

Multiple alignment, Phylogeny and Tree Building

aligning several sequences

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What's this lecture really about?

- ▶ Why align lots of sequences
- ▶ Types of homologous sequences
- ▶ How to solve a *difficult* problem
- ▶ Building trees from distances

Multiple Alignment. Why?

To some extent more natural than pairwise alignment since no reason to believe that similar sequences come in pairs.

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Paralogues Sequence groups that are homologous within species (i.e. several genes within a species that share an evolutionary origin).

There are usually lots of these. Hence multiple alignment is more 'natural' than simple pairwise alignment.

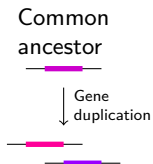
Homology, Orthology, Paralogy

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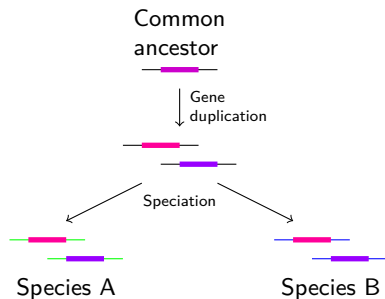
Common
ancestor



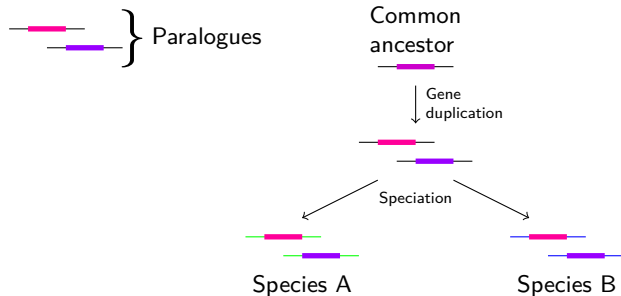
Homology, Orthology, Paralogy



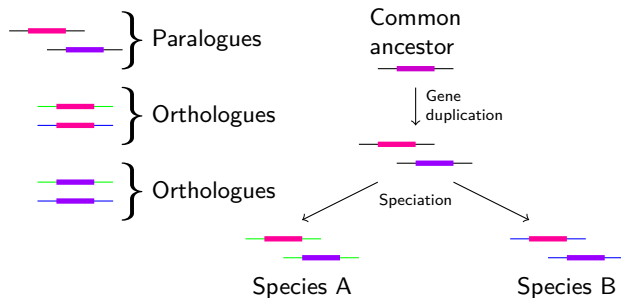
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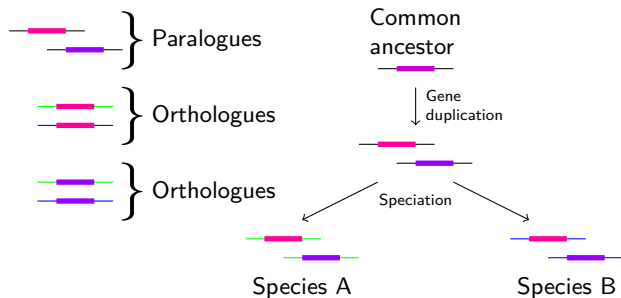
Homology, Orthology, Paralogy



Homology, Orthology, Paralogy



Homology, Orthology, Paralogy



All homologous to each other

Orthologues and Paralogues

Orthologues Arise from speciation

Paralogues Arise from gene duplication

- ▶ Paralogues can appear and disappear as a result of gene duplication followed by loss (esp. whole genome duplication).
- ▶ Gene duplication can be followed by functional specialisation of the paralogues
 - ▶ Change of regulatory environment (i.e. when the gene is expressed)
 - ▶ Change in coding other functional sequence

Multiple Alignment. More whys

Addresses many biological questions and technical issues:

- ▶ diagnostic patterns for protein families
- ▶ detect or demonstrate homology between sequences
- ▶ help predict secondary and tertiary structures
- ▶ to suggest oligonucleotide primers for PCR
- ▶ essential prelude to molecular evolutionary analysis (allows for ancestral state inference)
- ▶ ...

How to align many sequences?

