

## Basics

**Sylvester’s inequality:** If  $A$  is an  $M \times N$  matrix and  $B$  is an  $N \times K$  matrix, then  $\text{rank}(A) + \text{rank}(B) - N \leq \text{rank}(AB)$ .

**Cholesky decomposition:** Any  $A$  hermitian PD matrix can be expressed as  $A = LL^*$ , where where  $L$  is an invertible lower triangular matrix with real and positive diagonal entries.

**Vandermonde matrix:** A Vandermonde matrix  $V$  of size  $M \times N$  has entries of the form  $V_{i,j} = t_i^j$  for  $i = 0, \dots, M-1, j = 0, \dots, N-1$ , and it is of full rank if  $t_l \neq t_m$  when  $l \neq m$  since  $\det(V) = \prod_{0 \leq l < m \leq N-1} (t_l - t_m)$ .

**Circulant matrix:** Implement circular convolution. Its rows are circular rotations of a sequence. They are diagonalized by the DFT matrix.

**Toeplitz matrix:** Constant coefficients along diagonals. Implements linear convolution. In a matrix representation of a linear and shift-invariant system, the matrix will be Toeplitz.

**Unitary matrix:**  $U^{-1} = U^*$  and its eigenvalues satisfy  $|\lambda_j| = 1$ . Unitary matrices are isometries, i.e.  $\|U\mathbf{x}\| = \|\mathbf{x}\|$ .

**Normal matrices:**  $AA^* = A^*A$  (so need to be square).

**Eigenvalue decomposition:**  $A\mathbf{x} = \lambda\mathbf{x}$  (characteristic equation),  $\|\mathbf{x}\| = 1, \lambda \in \mathbb{F}$ .  $\lambda$  verifies  $\det(\lambda I - A) = 0$ . Square matrix diagonalizable if  $\exists S$   $n \times n$  and  $\Lambda$  diagonal matrix s.t.  $A = SAS^{-1}$ . For all square matrices,  $A = URU^*$  with  $R$  upper triangular matrix,  $U$  unitary (Schur decomposition). For normal matrices,  $R = \Lambda$ , with  $\Lambda$  diagonal.

**Singular value decomposition:**  $A = USV^*, U, V$  unitary,  $S$  (e.g. for tall matrix) hermitian diagonal block and 0 block,  $U = [U_1 \quad U_0]$ . This leads to thin SVD  $A = U_1 S_1 V^*, U_1, U_0$  form orthonormal basis for  $\mathcal{R}(A), \mathcal{N}(A^*)$  resp.  $U, V$  are resp- eigenvectors of normal matrices  $AA^*$  and  $A^*A$ .  $A^\dagger = VS^\dagger U^* (S^\dagger \text{ is } S^T \text{ with inverted singular values})$ .

**Determinant of a matrix:** Oriented volume of the hyper-parallelepiped defined by column vectors of  $A$ .  $|\det(U)| = 1$  for unitary matrices (e.g. rotation and reflection).

**Geometric series:** For  $r \neq 1, \sum_{k=n_0}^n r^k = \frac{r^{n_0} - r^{n+1}}{1-r}$ .

**Pseudo-inverse:** The pseudo-inverse of  $A$  denoted  $A^\dagger$  satisfies all the four statements:  $AA^\dagger A = A, A^\dagger AA^\dagger = A, (AA^\dagger)^* = AA^\dagger$  and  $(A^\dagger A)^* = A^\dagger A$ . If  $A$  has linearly independent columns (non-singular),  $A^*A$  invertible and  $A^\dagger = (A^*A)^{-1}A^*$  is a left-inverse. If  $A$  has linearly independent rows ( $A^*$  non-singular),  $AA^*$  invertible and  $A^\dagger = A^*(AA^*)^{-1}$  is a right-inverse.

**Building projections:**  $AA^\dagger$  orthogonal projection into  $\mathcal{R}(A)$ .  $A^\dagger A$  orthogonal projection into  $\mathcal{R}(A^*)$  (so orthogonal projection into  $\mathcal{N}(A)$  is  $I - A^\dagger A$ ). For  $U$  orthonormal set of vectors,  $U^\dagger = U^*$ .

**Spectral norm:**  $\|A\| := \sup_{\mathbf{x} \neq 0} \frac{\|A\mathbf{x}\|}{\|\mathbf{x}\|} = \sup_{\|\mathbf{x}\|=1} \|A\mathbf{x}\| = \sigma_{\max}(A) = \sqrt{\lambda_{\max}(A^*A)}$ . Spectral norm is submultiplicative:  $\|AB\| \leq \|A\| \|B\|$ .

**Change of coordinates:**  $\iint_C f(x,y) dx dy = \iint_P f(r,\phi) \left| \frac{\partial x}{\partial r} \frac{\partial y}{\partial r} \right| \frac{\partial x}{\partial \phi} \frac{\partial y}{\partial \phi} dr d\phi$ .

In case of switching from cartesian to polar,  $x = r \cos(\phi), y = r \sin(\phi), r = \sqrt{x^2 + y^2}, \phi = \arctan\left(\frac{y}{x}\right)$  and determinant becomes  $r$ .

## Trigonometric identities

**Even/odd:**  $\sin(-x) = -\sin(x), \cos(-x) = \cos(x), \tan(-x) = -\tan(x)$ .

**Cofunction:**  $\sin(\frac{\pi}{2} - x) = \cos(x), \cos(\frac{\pi}{2} - x) = \sin(x)$ .

**Sum and difference of angles:**  $\sin(x+y) = \sin x \cos y + \cos x \sin y, \sin(x-y) = \sin x \cos y - \cos x \sin y, \cos(x+y) = \cos x \cos y - \sin x \sin y, \cos(x-y) = \cos x \cos y + \sin x \sin y$ .

**Double angles:**  $\sin(2x) = 2 \sin x \cos x, \cos(2x) = \cos^2 x - \sin^2 x = 2 \cos^2 x - 1 = 1 - 2 \sin^2 x$ .

**Product to sum:**  $\sin x \sin y = \frac{1}{2} [\cos(x-y) - \cos(x+y)], \cos x \cos y = \frac{1}{2} [\cos(x-y) + \cos(x+y)], \sin x \cos y = \frac{1}{2} [\sin(x+y) + \sin(x-y)]$ .

**Derivatives:**  $\frac{d}{dx} \sin(x) = \cos(x), \frac{d}{dx} \cos(x) = -\sin(x)$ .

## Ranges, nullspaces and invertibility

For matrix  $A$  of size  $m \times n$ .

**Range:**  $\mathcal{R}(A) = \{y : y = A\mathbf{x}\}$  (linear subspace of  $\mathbb{F}^m$ ).

**Nullspace:**  $\mathcal{N}(A) = \{x : A\mathbf{x} = 0\}$  (linear subspace of  $\mathbb{F}^n$ ).

**Rank:** Cardinality of largest set of linearly independent columns (or rows) of  $A$ ;  $\text{rank}(A) = \dim(\mathcal{R}(A)) = \dim(\mathcal{R}(A^*)) \leq \min(m, n)$  (if equality, full rank).  $\text{rank}(A) = \text{rank}(A^*)$ .

**Rank-nullity theorem:**  $\dim(\mathcal{R}(A)) + \dim(\mathcal{N}(A^*)) = m$  (since  $\mathcal{R}(A) = \mathcal{N}(A^*)^\perp$ ),  $\dim(\mathcal{R}(A^*)) + \dim(\mathcal{N}(A)) = n$  (since  $\mathcal{R}(A^*) = \mathcal{N}(A)^\perp$ ).

