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1. (Soft margin SVM)

$$\begin{aligned} \text{(Primal)} \left\{ \begin{array}{ll} \min_{w, b, \xi} & \frac{1}{2} \|w\|_2^2 + C_1 \sum_{i=1}^n \xi_i + C_2 \sum_{i=1}^n \xi_i^2 \\ \text{s.t.} & y_i (x_i^T w + b) \geq 1 - \xi_i, \forall i \\ & \xi_i \geq 0, \forall i \end{array} \right. \quad (\text{Assume } C_1, C_2 > 0) \end{aligned}$$

The Lagrangian:

$$\begin{aligned} \mathcal{L}(w, \xi, b, \alpha) &= \frac{1}{2} \|w\|_2^2 + C_1 \sum_{i=1}^n \xi_i + C_2 \sum_{i=1}^n \xi_i^2 + \sum_{i=1}^n \alpha_i [1 - \xi_i - y_i (x_i^T w + b)] \\ &= \underbrace{\frac{1}{2} \|w - \sum_{i=1}^n \alpha_i y_i x_i\|_2^2}_{A_1} - \frac{1}{2} \sum_{i,j=1}^n \alpha_i \alpha_j y_i y_j x_i^T x_j + \sum_{i=1}^n \alpha_i - \underbrace{\sum_{i=1}^n \alpha_i y_i b}_{A_2} \\ &\quad + \underbrace{C_2 \sum_{i=1}^n \left(\xi_i + \frac{C_1 - \alpha_i}{2C_2} \right)^2 - \sum_{i=1}^n \frac{(C_1 - \alpha_i)^2}{4C_2}}_{A_3} \end{aligned}$$

$$\min_{w, \xi, b} \mathcal{L}(w, \xi, b, \alpha) \quad \text{with } \xi_i \geq 0 \forall i$$

$$\sum (C_1 - \alpha_i) \xi_i$$

$$\min_w A_1 \Rightarrow w = \sum_{i=1}^n \alpha_i y_i x_i \quad A_1 = 0$$

$$\min_b A_2 \Rightarrow \sum_{i=1}^n \alpha_i y_i = 0 \quad A_2 = 0$$

$$\min_{\xi_i} A_3 \Rightarrow \begin{cases} \xi_i = -\frac{C_1 - \alpha_i}{2C_2} & \alpha_i > C_1 \\ \xi_i = 0 & 0 \leq \alpha_i \leq C_1 \end{cases} \quad A_3 = -\sum_{i=1}^n \frac{(C_1 - \alpha_i)^2}{8C_2} [1 + \text{sign}(\alpha_i - C_1)]$$

$$\begin{aligned}
 \text{(Dual)} \quad & \left\{ \begin{aligned} & \max_{\alpha} \quad \sum_{i=1}^n \alpha_i - \frac{1}{2} \sum_{i,j=1}^n \alpha_i \alpha_j y_i y_j x_i^T x_j - \sum_{i=1}^n \frac{(C_1 - \alpha_i)^2}{8C_2} [1 + \text{sign}(\alpha_i - C_1)] \\ & \text{s.t.} \quad \sum_{i=1}^n \alpha_i y_i = 0 \\ & \quad 0 \leq \alpha_i, \forall i \end{aligned} \right.
 \end{aligned}$$

2. X is finite Set

(a). if $F: 2^X \mapsto \mathbb{R}_+$ is submodular,

that means $F(S) + F(T) \geq F(S \cup T) + F(S \cap T)$
 $\forall S, T \subseteq X$

then $F(S) + F(T) \geq F(S \cup T) + F(S \cap T) \geq F(S \cup T)$

so, F is subadditive.

(b).

① F is submodular

② F_S is subadditive

The statement ① \Leftrightarrow ② isn't true!

Counterexample: If $A \subseteq B$, $A \cup S \neq S$, but $F(A \cup S) < F(S)$

- If F_S is subadditive, $F_S(A) + F_S(B) \geq F_S(A \cup B) = F_S(B)$
 $\Rightarrow F(A \cup S) - F(S) \geq 0$ conflict!

With the addition condition F is monotone OR F_S is non-negative
 one can finish the proof.

