergoes an additive increase instead of an exponential one.
 - When the size of the congestion window reaches the slow-start threshold, the slow-start phase stops and the additive phase begins. In this algorithm, each time

the whole window of segments is acknowledged (one round), the size of the congestion window is increased by 1.

➤ Congestion Detection:

- **Multiplicative Decrease** If congestion occurs, the congestion window size must be decreased. The only way the sender can guess that congestion has occurred is by the need to retransmit a segment.
 - Most TCP implementations have two reactions:
- ➔ If a time-out occurs, there is a stronger possibility of congestion; a segment has probably been dropped in the network, and there is no news about the sent segments.
- In this case TCP reacts strongly:
 - It sets the value of the threshold to one-half of the current window size.
 - It sets cwnd to the size of one segment.
 - It starts the slow-start phase again.
- ➔ If three ACKs are received, there is a weaker possibility of congestion; a segment may have been dropped, but some segments after that may have arrived safely since three ACKs are received. This is called **fast transmission and fast recovery**.
- In this case, TCP has a weaker reaction:
 - It sets the value of the threshold to one-half of the current window size.
 - It sets cwnd to the value of the threshold (some implementations add three segment sizes to the threshold).
 - It starts the congestion avoidance phase.

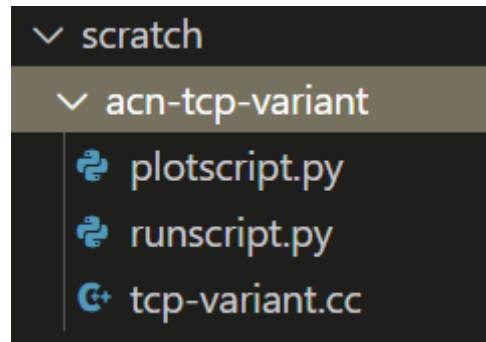


Requirements:

- Linux system
- Ns-3 network simulation.
- Python 3
- Numpy, Pandas, Matplotlib for graph plotting

Installation Steps:

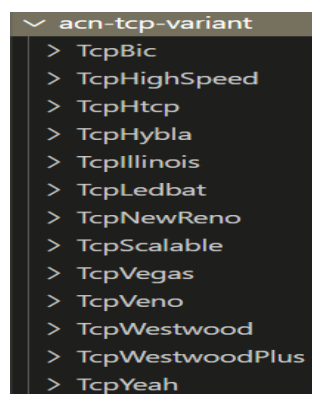
- Go to [link](#) and place the folder in scratch directory of ns3



- Navigate to ns3 folder /ns-3.30 and open terminal
- Run **Python3** `scratch/acn-tcp-variant/runscript.py`

```
baljot@BALJOT:~/myfolder/software/tar/ns-allinone-3.30/ns-3.30$ python3 scratch/acn-tcp-variant/runscript.py
Waf: Entering directory `/home/baljot/myfolder/software/tar/ns-allinone-3.30/ns-3.30/build'
Waf: Leaving directory `/home/baljot/myfolder/software/tar/ns-allinone-3.30/ns-3.30/build'
Build commands will be stored in build/compile_commands.json
'build' finished successfully (0.987s)
TcpNewReno Simulation
Waf: Entering directory `/home/baljot/myfolder/software/tar/ns-allinone-3.30/ns-3.30/build'
Waf: Leaving directory `/home/baljot/myfolder/software/tar/ns-allinone-3.30/ns-3.30/build'
Build commands will be stored in build/compile_commands.json
'build' finished successfully (1.038s)
TcpHybla Simulation
Waf: Entering directory `/home/baljot/myfolder/software/tar/ns-allinone-3.30/ns-3.30/build'
Waf: Leaving directory `/home/baljot/myfolder/software/tar/ns-allinone-3.30/ns-3.30/build'
```

- After this you can see various folder with name < tcp variant > in acn-tcp-variant folder



➤ **Content of each <tcp-variant> folder**

≡ ack.data	ack.data : acknowledged sequence no with time
🔥 anim.xml	ascii.tr : Ascii tracing file
≡ ascii.tr	
≡ cong-state.data	*.pcap : Packet capture file
≡ cwnd.data	
🔥 flow.xml	anim.xml : Animator file of Visualizing
≡ link1-0-1.pcap	Flow.xml : Flow Monitor File
≡ link1-0-2.pcap	
≡ link1-1-1.pcap	Cwnd.data : Congestion window data file
≡ link1-2-1.pcap	
≡ link2-0-1.pcap	Ssth.data : slow-start threshold data file
≡ link2-0-2.pcap	
≡ link2-1-1.pcap	Rtt.data : round-trip time data file
≡ link2-2-1.pcap	
≡ rtt.data	Cong-state.data : Congestion state data file
≡ ssth.data	

➤ **Navigate to ns-3.30 folder and Run `python3 scratch/acn-tcp-variant/plotscript.py`**

```
baljot@BALJOT:~/myfolder/software/tar/ns-allinone-3.30/ns-3.30$ python3 scratch/acn-tcp-variant/plotscript.py
```

- **Now you can see plot.png in each <tcp-variant> folder and ack.png in ack.png in acn-tcp-variant folder these file will be shown later**
- **To parse flow monitor file either open in netanim and change stats to flow monitor file or copy src/flow-monitor/examples/flowmon-parse-result.py into ns3 folder and Run command `python2 flowmon-parse-result.py scratch/acn-tcp-variant/<name>/flow.xml`**

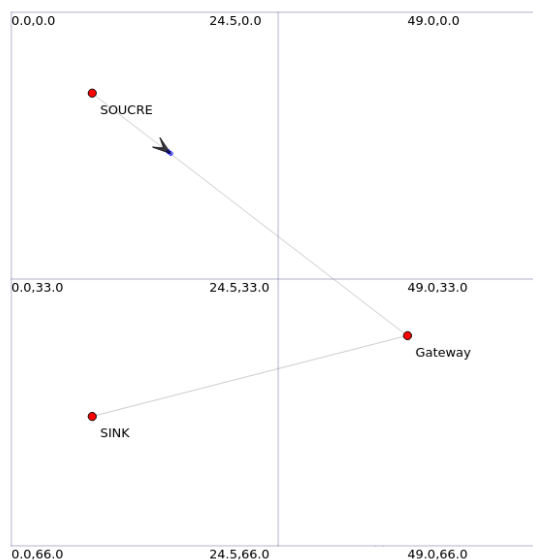
Using Terminal

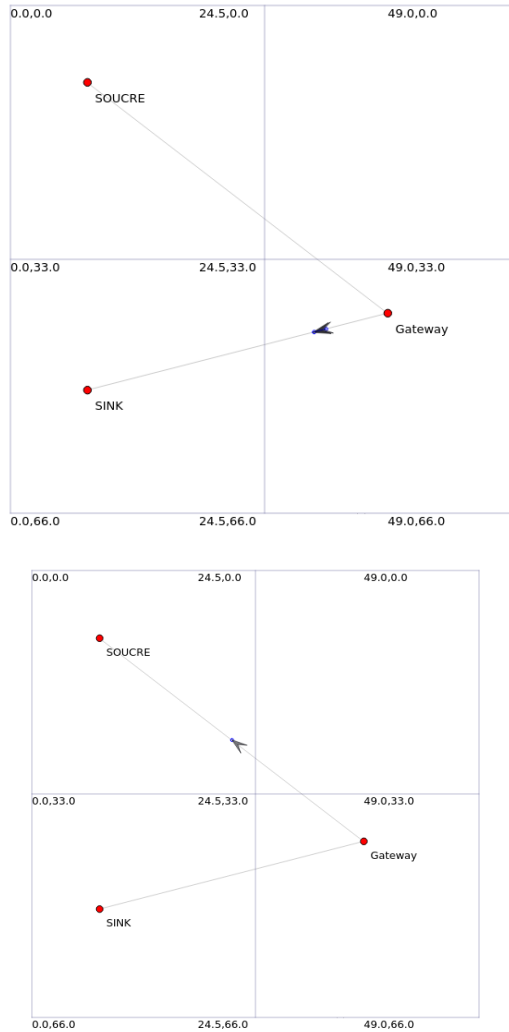
```
Reading XML file . done.
FlowID: 1 (TCP 10.1.1.1/49153 --> 10.1.2.2/2430)
TX bitrate: 189.90 kbit/s
RX bitrate: 186.75 kbit/s
Mean Delay: 93.07 ms
Packet Loss Ratio: 0.35 %
FlowID: 2 (TCP 10.1.2.2/2430 --> 10.1.1.1/49153)
TX bitrate: 14.93 kbit/s
RX bitrate: 14.67 kbit/s
Mean Delay: 90.27 ms
Packet Loss Ratio: 0.00 %
```

Using NetAnim

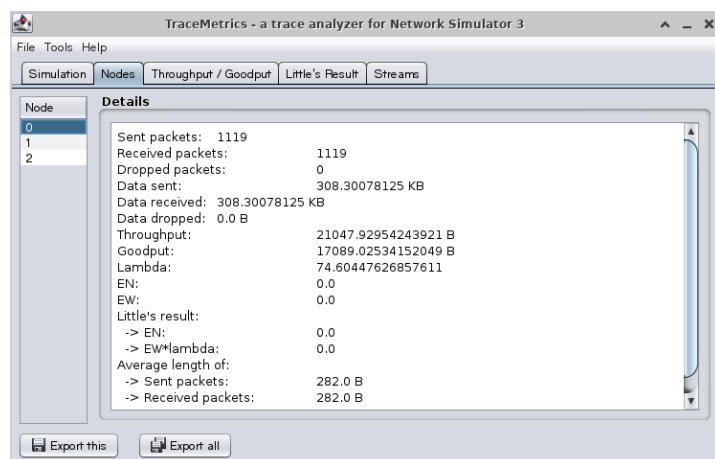


➤ To visualize Use netanim and open anim.xml file in NetAnim





➤ To Analyse Ascii trace download tracemetrics and open file in it



TraceMetrics - a trace analyzer for Network Simulator 3

File Tools Help

Simulation Nodes **Throughput / Goodput** Little's Result Streams

Node	Throughput	Goodput
0	21047.92954243921	17089.02534152049
1	19564.640545099373	17089.02534152049
2	1519.557840137075	0.0

TraceMetrics - a trace analyzer for Network Simulator 3

File Tools Help

Simulation Nodes **Throughput / Goodput** Little's Result Streams

Node

Details

Sent packets:	415
Received packets:	702
Dropped packets:	0
Data sent:	22.2578125 KB
Data received:	281.74609375 KB
Data dropped:	0.0 B
Throughput:	1519.557840137075 B
Goodput:	0.0 B
Lambda:	27.668326766272642
EN:	0.0
EW:	0.0
Little's result:	
-> EN:	0.0
-> EW*lambda:	0.0
Average lambda of:	
-> Sent packets:	54.0 B
-> Received packets:	410.0 B

