# Hashing

5 5

3×10+2×10+5×10

THE add 7 at the front.

3×10+2×10+5×10

Exappend 9 at the end  $x325 = xx10^3 + 3x10^7 + 2x10 + 5x10^9$ 

由 to avoid overflow use mod

Base and mod should be coppine

#### Use double hashing to avoid collison

base1 = 1827017

mod1 = 100888001

base2 = 5195977

mod2 = 133767331

```
\begin{aligned} \operatorname{hash}(s) &= s[0] + s[1] \cdot p + s[2] \cdot p^2 + \ldots + s[n-1] \cdot p^{n-1} \mod m \\ &= \sum_{i=0}^{n-1} s[i] \cdot p^i \mod m, \end{aligned}
```

```
long long compute_hash(string const& s) {
    const int p = 31;
    const int m = 1e9 + 9;
    long long hash_value = 0;
    long long p_pow = 1;
    for (char c : s) {
        hash_value = (hash_value + (c - 'a' + 1) * p_pow) % m;
        p_pow = (p_pow * p) % m;
    return hash_value;
```

### Find hash of a substring in O(1) time.

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$$H(f,i) = dH(f,0) - H(i-1,0)f/base^{i}$$

$$\frac{4}{3} \frac{3}{2} \frac{2}{9} \frac{1}{5} \frac{0}{4}$$

$$3 \frac{2}{9} \frac{9}{5} \frac{4}{9} \times \frac{9}{10} + \frac{9}{5} \times \frac{10}{10} + \frac{4}{10} \times \frac{10}{10}$$

#### Concatenate two strings

$$H = H_1 * boase^2 + H_2$$

$$S_1 = 25$$

$$H_1 = 2 \times 10^{1} + 5 \times 10^{1}$$

$$L_2 = 2$$

$$L_2 = 2$$

$$S_1 = 25$$
 $H_1 = 2 \times 10^{1} + 5 \times 10^{1}$ 
 $A_1 = 2$ 

$$S_2 = 3X$$
  
 $H_2 = 3X10 + 2X10$   
 $12 = 2$ 

```
struct Hash {
    11 \text{ vall} = 0, \text{ val2} = 0;
    11 \text{ base1} = 1827017, \text{ base2} = 5195977;
    11 \mod 1 = 1008888001, \mod 2 = 133767331;
    void push(int num) {
        vall *= basel; vall += num; vall %= modl;
        val2 *= base2; val2 += num; val2 %= mod2;
    bool operator < (const Hash& p) const {
        if (val1 == p.val1) return val2 < p.val2;
        return vall < p.vall;
    bool operator == (const Hash &p) const {
        return vall == p.vall && val2 == p.val2;
Hash x; // u can push many integers in x like x.push(tmp);
map<Hash,int>cnt; //can be used
```