

Lecture 2

What we have learnt

- What is system
- OS definition and its goals
- System components :processors, memory, storage , I/O subsystems, busses , files & network
- Abstract view of computing systems
- Operating system as a resource manager

User View of OS

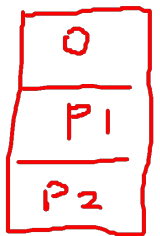
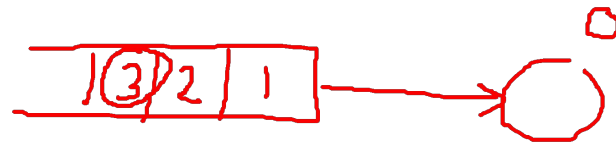
- Provides ease of operation
- Supports resource sharing
 - **Sequential sharing** : A resource is allocated for exclusive use of program eg. CPU
 - **Concurrent sharing**: Two or more program concurrently use same resource eg. Files, memory, disk array etc.

Systems View of OS

- OS acts as **Resource allocator** and is responsible for handling resource request, resource allocation and optimal utilization of system resources.
- Acts as **control program** that manages the execution of user program and prevents error and improper use of computer system.

What OS Does?

- Maintain a list of authorized users
- Construct list of all resources in the system
- Initiate execution of programs
- Maintain resource usage information by programs and current status of programs
- Maintain current status of all resources and allocate resources to programs when requested
- Perform scheduling
- Maintain information for protection



Goals Of Operating System

- Provide user interface for ***ease of operation***
 - Command line interface (CLI) eg. Unix, DOS
 - Graphical User Interface (GUI) eg. Windows
- ***Efficient*** use of hardware & software resources (system)
- Maximize ***System performance***
- Protection and access control
- ✓ • Ability to evolve and offer new services
- Footprint of OS should be small !!!

Can we measure System Efficiency, Performance and user convenience ?

- Efficiency---- CPU efficiency
- System performance ----- Throughput
- User services ---- Turn around time ,
Response time

Operating System and Computer Architecture are tightly bonded together.

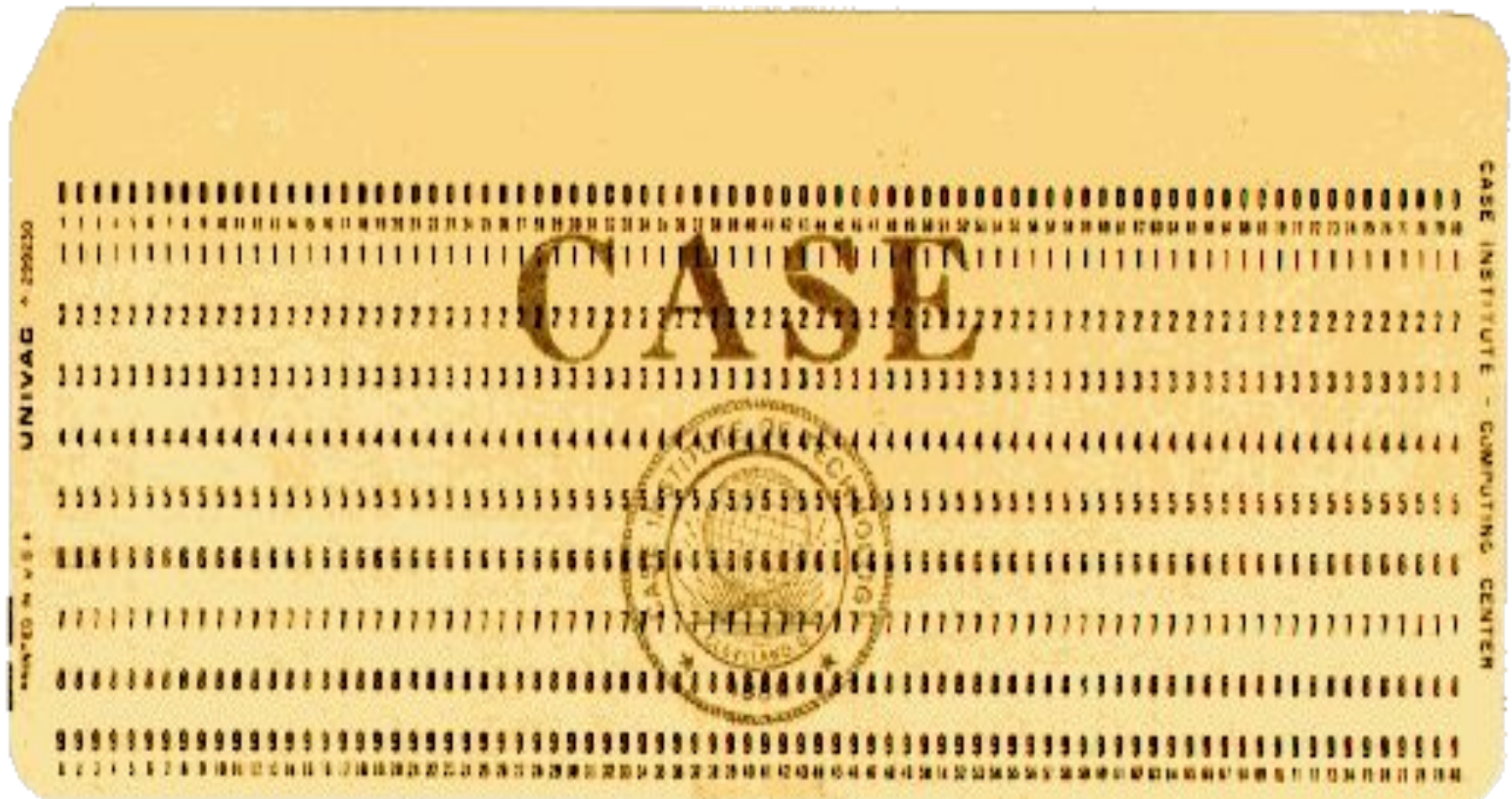
- Change in architecture usually leads to change in OS.
- The requirements of OS have led to several architectural changes.