- ▶ Complexity C of optimal alignment by dynamic programming

$$C = \prod_{i=1}^n l_i$$

where n = number of sequences and l_i = length of the i^{th} sequence.

- ▶ Requires too much:

- ▶ Memory
- ▶ computation (CPU cycles)

for more than a few sequences.

- ▶ *Heuristic* methods are used instead.

Do not guarantee an optimal result, but provide sufficient speed.

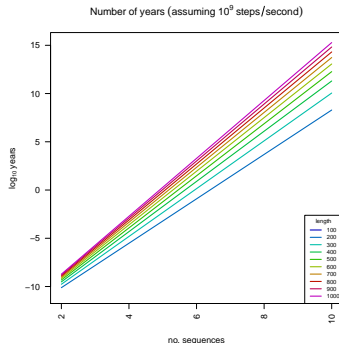
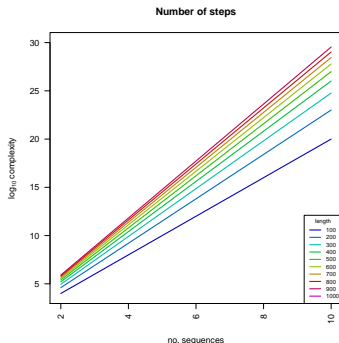
Many methods exist: we will look in detail at one of these.

Dynamic programming not possible!

Complexity (C) scales with length (l) and number (n) of sequences:

$$C = l^n$$

For $l = 100$ and $n = 2, 3$, this is 10000 and 1000000 steps respectively.



This is a huge underestimate as the complexity of each step scales with $2^n - 1$ steps

CLUSTAL W: improving the sensitivity of progressive multiple sequence alignment through sequence weighting, position-specific gap penalties and weight matrix choice

Julie D.Thompson, Desmond G.Higgins⁺ and Toby J.Gibson^{*}

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Received July 12, 1994; Revised and Accepted September 23, 1994

ABSTRACT

The sensitivity of the commonly used progressive multiple sequence alignment method has been greatly improved for the alignment of divergent protein sequences. Firstly, individual weights are assigned to each sequence in a partial alignment in order to down-weight near-duplicate sequences and up-weight the most divergent ones. Secondly, amino acid substitution matrices are varied at different alignment stages according to the divergence of the sequences to be aligned. Thirdly, residue-specific gap penalties and locally reduced gap penalties in hydrophilic regions encourage new gaps in potential loop regions rather than regular secondary structure. Fourthly, positions in early alignments where gaps have been opened receive locally reduced gap penalties to encourage the opening up of new gaps at these positions. These modifications are incorporated into a new program, CLUSTAL W which is freely available.

practical. The new methods are made available in a program called CLUSTAL W, which is freely available and portable to a wide variety of computers and operating systems.

In order to align just two sequences, it is standard practice to use dynamic programming (2). This guarantees a mathematically optimal alignment, given a table of scores for matches and mismatches between all amino acids or nucleotides [e.g. the PAM250 matrix (3) or BLOSUM62 matrix (4)] and penalties for insertions or deletions of different lengths. Attempts at generalising dynamic programming to multiple alignments are limited to small numbers of short sequences (5). For much more than eight or so proteins of average length, the problem is uncomputable given current computer power. Therefore, all of the methods capable of handling larger problems in practical timescales make use of heuristics. Currently, the most widely used approach is to exploit the fact that homologous sequences are evolutionarily related. One can build up a multiple alignment progressively by a series of pairwise alignments, following the

Why ClustalW

- ▶ One of the most widely used methods
- ▶ Easy to understand
- ▶ Includes phylogenetic analysis
- ▶ Paper describes the derivation and reasoning for the heuristics used nicely (eg. this can be a problem, so we tweaked this part of the method to give nicer results).
- ▶ The method extends naturally from pairwise alignment.

The clustal method

For a collection of sequences:

1. Align all pairs of sequences and calculate a distance matrix (table).
2. Use the distance matrix to calculate a guide tree.
3. Align the sequences progressively according to the branch order of the guide tree.