**Invertibility (equivalences):** For  $A$  square matrix,  $\mathcal{N}(A) = \{0\}$  (trivial nullspace), full-rank matrix, no 0 eigenvalue, determinant equal to 0.

## Hilbert spaces and projection operators

**Vector space:** Set of vectors ( $\mathbb{R}^N$ , functions,...), field of scalars (real, complex), vector addition, scalar multiplication. Satisfy  $x + y = y + x$  (commutativity),  $(x + y) + z = x + (y + z)$  and  $(\alpha\beta)x = \alpha(\beta x)$  (associativity),  $\alpha(x + y) = \alpha x + \alpha y$  and  $(\alpha + \beta)x = \alpha x + \beta x$  (distributivity),  $\exists 0$  s.t.  $x + 0 = x$  (additive identity),  $\exists -x$  s.t.  $x + (-x) = 0$  (additive inverse),  $1x = x$  (multiplicative identity).

**Subspace:**  $S \subseteq V, V$  vector space and also  $S$  itself.  $S \neq \emptyset$ , closed under vector addition  $(x + y \in S \forall x, y \in S)$  and scalar multiplication  $(\alpha x \in S \forall x \in S, \alpha \in \mathbb{F}, \text{ so } \{0\} \in S)$ .

**span( $S$ ):**  $\{\sum_{k=0}^N \alpha_k \phi_k : \alpha_k \in \mathbb{F}, \phi_k \in S, N \in \mathbb{N} \setminus \{\infty\}\}$ . Smallest vector space containing the set of vectors  $S$ . For  $S$  infinite,  $\overline{\text{span}}(S)$  (o.w. some vectors cannot be represented with **finite** linear combination of  $\phi_k$ ). Always a subspace.

**Inner product:**  $\langle \mathbf{x} + \mathbf{y}, \mathbf{z} \rangle = \langle \mathbf{x}, \mathbf{z} \rangle + \langle \mathbf{y}, \mathbf{z} \rangle$  (distributivity),  $\langle \alpha \mathbf{x}, \mathbf{y} \rangle = \alpha \langle \mathbf{x}, \mathbf{y} \rangle$  (linearity in the first argument),  $\langle \mathbf{x}, \mathbf{y} \rangle^* = \langle \mathbf{y}, \mathbf{x} \rangle$  (hermitian symmetry),  $\langle \mathbf{x}, \mathbf{x} \rangle \geq 0$  equality iff  $\mathbf{x} = 0$  (PD).

On  $\mathbb{C}^{\mathbb{R}}, \langle x, y \rangle = \int_{-\infty}^{\infty} x(t)y^*(t)dt$ .

**Norm:**  $\|\mathbf{x}\| \geq 0$  equality iff  $\mathbf{x} = 0$  (PD),  $\|\alpha \mathbf{x}\| = |\alpha| \|\mathbf{x}\|$  (positive scalability),  $\|\mathbf{x} + \mathbf{y}\| \leq \|\mathbf{x}\| + \|\mathbf{y}\|$  equality iff  $\mathbf{y} = \alpha \mathbf{x}$  (triangle inequality).  $\langle \mathbf{x}, \mathbf{y} \rangle = \|\mathbf{x}\| \|\mathbf{y}\| \cos \alpha$ .

**Pythagorean theorem:** For  $\{\mathbf{x}_k\}$  orthogonal,  $\|\sum_k \mathbf{x}_k\|^2 = \sum_k \|\mathbf{x}_k\|^2$ .

**Cauchy-Schwarz:**  $|\langle \mathbf{x}, \mathbf{y} \rangle| \leq \|\mathbf{x}\| \|\mathbf{y}\|$ . Equality when  $\mathbf{x} = \alpha \mathbf{y}$  (collinear).

**Parallelogram law:**  $\|x + y\|^2 + \|x - y\|^2 = 2(\|x\|^2 + \|y\|^2)$ .

**Linear operator:**  $A(x + y) = Ax + Ay$  (additivity),  $A(\alpha x) = \alpha(Ax)$  (scalability). Operator norm  $\|A\| = \sup_{\|\mathbf{x}\|=1} \|A\mathbf{x}\|$ .

**Adjoint:** For  $A : \mathcal{H}_0 \rightarrow \mathcal{H}_1, \langle A\mathbf{x}, \mathbf{y} \rangle_{\mathcal{H}_1} = \langle \mathbf{x}, A^*\mathbf{y} \rangle_{\mathcal{H}_0}$ . Exists and is unique.  $\|A^*\| = \|A\|$ . Operator is unitary iff  $A^{-1} = A^*$ . If a matrix is self-adjoint (hermitian), its eigenvectors define an orthogonal basis. If  $A$  invertible,  $(A^{-1})^* = (A^*)^{-1}$ .

**Projection:** Bounded linear operator ( $\|P\| < \infty$ ; linear operators with finite-dimensional domains are always bounded) that is idempotent  $P^2 = P$  (in such case  $\|P\| \geq 1$ ). If self-adjoint  $P^* = P$ , orthogonal projection and all eigenvalues are real valued (o.w. oblique). Bounded linear operator  $P$  satisfies  $\langle x - Px, Py \rangle = 0 \forall x, y \in H$  iff  $P$  orthogonal projection. If  $S, T$  closed subspaces s.t.  $H = S \oplus T \exists$  projection  $P$  on  $H$  s.t.  $S = \mathcal{R}(P), T = \mathcal{N}(P)$ . In an orthogonal projection, all eigenvalues are either 0 or 1, and we have that  $\|P\mathbf{x}\| \leq \|\mathbf{x}\|$  (orthogonal projection is a contraction). Moreover, if range space is not trivial,  $\|P\| = 1$ .

**Projection theorem:** For  $S$  closed subspace of Hilbert space  $H$  and  $x \in H, \|x - \hat{x}\| \leq \|x - s\| \forall s \in S$  iff  $x - \hat{x} \perp S, \hat{x} = Px, P$  orthogonal projection.

**Basis:**  $\Phi = \{\phi_k\}_{k \in \mathcal{K}} \subset V, V = \overline{\text{span}}(\Phi)$  so any  $x \in V$  can be expressed as  $x = \sum_{k \in \mathcal{K}} \alpha_k \phi_k$  and expansion coefficients  $\alpha_k$  are unique.

**Linear independence:**  $\sum_{k \in \mathcal{K}} \alpha_k \phi_k = 0$  iff  $\alpha_k = 0 \forall k$ .

**Riesz basis:**  $\Phi$  basis of Hilbert space  $H, \exists 0 < \lambda_{\min} \leq \lambda_{\max} < \infty$  s.t.  $\forall x \in H, \lambda_{\min} \|\mathbf{x}\|^2 \leq \sum_{k \in \mathcal{K}} |\alpha_k|^2 \leq \lambda_{\max} \|\mathbf{x}\|^2$ . Let  $G = \Phi^* \Phi$  the **Gram matrix**.  $\lambda_{\max}(G) = 1/\lambda_{\min}, \lambda_{\min}(G) = 1/\lambda_{\max}$ .

**Orthogonal projection:**  $\alpha_k = \langle x, \phi_k \rangle$ , i.e.  $\alpha = \Phi^* x$ . This gives orthogonal projection to  $\overline{\text{span}}\{\phi_k\}$  (also for orthonormal set, not necessarily basis).