Different Computer Systems

Mainframe Systems

- Are large systems
- Are very expensive
- Early mainframes were **Batch Systems**
- They used punched cards and tape for data and code input
- They used printer for results and error reporting

Punched Card(Hollerith's card)



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Mainframe Systems

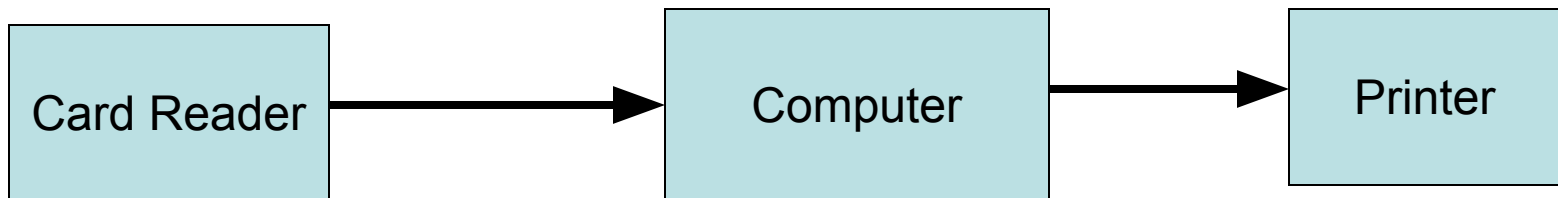
- Reduce setup time by batching similar jobs
- Automatic job sequencing – automatically transfers control from one job to another. **First rudimentary operating system.**
- Resident monitor
 - initial control in monitor
 - control transfers to job
 - when job completes control transfers back to monitor

Batch Systems

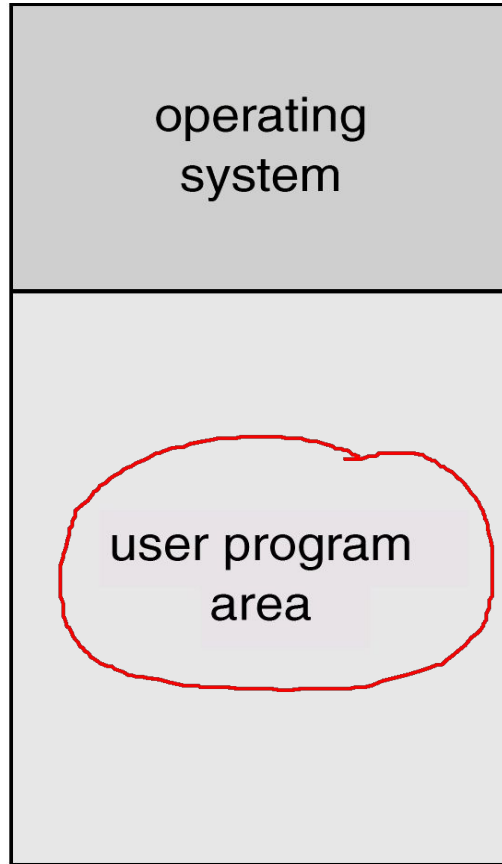
- Users submits job
- Operator forms a batch by bunching user jobs
- Loads job into the card reader
- Computer executes these jobs sequentially and produces results on printer
- Printed output is collected by operator and distributed to users.