Export this Export all

➤ To analyse Pcap file open file in wireshark

link1-0-1.pcap

File Edit View Go Capture Analyze Statistics Telephony Wireless Tools Help

Apply a display filter ... <Ctrl-F>

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000	10.1.1.1	10.1.2.2	TCP	58	49153 → 2430 [SYN] Seq=0 Win=65535 Len=0 TSval=90 TSecr=0 WS=4
2	0.000454	10.1.2.2	10.1.1.1	TCP	58	2430 → 49153 [ACK] Seq=0 Ack=1 Win=65535 Len=0 TSval=90 TSecr=90
3	0.180553	10.1.1.1	10.1.2.2	TCP	54	49153 → 2430 [ACK] Seq=1 Ack=1 Win=131072 Len=0 TSval=180 TSecr=90
4	0.180885	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=1 Ack=1 Win=131072 Len=360 TSval=180 TSecr=90
5	0.362235	10.1.2.2	10.1.1.1	TCP	54	2430 → 49153 [ACK] Seq=361 Ack=1 Win=131072 Len=0 TSval=362 TSecr=180
6	0.363131	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=721 Ack=1 Win=131072 Len=360 TSval=362 TSecr=362
7	0.363402	10.1.1.1	10.1.2.2	TCP	54	2430 → 49153 [ACK] Seq=1 Ack=1081 Win=131072 Len=0 TSval=456 TSecr=362
8	0.456659	10.1.2.2	10.1.1.1	TCP	414	49153 → 2430 [ACK] Seq=1081 Ack=1 Win=131072 Len=360 TSval=54 TSecr=456
9	0.547033	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=1441 Ack=1 Win=131072 Len=360 TSval=54 TSecr=54
10	0.547385	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=1801 Ack=1 Win=131072 Len=360 TSval=54 TSecr=54
11	0.547696	10.1.1.1	10.1.2.2	TCP	54	2430 → 49153 [ACK] Seq=1 Ack=1801 Win=131072 Len=0 TSval=640 TSecr=54
12	0.640551	10.1.2.2	10.1.1.1	TCP	414	49153 → 2430 [ACK] Seq=2161 Ack=1 Win=131072 Len=360 TSval=73 TSecr=640
13	0.730936	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=2521 Ack=1 Win=131072 Len=360 TSval=73 TSecr=73
14	0.731267	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=2881 Ack=1 Win=131072 Len=360 TSval=73 TSecr=73
15	0.731598	10.1.1.1	10.1.2.2	TCP	54	2430 → 49153 [ACK] Seq=1 Ack=2521 Win=131072 Len=0 TSval=822 TSecr=73
16	0.822808	10.1.2.2	10.1.1.1	TCP	414	49153 → 2430 [ACK] Seq=3241 Ack=1 Win=131072 Len=360 TSval=91 TSecr=822
17	0.913182	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=3601 Ack=1 Win=131072 Len=360 TSval=91 TSecr=91
18	0.913513	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=3961 Ack=1 Win=131072 Len=360 TSval=91 TSecr=91
19	0.913845	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=4321 Ack=1 Win=131072 Len=360 TSval=10 TSecr=91
20	1.005086	10.1.2.2	10.1.1.1	TCP	62	2430 → 49153 [ACK] Seq=1 Ack=2881 Win=131072 Len=0 TSval=1064 TSecr=91
21	1.006742	10.1.2.2	10.1.1.1	TCP	62	[TCP Dup ACK 20#1] 2430 → 49153 [ACK] Seq=1 Ack=2881 Win=131072 Len=0 TSval=10 TSecr=91
22	1.005467	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=4681 Ack=1 Win=131072 Len=360 TSval=10 TSecr=1064
23	1.005798	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=5041 Ack=1 Win=131072 Len=360 TSval=10 TSecr=1064
24	1.006129	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=5401 Ack=1 Win=131072 Len=360 TSval=10 TSecr=1064
25	1.007123	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=5761 Ack=1 Win=131072 Len=360 TSval=10 TSecr=1064
26	1.191403	10.1.2.2	10.1.1.1	TCP	70	[TCP Dup ACK 20#2] 2430 → 49153 [ACK] Seq=1 Ack=2881 Win=131072 Len=0 TSval=1310 TSecr=1064
27	1.189059	10.1.2.2	10.1.1.1	TCP	70	[TCP Dup ACK 20#3] 2430 → 49153 [ACK] Seq=1 Ack=2881 Win=131072 Len=0 TSval=1310 TSecr=1064
28	1.190715	10.1.2.2	10.1.1.1	TCP	70	[TCP Dup ACK 20#4] 2430 → 49153 [ACK] Seq=1 Ack=2881 Win=131072 Len=0 TSval=1310 TSecr=1064
29	1.192371	10.1.2.2	10.1.1.1	TCP	70	[TCP Dup ACK 20#5] 2430 → 49153 [ACK] Seq=1 Ack=2881 Win=131072 Len=0 TSval=1310 TSecr=1064
30	1.277790	10.1.1.1	10.1.2.2	TCP	414	[TCP Retransmission] 49153 → 2430 [ACK] Seq=2881 Ack=1 Win=131072 Len=360 TSval=14 TSecr=1310
31	1.282758	10.1.1.1	10.1.2.2	TCP	414	[TCP Retransmission] 49153 → 2430 [ACK] Seq=3241 Ack=1 Win=131072 Len=360 TSval=14 TSecr=1310
32	1.369094	10.1.2.2	10.1.1.1	TCP	62	2430 → 49153 [ACK] Seq=1 Ack=3961 Win=131072 Len=0 TSval=1369 TSecr=1310
33	1.374630	10.1.2.2	10.1.1.1	TCP	54	2430 → 49153 [ACK] Seq=1 Ack=5761 Win=131072 Len=0 TSval=1374 TSecr=1369
34	1.400075	10.1.1.1	10.1.2.2	TCP	414	[TCP Spurious Retransmission] 49153 → 2430 [ACK] Seq=3961 Ack=1 Win=131072 Len=360 TSval=14 TSecr=1374
35	1.460406	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=5761 Ack=1 Win=131072 Len=360 TSval=14 TSecr=1374
36	1.465005	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=6121 Ack=1 Win=131072 Len=360 TSval=14 TSecr=1374
37	1.551947	10.1.2.2	10.1.1.1	TCP	94	[TCP Dup ACK 33#1] 2430 → 49153 [ACK] Seq=1 Ack=5761 Win=131072 Len=0 TSval=1556 TSecr=1374
38	1.556877	10.1.2.2	10.1.1.1	TCP	54	2430 → 49153 [ACK] Seq=1 Ack=6481 Win=131072 Len=0 TSval=1556 TSecr=1374
39	1.647251	10.1.1.1	10.1.2.2	TCP	414	49153 → 2430 [ACK] Seq=6481 Ack=1 Win=131072 Len=360 TSval=16 TSecr=1556

Frame 5: 54 bytes on wire (432 bits), 54 bytes captured (432 bits)

```

0000  60 21 45 00 00 34 00 01 00 00 3f 06 00 00 0a 01  :E: 4...?.....
0010  02 02 0a 01 01 01 00 7e c0 01 00 00 00 01 00 00  :.....1.....
0020  01 69 00 10 00 00 00 00 00 00 08 0a 00 00 01 10  :.....
0030  00 00 00 b4 00 00  :.....

```


Code:

Tcp-variant.cc

```
/*
```

Topology used



```
*/
```

```
#include <iostream>
#include <fstream>
#include <string>
#include <sys/stat.h>
#include <sys/types.h>
#include "ns3/core-module.h"
#include "ns3/network-module.h"
#include "ns3/internet-module.h"
#include "ns3/point-to-point-module.h"
#include "ns3/applications-module.h"
#include "ns3/error-model.h"
#include "ns3/tcp-header.h"
#include "ns3/udp-header.h"
#include "ns3/enum.h"
#include "ns3/event-id.h"
#include "ns3/flow-monitor-helper.h"
#include "ns3/ipv4-global-routing-helper.h"
#include "ns3/netanim-module.h"
using namespace ns3;
```

```
NS_LOG_COMPONENT_DEFINE ("TcpVariantsComparison");
```

```
bool firstCwnd = true;
bool firstSshThr = true;
bool firstRtt=true;
// output streams
Ptr<OutputStreamWrapper> cWndStream;
Ptr<OutputStreamWrapper> ssThreshStream;
Ptr<OutputStreamWrapper> ackStream;
Ptr<OutputStreamWrapper> congStateStream;
Ptr<OutputStreamWrapper> rttStream;
```

```
uint32_t cWndValue = 0;
uint32_t ssThreshValue = 0;
uint32_t rtovalue=0;
```

```
static void
CwndTracer (uint32_t oldval, uint32_t newval)
{

    if (firstCwnd)
```

```

{
    *cWndStream->GetStream () << "0.0 " << oldval << std::endl;
    firstCwnd = false;
}
*cWndStream->GetStream () << Simulator::Now ().GetSeconds () << " " << newval << std::endl;
cWndValue = newval;

if (!firstSshThr)
{
    *ssThreshStream->GetStream () << Simulator::Now ().GetSeconds () << " " << ssThreshValue << std::endl;
}
}

static void
SsThreshTracer (uint32_t oldval, uint32_t newval)
{
    if (firstSshThr)
    {
        *ssThreshStream->GetStream () << "0.0 " << oldval << std::endl;
        firstSshThr = false;
    }
    *ssThreshStream->GetStream () << Simulator::Now ().GetSeconds () << " " << newval << std::endl;
    ssThreshValue = newval;

    if (!firstCwnd)
    {
        *cWndStream->GetStream () << Simulator::Now ().GetSeconds () << " " << cWndValue << std::endl;
    }
}

static void
RttTracer (Time oldval, Time newval)
{
    if (firstRtt)
    {
        *rttStream->GetStream () << "0.0 " << oldval.GetSeconds () << std::endl;
        firstRtt = false;
    }
    *rttStream->GetStream () << Simulator::Now ().GetSeconds () << " " << newval.GetSeconds () << std::endl;
}

static void
CongStateTracer (TcpSocketState::TcpCongState_t old, TcpSocketState::TcpCongState_t newState)
{
    *congStateStream->GetStream () << Simulator::Now ().GetSeconds () << " " << newState << std::endl;
}

static void
AckTracer (SequenceNumber32 old, SequenceNumber32 newAck)
{
    *ackStream->GetStream () << Simulator::Now ().GetSeconds () << " " << newAck << std::endl;
}

static void
TraceCwnd (std::string cwnd_tr_file_name)
{
    AsciiTraceHelper ascii;
    cWndStream = ascii.CreateFileStream (cwnd_tr_file_name.c_str ());

```

```

    Config::ConnectWithoutContext ("/NodeList/1/$ns3::TcpL4Protocol/SocketList/0/CongestionWindow", MakeCallback (&
CwndTracer));
}

```

static void

```

TraceSsThresh (std::string ssthresh_tr_file_name)
{
    AsciiTraceHelper ascii;
    ssThreshStream = ascii.CreateFileStream (ssthresh_tr_file_name.c_str ());
    Config::ConnectWithoutContext ("/NodeList/1/$ns3::TcpL4Protocol/SocketList/0/SlowStartThreshold", MakeCallback (&
SsThreshTracer));
}

```

static void

```

TraceRtt (std::string rtt_tr_file_name)
{
    AsciiTraceHelper ascii;
    rttStream = ascii.CreateFileStream (rtt_tr_file_name.c_str ());
    Config::ConnectWithoutContext ("/NodeList/1/$ns3::TcpL4Protocol/SocketList/0/RTT", MakeCallback (&RttTracer));
}

```

static void

```

TraceAck (std::string &ack_file_name)
{
    AsciiTraceHelper ascii;
    ackStream = ascii.CreateFileStream (ack_file_name.c_str ());
    Config::ConnectWithoutContext ("/NodeList/1/$ns3::TcpL4Protocol/SocketList/0/HighestRxAck", MakeCallback (&AckT
racer));
}

```

static void

```

TraceCongState (std::string &cong_state_file_name)
{
    AsciiTraceHelper ascii;
    congStateStream = ascii.CreateFileStream (cong_state_file_name.c_str ());
    Config::ConnectWithoutContext ("/NodeList/1/$ns3::TcpL4Protocol/SocketList/0/CongState", MakeCallback (&CongStat
eTracer));
}

```

int main (int argc, char *argv[])

```

{
    std::string transport_prot = "TcpWestwood";
    double error_p = 0.0;
    std::string bandwidth = "2Mbps";
    std::string delay = "0.01ms";
    std::string access_bandwidth = "10Mbps";
    std::string access_delay = "45ms";

```

```

    double data_mbytes = 0;
    uint32_t mtu_bytes = 400;
    float duration = 15;
    uint32_t run = 0;
    uint32_t cbr=0;

```

CommandLine cmd;

```

cmd.AddValue ("transport_prot", "Transport protocol to use: TcpNewReno, TcpLinuxReno, "