Pairwise alignment

- ▶ Global alignment of all pairs using a modification of Needleman-Wunsch, or a faster k-tuple based heuristic method.
- ▶ Scores are calculated as: number of identities / number of residues compared (gap positions are excluded).
- ▶ Distances are simply $(1 - \text{score})$

This gives an n by n distance matrix which is then used to make a guiding tree.

Pairwise alignment

Alignments

Seq_0

Seq_0

Seq_1

Seq_1

Seq_2

Seq_2

Seq_3

Seq_3

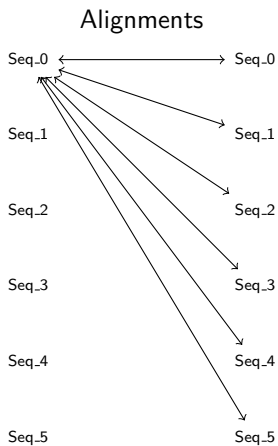
Seq_4

Seq_4

Seq_5

Seq_5

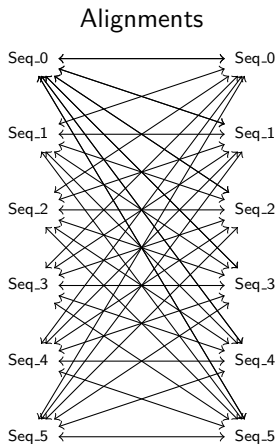
Pairwise alignment



Distance table

	S_0	S_1	S_2	S_3	S_4	S_5
S_0	$d_{0,0}$	$d_{0,1}$	$d_{0,2}$	$d_{0,3}$	$d_{0,4}$	$d_{0,5}$
S_1						
S_2						
S_3						
S_4						
S_5						

Pairwise alignment



Distance table

	S_0	S_1	S_2	S_3	S_4	S_5
S_0	$d_{0,0}$	$d_{0,1}$	$d_{0,2}$	$d_{0,3}$	$d_{0,4}$	$d_{0,5}$
S_1	$d_{1,0}$	$d_{1,1}$	$d_{1,2}$	$d_{1,3}$	$d_{1,4}$	$d_{1,5}$
S_2	$d_{2,0}$	$d_{2,1}$	$d_{2,2}$	$d_{2,3}$	$d_{2,4}$	$d_{2,5}$
S_3	$d_{3,0}$	$d_{3,1}$	$d_{3,2}$	$d_{3,3}$	$d_{3,4}$	$d_{3,5}$
S_4	$d_{4,0}$	$d_{4,1}$	$d_{4,2}$	$d_{4,3}$	$d_{4,4}$	$d_{4,5}$
S_5	$d_{5,0}$	$d_{5,1}$	$d_{5,2}$	$d_{5,3}$	$d_{5,4}$	$d_{5,5}$

The guide tree

- ▶ Tree created from the distances to represent the similarities between the sequences and to suggest an order for the progressive alignment.
- ▶ Earlier versions used UPGMA. Newer version uses Neighbor joining algorithm.

What is a tree

- ▶ A way to represent a set of relationships (commonly distances or dis-similarities).
- ▶ Often obtained by hierarchical clustering methods from distances matrices (see below).
- ▶ Developed to represent evolutionary relationships (i.e. phylogenetic trees).
- ▶ Can be *evaluated* by maximum parsimony and likelihood methods.
- ▶ Can summarise N-dimensional data sets in general (eg. gene expression data)

A phylogenetic tree represents a *hypothesis* about how a set of species or sequences evolved.