Characteristics of Batch System

- Once a job starts executing , until finished no other job can execute
- **Low CPU utilization** because when program is reading data or producing output CPU is idle
- **No user interaction**
- Batch systems are appropriate for executing large jobs that require little interaction

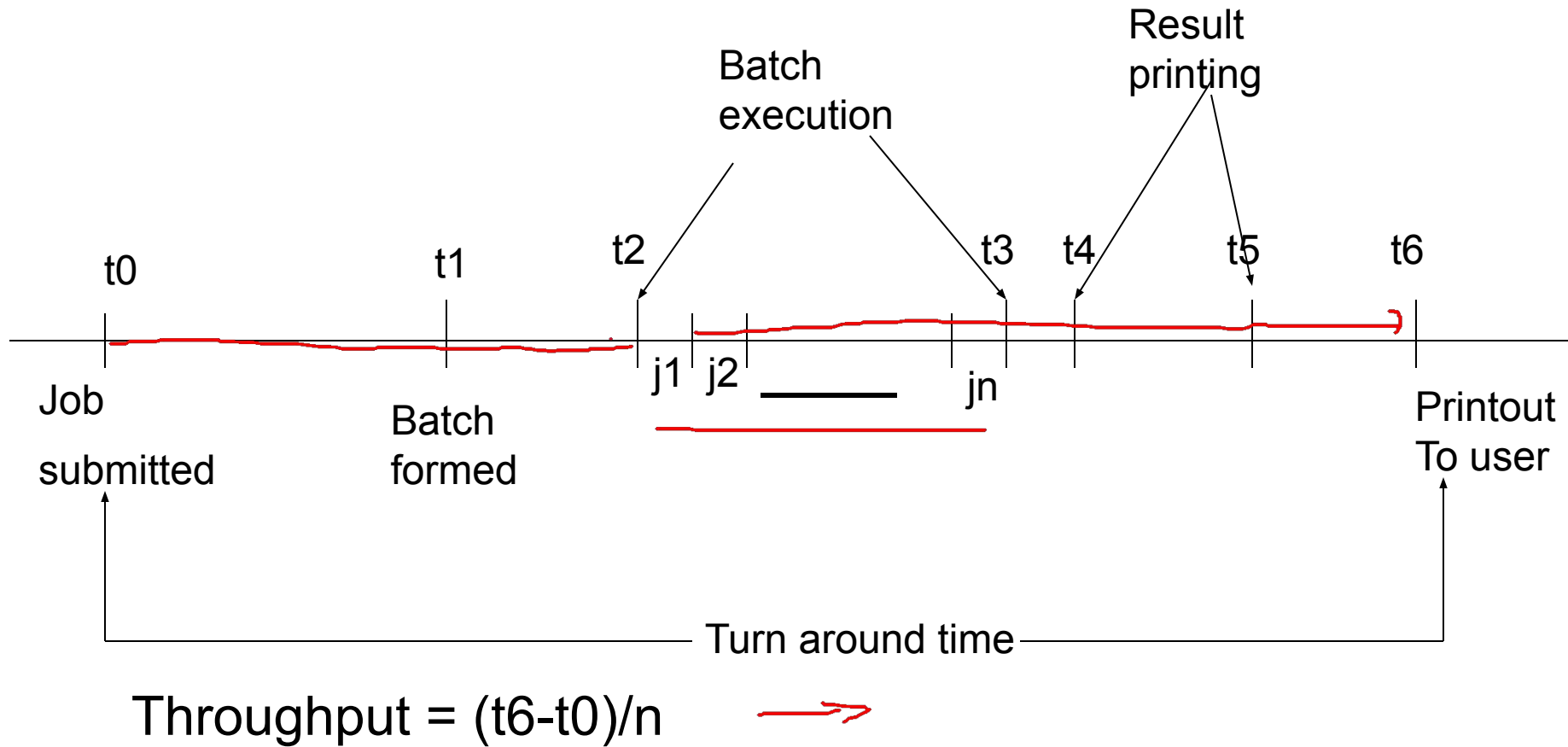


Batch system Memory Layout

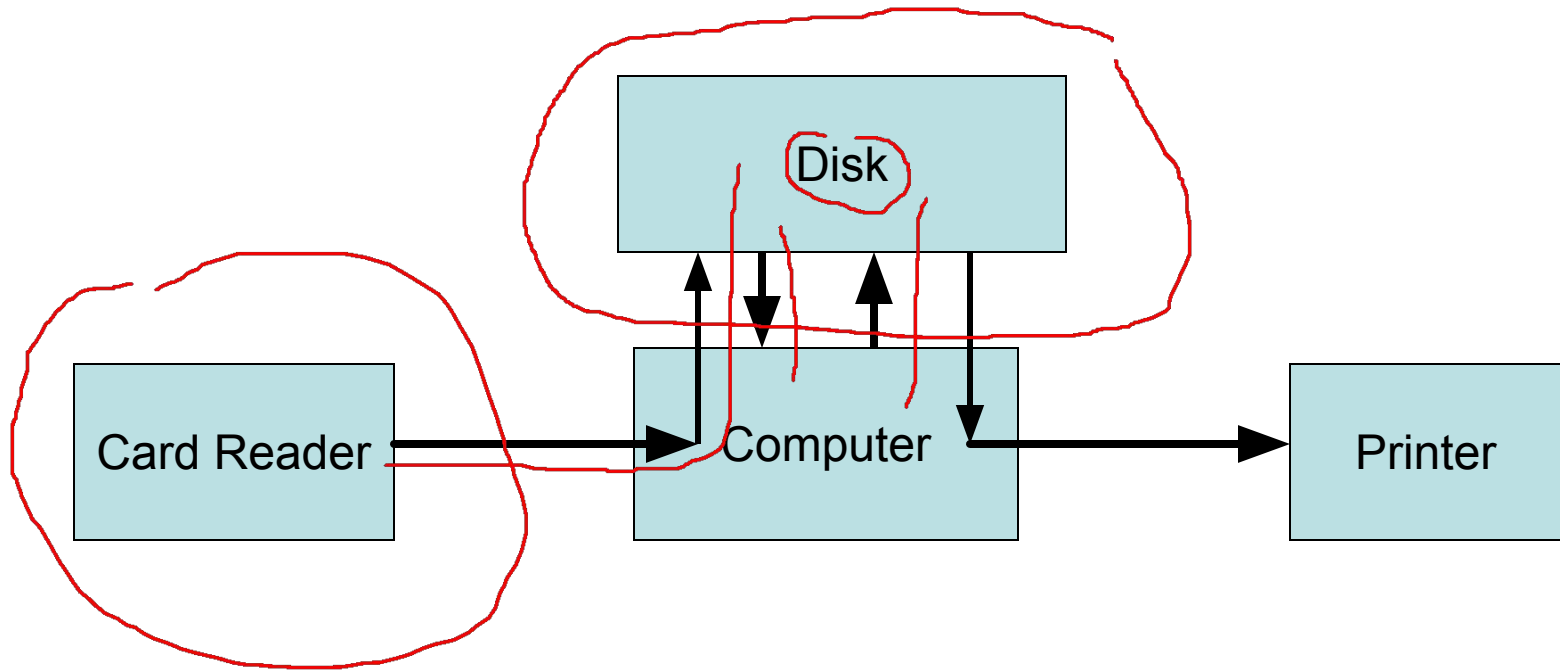


- Main task of the Batch processing OS was to transfer control from one job to another.
 - Problems
 1. How does the monitor know about the nature of the job (e.g., Fortran versus Assembly) How does the monitor distinguish
 - (a) one job from other job?
 - (b) data from program?
 - Solution
 - Job Control Language

Turn around time in batch system



Batch system with spooling



- Every executing program ~~usually~~ will have CPU Burst then I/O followed by CPU Burst and so on. The last burst is always CPU Burst and then program terminates
- c1 io1 c2 io2 c3-----cn