```

```

        "TcpHybla, TcpHighSpeed, TcpHtcp, TcpVegas, TcpScalable, TcpVeno, "
        "TcpBic, TcpYeah, TcpIllinois, TcpWestwood, TcpWestwoodPlus, TcpLedbat", transport_prot);
cmd.AddValue ("error_p", "Packet error rate", error_p);
cmd.AddValue ("bandwidth", "Bottleneck bandwidth", bandwidth);
cmd.AddValue ("delay", "Bottleneck bandwidth", delay);
cmd.AddValue ("access_bandwidth", "Access link bandwidth", access_bandwidth);
cmd.AddValue ("access_delay", "Access link delay", access_delay);
cmd.AddValue ("data", "Number of Megabytes of data to transmit", data_mbytes);
cmd.AddValue ("mtu", "Size of IP packets to send in bytes", mtu_bytes);
cmd.AddValue ("cbr", "No of cbr traffic genrator", cbr);

cmd.Parse (argc, argv);
NS_LOG_UNCOND(transport_prot+" Simulation");
std::string file="scratch/acn-tcp-variant/"+transport_prot;
mkdir(file.c_str(), 0777);
SeedManager::SetSeed (1);
SeedManager::SetRun (run);

// User may find it convenient to enable logging
LogComponentEnable("TcpVariantsComparison", LOG_LEVEL_ALL);

// Calculate the ADU size
Header* temp_header = new Ipv4Header ();
uint32_t ip_header = temp_header->GetSerializedSize ();
//NS_LOG_LOGIC ("IP Header size is: " << ip_header);
delete temp_header;
temp_header = new TcpHeader ();
uint32_t tcp_header = temp_header->GetSerializedSize ();
//NS_LOG_LOGIC ("TCP Header size is: " << tcp_header);
delete temp_header;
uint32_t tcp_adu_size = mtu_bytes - (ip_header + tcp_header);
//NS_LOG_LOGIC ("TCP ADU size is: " << tcp_adu_size);

Config::SetDefault ("ns3::TcpSocket::SegmentSize", UIntegerValue (tcp_adu_size));

// Set the simulation start and stop time
float start_time = 0.0;
float stop_time = start_time + duration;
transport_prot = std::string ("ns3::") + transport_prot;

// Select TCP variant
if (transport_prot.compare ("ns3::TcpWestwoodPlus") == 0)
{
    // TcpWestwoodPlus is not an actual TypeId name; we need TcpWestwood here
    Config::SetDefault ("ns3::TcpL4Protocol::SocketType", TypeIdValue (TcpWestwood::GetTypeId ()));
    // the default protocol type in ns3::TcpWestwood is WESTWOOD
    Config::SetDefault ("ns3::TcpWestwood::ProtocolType", EnumValue (TcpWestwood::WESTWOODPLUS));
}
else
{
    TypeId tcpTid;
    Config::SetDefault ("ns3::TcpL4Protocol::SocketType", TypeIdValue (TypeId::LookupByName (transport_prot)));
}

// Create gateways, sources, and sinks
NodeContainer gateways;

```

```

gateways.Create (1);
NodeContainer sources;
sources.Create (1);
NodeContainer sinks;
sinks.Create (1);

// Configure the error model
// Here we use RateErrorModel with packet error rate
Ptr<UniformRandomVariable> uv = CreateObject<UniformRandomVariable> ();
uv->SetStream (50);
RateErrorModel error_model;
error_model.SetRandomVariable (uv);
error_model.SetUnit (RateErrorModel::ERROR_UNIT_PACKET);
error_model.SetRate (error_p);

InternetStackHelper stack;
stack.InstallAll ();

// Configure the sources and sinks net devices
// and the channels between the sources/sinks and the gateways
PointToPointHelper link1;
link1.SetDeviceAttribute ("DataRate", StringValue (access_bandwidth));
link1.SetChannelAttribute ("Delay", StringValue (access_delay));

PointToPointHelper link2;
link2.SetDeviceAttribute ("DataRate", StringValue (bandwidth));
link2.SetChannelAttribute ("Delay", StringValue (access_delay));
link2.SetDeviceAttribute ("ReceiveErrorModel", PointerValue (&error_model));

NetDeviceContainer devices1,devices2;
devices1 = link1.Install (sources.Get (0), gateways.Get (0));
devices2=link2.Install(gateways.Get (0), sinks.Get (0));

Ipv4AddressHelper address;

Ipv4InterfaceContainer interfaces1,interfaces2;

address.SetBase ("10.1.1.0", "255.255.255.0");
interfaces1 = address.Assign (devices1);
address.SetBase ("10.1.2.0", "255.255.255.0");
interfaces2= address.Assign (devices2);

//NS_LOG_INFO ("Initialize Global Routing.");
Ipv4GlobalRoutingHelper::PopulateRoutingTables ();

```

```

uint16_t port = 2430;

Address sinkLocalAddress (InetSocketAddress (Ipv4Address::GetAny (), port));
PacketSinkHelper sinkHelper ("ns3::TcpSocketFactory", sinkLocalAddress);
sinkHelper.SetAttribute ("Protocol", TypeIdValue (TcpSocketFactory::GetTypeId ()));

AddressValue remoteAddress (InetSocketAddress (interfaces2.GetAddress (1), port));

// source traffic generator
BulkSendHelper source ("ns3::TcpSocketFactory", Address ());

source.SetAttribute ("Remote", remoteAddress);
source.SetAttribute ("SendSize", UintegerValue (tcp_adu_size));
// 0 means unlimited
source.SetAttribute ("MaxBytes", UintegerValue (int(data_mbytes * 1000000)));

ApplicationContainer sourceApp = source.Install (sources.Get (0));

sourceApp.Start (Seconds (start_time));
sourceApp.Stop (Seconds (stop_time));

ApplicationContainer sinkApp = sinkHelper.Install (sinks);
sinkApp.Start (Seconds (start_time));
sinkApp.Stop (Seconds (stop_time));

// traffic generator installed in node 0
ApplicationContainer cbrApps;
ApplicationContainer cbrSinkApps;
uint32_t cbrPort=6759;

for(uint32_t i=0; i<cbr; i++)
{
    OnOffHelper onOffHelper ("ns3::UdpSocketFactory", InetSocketAddress (interfaces2.GetAddress (1), cbrPort));
    onOffHelper.SetAttribute ("PacketSize", UintegerValue (1024));
    onOffHelper.SetAttribute ("DataRate", StringValue ("1Mbps"));
    cbrApps.Add (onOffHelper.Install (gateways.Get (0)));
    cbrApps.Start (Seconds (7));
    cbrApps.Stop (Seconds (10));

}

PacketSinkHelper sink ("ns3::UdpSocketFactory", InetSocketAddress (Ipv4Address::GetAny (), cbrPort));
cbrSinkApps.Add(sink.Install (sinks.Get (0)));
cbrSinkApps.Start (Seconds (5));
cbrSinkApps.Stop (Seconds (10));

// Set up tracing
AsciiTraceHelper ascii;
stack.EnableAsciiIpv4All (ascii.CreateFileStream(file+"/ascii.tr"));
Simulator::Schedule (Seconds (0.00001), &TraceCwnd, file + "/cwnd.data");

```

```

Simulator::Schedule (Seconds (0.00001), &TraceSsThresh, file + "/ss.th.data");
Simulator::Schedule (Seconds (0.00001), &TraceAck, file + "/ack.data");
Simulator::Schedule (Seconds (0.00001), &TraceCongState, file + "/cong-state.data");
Simulator::Schedule (Seconds (0.00001), &TraceRtt, file + "/rtt.data");

link2.EnablePcapAll (file+"/link2", true);
link1.EnablePcapAll (file+"/link1", true);

// Flow monitor
FlowMonitorHelper flowHelper;
flowHelper.InstallAll ();

//net anim
AnimationInterface anim(file+"/anim.xml"); //to save file for netAnim
anim.UpdateNodeDescription(sources.Get(0), "SOUCRE");

anim.UpdateNodeDescription(gateways.Get(0), "Gateway");

anim.UpdateNodeDescription(sinks.Get(0), "SINK");

Simulator::Stop (Seconds (stop_time));
Simulator::Run ();
flowHelper.SerializeToXmlFile (file+"/flow.xml", true, true);
Simulator::Destroy ();
return 0;
}

```

PLOTSRIPT.PY

```

import pandas as pd
import numpy as np
import matplotlib.pyplot as plt

tcp_variants=[ "TcpNewReno", "TcpWestwood", "TcpVeno", "TcpHybla", "TcpHighSpeed", "TcpHtcp", "TcpVegas", "TcpScalable", "TcpBic", "TcpYeah", "TcpLedbat", "TcpIllinois", "TcpWestwoodPlus" ]

segment = 360

def get_data(type, protocol):
    data=pd.read_csv(f"scratch/acn-tcp-variant/{protocol}/{type}.data", sep="\s+", names=[0,1], engine='python')

    return data.T.to_numpy()

def wrappper(cwnd, statedata):
    xaxis=[]
    yaxis=[]
    counter=0
    state=0;
    for x in cwnd[0]:
        if counter!=len(statedata[0]) and statedata[0][counter]>=x:
            state=statedata[1][counter]
            counter+=1

```

```

    xaxis.append(x)
    yaxis.append(state)

```

```

return statedata

```

```

def plot(protocol):
    plt.figure(1,figsize=(12,12))
    plt.tight_layout()

    plt.subplot(2,1,1)
    cwnddata=get_data("cwnd",protocol)
    ssthdata=get_data("ssth",protocol)
    statedata=get_data("cong-state",protocol)
    plt.plot(cwnddata[0], cwnddata[1] / segment,color='darkblue', label='cwnd',linewidth=2)
    plt.plot(ssthdata[0], ssthdata[1] / segment,color='red', linestyle='dotted', label='ssth',linewidth=2)

    plt.ylim(0, cwnddata[1] .max()/segment + 20)
    plt.ylabel('segment(1seg=360bytes)')
    plt.xlabel('Time(sec)')

    fill_data=wrappper(cwnddata,statedata)
    plt.axvspan(0,15,facecolor='lightblue', alpha=0.1)
    for i in range(0,len(fill_data[0])-1):
        if fill_data[1][i] == 1: #disorder
            plt.axvspan(fill_data[0][i],fill_data[0][i+1]+0.1,facecolor='green', alpha=0.5)
        elif fill_data[1][i]== 3: #recovery
            plt.axvspan(fill_data[0][i],fill_data[0][i+1]+0.1,facecolor='yellow', alpha=0.5)
        elif fill_data[1][i]== 4: #Loss
            plt.axvspan(fill_data[0][i],fill_data[0][i+1]+0.1,facecolor='red', alpha=0.5)

    plt.plot([],[],color='lightblue',label='open')
    plt.plot([],[],color='yellow',label='recovery')
    plt.plot([],[],color='green',label='disorder')
    plt.plot([],[],color='red',label='loss')
    plt.title("CONGESTION WINDOW VS SLOW START THRESHOLD")
    plt.legend()

    plt.subplot(2,1,2)
    plt.subplots_adjust(wspace=2)
    rttdata=get_data("rtt",protocol)
    plt.plot(rttdata[0],rttdata[1])
    plt.title("ROUND TRIP TIME")
    plt.ylabel('RTT')
    plt.xlabel('Time(sec)')

    plt.suptitle(protocol)

    plt.savefig(f'scratch/acn-tcp-variant/{protocol}/plot.png')
    plt.close()

```

```

def plotack(protocol):
    plt.figure(500)
    data=get_data("ack",protocol)

```



```
plt.plot(data[0],data[1] / segment,label=protocol)
```

```
if __name__ == '__main__':

    for protocol in tcp_variants:
        plot(protocol)
        plotack(protocol)
    plt.figure(500,figsize=(10,10))
    plt.title("ACKNOWLEDGEMENT")
    plt.xlabel("TIME(ns)")
    plt.ylabel("SEGMENT(360bytes)")
    plt.legend()
    plt.savefig(f'scratch/acn-tcp-variant/ack-plot.png')
    plt.close()
    print("Potting done")
```

RUNSCRIPT.PY

```
#running tcp variants one by one
import os

tcp_variants=["TcpNewReno", "TcpHybla", "TcpHighSpeed","TcpHtcp","TcpVegas", "TcpScalable","TcpVeno", "TcpBic",
"TcpYeah", "TcpIllinois", "TcpWestwood", "TcpWestwoodPlus", "TcpLedbat"]
for protocol in tcp_variants:
    os.system(f'./waf --run "scratch/acn-tcp-variant/acn-tcp-variant --transport_prot={protocol} --error_p=0.00 --cbr=0" ')
```

NS3-CONGESTION-STATE

- **OPEN** : This is the so-called normal state. Some algorithms have two phases: Slow start (SS) and Congestion avoidance (CA).if cwnd is smaller than ssthresh, it corresponds to the **Slow start phase**, and if it is larger, it corresponds to the **Congestion avoidance phase**.
- **DISORDER** : A state in which a duplicate ACK (Duplicate ACK) has been received. There is a possibility of packet loss or disorder of reception order.
- **RECOVERY** : A state in which a triple duplicate ACK has been received. We are confident of packet loss and start resending.
- **LOSS** : RTT becomes larger than the retransmission time out (RTO), that is, the ACK timeout is detected. There may be serious congestion.