UPGMA: the simplest tree

Unweighted Pair Group Method with Arithmetic Mean

Distances

	A	B	C	D	E
B	2				
C	4	4			
D	6	6	6		
E	6	6	6	4	
F	8	8	8	8	8

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→
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Calculate distances between merged node (AB) and other nodes (using arithmetic mean)

$$\text{dist}(AB, C) \leftarrow (\text{dist}(A, C) + \text{dist}(B, C)) / 2$$

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1	A
1	B

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F	8	8	8	8

repeat

2	D
2	E

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Unweighted Pair Group Method with Arithmetic Mean

Distances

	A	B	C	D	E
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join two
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1	B

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	AB	C	D	E
C	4			
D	6	6		
E	6	6	4	
F	8	8	8	8

repeat

2	D
2	E

	AB	C	DE
C	4		
DE	6	6	
F	8	8	8

UPGMA: the simplest tree

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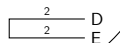


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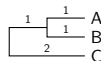
	AB	C	D	E
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repeat



	AB	C	DE
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and again



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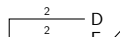


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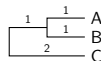
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repeat

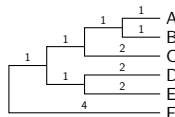


	AB	C	DE
C	4		
DE	6	6	
F	8	8	8

and again



eventually



Neighbor joining algorithm

- ▶ Underlying algorithm method similar to UPGMA (i.e. progressively merge neighboring nodes until a single tree is obtained).
- ▶ Modified distance matrix used to find nearest nodes to join.
- ▶ Distances of pair members to joins are influenced by distances to external nodes.
- ▶ Does not assume equal rate of evolution
⇒ neighbours have differing distances to their joining nodes.
- ▶ Better than UPGMA (?)

Neighbor joining (1)

Nodes to be joined (i.e. neighbors) are chosen from a Q matrix:

$$Q_{i,j} = (n - 2)d_{i,j} - \sum_{k=1}^n d_{i,k} - \sum_{k=1}^n d_{j,k}$$

$d_{i,j}$ distance between nodes i and j

n the number of nodes

Q is *only* used to rank node pairs

Outlier pairs have low Q and are joined first (i.e. pairs of nodes which are distant from the larger set).

It's OK to be confused here

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From 'Neighbour-Joining Revealed', published almost 20 years after Saitou and Nei:

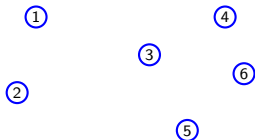
'Yet the question "what does the NJ method seek to do?" has until recently proved somewhat elusive, leading to some imprecise claims and misunderstanding'

eh??

$$Q_{i,j} = (n-2)d_{i,j} - \sum_{k=1}^n d_{i,k} - \sum_{k=1}^n d_{j,k}$$

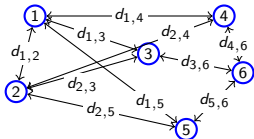
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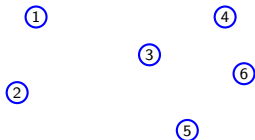
$$Q_{i,j} = (n-2)d_{i,j} - \sum_{k=1}^n d_{i,k} - \sum_{k=1}^n d_{j,k}$$



$d_{i,j}$ is the distance between nodes i and j

eh??

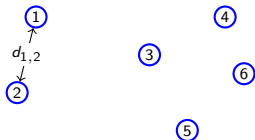
$$Q_{i,j} = (n-2)d_{i,j} - \sum_{k=1}^n d_{i,k} - \sum_{k=1}^n d_{j,k}$$



$$Q_{1,2} = (n-2)d_{1,2} - \sum_{k=1}^n d_{1,k} - \sum_{k=1}^n d_{2,k}$$

eh??

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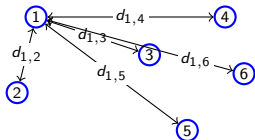


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$$(6-2) \times d_{1,2}$$

eh??

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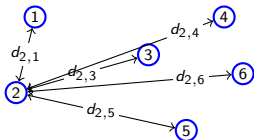
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$$-(d_{1,1} + d_{1,2} + d_{1,3} + d_{1,4} + d_{1,5} + d_{1,6})$$

eh??

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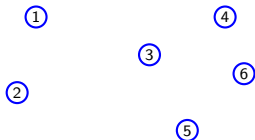
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$$(6-2) \times d_{1,2}$$

$$-(d_{1,1} + d_{1,2} + d_{1,3} + d_{1,4} + d_{1,5} + d_{1,6})$$

$$-(d_{2,1} + d_{2,2} + d_{2,3} + d_{2,4} + d_{2,5} + d_{2,6})$$

This favours nodes close to each other, but which are far from the others; i.e. pairs of outliers.