Problems of mainframe

- Early days Mainframes were inefficient as execution and I/O was done sequentially
- AS I/O devices are much slower than CPU execution , CPU utilization was poor

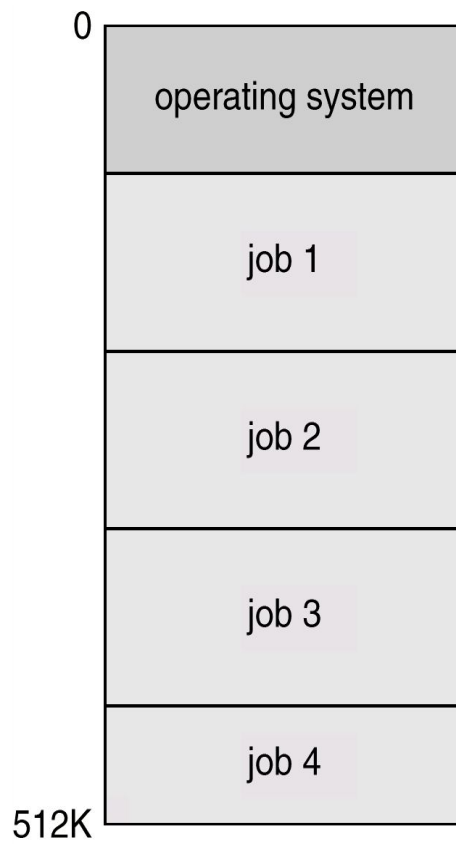
Solution

- Perform execution and I/O in concurrent manner
- Multiprogramming & Time sharing

Multiprogramming

- In this environment, multiple programs are kept in memory and one program is executed at a time. When a program terminates or goes for I/O resources, another program starts executing
- Multiprogramming improves CPU utilization
- Improves throughput

Memory Layout for multiprogramming



OS Features Needed for Multiprogramming

- **Memory management** – the system must allocate the memory to several jobs.
- **CPU scheduling** – the system must choose amongst several jobs ready to run.
- **Allocation of devices.**

Time sharing

- A time unit is divided into small slice and each user can be allocated CPU in round robin manner
- A time shared OS allows many users to share the computer simultaneously
- Time shared OS uses multiprogramming and CPU scheduling to provide each user with a small portion of time shared computer
- Time sharing machine with Interactive I/O devices improves user response time
- Gives illusion that each user has his own machine

Requirements of time shared system

- Needs to keep track of time
- Requires Dynamic Resource allocation
 - To manage memory processes can be swapped out requiring backing store
- Dynamic resource allocation requires elaborate and complicated resource management and access control technique

Problems with early systems and change drivers

- High cost of resources
 - This necessitated concept of resource sharing and eventually led to the development of computer network
 - Technological development and IC fabrication has been main driver for development of low cost computers and variety of peripheral devices
 - Applications are main drivers for development of faster processor and computing system

Desktop systems

- Provided interaction
- Main concern is response time
- Since it serves single user, optimal resource utilization and protection are not a major concern
- Present day PCs support time sharing , multiprogramming and are network enabled

Parallel & Distributed System

- Parallel processing improves performance
- Approaches for parallel processing
 - Pipelining eg. Instruction pipelining
 - Multiplicity of components (Multiprocessor Systems)
 - Symmetric Multiprocessing
 - Asymmetric Multiprocessing
 - Multiplicity of Systems
 - Cluster
 - Grid

- Multiprocessor System Offers
 - Increased throughput
 - Economy
 - Increased Reliability
 - Graceful Degradation
 - Fault tolerant
- Symmetric Multiprocessing :Each processor runs identical copy of OS
- Asymmetric multiprocessing: master slave relation exist.
- These are tightly coupled systems

Distributed Systems

- These systems depend on network for their functionality
 - Client Server Systems
 - Peer to peer systems
- Each processor has its own local memory
- Processors communicate with each other through communication links
- Loosely coupled Systems

Cluster System

- Cluster uses multiple CPUs to accomplish computational work
- Provide High availability
- Asymmetric Cluster (Hot Standby)
 - Hot standby machine monitors the active server and in case of failure of active server, hot standby becomes active server.
- Symmetric

Real Time system

- Hard Real time system
- Soft Real Time System