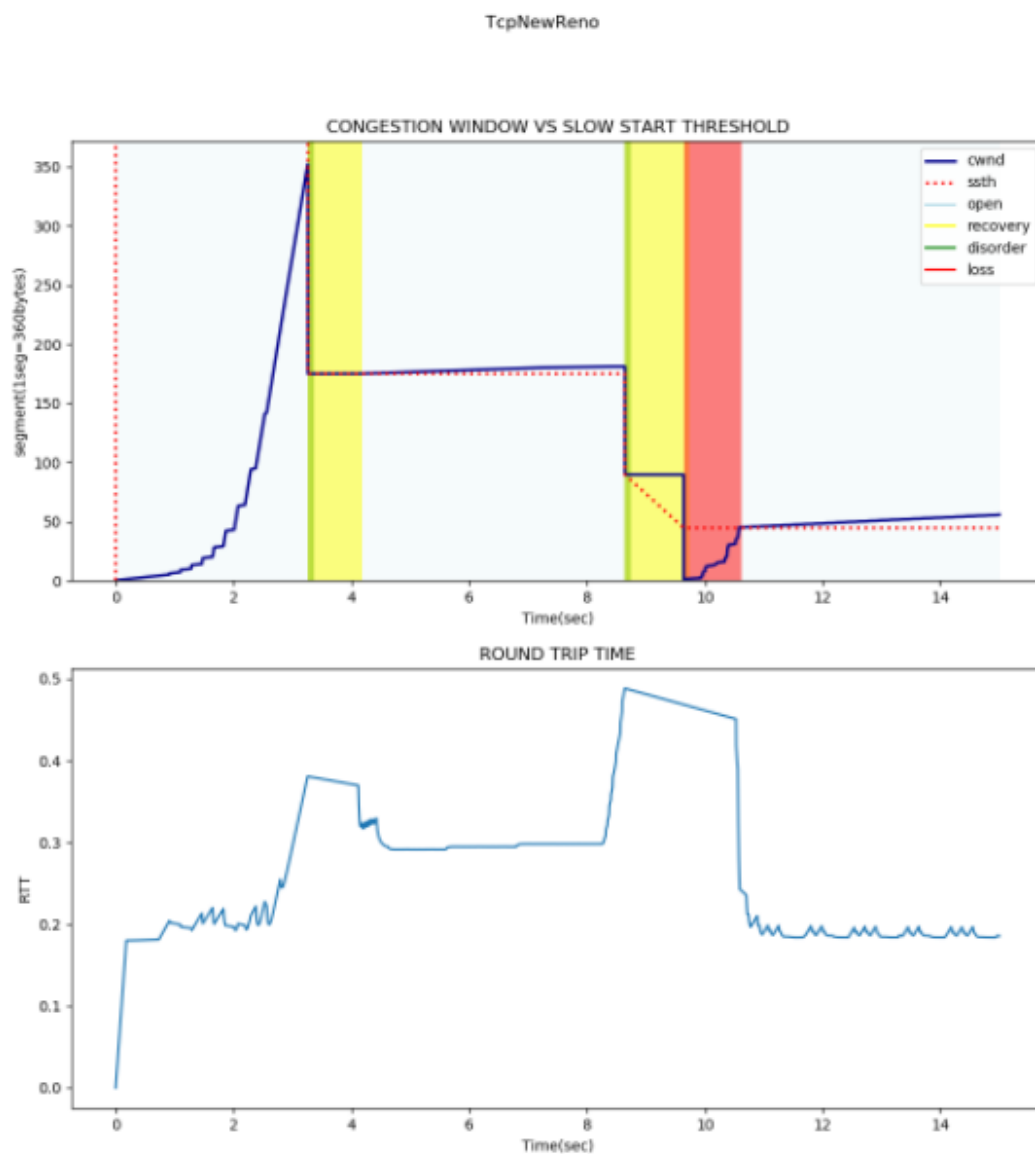
TCP VARIANTS USED IN NS3 SIMULATION

Note: To see every congestion state we have used 5 constant bit rate traffic generator between 7 to 10 seconds to cause serious congestion state in link2

➤ NewReno

TCP New Reno is efficient as compare to Tahoe and Reno. In Reno Fast Recovery exits when three duplicate acknowledgements are received but in New Reno it does not exist until all outstanding data is acknowledged or ends when retransmission timeout occurs. In this way it avoids unnecessary multiple fast retransmit from single window data.

Plot:

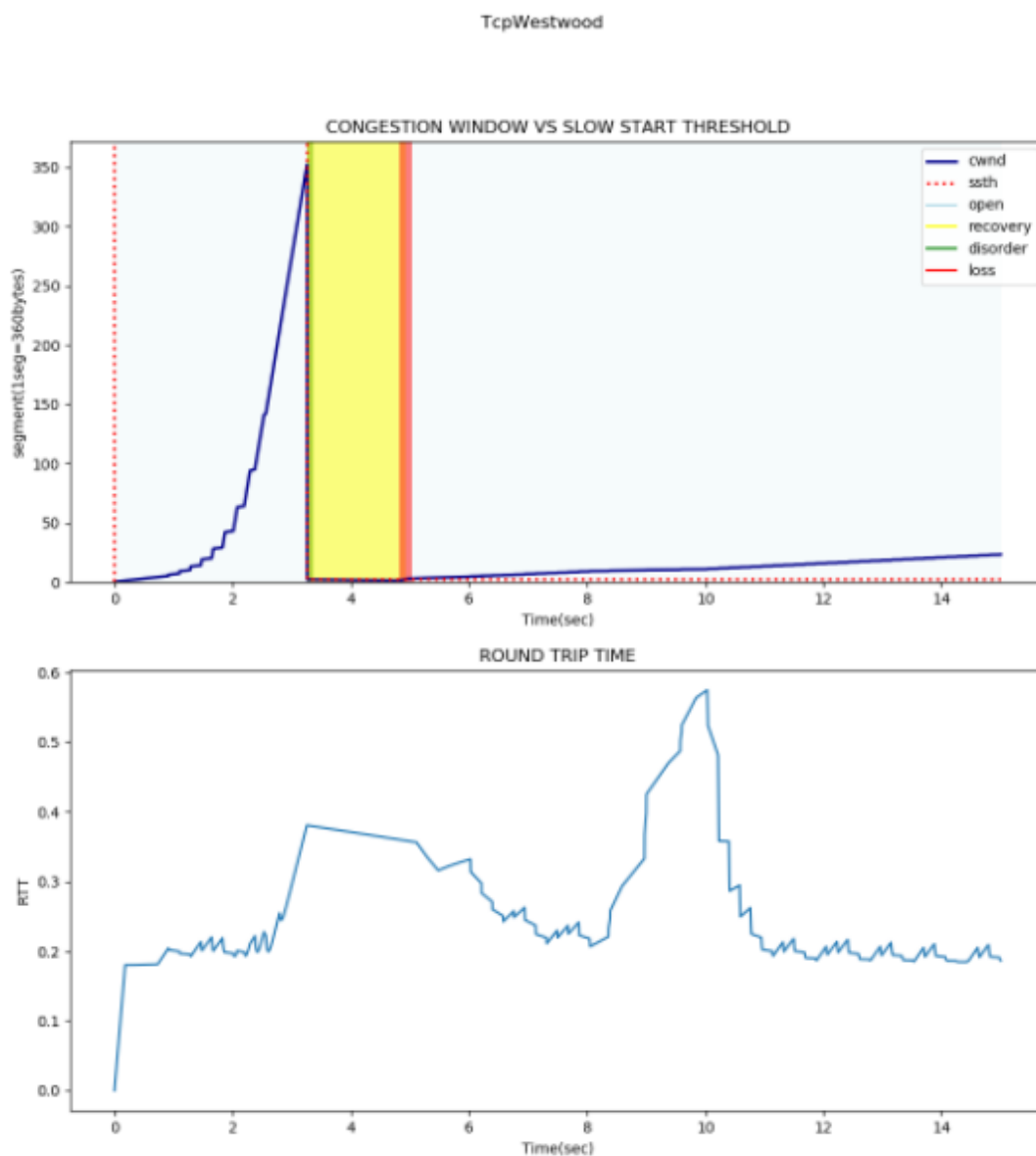


➤ Westwood

TCP Westwood (TCPW) is a sender-side-only modification to TCP New Reno that is intended to better handle large bandwidth-delay product paths (large pipes), with potential packet loss due to transmission or other errors (leaky pipes), and with dynamic load (dynamic pipes).^[1]

TCP Westwood relies on mining the ACK stream for information to help it better set the congestion control parameters: Slow Start Threshold (ssthresh), and Congestion Window (cwin).

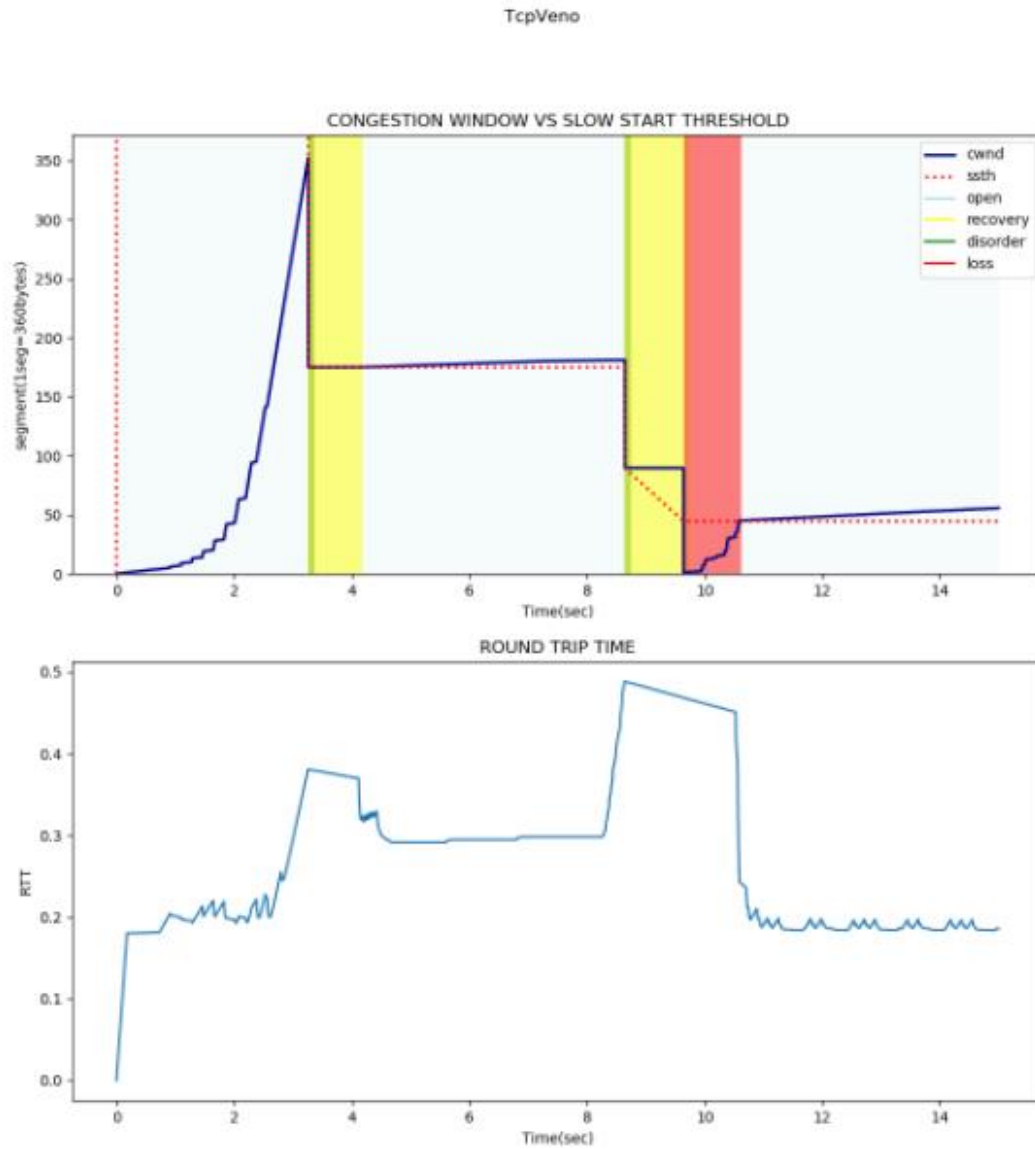
Plot:



➤ Veno

TCP Veno that is simple and effective for dealing with random packet loss. A key ingredient of Veno is that it monitors the network congestion level and uses that information to decide whether packet losses are likely to be due to congestion or random bit errors.