Neighbor joining (2)

The pair of nodes (f , g) with the lowest Q value are joined through a new node u .

The distances between the new node u and f and g are:

$$\delta_{f,u} = \frac{1}{2}d_{f,g} + \frac{1}{2(n-2)} \left[\sum_{k=1}^n d_{f,k} - \sum_{k=1}^n d_{g,k} \right] \quad (1)$$

$$\delta_{g,u} = d_{f,g} - \delta_{f,u}$$


$\delta_{f,u}$ and $\delta_{g,u}$ are adjusted such that they are proportional to their respective distances to the remaining nodes.

ehh again?

$$\delta_{f,u} = \frac{1}{2}d_{f,g} + \frac{1}{2(n-2)} \left[\sum_{k=1}^n d_{f,k} - \sum_{k=1}^n d_{g,k} \right] \quad (2)$$

$$\delta_{g,u} = d_{f,g} - \delta_{f,u}$$

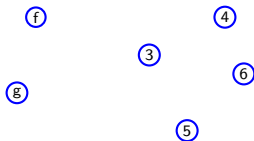
Here f and g are equivalent to 1 and 2

Note that, $d_{f,1}$ and $d_{g,2}$ are 0 and that $d_{f,2} = d_{g,1}$ and so cancel each other. 

ehh again?

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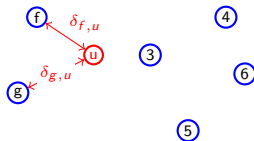
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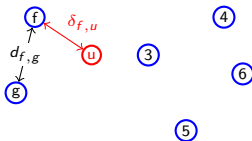
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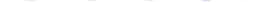
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$$\delta_{f,u} = \frac{1}{2} d_{f,g} +$$

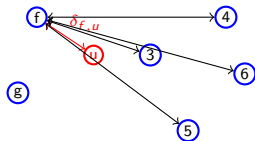
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Note that, $d_{f,1}$ and $d_{g,2}$ are 0 and that $d_{f,2} = d_{g,1}$ and so cancel each other. 

ehh again?

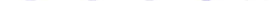
$$\delta_{f,u} = \frac{1}{2}d_{f,g} + \frac{1}{2(n-2)} \left[\sum_{k=1}^n d_{f,k} - \sum_{k=1}^n d_{g,k} \right] \quad (2)$$

$$\delta_{g,u} = d_{f,g} - \delta_{f,u}$$



$$\delta_{f,u} = \frac{1}{2}d_{f,g} + \frac{1}{2 \times (6-2)} \times ((d_{f,1} + d_{f,2} + d_{f,3} + d_{f,4} + d_{f,5} + d_{f,6}))$$

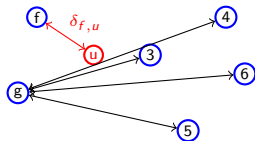
Here f and g are equivalent to 1 and 2

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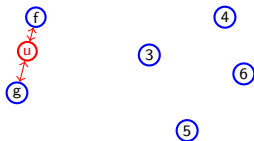
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Note that, $d_{f,1}$ and $d_{g,2}$ are 0 and that $d_{f,2} = d_{g,1}$ and so cancel each other.

Neighbor joining (3)

The distances of the remaining nodes to the joining node u are set as:

$$\delta_{u,k} = \frac{1}{2}[d_{f,k} + d_{g,k} - d_{f,g}]$$

u joining node

k a remaining node

f, g the joined nodes

This assures that the total distance within the tree is consistent.

should be: $k \neq f, g$

These equations taken from the Wikipedia page:

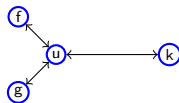
https://en.wikipedia.org/wiki/Neighbor_joining

They differ slightly in their form, but not substance from the original publication:

Neighbor joining (3)

The total distance of the tree is consistent:

