

# Assignment 1

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## Question 1

Proof that  $\mathcal{R} = \{(A, B) \mid A \supseteq B\} \subseteq \mathcal{P}(S) \times \mathcal{P}(S)$  is a partial order

We prove that  $\mathcal{R}$  is a partial order by verifying that it satisfies reflexivity, antisymmetry, and transitivity.

### 1. reflexivity

For any  $A \in \mathcal{P}(S)$ , clearly  $A \supseteq A$  since every set contains itself. Therefore,  $(A, A) \in \mathcal{R}$ , and so  $\mathcal{R}$  is reflexive.

### 2. antisymmetry

Suppose  $(A, B) \in \mathcal{R}$  and  $(B, A) \in \mathcal{R}$ . This means  $A \supseteq B$  and  $B \supseteq A$ . This means they are equal. Therefore,  $A = B$ , and  $\mathcal{R}$  is antisymmetric.

### 3. Transitivity

Suppose  $(A, B) \in \mathcal{R}$  and  $(B, C) \in \mathcal{R}$ . Then  $A \supseteq B$  and  $B \supseteq C$ . By transitivity of the superset relation,  $A \supseteq C$ , so  $(A, C) \in \mathcal{R}$ . Therefore,  $\mathcal{R}$  is transitive.

## Conclusion

Since  $\mathcal{R}$  is reflexive, antisymmetric, and transitive, it follows that  $\mathcal{R}$  is a **partial order** on  $\mathcal{P}(S)$ . ■

## Question 2

To run the program implemented in Python, follow the steps below.

### Prerequisites

You must have Python 3 installed on your system.

## Steps

```
git clone https://github.com/BangruiW/transaction_systems_ss25.git
cd transaction_systems_ss25/assignment_1
python3 serializability_checker.py
```

The input should be a space-separated list of operations, for example:

```
w 1 x r 2 x w 2 y r 3 y w 3 z r 1 z
```

## Question 3

Our algorithm constructs a **precedence graph** based on conflicting operations between transactions. According to serializability theory, a schedule is **conflict-serializable** if and only if this graph has **no cycles**. We apply the algorithm to a theoretical instance:

```
w 1 x    r 2 x    w 2 y    r 3 y    w 3 z    r 1 z
```

The algorithm detects WR conflicts:  $1 \rightarrow 2 \rightarrow 3 \rightarrow 1$ , forming a cycle. Therefore, the algorithm returns **false**, providing that the history is not conflict-serializable. This matches the theoretical prediction and confirms the algorithm's correctness for this instance.