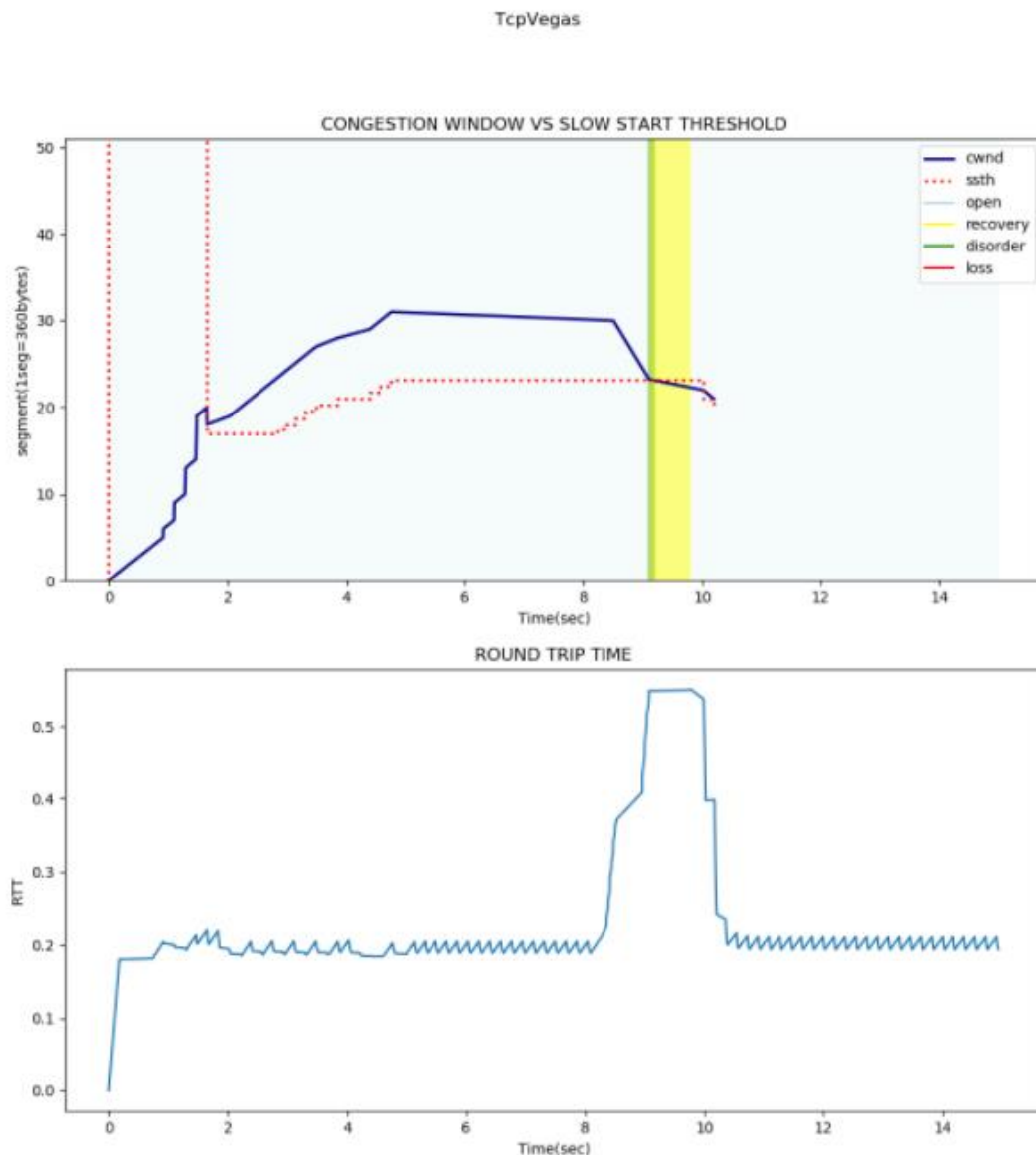
Plot:



➤ Vegas

TCP Vegas detects congestion at an incipient stage based on increasing Round-Trip Time (RTT) values of the packets in the connection unlike other flavors such as Reno, New Reno, etc., which detect congestion only after it has actually happened via packet loss. The algorithm depends heavily on accurate calculation of the Base RTT value. If it is too small then throughput of the connection will be less than the bandwidth available while if the value is too large then it will overrun the connection.

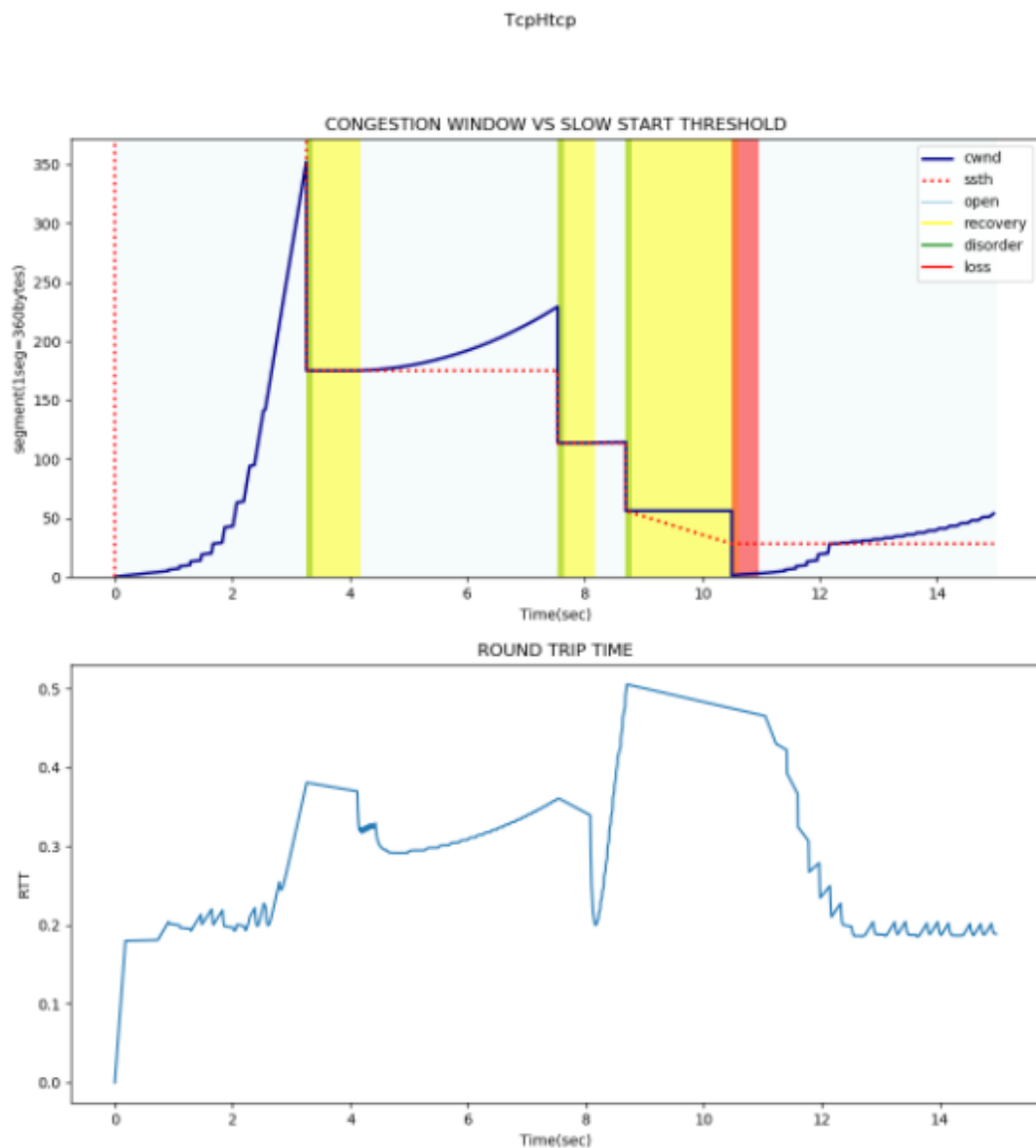
Plot:



➤ HTCP

H-TCP is another implementation of TCP with an optimized congestion control algorithm for high speed networks with high latency. H-TCP is a loss-based algorithm, using additive-increase/multiplicative-decrease (AIMD) to control TCP's congestion window. H-TCP increases its aggressiveness (in particular, the rate of additive increase) as the time since the previous loss increases.

Plot:

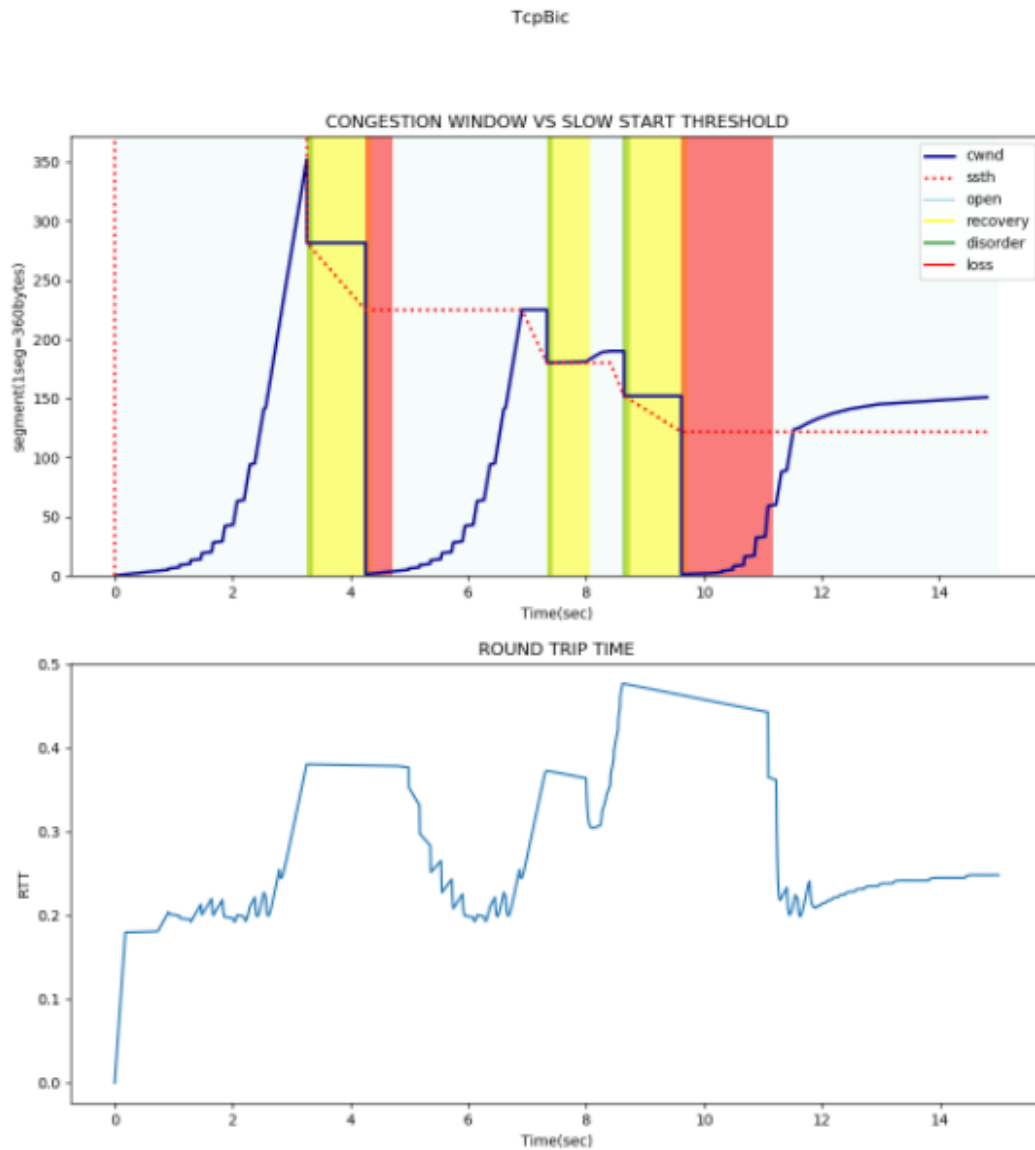


➤ BIC

Binary Increase Congestion control (BIC) is a TCP implementation with an optimized CCA for high speed networks with high latency,

