

## Exercise 1 – Herbrand Semantics

Given the schedule:

$$s = r_1[x] \ r_3[x] \ w_3[y] \ w_2[x] \ c_3 \ r_4[y] \ w_4[x] \ c_2 \ r_5[x] \ c_4 \ w_5[z] \ w_1[z] \ c_1 \ c_5$$

We are told that transactions  $t_3$  and  $t_4$  are **copiers**, i.e., the write functions are identity functions on one of their inputs:

$$f_{ix}(v_1, \dots, v_m) = v_j \quad \text{for some } j$$

### Step 1: Compute values of each write using Herbrand semantics

- $H_s[r_3[x]] = H_s[w_0[x]] = f_{0x}()$
- $H_s[w_3[y]] = H_s[r_3[x]] = f_{0x}()$  (copier)
- $H_s[w_2[x]] = f_{2x}()$  (no read in  $t_2$ )
- $H_s[r_4[y]] = H_s[w_3[y]] = f_{0x}()$
- $H_s[w_4[x]] = H_s[r_4[y]] = f_{0x}()$  (copier)
- $H_s[r_5[x]] = H_s[w_4[x]] = f_{0x}()$
- $H_s[w_5[z]] = f_{5z}(H_s[r_5[x]]) = f_{5z}(f_{0x}())$
- $H_s[r_1[x]] = f_{0x}()$
- $H_s[w_1[z]] = f_{1z}(H_s[r_1[x]]) = f_{1z}(f_{0x}())$

### Step 2: Final writes to each data item

The last writes in the schedule are:

- $x$ : last write is  $w_4[x] \Rightarrow H[s](x) = H_s[w_4[x]] = f_{0x}()$
- $y$ : last write is  $w_3[y] \Rightarrow H[s](y) = H_s[w_3[y]] = f_{0x}()$
- $z$ : last write is  $w_1[z] \Rightarrow H[s](z) = H_s[w_1[z]] = f_{1z}(f_{0x}())$

### Answer: Herbrand Semantics

$H[s](x) = f_{0x}(), \quad H[s](y) = f_{0x}(), \quad H[s](z) = f_{1z}(f_{0x}())$
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