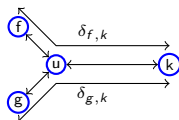
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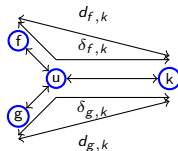
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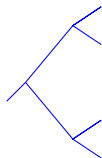


$$d_{f,k} + d_{g,k} = \delta_{f,k} + \delta_{g,k}$$

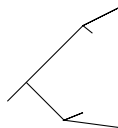
Neighbor joining: putting it together

1. Determine the Q matrix based on the current distance matrix.
2. Find the pair of nodes with the smallest Q value.
3. Create a new node that connects this pair.
4. Determine the distances of all the nodes to this new joining node.
5. Replace the neighbour pair with the new node and update the distance matrix.
6. Repeat from (1) until the tree is fully connected.

UPGMA vs Neighbor joining



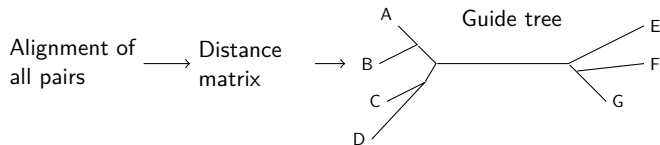
UPGMA



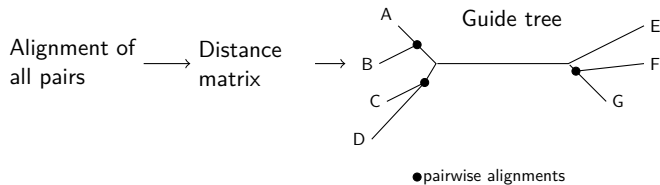
Neighbor Joining

Neighbor joining does not assume equal rate of evolution when joining nodes.

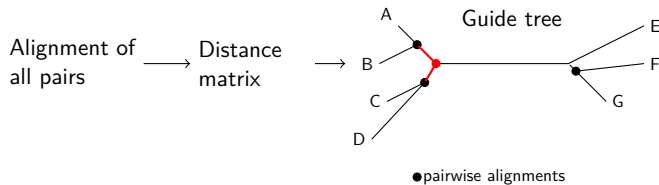
Progressive alignment



Progressive alignment

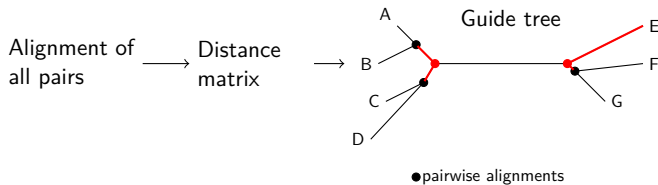


Progressive alignment



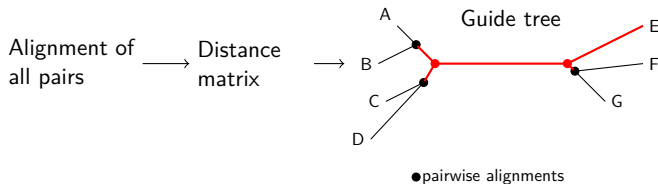
1. Align AB to CD

Progressive alignment



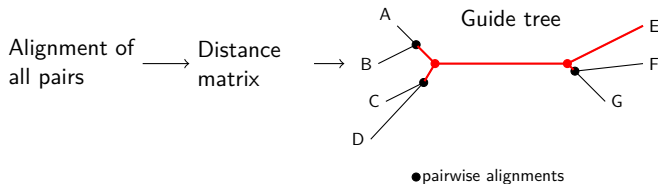
1. Align AB to CD
2. Align E to FG

Progressive alignment



1. Align AB to CD
2. Align E to FG
2. Align ABCD to EFG

Progressive alignment



1. Align AB to CD
2. Align E to FG
2. Align ABCD to EFG

How to align two alignments?

Aligning alignments

Modify the scoring function to use several sequences.