BIC implements a unique congestion window (cwnd) algorithm. This algorithm tries to find the maximum cwnd by searching in three parts: binary search increase, additive increase, and slow start. When a network failure occurs, the BIC uses multiplicative decrease in correcting the cwnd

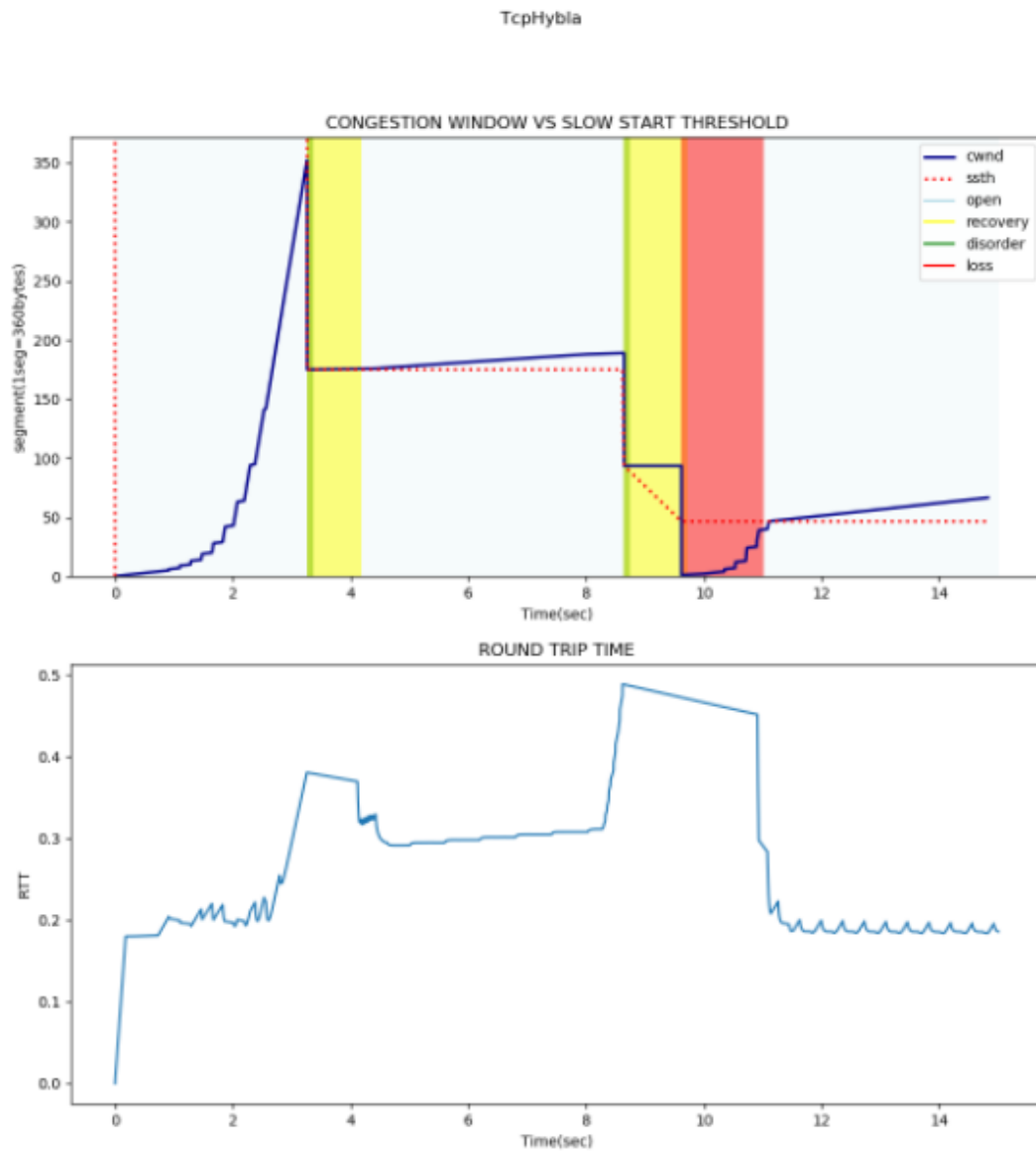
Plot:



➤ Hybla

TCP Hybla aims to eliminate penalties to TCP connections that incorporate a high-latency terrestrial or satellite radio links. Hybla improvements are based on analytical evaluation of the congestion window dynamics.

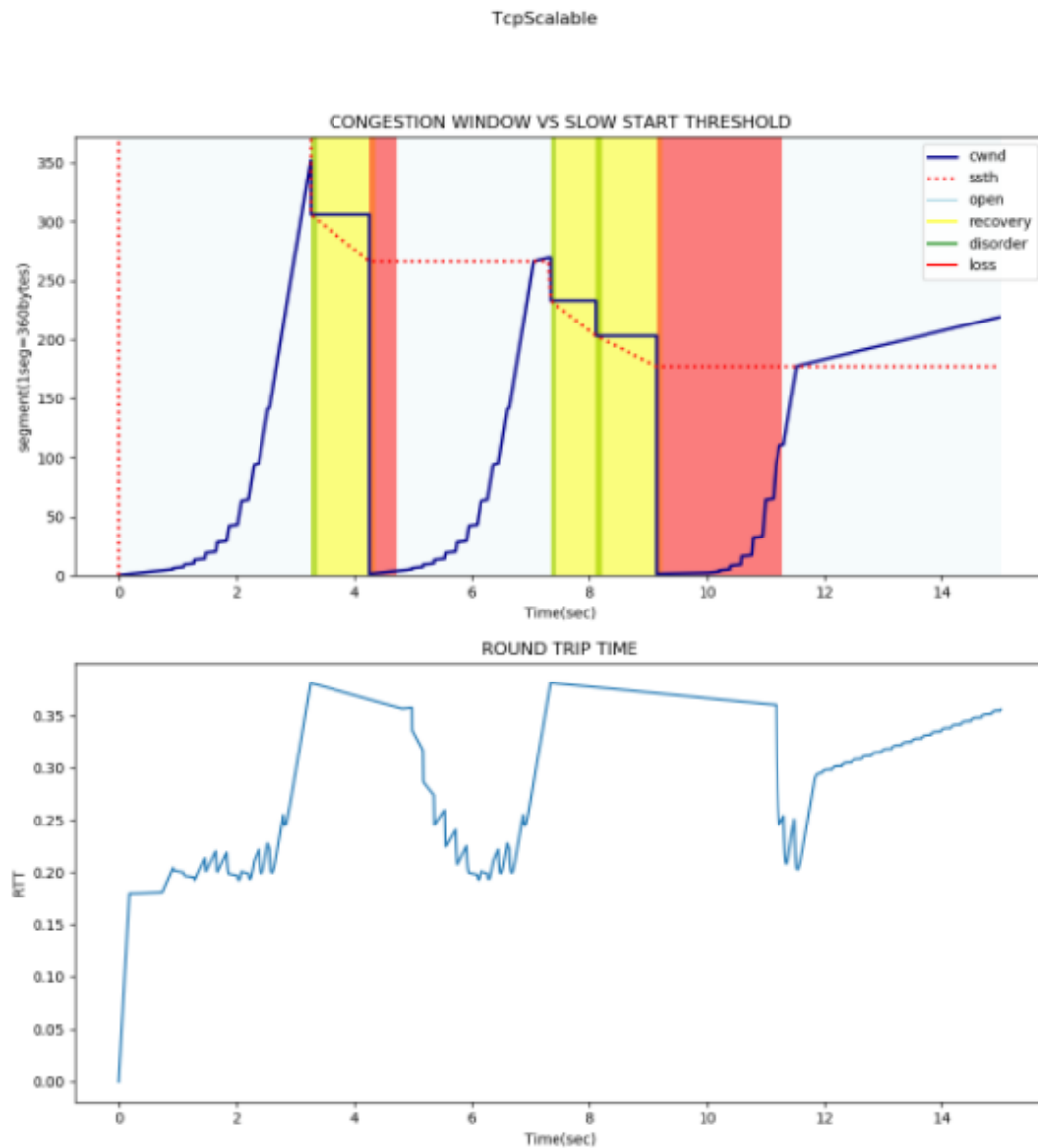
Plot:



➤ Scalable

Instead of halving the congestion window size, each packet loss decreases the congestion window by a small fraction (a factor of $1/8$ instead of Standard TCP's $1/2$) until packet loss stops. When packet loss stops, the rate is ramped up at a slow fixed rate instead of the Standard TCP rate that's the inverse of the congestion window size.

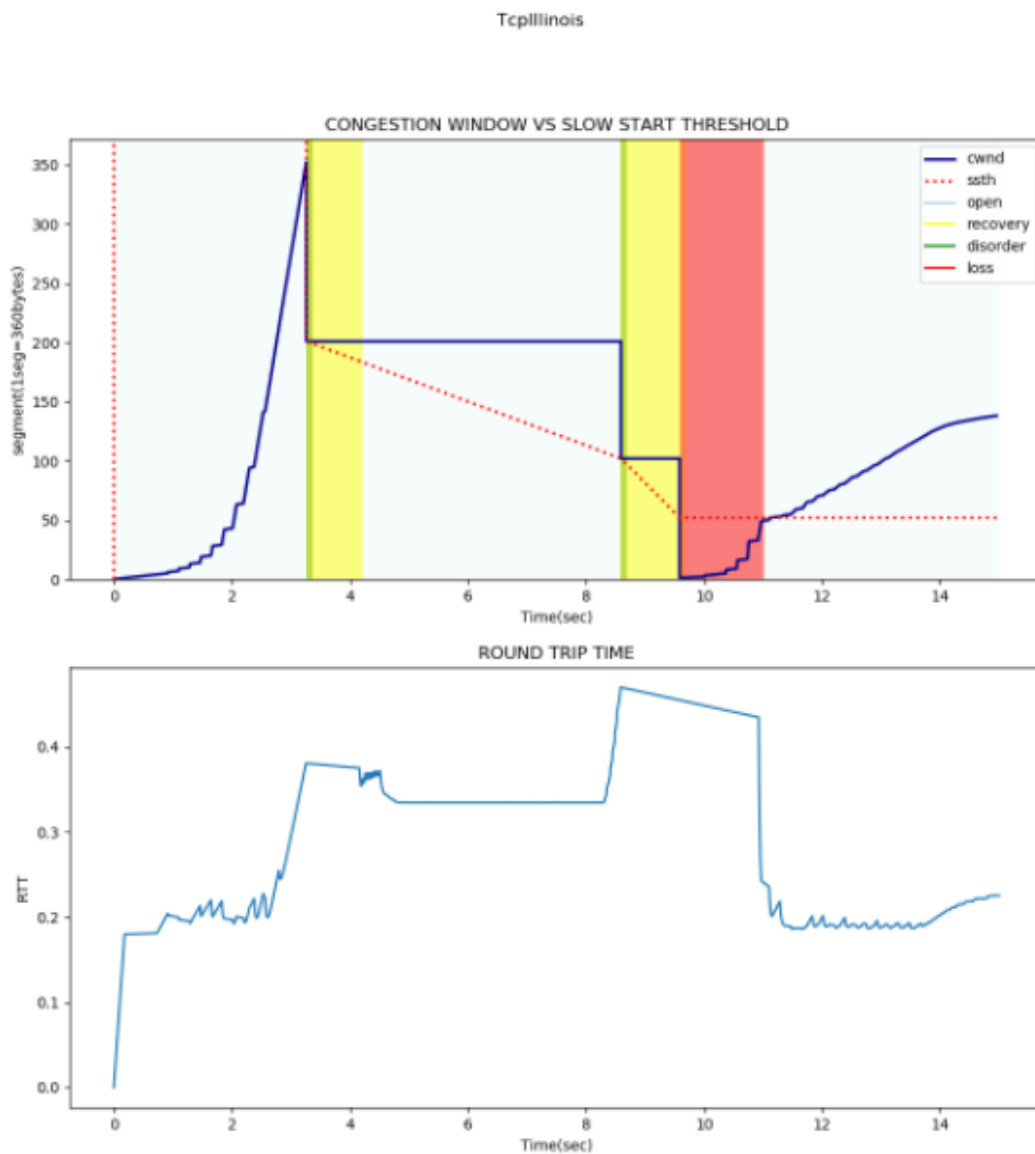
Plot:



➤ Illinois

A sender side modification to the standard TCP congestion control algorithm, it achieves a higher average throughput than the standard TCP, allocates the network resource fairly as the standard TCP, is compatible with the standard TCP. It is a loss-delay based algorithm, which uses packet loss as the primary congestion signal to determine the direction of window size change, and uses queuing delay as the secondary congestion signal to adjust the pace of window size change's-Illinois increases the window size W by α/W for each acknowledgment, and decreases W by $\beta \cdot W$ for each loss event, where α, β are not constants.

