

1 . I created the hash table with linear probing.

```
50 class HashTable
51 {
52 public:
53     explicit HashTable(int size = 10000 ); // constructor
54     returnType contains( const int & x ) const; // tries to find and returns probe number
55
56     void makeEmpty( ); // make the table empty
57     returnType insert( const int & x ); // inserting new number to the table
58     returnType remove( const int & x ); // removing from table
59     int getCurrentSize(){ // returns currentSize
60         return currentSize;
61     };
62     int getTotalSize(){ // returns totalSize
63         return totalSize;
64     };
65     enum EntryType { ACTIVE, EMPTY, DELETED }; // types to check if its active, empty or deleted
66
67 private:
68     struct HashEntry
69     {
70         int element;
71         EntryType info;
72
73         HashEntry( const int & e = int( ), EntryType i = EMPTY )
74             : element( e ), info( i ) { }
75     };
76
77     vector<HashEntry> array;
78     int currentSize;
79     int totalSize;
80
81     bool isActive( int currentPos ) const{ // returns if currentPos is active or not
82         return array[currentPos].info == ACTIVE;
83     };
84
85     int findPos( const int & x, int & probe ) const; // finds position of x and returns it, increments probe count too.
86     int myhash(const int & x) const{ // x mod M
87         return x % totalSize;
88     };
89 };
```

2. I created a function for 3 combination(6-1-1, 4-2-2, 2-1-5)

```
int chooseSelect(const int & number , const int & combination){ // 0 for 6-1-1 ~~ 1 for 4-2-2 ~~ 2 for 2-1-5
    if(combination == 0){
        if(number < 6)
            return 0; // insert
        else if(number < 7)
            return 1; // delete
        else
            return 2; // find
    }
    else if(combination == 1){
        if(number < 4)
            return 0; // insert
        else if(number < 6)
            return 1; // delete
        else
            return 2; // find
    }else{
        if(number < 2)
            return 0; // insert
        else if(number < 3)
            return 1; // delete
        else
            return 2; // find
    }
}
```

3. I created 6 [size+1] arrays then fill them up with # of probes and # of transactions.

```
element* insertionSuccess = new element[size+1];
element* insertionFailure = new element[size+1];
element* removeSuccess = new element[size+1];
element* removeFailure = new element[size+1];
element* findSuccess = new element[size+1];
element* findFailure = new element[size+1];
```

Then computed the averages of the transactions after that pushed that information to the 2D arrays.

```
if(insertionSuccess[j].transactions != 0)
    insertionSuccess2D[i][j] = insertionSuccess[j].probe / insertionSuccess[j].transactions;
```

4. I limited the range of the integers to 10M so that most integers have a decent chance of being deleted or searched. I generated such integers by generating a random number and taking mod 10M.

```
int toTransaction = rand() % (size * 10);
```

