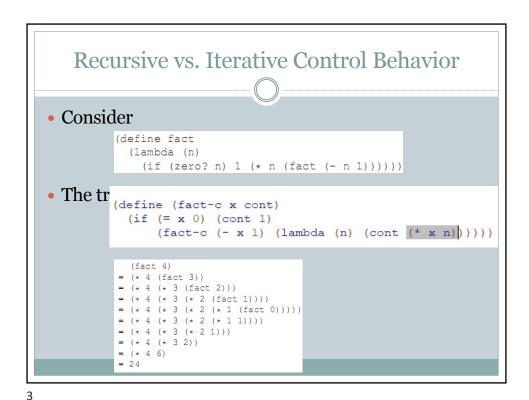
Lecture 26 Continuation Passing Style

T. METIN SEZGIN

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Learning outcomes of this lecture

- A student attending this lecture should be able to:
 - 1. Understand the difference between tail-recursion & recursion
 - 2. Understand the concept of continuation
 - 3. Implement simple procedures using continuations



Recursive vs. Iterative Control Behavior Consider (define fact-iter (lambda (n) (fact-iter-acc n 1))) (define fact-iter-acc (lambda (n a) (if (zero? n) a (fact-iter-acc (- n 1) (* n a))))) • The trace (fact-iter 4) = (fact-iter-acc 4 1) = (fact-iter-acc 3 4) = (fact-iter-acc 2 12) = (fact-iter-acc 1 24) = (fact-iter-acc 0 24) = 24

Alternative implementation

Consider

The trace

```
> (fact 4)
>(fact-c 4 #<procedure:...33/fact-cont.rkt:6:27>)
>(fact-c 3 #<procedure:...33/fact-cont.rkt:11:22>)
>(fact-c 2 #<procedure:...33/fact-cont.rkt:11:22>)
>(fact-c 1 #<procedure:...33/fact-cont.rkt:11:22>)
>(fact-c 0 #<procedure:...33/fact-cont.rkt:11:22>)
<24
24</pre>
```

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What is the key difference between the recursive and tail recursive versions?



- What do we do after each call?
- How does the control context grow?
- Continuation:
 - Captures the control context
 - O Describes what needs to be done next!

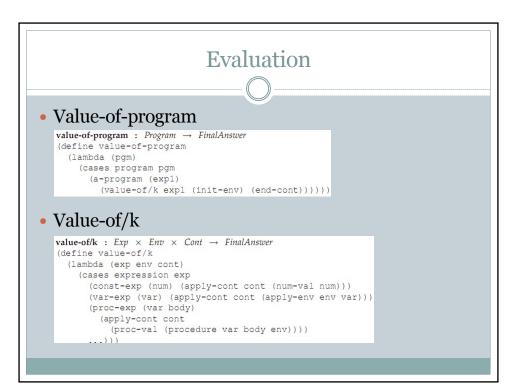
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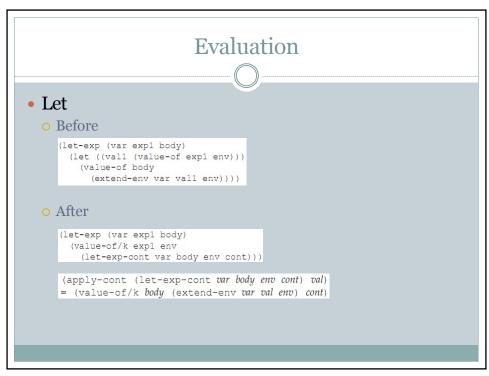
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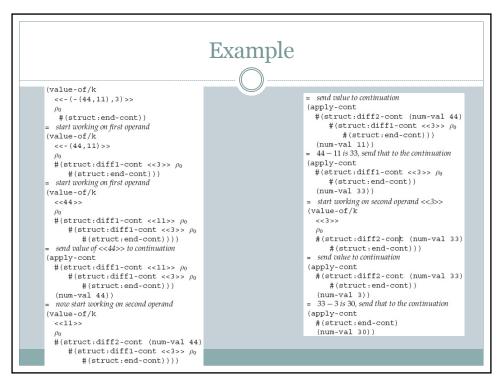
A CPS Interpreter

- The environment grows as we evaluate expressions
- Now we need to keep around a list of things to do after the evaluation of each expression.
- Introduce apply-cont
 - o Example:





```
Example
(value-of/k << letrec p(x) = x in if b then 3 else 4>>
 \rho_0 \ cont_0)
= letting \rho_1 be (extend-env-rec ... \rho_0)
(value-of/k <<if b then 3 else 4>> \rho_1 cont<sub>0</sub>)
= next, evaluate the test expression
(value-of/k <<br/>b>> \rho_1 (test-cont <<3>> <<4>> \rho_1 cont_0))
= send the value of b to the continuation
(apply-cont (test-cont <<3>> <<4>> \rho_1 \ cont_0)
               (bool-val #t))
= evaluate the then-expression
(value-of/k <<3>> \rho_1 \ cont_0)
= send the value of the expression to the continuation
(apply-cont cont_0 (num-val 3))
= invoke the final continuation with the final answer
(begin (eopl:printf ...) (num-val 3))
```



```
Evaluation
• Procedure application
       (call-exp (rator rand)
         (let ((proc1 (expval->proc (value-of rator env)))
                (val (value-of rand env)))
            (apply-procedure proc1 val)))
   After
        (call-exp (rator rand)
  (value-of/k rator env
        (apply-cont (rator-cont rand env cont) val1) (apply-cont (rand-cont val1 cont) val2)
         (value-of/k rand env
                                                         = (let ((proc1 (expval->proc val1)))
                                                             (apply-procedure/k proc1 val2 cont))
             (rand-cont val1 cont))
                      apply-procedure/k : Proc \times ExpVal \times Cont \rightarrow FinalAnswer
                       (define apply-procedure/k
                         (lambda (procl val cont)
  (cases proc procl
    (procedure (var body saved-env))
                                (value-of/k body
                                  (extend-env var val saved-env)
                                  cont)))))
```