**Gram matrix:**  $G = \Phi^* \Phi, G_{i,k} = \langle \phi_k, \phi_i \rangle$ .

**Biorthogonal pairs of bases:**  $\Phi, \tilde{\Phi}$  both bases for  $H$  and biorthogonal:  $\langle \phi_i, \tilde{\phi}_k \rangle = \delta_{i-k}$ . In this case  $\alpha_k = \langle x, \tilde{\phi}_k \rangle$ , i.e.  $\alpha = \tilde{\Phi}^* x$ . Residual of projection on biorthogonal pairs of sets, satisfies  $x - Px \perp \overline{\text{span}}(\{\tilde{\phi}_k\}_{k \in \mathcal{I}})$   $\tilde{\Phi} = \Phi G^{-1}$  if  $\Phi$  Riesz basis.

**Parseval equality:** For  $\Phi$  orthonormal basis for  $H, \langle x, y \rangle = \langle \Phi^* x, \Phi^* y \rangle = \langle \alpha, \beta \rangle$  since  $\Phi^* \Phi = I$  on  $\ell^2(\mathcal{K})$ ,  $\Phi \Phi^* = I$  on  $H$  (unitary operator). For biorthogonal pairs of bases,  $\langle x, y \rangle = \langle \tilde{\alpha}, \beta \rangle$  since  $\tilde{\Phi}^* \Phi = I$  on  $\ell^2(\mathcal{K})$ ,  $\Phi P \tilde{H}^* = I$  on  $H$ .

**Bessel’s inequality:** For orthonormal sets,  $\Phi^* \Phi = I$  but  $\Phi \Phi^* \neq I$  in general.  $\|\mathbf{x}\|^2 \geq \|\Phi_T^* \mathbf{x}\|^2$  with equality when  $\Phi$  is a basis.

**Inverse problem:** Find  $x$  s.t.  $y = Ax$ . Solution  $\hat{x}$  is consistent with measurements if  $\hat{x} : \hat{x} = x + \tilde{x}, \tilde{x} \in \mathcal{N}(A)$ , i.e.  $A\tilde{x} = y$ .

**Dual basis;**  $\tilde{\Phi} = \Phi(\Phi\Phi^*)^{-1} = \Phi G^{-1}, G$  the Gram matrix.

**Change of basis:**  $x, y \in H, y = Ax, x = \Phi\alpha, y = \Psi\beta, \beta = \Gamma\alpha. \Gamma_{i,j} = \langle A\phi_j, \tilde{\psi}_i \rangle$ .

**Gram-Schmidt procedure:** Want to convert original set  $\{s^{(k)}\}$  onto orthonormal set  $\{u^{(k)}\}$ . At each step  $k, p^{(k)} = s^{(k)} - \sum_{n=0}^{k-1} \langle u^{(n)}, s^{(k)} \rangle u^{(n)}$ .

$u^{(k)} = p^{(k)} / \left\| p^{(k)} \right\|$ .

## Discrete Systems

**Delta function:**  $\int_{-\infty}^{\infty} \delta(t)dt = 1, \delta(t) = 0 \forall t \neq 0, \int_{-\infty}^{\infty} x(t)\delta(t)dt = x(0), \int_{-\infty}^{\infty} x(\tau)\delta(t - \tau)d\tau = x(t). \delta(at) = \delta(t)/|a|$ .

**Impulse response:** Let  $y_k(t) = A(x_k(t))$ . If a linear  $(A(\sum_k \alpha_k x_k(t)) = \sum_k \alpha_k y_k(t))$  shift-invariant  $(A(x_k(t - \tau)) = y_k(t - \tau))$  system, can express  $A$  as convolution with impulse response  $h(t) = A(\delta(t)); A(x(t)) = (x * h)(t)$ .

**Filter as projection:** A filter is said to be a projection if  $h_n = (h * h)_n$ , or equivalently  $H(e^{j\omega}) = H^2(e^{j\omega})$ . In case of being a projection, orthogonal projection if  $h_n = h_{-n}^*$ , or equivalently  $H(e^{j\omega}) = H^*(e^{j\omega})$ .

## Transforms

	Time domain	Frequency domain
Fourier Transform	Continuous aperiodic	Continuous aperiodic
Fourier Series	Continuous periodic	Discrete aperiodic
DTFT	Discrete aperiodic	Continuous aperiodic
DFT	Discrete periodic	Discrete periodic

## Fourier Transform

**Definition:**  $X(\omega) = \int_{-\infty}^{\infty} x(t)e^{-j\omega t}dt, x(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} X(\omega)e^{j\omega t}d\omega$ .

**Continuous-time convolution:**  $(x * y)(t) = \int_{-\infty}^{\infty} x(\tau)y(t - \tau)d\tau$ . Eigenfunctions  $e^{j\omega t}$ .

**Parseval:**  $\langle x(t), y(t) \rangle = \frac{1}{2\pi} \langle X(\omega), Y(\omega) \rangle$ .

**Properties:**  $x(\alpha t) \leftrightarrow \frac{1}{\alpha} X(\frac{\omega}{\alpha}), (x * y)(t) \leftrightarrow X(\omega)Y(\omega), x(t)y(t) \leftrightarrow \frac{1}{2\pi} (X * Y)(\omega), x(t - \tau) \leftrightarrow e^{-j\omega\tau} X(\omega), e^{j\omega_0 t} x(t) \leftrightarrow X(\omega - \omega_0), x(-t) \leftrightarrow X(-\omega), x^*(t) \leftrightarrow X^*(-\omega), \frac{d^n x(t)}{dt^n} \leftrightarrow (j\omega)^n X(\omega), (-j t)^n x(t) \leftrightarrow \frac{d^n X(\omega)}{d\omega^n}, \int_{-\infty}^t x(\tau)d\tau \leftrightarrow \frac{X(\omega)}{j\omega}, X(0) = 0$ .

**Poisson sum formula:**  $\sum_{n \in \mathbb{Z}} x(t - nT) = \frac{1}{T} \sum_{m \in \mathbb{Z}} X(\frac{2\pi m}{T})e^{j\frac{2\pi m t}{T}}$ . Proof: LHS is  $T$ –periodic so compute Fourier Series and anti-transform result.

**Common transforms:**  $\text{sinc}(\omega_0 t) = \frac{\sin \omega_0 t}{\omega_0 t} \leftrightarrow \frac{\pi}{\omega_0} \mathbb{1}_{\{\omega \in [-\omega_0, \omega_0]\}}$ .  $\sum_{n \in \mathbb{Z}} \delta(t - nT) \leftrightarrow \frac{2\pi}{T} \sum_{k \in \mathbb{Z}} \delta(\omega - \frac{2k\pi}{T}). \frac{1}{t_0} \mathbb{1}_{\{t \in [-t_0/2, t_0/2]\}} \leftrightarrow \text{sinc}(\frac{t_0 \omega}{2}). 1 - |t|, |t| < 1 \leftrightarrow \text{sinc}^2(\frac{\omega}{2}). 1 \leftrightarrow 2\pi \delta(\omega). \delta(t) \leftrightarrow 1$ .

## Fourier Series (signal of period T)

**Definition:**  $X_k = \frac{1}{T} \int_{-\frac{T}{2}}^{\frac{T}{2}} x(t)e^{-j\frac{2\pi k t}{T}}dt, x(t) = \sum_{k \in \mathbb{Z}} X_k e^{j\frac{2\pi k t}{T}}$ .

**Circular continuous-time convolution:** For  $x, y T$ –periodic.  $(x \otimes y)(t) = \int_{-\frac{T}{2}}^{\frac{T}{2}} x(\tau)y(t - \tau)d\tau$ . Eigenfunctions  $e^{j\frac{2\pi k t}{T}}$ .