5. I generated transactions until the either the table becomes full or I have a total of 1,000,000 transactions.

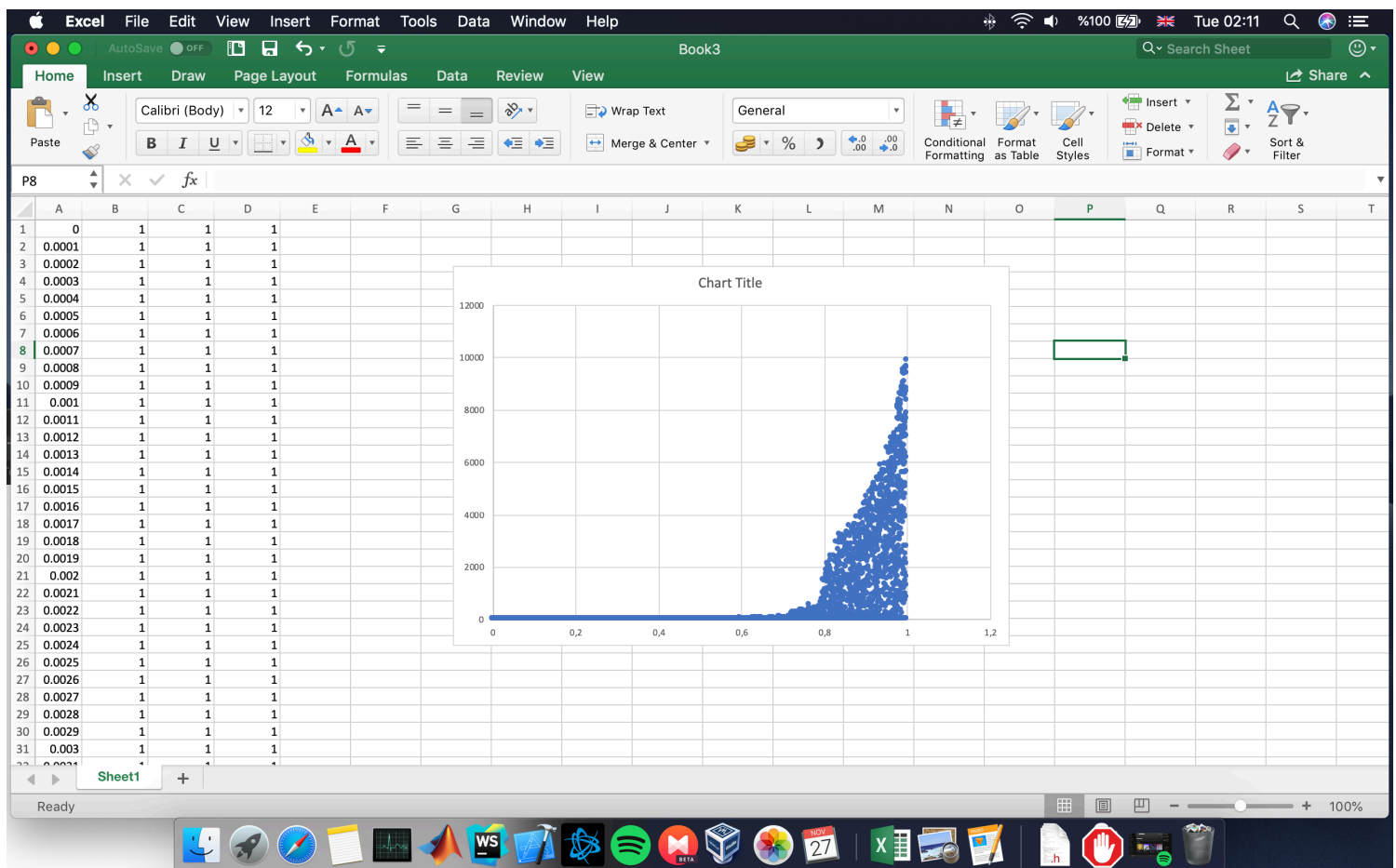
```
int transactions = 1000000;
while(transactions && table.getCurrentSize() < size){
```

6. I put comma after each printing to file and i saved file as .csv so excel opens it perfectly and used setprecision(4) to get some observable values.

```
insertSuccessOS.open("insertSuccess.csv");
insertFailureOS.open("insertFailure.csv");
removeSuccessOS.open("removeSuccess.csv");
removeFailureOS.open("removeFailure.csv");
findSuccessOS.open("findSuccess.csv");
findFailureOS.open("findFailure.csv");

insertSuccessOS << fixed << setprecision(4);
insertFailureOS << fixed << setprecision(4);
removeSuccessOS << fixed << setprecision(4);
removeFailureOS << fixed << setprecision(4);
findSuccessOS << fixed << setprecision(4);
findFailureOS << fixed << setprecision(4);
```

7. For example successful insertion excel file looks like this.



8. My hash function is $x \bmod M$

```
int myhash(const int & x) const{ // x mod M
    return x % totalSize;
};
```

9. I took all the functions from book and changed couple of things.

=> findPos

```
int HashTable::findPos(const int & x, int & probe ) const{
    int pos = myhash(x);
    while(array[pos].info == ACTIVE && array[pos].element != x){
        pos++;
        probe++;

        if( pos >= array.size( ) )
            pos -= array.size( );
    }
    return pos;
}
```

