

Proximal Methods

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May 18, 2021

Chapter 1

Proximal Algorithms

1.1 Motivation

In this chapter, we briefly outline the problem of minimizing functions that are not necessarily differentiable. A typical example is the l_1 -regularized problem. For example, the object might look like

$$\min_{\beta} \sum_{i=1}^N (y_i - x_i^T \beta)^2 + \lambda \|\beta\|_1.$$

Here, β is the parameter we want to find. Should we not have $\lambda \|\beta\|_1$, then everything is differentiable and can be solved using quasi-Newton methods, among other things. However, the absolute value function is not differentiable everywhere, which causes problems.

The first solution is to consider *sub-differentials*. Sub-differentials are defined as

$$\partial f(x) = \{y \mid f(z) \geq f(x) + y^T(z - x) \text{ for all } z \in \text{dom } f\},$$

where $\text{dom } f$ is the domain of the function. Note that if a function is differentiable then $\partial f = \{\nabla f\}$. However, in general case, the sub-differential is not a singleton.

For simplicity, we assume all the functions we discuss are subdifferentiable.

1.2 Proximal Algorithms

The proximal operator is defined as

$$\text{prox}_f(v) = \underset{x}{\operatorname{argmin}} (f(x) + (1/2)\|x - v\|_2^2).$$

As simple as the definition might look like, it has quite some nice results. The first one is a fixed-point properties. That is, the point x^* minimizes f if and only if

$$x^* = \text{prox}_f(x^*).$$

Proof. First we show that if x^* is the minimizer, then $x^* = \text{prox}_f(x^*)$. Note that for any x ,

$$f(x) + (1/2) \|x - x^*\|_2^2 \geq f(x^*) = f(x^*) + (1/2) \|x^* - x^*\|_2^2,$$

and thus by definition, $x^* = \text{prox}_f(x^*)$.

Now consider the reverse case, let $\tilde{x} = \text{prox}_f(v)$. Take the subdifferential operator, we see that this is equivalent to

$$0 \in \partial f(\tilde{x}) + (\tilde{x} - v).$$

Taking $\tilde{x} = v = x^*$, it follows that $0 \in \partial f(x^*)$, so x^* minimizes f . \square

The second interesting, and rather surprising fact is that, the proximal operator is actually the resolvent of subdifferential operator. More specifically,

$$\text{prox}_{\lambda f} = (I + \lambda \partial f)^{-1}.$$

Proof. If $z \in (I + \lambda \partial f)^{-1}(x)$, then $0 \in \partial f(z) + (1/\lambda)(z - x)$. This implies $0 \in \partial_z (f(z) + (1/2\lambda)\|z - x\|_2^2)$. Now, since we can prove that $f(z) + (1/2\lambda)\|z - x\|_2^2$ is strongly convex, we can deduce that $z = \underset{u}{\text{argmin}} (f(u) + (1/2\lambda)\|u - x\|_2^2)$. \square

Finally, let us look at the case of minimizing $f + g$ where f is differentiable and g is not. A famous algorithms goes,

$$x^{k+1} := \text{prox}_{\lambda^k g} (x^k - \lambda^k \nabla f(x^k)).$$

To see why, consider the fixed point version that is $x^* = \text{prox}_{\lambda g} (x^* - \lambda \nabla f(x^*))$, we show that if this is true then x is indeed the solution to

$$x^* = \underset{x}{\text{argmin}} f(x) + g(x).$$

Proof. Note that x^* is the minimizer if and only if $0 \in \nabla f(x^*) + \partial g(x^*)$. Now with some straight forward computation,

$$\begin{aligned} 0 &\in \lambda \nabla f(x^*) + \lambda \partial g(x^*) \\ 0 &\in \lambda \nabla f(x^*) - x^* + x^* + \lambda \partial g(x^*) \\ (I + \lambda \partial g)(x^*) &\ni (I - \lambda \nabla f)(x^*) \\ x^* &= (I + \lambda \partial g)^{-1}(I - \lambda \nabla f)(x^*) \\ x^* &= \text{prox}_{\lambda g}(x^* - \lambda \nabla f(x^*)) \end{aligned}$$

□

1.3 Applications

Now let us return to the previous topics. Let us assume that the **negative** log-likelihood function is $f_\beta(X)$, here β is the parameter we want to estimate and X is the data, which include both the dependent and independent variables. The goal now is to minimize the following quantities.

$$f_\beta(X) + \lambda \|\beta\|_1,$$

where $\lambda > 0$ is a positive hyperparameter, $\|\cdot\|$ is the l_1 norm, namely for any $x \in \mathbb{R}^N$, $\|x\|_1 = \sum_{i=1}^N |x_i|$. Since the latter term is not differentiable at 0, we are in position to work out the solution as is indicated in the previous section.

Now let $w := \beta^k - \lambda^k \nabla f_{\beta^k}(X)$, where λ^k is the step-size, and β^k is the value of β at step k , and let $g : x \mapsto \lambda \|x\|_1$, we have

$$\begin{aligned} \beta^{k+1} &:= \text{prox}_{\lambda^k g}(w) \\ &= \underset{x}{\text{argmin}} \left(\lambda^k g(x) + \frac{1}{2} \|x - w\|_2^2 \right) \\ &= \underset{x}{\text{argmin}} \left(\lambda \|x\|_1 + \frac{1}{2\lambda_k} \|x - w\|_2^2 \right) \end{aligned}$$

Now it suffices to calculate the for each i , since none of i depends on others. For that, note that (**exercises!**) we have

$$\underset{z}{\text{argmin}} \frac{1}{2} \|\beta - z\|_2^2 + \lambda t \|z\|_1 = S_{\lambda t}(\beta)$$

where we have

$$[S_\lambda(\beta)]_i = \begin{cases} \beta_i - \lambda & \text{if } \beta_i > \lambda \\ 0 & \text{if } -\lambda \leq \beta_i \leq \lambda \\ \beta_i + \lambda & \text{if } \beta_i < -\lambda \end{cases}$$

Therefore we have

$$[\beta^{k+1}]_i = \begin{cases} w_i - \lambda \lambda_k & \text{if } w_i > \lambda \lambda_k \\ 0 & \text{if } -\lambda \lambda_k \leq w_i \leq \lambda \lambda_k \\ w_i + \lambda \lambda_k & \text{if } w_i < -\lambda \lambda_k \end{cases}$$

With that, one can easily implement the algorithm.

1.4 Improving the algorithm

Unfortunately, the algorithm itself isn't the fastest possible. One way to implement the algorithm is by backtracking line search. Specifically, let us define the following quantities: $G_t(x) = \frac{x - \text{prox}_t(x - t\nabla g(x))}{t}$.

Now, fix an shrinkage parameter $\beta \in (0, 1)$, let the initial step size $t = 1$, then while

$$g(x - tG_t(x)) > g(x) - t\nabla g(x)^T G_t(x) + \frac{t}{2} \|G_t(x)\|_2^2.$$

Shrink t by β , i.e. $t := \beta t$.