**Parseval:**  $\langle x(t), y(t) \rangle = T \langle X_k, Y_k \rangle$ .

**Properties:**  $x(t - t_0) \leftrightarrow e^{-j\frac{2\pi k t_0}{T}} X_k, e^{j\frac{2\pi k_0 t}{T}} x(t) \leftrightarrow X(k - k_0), (h \otimes x)(t) \leftrightarrow T H_k X_k, h(t)x(t) \leftrightarrow (H * X)_k$ .

**Common transforms:**  $\sum_{n \in \mathbb{Z}} \delta(t - nT) \leftrightarrow 1/T$ .

## DFTF

**Definition:**  $X(e^{j\omega}) = \sum_{n \in \mathbb{Z}} x_n e^{-j\omega n}, x_n = \frac{1}{2\pi} \int_{-\pi}^{\pi} X(e^{j\omega})e^{j\omega n}d\omega$ .

**Discrete convolution:**  $(x * y)_n = \sum_{k \in \mathbb{Z}} x_k y_{n-k}$ . Eigenfunctions  $e^{j\omega n}$ .

**Properties:** Transform is  $2\pi$ –periodic.  $(h * x)_n \leftrightarrow H(e^{j\omega})X(e^{j\omega}), h_n x_n \leftrightarrow \frac{1}{2\pi} (H \otimes X)(e^{j\omega})$ .

## DFT

**Definition:**  $X_k = \sum_{n=0}^{N-1} x_n W_N^{kn}$ ,  $x_n = \sum_{k=0}^{N-1} X_k W_N^{-kn}$ ,  $W_N^{kn} = e^{-j \frac{2\pi kn}{N}}$ .

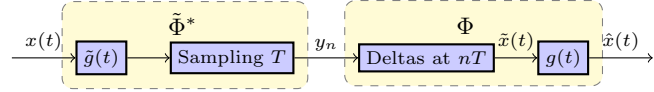
**Discrete circular convolution:** For  $x, y$  sequences of length  $N$ .  $(x \otimes y)_n = \sum_{k=0}^{N-1} x_k y_{n-k}$ . Eigenfunctions  $e^{j \frac{2\pi kn}{N}}$ . Same as discrete convolution if  $x$  has support  $M$ ,  $y$  has support  $L$  and  $N \geq L + M - 1$ .

**Properties:**  $x_{(n-n_0) \bmod N} \leftrightarrow W_N^{kn_0} X_k$ ,  $W_n^{-k_0 n} \leftrightarrow X_{(k-k_0) \bmod N}$ ,  $(h \otimes x)_n \leftrightarrow H_k X_k$ ,  $h_{nN} \leftrightarrow \frac{1}{N} (H \otimes X)_k$ ,  $x_n^* = X_{-k}^* \bmod N$ ,  $x_{-n}^* \bmod N \leftrightarrow X_k^*$ .

## Sampling and Interpolation

Interpolation  $\tilde{\Phi}$ , sampling  $\tilde{\Phi}^*$ .  $\tilde{S} = \mathcal{N}(\tilde{\Phi}^*)^\perp = \mathcal{R}(\tilde{\Phi})$  (what can be measured).  $S = \mathcal{R}(\Phi)$  (what can be reproduced).  $P = \Phi \tilde{\Phi}^*$  projection (with range  $S$  and  $x - \hat{x} \perp \tilde{S}$ , where  $\hat{x} = Px$  iff  $\tilde{\Phi}^* \Phi = I$  (or equivalently  $\langle g(t - kT), \tilde{g}^*(nT - T) \rangle_t = \delta_{k-n}$ ). In this case, we say that sampling and interpolation are consistent (and  $\hat{x}$  is the best least-squares approximation of  $x$  in  $S$ ). When  $\Phi = (\tilde{\Phi}^*)^\dagger = \tilde{\Phi}(\tilde{\Phi}^* \tilde{\Phi})^{-1}$  (for orthogonal vectors, pseudo-inverse is the adjoint), they form a biorthogonal pair of bases for  $S$  and hence  $S = \tilde{S}$  and we say that operators are ideally matched (orthogonal projection).

To prove  $S = \tilde{S}$  when  $\Phi = (\tilde{\Phi}^*)^\dagger$ , note that  $\Phi$  is a linear combination of columns of  $\tilde{\Phi}$  with coefficients given by the corresponding column of  $(\tilde{\Phi}^* \tilde{\Phi})^{-1}$  (full rank matrix). Hence, columns of  $\Phi, \tilde{\Phi}$  must span the same space.



$$y_n = \int_{-\infty}^{\infty} x(\tau) g^*(\tau - t) d\tau \Big|_{t=nT} \quad ; \quad \hat{x}(t) = \sum_{n \in \mathbb{Z}} y_n \delta(t - nT)$$

$$\hat{x}(t) = \sum_{k \in \mathbb{Z}} y_k g(t - kT) = \int_{-\infty}^{\infty} \tilde{x}(\tau) g(t - \tau) d\tau = \Phi \tilde{\Phi}^* x = Px.$$

In the orthogonal case,  $\tilde{\Phi}^* = \Phi^*$  and  $\tilde{g}(t) = g^*(-t)$ .

**Shift-invariant subspace:** A subspace  $S \subset \mathcal{L}^2(\mathbb{R})$  is a shift-invariant subspace with respect to shift  $T \in \mathbb{R}^+$  if  $x(t) \in S \implies x(t - kT) \in S \forall k \in \mathbb{Z}$ .  $s \in \mathcal{L}^2(\mathbb{R})$  is called a generator of  $S$  when  $S = \overline{\text{span}}(\{s(t - kT)\}_{k \in \mathbb{Z}})$ . To check for latter, have to show that  $\forall x \in S$ ,  $\exists \{\alpha_k\}_{k \in \mathbb{Z}}$  unique s.t.  $x(t) = \sum_{k \in \mathbb{Z}} \alpha_k s(t - kT)$ .

**Subspace of bandlimited functions:** A function  $x(t) \in \mathcal{L}^2(\mathbb{R})$  has bandwidth  $\omega_0 \in [0, \infty)$  if its FT satisfies  $X(\omega) = 0 \forall |\omega| > \frac{\omega_0}{2}$ . This defines the space of  $\omega_0$ -bandlimited functions,  $\text{BL}[-\omega_0/2, \omega_0/2]$  which is a shift-invariant subspace (proof by delay property of FT).

**Sampling theorem:** If  $x \in \text{BL}[-\pi/T, \pi/T]$ ,  $x(t) = \sum_{n \in \mathbb{Z}} x(nT) \text{sinc}(\frac{\pi}{T}(t - nT))$ . This means that sinc is a generator for  $\frac{2\pi}{T}$ -band limited functions. In this case, choose  $g(t) = \frac{1}{\sqrt{T}} \text{sinc}(\frac{\pi}{T}t) = g^*(-t)$  since it is a generator with shift  $T$  of  $\text{BL}[-\pi/T, \pi/T]$  and  $\langle g(t - kT), g(t - lT) \rangle = \delta_{k-l}$  are orthonormal. If  $x \in \text{BL}[-\omega_0/2, \omega_0/2]$ , we need  $T < 2\pi/\omega_0$  (Nyquist interval).  $\omega_0/2\pi$  is called the Nyquist rate.

**Continuous-time convolution via DSP:** For  $x \in \text{BL}[-\pi/T, \pi/T]$ ,  $y = h * x$  can be computed by sampling, filtering the resulting sequence and interpolating the result of the convolution.  $\hat{h}_n = \langle h(t), \text{sinc}(\frac{\pi}{T}(t - nT)) \rangle$  without pre-filter (first filter in sampling is multiplying by  $\sqrt{T}$ ).