Match score is set to the mean of all independent pairs:

1 peeksavtal

2 geekaavlal

3 padktnvkaa

4 aadktnvkaa

5 egewqlvlhv

6 aaektkirsa

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```

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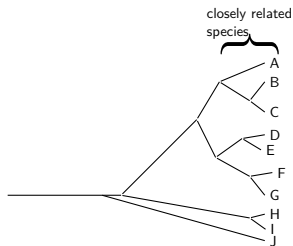
1	peeksavtal	Score	=	$M(t,v)$
2	geekaavllal		+	$M(t,i)$
3	padktnvkaa		+	$M(l,v)$
4	aadktnvkaa		+	$M(l,i)$
			+	$M(k,v)$
			+	$M(k,i)$
5	egewqlvlhv		+	$M(k,v)$
6	aaektkirsa		+	$M(k,i)$
				<hr/>
				8

Where $M(i,j)$ is are values taken from the given substitution matrix

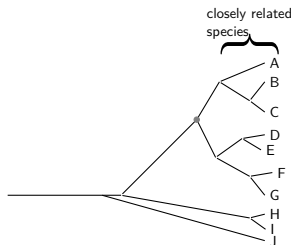
Refinements

- ▶ Weighting of sequences to correct for unequal sampling across evolutionary distances in the data set (greater weight to outlier sequences)
- ▶ Dynamic variation of gap penalties (to mimic known tendencies in proteins)
 - ▶ Increase gap opening penalty within 8 amino acid of a gap opening
 - ▶ Decrease gap opening penalty in hydrophilic stretches (associated with loops)
 - ▶ Decreased gap opening penalties at positions of gaps in early alignments.
- ▶ Dynamic use of substitution matrices: starting with substitution matrices suitable for closely related sequences and moving to divergent matrices.

Why weight sequences

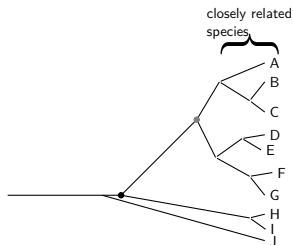


Why weight sequences



- A. GKSWKALTPP
- B. GKSWKSLSPS
- C. GKSWKSLSTS
- D. GKSWKSLSPS
- E. GKSWKSLSPS
- F. GKSWRALSPS
- G. GRSWKSLSPS

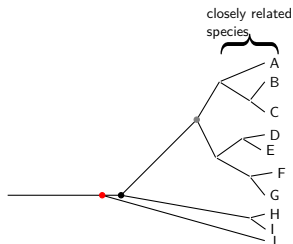
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E. GKSWKSLSPS
F. GKSWRALSPS
G. GRSWKSLSPS
H. PKSWRASSPS
I. PRSWKSSSPS

$$S = \frac{7 \times M_{P,G} + 7 \times M_{P,G}}{14}$$

Why weight sequences



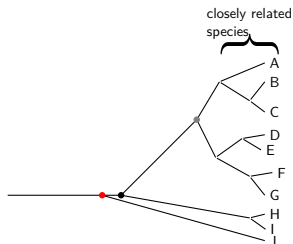
$$S = \frac{7 \times M_{P,G} + 2 \times M_{P,P}}{9}$$

A. GKSWKALTPP
 B. GKSWKSLSPS
 C. GKSWKSLSTS
 D. GKSWKSLSPS
 E. GKSWKSLSPS
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 H. PKSWRASSPS
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[A – G] and [H – I] represent single lineages. But:
 S Dominated by $M_{P,G}$

Why weight sequences



$$S = \frac{7 \times M_{P,G} + 2 \times M_{P,P}}{9}$$

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[A – G] and [H – I] represent single lineages. But:
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May also wish to weight by branch length
 (leads to maximum likelihood)

Problems?

- ▶ All alignments are global alignments and it may be necessary to trim sequences to give reasonable alignments.
- ▶ The guide tree is based on a matrix of distances of separately aligned sequences and may not be reliable. This may lead to mistakes early in the merging process that cannot be corrected later.

works ok!