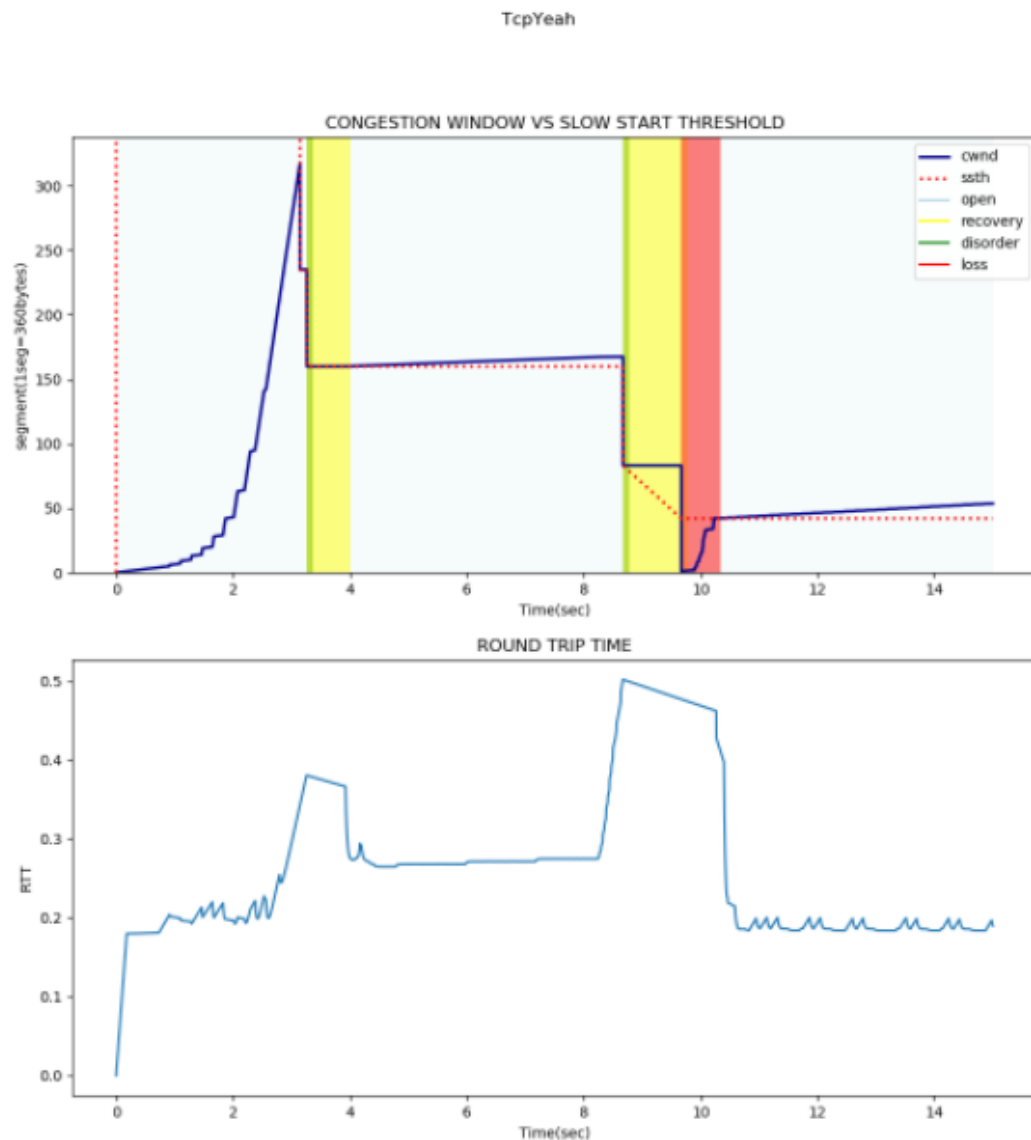
Plot:



➤ Yeah

TCP is a high speed TCP congestion control algorithm which uses a mixed loss/delay approach to calculate congestion windows. Its purpose is to target high efficiency, fairness, and minimizing link loss while keeping network elements load as low as possible. Through delay measures, when the network is sensed unloaded it will quickly exploit the available capacity, trying to keep the network buffer utilization always lower than a threshold

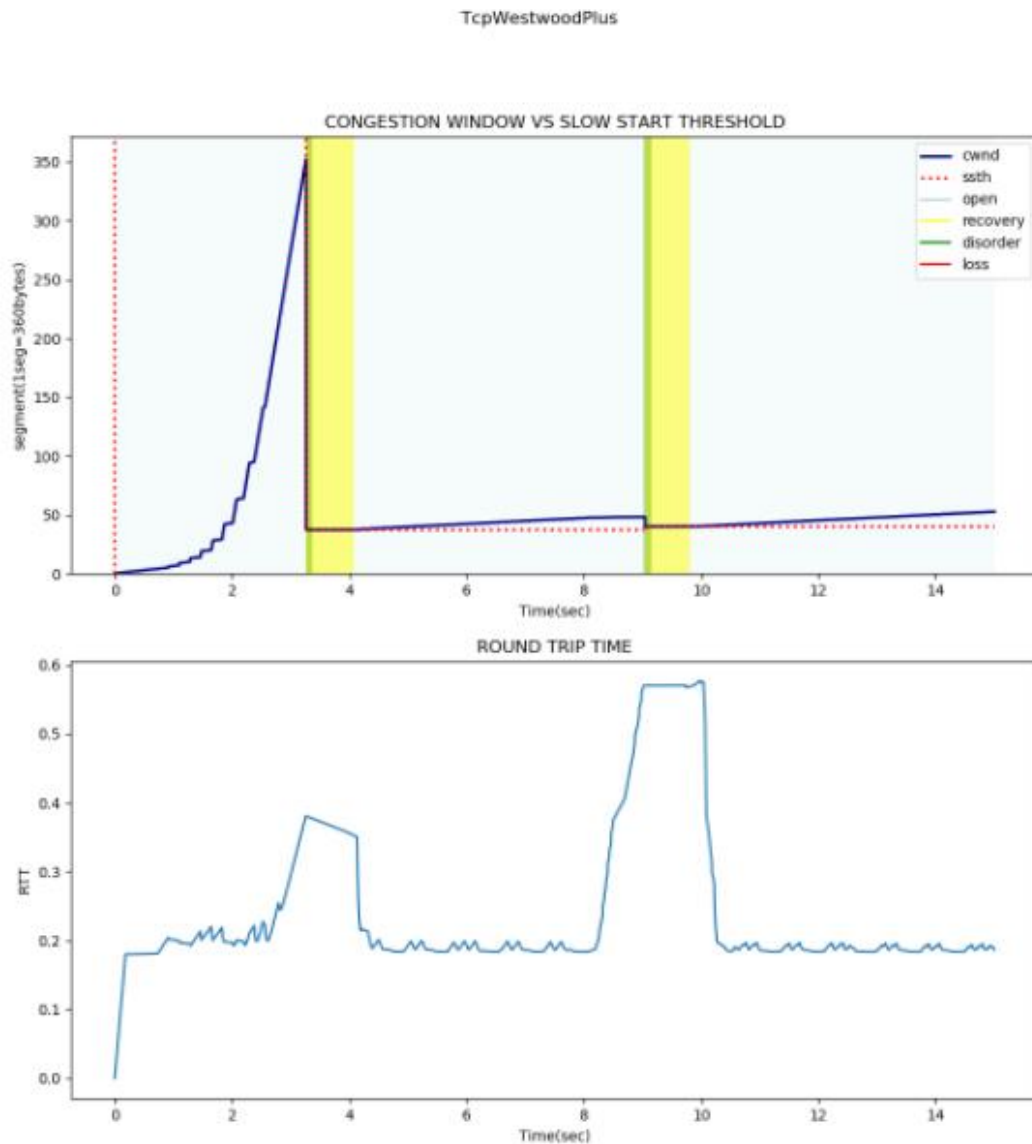
Plot:



➤ Westwood Plus

TCP Westwood+ is an evolution of *TCP Westwood*. TCP Westwood+ is a sender-side only modification of the TCP Reno protocol stack that optimizes the performance of TCP congestion control over both wireline and networks. Westwood+ is based on end-to-end bandwidth estimation to set congestion window and slow start threshold after a congestion episode, that is, after three duplicate acknowledgments or a timeout. The bandwidth is estimated by properly low-pass filtering the rate of returning acknowledgment packets.

Plot:

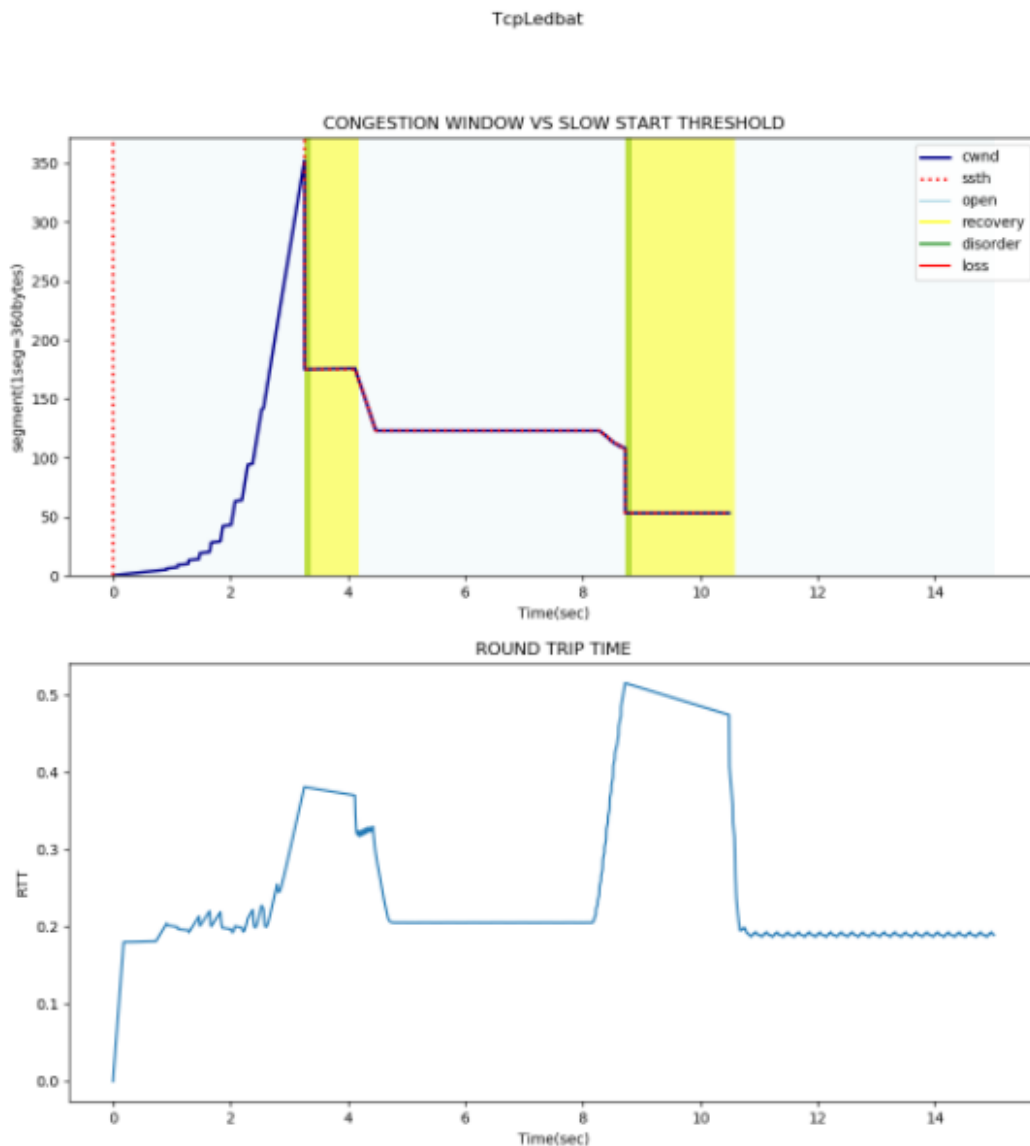


➤ Tcp Ledbat

Low Extra Delay Background Transport (LEDBAT) is an experimental delay-based congestion control algorithm that seeks to utilize the available bandwidth on an end-to-end path while limiting the consequent increase in queueing delay on that path.