## Computational Tomography

Parametrize line with angle  $\theta \in [0, \pi)$  of the line's normal vector and signed distance  $t \in \mathbb{R}$  from the origin. Cannot use  $y = mx + n$  since vertical lines cannot be described.

**Explicit:**  $x, y$  parametrized by  $s \in \mathbb{R}$ , where  $L_{\theta, t} = \{(x(s), y(s))\}$ .  $x(s) = t \cos \theta - s \sin \theta$ ,  $y(s) = t \sin \theta + s \cos \theta$ .

**Implicit:**  $L_{\theta, t} = \{(x, y) : x \cos \theta + y \sin \theta = t\}$ .

**Radon transform:** Gives sinogram.  $\mathcal{R}[f](\theta, t) = p(\theta, t) = \int_{L_{\theta, t}} f(x, y) ds = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x, y) \delta(x \cos \theta + y \sin \theta - t) dx dy$ .

**Laminogram or Backprojection Summation:** Adjoint of Radon transform. Assign every point in the image along  $L_{\theta, t}$  the value  $p(\theta, t)$ . Gives blurred reconstruction.  $\mathcal{B}[p](x, y) = f_b(x, y) = \int_0^\pi p(\theta, x \cos \theta + y \sin \theta) d\theta$ .

**Fourier Slice Theorem:** The Fourier transform of a parallel projection of

an image  $f(x, y)$  taken at angle  $\theta$  gives a slice of its two-dimensional Fourier transform,  $F(u, v)$ , that subtends an angle  $\theta$  with the  $u$ -axis.  $P(\theta, \omega) = F(u, v)|_{u=\omega \cos \theta, v=\omega \sin \theta} = F(\omega \cos \theta, \omega \sin \theta)$ , where  $P(\theta, \omega)$  is the 1D-FT of  $p(\theta, t)$  w.r.t  $t$ .

**Fourier Reconstruction Method:** Take the 1D Fourier transform of each projection (w.r.t.  $t$ ). Insert the results in the appropriate slices in the  $(u, v)$ -plane. Resample on a rectangular grid in the  $(u, v)$ -plane. Take the 2D IFT of the formed 2D spectrum.

## Filtered Backprojection (FBP)

Gives images very close to the original. FBP does not work easily with constraints. With limited data and/or non-uniform distribution of projection angles, reconstruction with FBP contains artifacts. In theory,  $f(x, y) = \mathcal{B}[q](x, y)$  where  $q(\theta, t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} P(\theta, \omega) \frac{|\omega|}{2\pi} e^{j\omega t} d\omega$ . To prove this, do inverse 2D-IFT of  $F(u, v)$ . Change Cartesian to polar coordinates with  $\omega \in \mathbb{R}_+, \theta \in [0, 2\pi)$  (Recall that  $dudv = \omega d\theta d\omega$ ). Split integral in  $\theta$  from 0 to  $\pi$  for  $\theta$  and  $\theta + \pi$ . Note that  $\tilde{F}(\omega, \theta + \pi) = \tilde{F}(-\omega, \theta)$  since  $\cos(\theta + \pi) = -\cos \theta$  and  $\sin(\theta + \pi) = -\sin \theta$ . Apply Fourier Slice Theorem.

**Convolution backprojection:** Can also write  $f(x, y) = \int_{-\infty}^{\infty} p(\theta, t) h(x \cos \theta + y \sin \theta - t) dt d\theta$ , where  $h(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} \frac{|\omega|}{\pi} e^{j\omega t} d\omega$ . Given that  $|\omega| \notin \mathcal{L}^2(\mathbb{R})$ , its FT doesn't exist, so in practice  $H(\omega) = \frac{|\omega|}{2\pi} W(\omega)$  for  $W$  a suitable window. IFT of ramp filter  $\frac{|\omega|}{2\pi}$  is sinc - sinc<sup>2</sup> (rectangle - triangle).

**Reconstruct image from sinogram:** Take 1D FT of projections (w.r.t.  $t$ ) apply filter  $H(\omega)$  and take 1D IFT (if  $h(t)$  short, do convolution in spatial domain). Backproject filtered projections and sum backprojected images.

## Algebraic Reconstruction

Suppose  $f(x, y) = \sum_{i=0}^{N-1} f_i \phi_i(x, y)$ , where  $\phi_k$  are the basis functions. Beam shapes are  $h_i$ . Hence,  $b_i = \langle \mathbf{f}, h_i \rangle$ . With spline surface model,  $\mathbf{b}$  (the measurement) results from  $\mathbf{b} = \mathbf{A}\mathbf{f}$ , where  $A$  is the measurement matrix or forward operator with  $A_{i,k} = \langle \phi_k, h_i \rangle$  of size  $M \times N$  and  $\mathbf{f}$  the unknown pixel weight vector. Usually, number of non-zero coefficients in each row of  $A$  is  $\mathcal{O}(\sqrt{N})$ , so  $A$  sparse (so avoid computing pseudo-inverse, usually dense). Can view  $\mathbf{A}\mathbf{f}$  as forward projection,  $A^T \mathbf{b}$  as backprojection. In presence of noise system can be inconsistent: Usually  $\hat{\mathbf{f}} = \arg \min_{\mathbf{f}} \|\mathbf{b} - \mathbf{A}\mathbf{f}\|_2^2$ . If  $\text{rank}(A) = N$ ,  $\hat{\mathbf{f}} = (A^T A)^{-1} A^T \mathbf{b}$  unique. If  $\text{rank}(A) = M < N$ , take solution with minimum norm  $\hat{\mathbf{f}} = A^T (A A^T)^{-1} \mathbf{b}$  (both interpreted as FBP). With this approach, handle constraints of scanning topology.

## Kaczmarz's algorithm

Let  $\{\mathbf{r}_1^T, \dots, \mathbf{r}_{M-1}^T\}$  the rows of  $A$ . Hence  $\langle \mathbf{f}, \mathbf{r}_i \rangle = b_i$ , which define affine hyperplanes in  $\mathbb{R}^N$ . If solution to  $\mathbf{A}\mathbf{f} = \mathbf{b}$  exists and is unique, it's intersection of  $M$  affine hyperplanes. Kaczmarz's satisfies constraints iteratively starting from initial guess  $\mathbf{f}^{(-1)}$ .  $\mathbf{f}^{(k)} = \mathbf{f}^{(k-1)} + \frac{b_n - \langle \mathbf{f}^{(k-1)}, \mathbf{r}_n \rangle}{\|\mathbf{r}_n\|_2^2} \mathbf{r}_n$

for  $n = k \bmod M$ . That is, apply orthogonal projection of  $\mathbf{f}^{(k-1)}$  onto  $\mathcal{H}_n = \{\mathbf{f}^{(-1)} : \langle \mathbf{f}^{(k-1)}, \mathbf{r}_n \rangle = b_n\}$ . If there's a unique solution, found with  $k \rightarrow \infty$ . Ordering of rows influence convergence rate (can order to increase angles between consecutive rows or select randomly with e.g.  $p_n = \frac{\|\mathbf{r}_n\|_2^2}{\|\mathbf{A}\|_F^2}$ ).

Can incorporate constraints such as box constraint ( $0 \leq f_i \leq 1$ ) by projecting each step onto this convex set.