ASV_v5RC	tfvallydyearie	tdlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
RSV_v5RC	tfvallydyearie	tdlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_cSRCL	tfvallydyearie	tdlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
X1_cSRCL	tfvallydyearie	tdlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
ASV_v5RC	tfvallydyearie	tdlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
X1_cSRCL	tfvallydyearie	tdlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
ASV_YYES	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
C_YYES	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_cYEB1	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
X1_YYES	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
X1_YYES	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_cYFN	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
M_cPGR	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_cPGR	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
HA_STK	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_cYFN	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
M_HCK	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_LYN	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
M_BLK	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
M_LSKT	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
SV_YABL	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
Dm_ABL1	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
C_cTKL	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
Ce_ses5/1	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
Ces_ses5/2	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
Dm_SRC2	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
SV_GAOCRK	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
C_SPCA	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
Dm_SPCA	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
Dm_SPCB	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
B_C2C	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
B_PLIC1	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
B_PLIC2	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_PLIC1	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
I_RASA/GAP	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
Ac_MITC	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_H81	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_VAV	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
Dm_SRC2	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/1	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/2	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/3	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/4	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/5	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/6	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/7	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/8	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/9	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/10	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/11	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/12	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/13	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/14	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/15	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/16	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/17	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/18	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/19	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1/20	tfvallydyearie	ddlrfk	qgqrclvnmv	ggdwlahahltg	qtrvpanyapad
H_NCK1					

Other methods

Multiple Sequence Alignment

Tools > Multiple Sequence Alignment

Multiple Sequence Alignment (MSA) is generally the alignment of three or more biological sequences (protein or nucleic acid) of similar length. From the output, homology can be inferred and the evolutionary relationships between the sequences studied.

By contrast, Pairwise Sequence Alignment tools are used to identify regions of similarity that may indicate functional, structural and/or evolutionary relationships between two biological sequences.

Clustal Omega

New MSA tool that uses seeded guide trees and HMM profile-profile techniques to generate alignments. Suitable for medium-large alignments.

[Launch Clustal Omega](#)

EMBOSS Cons

EMBOSS Cons creates a consensus sequence from a protein or nucleotide multiple alignment.

[Launch EMBOSS Cons](#)

Kalign

Very fast MSA tool that concentrates on local regions. Suitable for large alignments.

[Launch Kalign](#)

MAFFT

MSA tool that uses Fast Fourier Transforms. Suitable for medium-large alignments.

[Launch MAFFT](#)

MUSCLE

Accurate MSA tool, especially good with proteins. Suitable for medium alignments.

[Launch MUSCLE](#)

MView

Transforms a Sequence Similarity Search result into a Multiple Sequence Alignment or reformat a Multiple Sequence Alignment using the MView program.

[Launch MView](#)

T-Coffee

Consistency-based MSA tool that attempts to mitigate the pitfalls of progressive alignment methods. Suitable for small alignments.

[Launch T-Coffee](#)

WebPRANK

The EBI has a new phylogeny-aware multiple sequence alignment program which makes use of evolutionary information to help place insertions and deletions. Try it out at [WebPRANK](#).

- ▶ Clustal Omega
- ▶ EMBOSS Cons
- ▶ Kalign
- ▶ MAFFT
- ▶ MUSCLE
- ▶ MView
- ▶ T-Coffee
- ▶ WebPRANK

<https://www.ebi.ac.uk/Tools/msa/>

Other methods

- ▶ DCA. Semi-exhaustive, divide and conquer algorithm. Breaks sequences into segments based on local similarity. Segments are aligned by dynamic programming and then joined.
bibiserv.techfak.uni-bielefeld.de/dca?id=dca_view_webservice
- ▶ Poa (Partial order alignments). Uses a graph representation of the multiple alignment that can be aligned by dynamic programming.
bioinformatics.oxfordjournals.org/content/18/3/452.short
- ▶ Dialign (and Dialign2, Dialign-TX). Compares segments of sequences rather than individual residues without using gap penalties. Good for comparing sequences with only local similarity.
mobyle.pasteur.fr/cgi-bin/portal.py?#forms::dialign
- ▶ Can also be achieved by probabilistic methods (eg. Hidden Markov Models), though these are normally used to model (describe) sequences rather than to perform the actual alignment. But see hmmbuild:
hmmer.janelia.org/

Many more available. Method usually determined by the sheep algorithm (i.e. use what others in the field are using).