LEDBAT uses changes in one-way delay measurements to limit congestion that the flow itself induces in the network. LEDBAT is designed for use by background bulk-transfer applications to be no more aggressive than standard TCP congestion control (as specified in RFC 5681) and to yield in the presence of competing flows, thus limiting interference with the network performance of competing flows.

Plot:



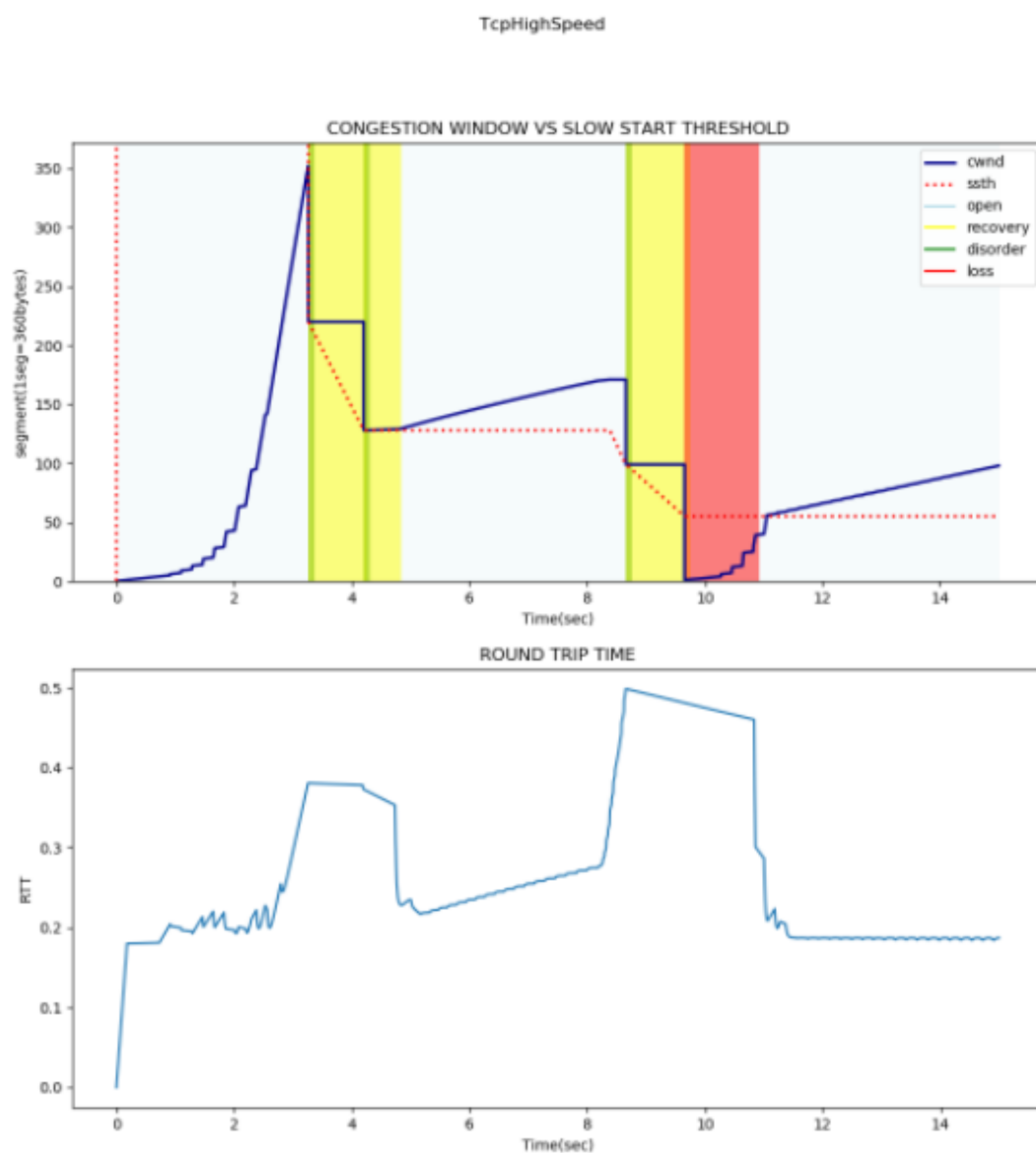
➤ Tcp HighSpeed

HighSpeed TCP (HSTCP) is a congestion control algorithm protocol defined in RFC 3649 for Transport Control Protocol (TCP). Standard TCP performs poorly in networks with a large bandwidth-delay product. It is unable to fully utilize available bandwidth. HSTCP makes minor modifications to standard TCP's congestion control mechanism to overcome this limitation.

When an ACK is received (in congestion avoidance), the window is increased by $a(w)/w$ and when a loss is detected through triple duplicate acknowledgments, the window equals $(1-b(w))w$, where w is the current window size

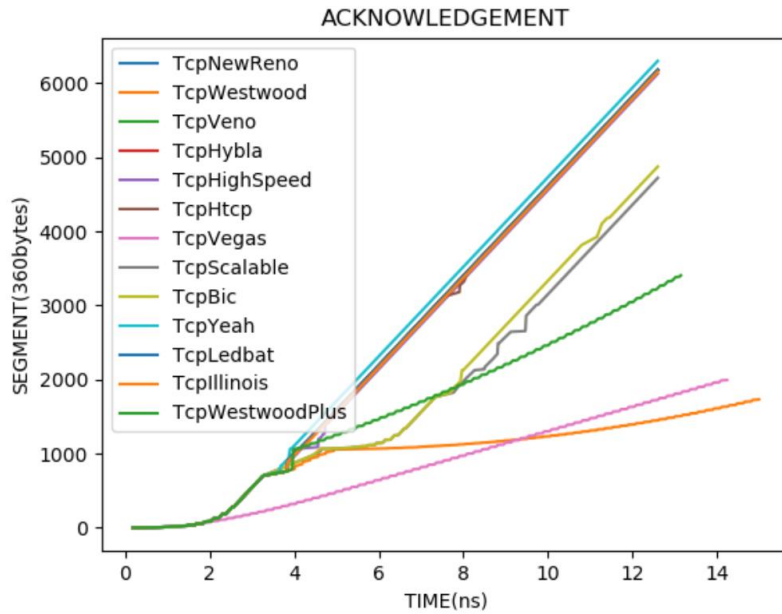
This means that HSTCP's window will grow faster than standard TCP and also recover from losses more quickly

Plot:

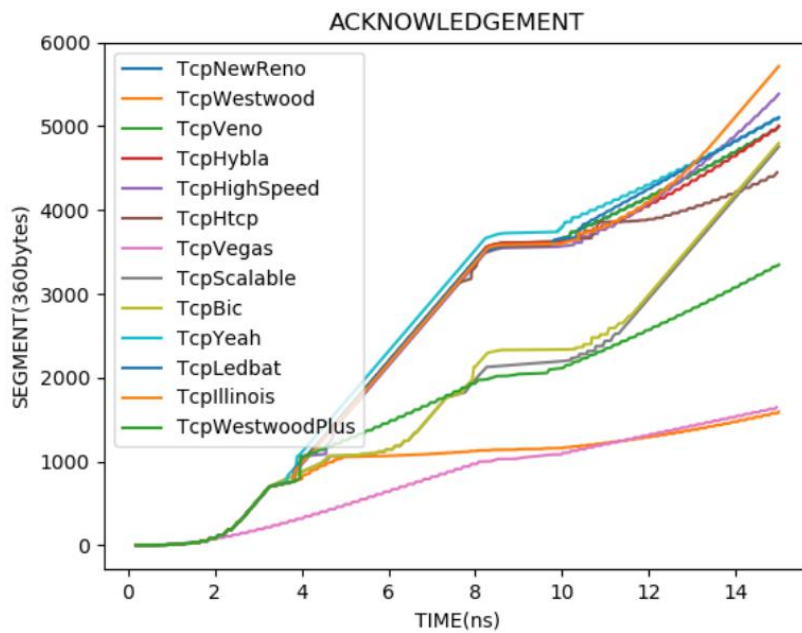


ACKNOWLEDGE VS TIME COMPARISON OF TCP VARIANT WITH VARIANTION IN ERROR MODEL OF NS3

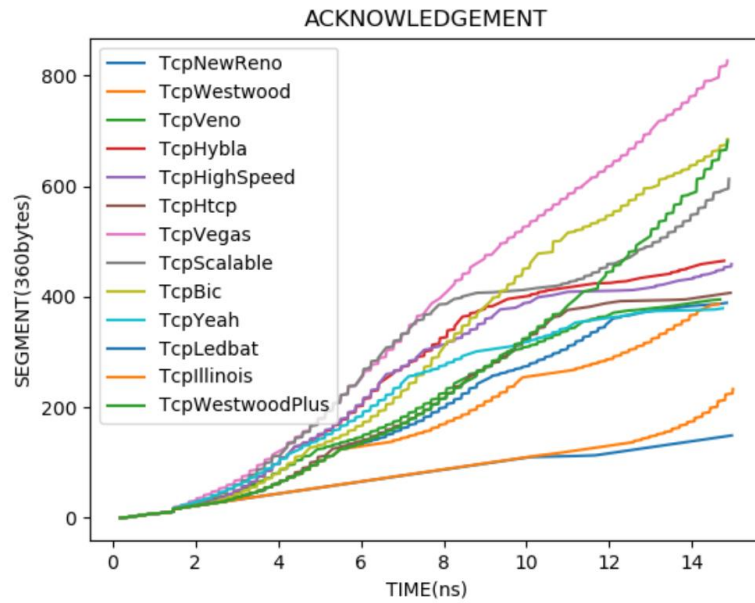
0 CBR(constant bit rate) generator used



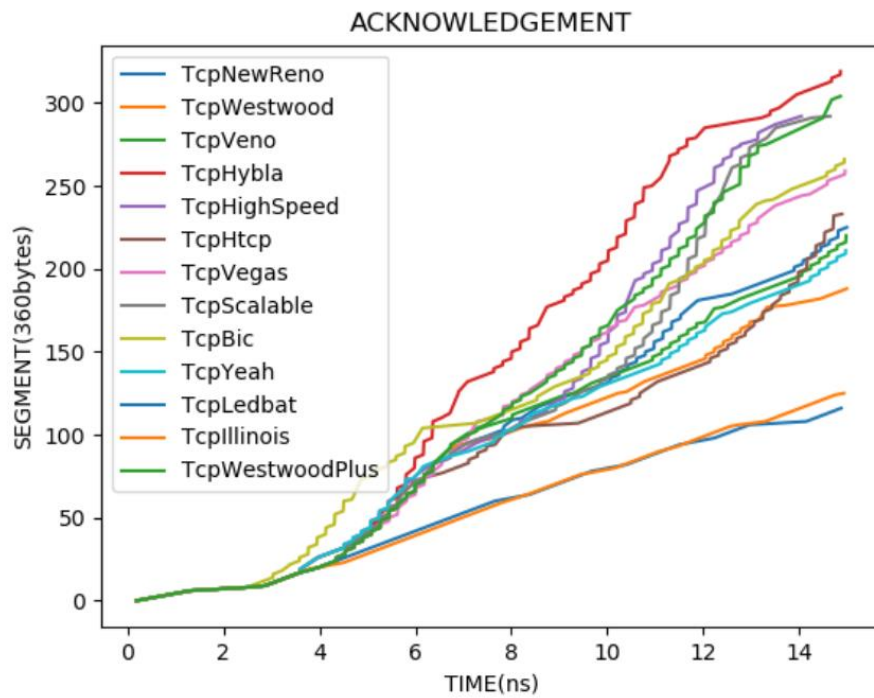
5 CBR(constant bit rate) generator used



At packet error rate=2%



At packet error rate=5%



Github link: <https://github.com/Baljotsinghchoudhary/ns3-Tcp-comparison>