## Cimmino's method

Instead of updating one row at a time, update once per sweep with average of all projections.  $\mathbf{f}^{(k)} = \mathbf{f}^{(k-1)} + A^T W^{-1} (\mathbf{b} - \mathbf{A}\mathbf{f}^{(k-1)})$  where  $W = \text{diag}(M \|\mathbf{r}_0\|^2, \dots, M \|\mathbf{r}_{M-1}\|^2)$ .

## Stochastic Processes

Covariance  $\Sigma_x = \mathbb{E}[(x - \mu_x)(x - \mu_x)^*]$ , which is PSD since  $u^* \Sigma_x u = \mathbb{E}[u^*(x - \mu_x)(x - \mu_x)^* u] = \mathbb{E}[|u^*(x - \mu_x)|^2]$ .  $\langle x, y \rangle = \sum_{n=0}^{N-1} \mathbb{E}[x_n y_n^*]$ . Autocorrelation  $a_{x,n,k} = \mathbb{E}[x_n x_{n-k}^*]$ , crosscorrelation  $c_{x,y,n,k} = \mathbb{E}[x_n y_{n-k}^*]$ . For i.i.d. process  $a_{x,n,k} = |\mu_x|^2 + \sigma_x^2 \delta_k$ .

**Stationary process:** Joint distribution of  $(x_{n_0}, \dots, x_{n_L})$  and  $((x_{n_0+k}, \dots, x_{n_L+k}))$  are identical  $\forall \{n_0, \dots, n_L\} \subset \mathbb{Z}, \forall k \in \mathbb{Z}$  and  $L$  finite.

**WSS:**  $\mu_{x,n} = \mu_x$ ,  $a_{x,n,k} = a_{x,k}$ ,  $n, k \in \mathbb{Z}$ .  $x$  and  $y$  jointly WSS if each is WSS and  $c_{x,y,n,k} = c_{x,y,k}$ . With WSS, we have  $\sigma_{x,n}^2 = a_{x,0} - |\mu_x|^2 = \sigma_x^2$ ,  $a_{x,k} = a_{x,-k}^*$ . With joint WSS,  $c_{y,x,k} = c_{y,x,-k}^*$ .

**White noise:**  $\mu_{x,n} = 0$ ,  $\sigma_{x,n}^2 = \sigma_x^2$  and  $a_{x,k} = \sigma_x^2 \delta_k$  (uncorrelated). Gaussian rv's uncorrelated iff independent, so white Gaussian process is i.i.d.

**Independent vs. uncorrelated:** If  $X, Y$  independent,  $\mathbb{E}[XY] = \mathbb{E}[X] \mathbb{E}[Y]$ , which implies that they are uncorrelated, i.e.  $\mathbb{E}[(X - \mathbb{E}[X])(Y - \mathbb{E}[Y])] = 0$ . Converse is only true for zero-mean random variables and jointly Gaussian random variables.

**Whitening or decorrelation:** Processing that results in white noise process. Diagonalization of covariance matrix.

**Filtering WSS processes:**  $y = x * h$  with  $h$  the impulse response of BIBO-stable LSI system (BIBO means L1 norm of  $h$  finite).  $\mu_{y,n} = \mu_x H(e^{j\omega}) = \mu_y$ ,  $a_{y,n,k} = \sum_{p \in \mathbb{Z}} a_{h,p} a_{x,k-p} = a_{y,k}$  so if x WSS, y WSS and x and y jointly WSS.  $C_{x,y}(e^{j\omega}) = H^*(e^{j\omega}) A_x(e^{j\omega})$  and  $C_{y,w}(e^{j\omega}) = H(e^{j\omega}) A_x(e^{j\omega})$ .

**Power spectral density:** x WSS. DFT of its autocorrelation (FT in case of continuous time).  $A_x(e^{j\omega}) = \sum_{k \in \mathbb{Z}} a_{x,k} e^{-j\omega k}$ .  $A_y(e^{j\omega}) = |H(e^{j\omega})|^2 A_x(e^{j\omega})$ .

**Power:**  $P_x = \frac{1}{2\pi} \int_{-\pi}^{\pi} A_x(e^{j\omega}) d\omega = a_{x,0}$ .

**Orthogonal Stochastic Processes:**  $c_{x,y,k,n} = \mathbb{E}[x_k y_n^*] = 0 \forall k, n \in \mathbb{Z}$ . For jointly WSS processes, equivalent to  $C_{x,y}(e^{j\omega}) = 0 \forall \omega \in \mathbb{R}$ .

**Wiener filtering:**  $\hat{x} = h * y$ ,  $h = \arg \min_h \mathbb{E}[|e_n := x_n - \hat{x}_n|^2]$ . Assume  $x, w$  uncorrelated, WSS, zero-mean.  $\hat{x} \in \mathcal{S} := \overline{\text{span}}(\{y_{n-k}\}_{k \in \mathbb{Z}})$ , so best estimator  $e \perp \mathcal{S}$ . This gives  $H(e^{j\omega}) = \frac{C_{x,y}(e^{j\omega})}{A_y(e^{j\omega})}$ .

## Beamforming

In narrowband beamforming, output is  $y_n = \sum_{k=0}^{M-1} h_k^* x_{k,n}$ ,  $M$  number of array elements. **Array response vector:**  $\mathbf{a}(\theta) = [H_0(\omega) e^{-j\omega \tau_0(\theta)} \dots H_{M-1}(\omega) e^{-j\omega \tau_{M-1}(\theta)}]^T$ ,  $H_k$  the transfer function of the  $k$ -th sensor,  $\tau_k(\theta)$  the time needed for the wave to travel from reference point to sensor  $k$ . If sensors ideal ( $H_k(\omega) = 1 \forall k$ ) and sensor 0 taken as reference,  $\mathbf{a}(\theta) = [1 \ e^{-j\omega \tau_1(\theta)} \dots (\omega) e^{-j\omega \tau_{M-1}(\theta)}]^T$ .  $x_n = \mathbf{a}(\theta) s_n + e_n$ . For multiple sources  $\mathbf{x} = [\mathbf{a}(\theta_0) \dots \mathbf{a}(\theta_{N-1})][s_0 \dots s_{N-1}]^T + \mathbf{e} = \mathbf{A}\mathbf{s} + \mathbf{e}$ .

**Uniform linear array:** Equispaced sensors in same line. Under plane-wave assumption,  $\tau_k = k \frac{d \sin \theta}{c}$ ,  $\theta \in [-\frac{\pi}{2}, \frac{\pi}{2})$  (front-back ambiguity),  $c$  velocity of wave propagation,  $d$  distance between consecutive sensors.

**Beamformer's response:**  $r(\theta) = \mathbf{h}^* \mathbf{a}(\theta)$ . Beampattern is  $|r(\theta)|^2$ . If  $\mathbf{a}(\theta_1) = \mathbf{a}(\theta_2)$  and  $\theta_1 \neq \theta_2$ , spatial aliasing, i.e. spacing between sensors is too large.

**Signal model:** For narrowband source with DOA  $\theta$  and power  $\sigma_s^2$ , if there's no noise  $\mathbf{x}_n = \mathbf{a}(\theta) s_n$ , so  $R_x = \sigma_s^2 \mathbf{a}(\theta) \mathbf{a}^*(\theta)$ .

## Data independent beamforming

**Phased array (delay-and-sum) beamformer:**  $\mathbf{h} = \mathbf{a}(\theta_0)$ , where signal comes from single location  $\theta_0$ . Can control beam width of main lobe and height of side lobes with tapering function  $\mathbf{h} = T \mathbf{a}(\theta_0)$ , with  $T$  diagonal  $M \times M$  with tapering weights.