Pf:

① \Rightarrow ② F is submodular, then let $T \subseteq M$, $e \in X/M$

$$\begin{aligned}
 F_S(T \cup \{e\}) - F_S(T) &= F(S \cup T \cup \{e\}) - F(S \cup T) \\
 &\geq F(S \cup M \cup \{e\}) - F(S \cup M) \\
 &= F_S(M \cup \{e\}) - F_S(M)
 \end{aligned}$$

So, F_S is submodular, with the addition condition
 F_S is subadditive

② \Rightarrow ① Let $S \subseteq T \subseteq X$, $e \in X \setminus T$, $T = S \cup T'$

$$\begin{aligned}
 F(T \cup \{e\}) - F(T) &= F(S \cup T' \cup \{e\}) - F(S \cup T') \\
 &= F_S(T' \cup \{e\}) - F_S(T') \\
 &\leq F_S(T') + F_S(\{e\}) - F_S(T') \\
 &= F_S(\{e\}) = F(S \cup \{e\}) - F(S)
 \end{aligned}$$

$$\text{so, } F(T \cup \{e\}) - F(T) \leq F(S \cup \{e\}) - F(S)$$

$\Rightarrow F$ is submodular

Q.E.D.

3. $f(S) = \# \text{ of edge } e = (u, v) \in E \text{ } u \in S, v \in V/S.$

Def: $E(X, Y) = \{(x, y) \in E \mid x \in X, y \in Y\}$

then $f(S) = |E(S, V/S)|$, $|\cdot|$ means the number of elements

★ we have $E(X \cup \{a\}, Y) = E(X, Y) \cup E(\{a\}, Y)$
 $E(X, Y) = E(X, Y/\{b\}) \cup E(X, \{b\})$

$$\text{if } a \notin X, |E(X \cup \{a\}, Y)| = |E(X, Y)| + |E(\{a\}, Y)|$$

$$\text{if } b \in Y, |E(X, Y)| = |E(X, Y/\{b\})| + |E(X, \{b\})|$$

$$\text{What's more, if } A \cap X = \emptyset, |E(X \cup A, Y)| = |E(X, Y)| + |E(A, Y)|$$

$$\text{if } B \subseteq Y \quad |E(X, Y)| = |E(X, Y/B)| + |E(X, B)|$$

$$\text{Let } A \subset B \subseteq V, x \in V/B, B = A \cup S, \hat{B} = B \cup \{x\}, \hat{A} = A \cup \{x\}, \hat{S} = S \cup \{x\},$$

$$\begin{aligned} f(\hat{B}) - f(B) &= |E(\hat{B}, V/\hat{B})| - |E(B, V/B)| \\ &= |E(\hat{A}, V/\hat{B})| + |E(\hat{S}, V/\hat{B})| - |E(A, V/B)| - |E(S, V/B)| \\ &= |E(\hat{A}, V/\hat{A})| - |E(\hat{A}, S)| - |E(A, V/A)| + |E(A, S)| \\ &\quad + |E(\hat{S}, V/\hat{B})| - |E(S, V/B)| \\ &= f(\hat{A}) - f(A) - 2|E(\{x\}, S)| \\ &\leq f(\hat{A}) - f(A) \end{aligned}$$

$$\Rightarrow f(B \cup \{x\}) - f(B) \leq f(A \cup \{x\}) - f(A)$$

f is submodular

4.

s	a	s'	r	$p(s', r s, a)$
high	search	high	r_{search}	α
high	search	low	r_{search}	$1 - \alpha$
high	wait	high	r_{wait}	1.
high	wait	low	-	0

low	search	high	-3	$1-\beta$
low	search	low	γ_{search}	β
low	wait	low	γ_{wait}	1
low	wait	high	-	0
low	recharge	high	0	1
low	recharge	low	-	0

5. The Bellman equation for q_x

$$q_x(s|a) = \sum_{s', r'} p(s', r' | s, a) [r + \gamma \max_a q_x(s', a')].$$

For cycling robot, state = $\{h, l\}$, action = $\{s, w, re\}$.

$$\left\{ \begin{array}{l} q_x(l, s) = \beta (\gamma_{search} + \gamma \max_{a'} q_x(l, a')) + (1-\beta) (-3 + \gamma \max_{a'} q_x(h, a')) \\ q_x(l, w) = (\gamma_{wait} + \gamma \max_{a'} q_x(l, a')) \\ q_x(l, re) = \gamma \max_{a'} q_x(h, a') \\ q_x(h, s) = \alpha (\gamma_{search} + \gamma \max_{a'} q_x(h, a')) + (1-\alpha) (\gamma_{search} + \gamma \max_{a'} q_x(l, a')) \\ q_x(h, w) = \gamma_{wait} + \gamma \max_{a'} q_x(h, a') \end{array} \right.$$