**Response design:** To find  $\mathbf{h}$  giving response similar to  $r_d(\theta)$ , pose as overdetermined problem  $\arg \min_{\mathbf{h}} \|\mathbf{A}^* \mathbf{h} - \mathbf{r}_d\|_2^2$ ,  $A = [\mathbf{a}(\theta_0) \dots \mathbf{a}(\theta_{P-1})]$ ,  $\mathbf{r}_d = [r_d(\theta_0) \dots r_d(\theta_{P-1})]^*$ . If  $A$  full rank  $\mathbf{h} = (A A^*)^{-1} A \mathbf{r}_d$ .

**White noise gain:** Output power due to white noise of unit power, i.e.  $\mathbf{h}^* \mathbf{h}$ . If this is high, beamformer output could have poor SNR. Control it by low-rank approximation of  $A$  or solving previous problem with regularization  $+\lambda \|\mathbf{h}\|_2^2$ , which gives  $\mathbf{h} = (A A^* + \lambda I)^{-1} A \mathbf{r}_d$ .

## Data dependent beamforming

Output of beamformer approximates desired  $\mathbf{y}_d$ . Usually pose as  $\arg \min_{\mathbf{h}} \mathbb{E}[|y - \mathbf{y}_d|^2]$ , which gives  $\mathbf{h} = R_x^{-1} \mathbf{r}_d$ .

**LCMV:** Constrain so that signals from desired directions have specified gain with  $C^* \mathbf{h} = \mathbf{f}$ . Solve  $\min_{\mathbf{h}} \mathbf{h}^* R_x \mathbf{h}$  (minimize power at beamformer's output) subject to constraints. This gives  $\mathbf{h} = R_x^{-1} C (C^* R_x^{-1} C)^{-1} \mathbf{f}$ .

**GSC:** Transform LCMV to unconstrained. Decompose  $\mathbf{h} = \mathbf{h}_0 - \mathbf{g}$ ,  $\mathbf{h}_0 \in \mathcal{R}(C)$ ,  $\mathbf{g} \in \mathcal{N}(C^*)$ .  $\mathbf{h}_0 = C(C^* C)^{-1} \mathbf{f}$ , the min norm solution

satisfying constraints (data independent),  $\mathbf{g} = C_n \mathbf{h}_n$ , where  $C_n$  basis of  $\mathcal{N}(C^*)$  (has no contribution of satisfaction of constraint bt allows to minimize objective). Solve  $\min_{\mathbf{h}_n} (\mathbf{h}_0 - C_n \mathbf{h}_n)^* R_x (\mathbf{h}_0 - C_n \mathbf{h}_n)$ , which gives  $\mathbf{h}_n = (C_n^* R_x C_n)^{-1} C_n^* R_x \mathbf{h}_0$ . The data-dependent beamformer vector has nulls in the directions of the constraints, which is ensured by the signal blocking matrix  $C_n$ .

Approximation Theory

Polynomials (finite interval)

Approximate  $x(t)$  in finite interval  $[a, b]$  by polynomial of order  $K$   $p_K(t) = \sum_{k=0}^K a_k t^k$ . Approximation error  $e_K(t) = x(t) - p_K(t)$ ,  $t \in [a, b]$ . **Smooth approximation.** Can approximate continuous functions arbitrarily well over finite intervals (Weierstrass theorem). **Polynomials are infinitely differentiable.** Approximating continuous functions with high degree polynomials tends to be problematic (e.g. at ending points of interval). Cannot approximate discontinuous functions or over infinite intervals well.

Least square minimization

Minimize  $\|e_K\|_2^2 = \int_a^b |x(t) - p_K(t)|^2 dt$ . Since  $\mathcal{P}_K([a, b]) = \text{span}(\{1, t, \dots, t^K\}) \subset \mathcal{L}^2([a, b])$ , solution (by projection theorem) is  $p_K(t) = \sum_{k=0}^K \langle x, \phi_k \rangle \phi_k(t)$ ,  $\{\phi_k\}_{k=0}^K$  orthonormal basis of  $\mathcal{P}_K([a, b])$ . These inner products with basis functions are not always easy to obtain (e.g. if we only have samples). **Gibbs phenomenon:** No matter how large  $K$  is, absolute error stays the same at the ripple near to the boundary.

**Legendre polynomials:** Orthonormal basis of  $\mathcal{P}_K([-1, 1])$ . The Legendre polynomial of degree  $n$ ,  $L_n(t) = \frac{1}{2^n n!} \frac{d^n}{dt^n} (t^2 - 1)^n$ , has  $n$  real distinct zeros in the interior of the interval  $[-1, 1]$

Lagrange interpolation

Observe function  $x(t)$  at points (nodes)  $t_0, \dots, t_K$ . Constrain  $p_K(t_i) = x(t_i) \forall i$ . Leads to Vandermonde system, invertible iff  $\{t_i\}$  distinct with solution  $p_K(t) = \sum_{k=0}^K x(t_k) \prod_{i=0, i \neq k}^K \frac{t - t_i}{t_k - t_i}$ .

**Error:** For  $x(t) \in CK + 1([a, b])$  and  $\{t_i\}$  distinct,  $|e_K(t)| \leq \frac{\prod_{k=0}^K |t - t_k|}{(K+1)!} \max_{\eta \in [a, b]} |x^{(K+1)}(\eta)|$ . So error increases at the boundaries. Last term means error higher for wigglier functions.

Taylor series expansion

Assume  $x(t) \in C^K([a, b])$  and find degree  $K$  polynomial with matching derivatives at  $t_0 \in [a, b]$ . Solution  $p_K(t) = \sum_{k=0}^K x(t_k) \frac{(t - t_0)^k}{k!} x^{(k)}(t_0)$ .

**Error:** For  $x(t) \in CK + 1([a, b])$ ,  $|e_K(t)| \leq \frac{|t - t_0|^{K+1}}{(K+1)!} \max_{\eta \in [a, b]} |x^{(K+1)}(\eta)|$ .

Minimax approximation

Minimize  $\|e_K\|_\infty = \max_{t \in [a, b]} |e_K(t)|$ . Non trivial since not Hilbert space. With polynomial of degree  $\leq K$ , minimax approximation  $p_K$  unique and determined by at least  $K + 2$  points  $a \leq s_0 < s_1 < \dots < s_{K+1} \leq b$  for which  $e_K(s_k) = \pm 1(-1)^k \|e_K\|_\infty$  (Chebyshev equioscillation theorem). So expect to reach maximum error in  $K + 2$  points with alternating sign. Nearly solve by taking nodes minimizing maximum error for Lagrange interpolation. These optimal  $K + 1$  nodes given by roots of  $K + 1$ -degree Chebyshev polynomial (its scaled version  $2^{-K} T_{K+1}$  have the minimum  $\ell_\infty$  norm among all  $K + 1$ -degree polynomials):  $t_k = \cos\left(\frac{2k+1}{2(K+1)}\pi\right)$ .

Splines

Can approximate discontinuous functions well and over the entire real line. Spline of degree  $K$  with knots  $\tau$  (countable strictly-increasing sequence) is a polynomial of degree  $\leq K$  on  $[\tau_n, \tau_{n+1})$  and its derivatives of order  $0, \dots, K - 1$  are continuous.  $S_{K, \tau} \subset \mathcal{L}^2(\mathbb{R})$  is the spline space of degree  $K$  with knots  $\tau$ . When  $\tau$  evenly spaced and doubly infinite, spline space called uniform. For spline of degree  $K$  with  $L + 1$  knots,  $K - 1$  degrees of freedom. In practice for  $K = 3$ , specifies derivatives at end-points or make sure 3rd derivative continuous at second and penultimate knots (not-a-knot condition).

B-Splines

Elementary B-Spline of degree 0 is  $\beta^{(0)}(t) = 1$  for  $t \in [-0.5, 0.5)$  (0 o.w.) and of degree  $K$ ,  $\beta^{(K)} = \beta^{(K-1)} * \beta^{(0)}$  (supported on  $[-\frac{K+1}{2}, \frac{K+1}{2})$ ). Shifts of this are called B-splines of degree  $K$ . FT is  $B^{(K)}(\omega) = \text{sinc}^{K+1}(\omega/2)$ . For even  $K$ ,  $\beta^{(K)}$  smooth on  $(z - 0.5, z + 0.5)$  and if odd on  $(z, z + 1)$   $z \in \mathbb{Z}$ .

**Causal B-Splines:** Causal elementary B-Spline of degree 0 is  $\beta_+^{(0)}(t) = 1$  for  $t \in [0, 1)$  (0 o.w.) and of degree  $K$ ,  $\beta_+^{(K)} = \beta_+^{(K)}(t - 0.5(K + 1))$ .  $\beta_+^{(K)}$  is a generator of shift-invariant subspace  $S_{K, \mathbb{Z}}$ , in fact the one with shortest support.  $\overline{\text{span}}(\{\beta^{(K)}(t - k)\}_{k \in \mathbb{Z}}) = S_{K, \mathbb{Z}}$  for odd  $K$  and  $S_{K, \mathbb{Z} + 0.5}$  for even  $K$ .

**Differentiation:**  $\frac{d}{dt} x(t) = \sum_{k \in \mathbb{Z}} \alpha'_k \beta_+^{(K-1)}(t - k)$ ,  $\alpha'_k = \alpha_k - \alpha_{k-1}$ .

**Integration:**  $\int_{-\infty}^\tau x(\tau) d\tau = \sum_{k \in \mathbb{Z}} \alpha_k^{(1)} \beta_+^{(K+1)}(t - k)$ ,  $\alpha_k^{(1)} = \sum_{m=-\infty}^k \alpha_m$ .

**Canonical Dual Spline Basis:** For dual basis of  $\{\beta_+^{(1)}(t - k)\}_{k \in \mathbb{Z}}$ , need  $\langle \tilde{\beta}_+^{(1)}(t - i), \beta_+^{(1)}(t - k) \rangle = \delta_{i-k}$ . For canonical dual, need  $\tilde{\beta}_+^{(1)} \in S_{1, \mathbb{Z}}$ .  $\hat{x}(t) = \sum_{k \in \mathbb{Z}} \langle x(t), \tilde{\beta}_+^{(1)}(t - k) \rangle \beta_+^{(1)}(t - k)$ . Dual spline of degree 1 has infinite support but decays exponentially.

Polynomial reproduction (Strang-Fix theorem)

If  $\int_{-\infty}^\infty (1 + |t|^K) |\phi(t)| dt < \infty$  for  $k \in \mathbb{N}$  and  $\phi$  function with FT  $\Phi$ , following are equivalent:

- (i)  $p_K(t) = \sum_{k \in \mathbb{Z}} \alpha_k \phi(t - k)$  for  $p_K$  a polynomial of degree  $\leq K$ .
- (ii)  $\Phi$  and its first  $K$  derivatives satisfy  $\Phi(0) \neq 0$  and  $\Phi^{(k)}(2\phi l) = 0$  for  $k = 1, \dots, K$ ,  $l \in \mathbb{Z} \setminus \{0\}$ .

**Partition of unity (case of Strang-Fix):**  $\phi_1(t) = \sum_{n \in \mathbb{Z}} \phi(t - n) = 1$  (periodized version with period 1 of  $\phi \in \mathcal{L}^1(\mathbb{R})$ ) iff  $\Phi(2\phi k) = \delta_k$   $k \in \mathbb{Z}$ .

Series Truncation

Given orthonormal expansion in infinite Hilbert space  $x = \sum_{k \in \mathbb{Z}} c_k \phi_k$ ,  $c_k = \langle x, \phi_k \rangle$ . Cannot store all coefficients. **Linear approximation:** Retain coefficients with a priori fixed set of indices. **Don't depend on x. Linear. Non optimal in error.** **Nonlinear approximation:** Retain  $M$  largest coefficients in absolute value. **Depend on x. Non linear. Store index of coefficients. Optimal in error.** To proof latter select indices in  $\mathcal{I}_M$  and apply Parseval:  $\|x - \hat{x}\|^2 = \left\| \sum_{m \in \mathbb{Z} \setminus \mathcal{I}_M} c_m \phi_m \right\|^2 = \sum_{m \in \mathbb{Z} \setminus \mathcal{I}_M} |c_m|^2$ .

Uncertainty Principles

$\mu_t := \int_{-\infty}^\infty t \frac{|x(t)|^2}{\|x\|^2} dt$ ,  $\Delta_t^2 := \int_{-\infty}^\infty (t - \mu_t)^2 \frac{|x(t)|^2}{\|x\|^2} dt$ .  $\mu_f := \int_{-\infty}^\infty \omega \frac{|X(\omega)|^2}{2\pi \|x\|^2} d\omega$ ,

$\Delta_f^2 := \int_{-\infty}^\infty (\omega - \mu_t)^2 \frac{|X(\omega)|^2}{2\pi \|x\|^2} d\omega$ .

Heisenberg principle

For  $x \in \mathcal{L}^2(\mathbb{R})$   $\Delta_t^2 \Delta_f^2 \geq \frac{1}{4}$  with equality for Gaussian functions  $x(t) = \gamma e^{-\alpha t^2}$ ,  $\alpha > 0$ . To proof, suppose w.l.o.g.  $\mu_t = \mu_f = 0$ , use Cauchy-Schwarz, Parseval, integration by parts with  $(|x(t)|^2)$  and  $\lim_{t \rightarrow \infty} t x^2(t) = 0$  (since  $x \in \mathcal{L}^2(\mathbb{R})$ , i.e. decays faster than  $1/t$ ).

**Shifts and scalings:** Shifting changes  $\mu$  in domain of shift.  $\sqrt{a} x(at) \leftrightarrow X(\omega/a)/\sqrt{a}$  gives  $\mu_t/a, \Delta_t/a, a\mu_f, a\Delta_f$ .

Heisenberg principle for infinite sequences

Same definitions with discrete sum in time and DTFT (so integrals in  $[-\pi, \pi]$ ). For  $x \in \ell^2(\mathbb{Z})$  and  $X(e^{j\pi}) = 0$  (necessary extra condition),  $\Delta_n^2 \Delta_f^2 > \frac{1}{4}$  and lower bound cannot be achieved. **Shifts and scalings:** Shifting in frequency gives  $\mu_f + \omega_0$  if signal still  $X(e^{j\pi}) = 0$  ("signal not splitted"). Upsample and postfilter gives  $N\mu_n, N\Delta_n, \mu_f/N, \Delta_f/N$ . Prefilter and downsample  $\mu_n/N, \Delta_n/N, N\mu_f, N\Delta_f$  if  $x \in \text{BL}[-\pi/N, \pi/N]$ .

Uncertainty principle for finite-length sequences

$x_n \in \mathbb{C}^N$  and  $X_k$  its DFT with  $N_n$  and  $N_k$  nonzero coefficients respectively.  $N_n N_k \geq N$ . So cannot be sparse in both domains.