=> insert

```
124
125 returnType HashTable::insert( const int & x ){
126     int probe = 1;
127     int currentPos = findPos( x , probe );
128     if(isActive( currentPos ) )
129         return returnType(probe,false);
130
131     array[ currentPos ].element = x;
132     array[ currentPos ].info = ACTIVE;
133     currentSize++;
134
135     return returnType(probe,true);
136 }
137
```

=> remove

```
137
138  returnType HashTable::remove(const int & x){
139
140      int probe = 1;
141      int currentPos = findPos( x, probe );
142      if(!isActive(currentPos))
143          return returnType(probe, false);
144
145      array[currentPos].info = DELETED;
146      currentSize--;
147      return returnType(probe, true);
148  }
149
```

=> contains

```
116
117 returnType HashTable::contains( const int & x) const{
118     int probe = 1;
119     if(isActive( findPos( x, probe ) ))
120         return returnType(probe,true);
121     else
122         return returnType(probe,false);
123 }
124
```

=> To choose prime size I implemented some prime functions

```
bool isPrime ( int x ){ // returns if x is a prime number
    for(int i = 2; i <= x / 2; ++i)
    {
        if(x % i == 0)
        {
            return false;
        }
    }
    return true;
}

int nextPrime( int n ) // finds smallest next prime number
{
    if( n % 2 == 0 )
        n++;

    for( ; !isPrime( n ); n += 2 )
        ;

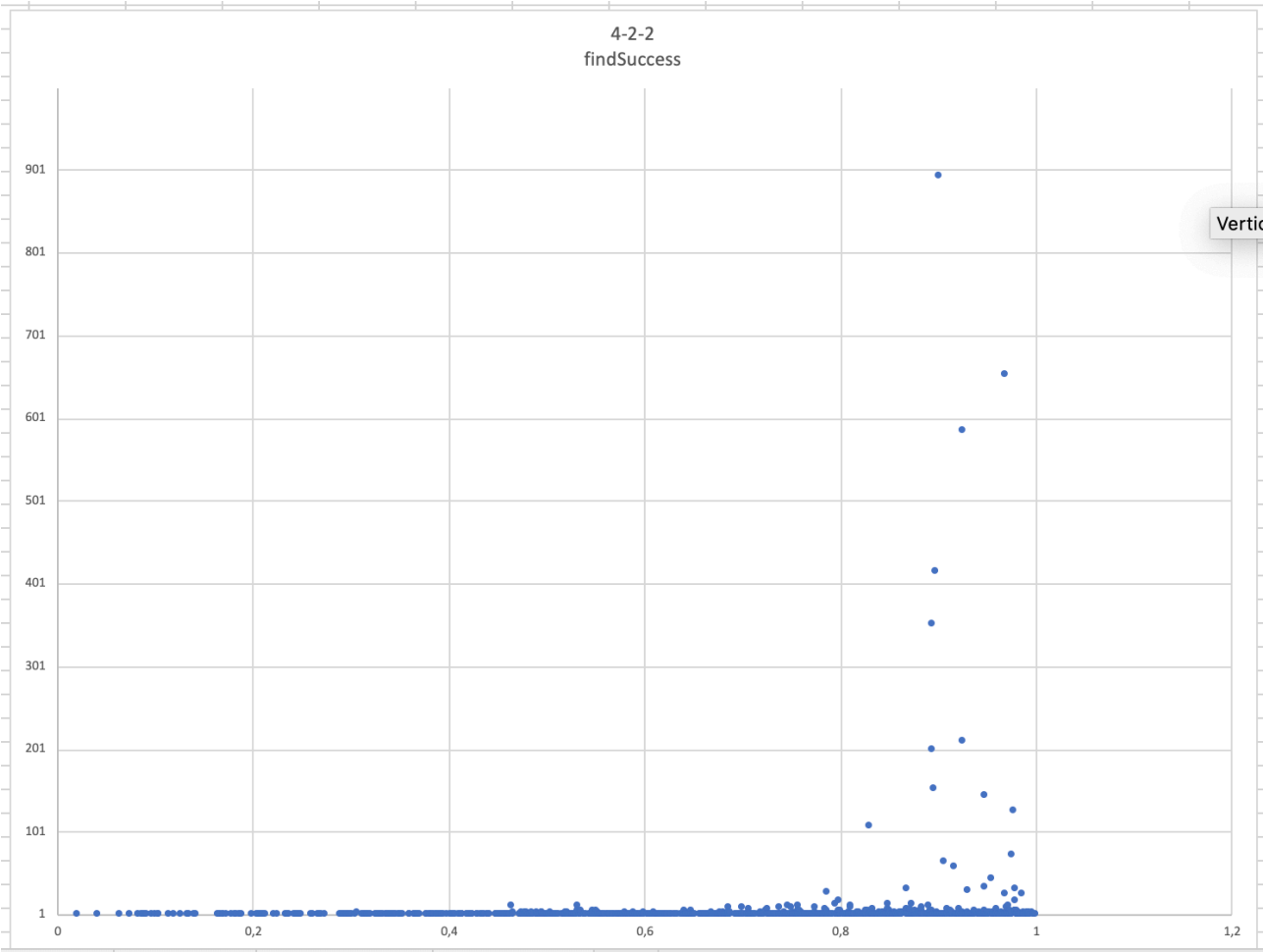
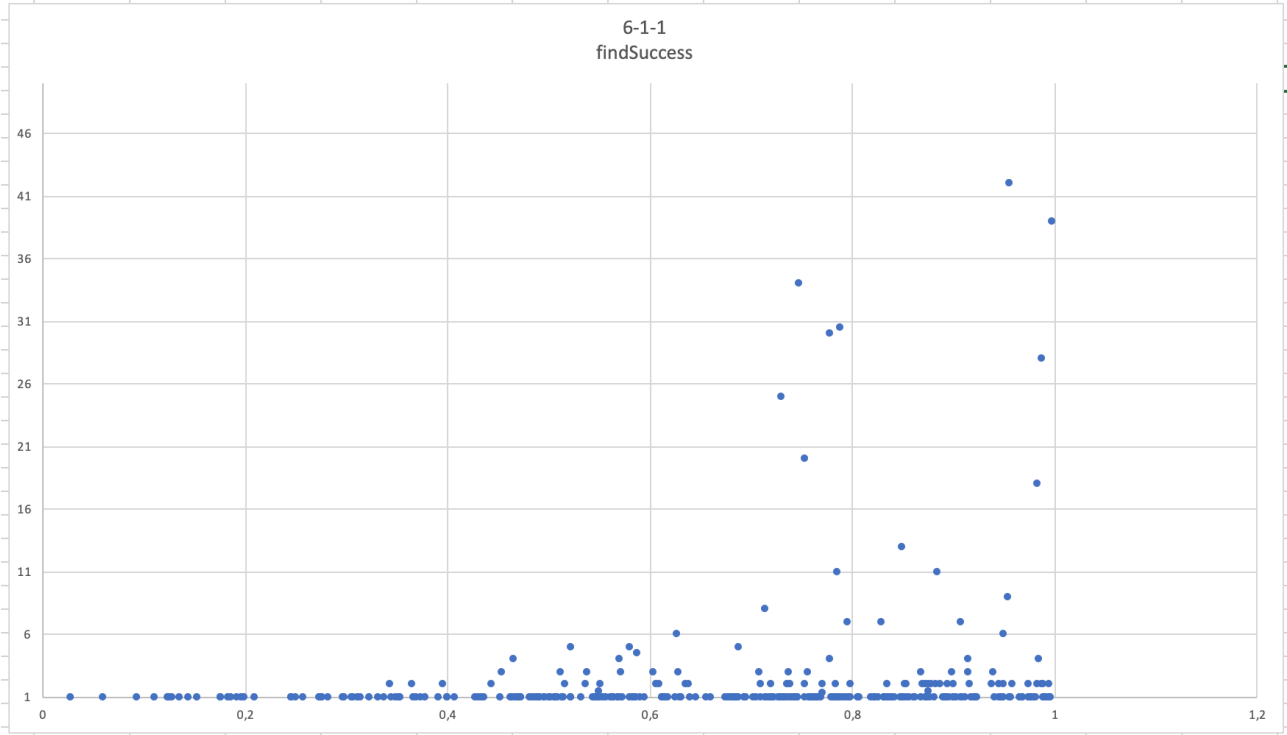
    return n;
}
```

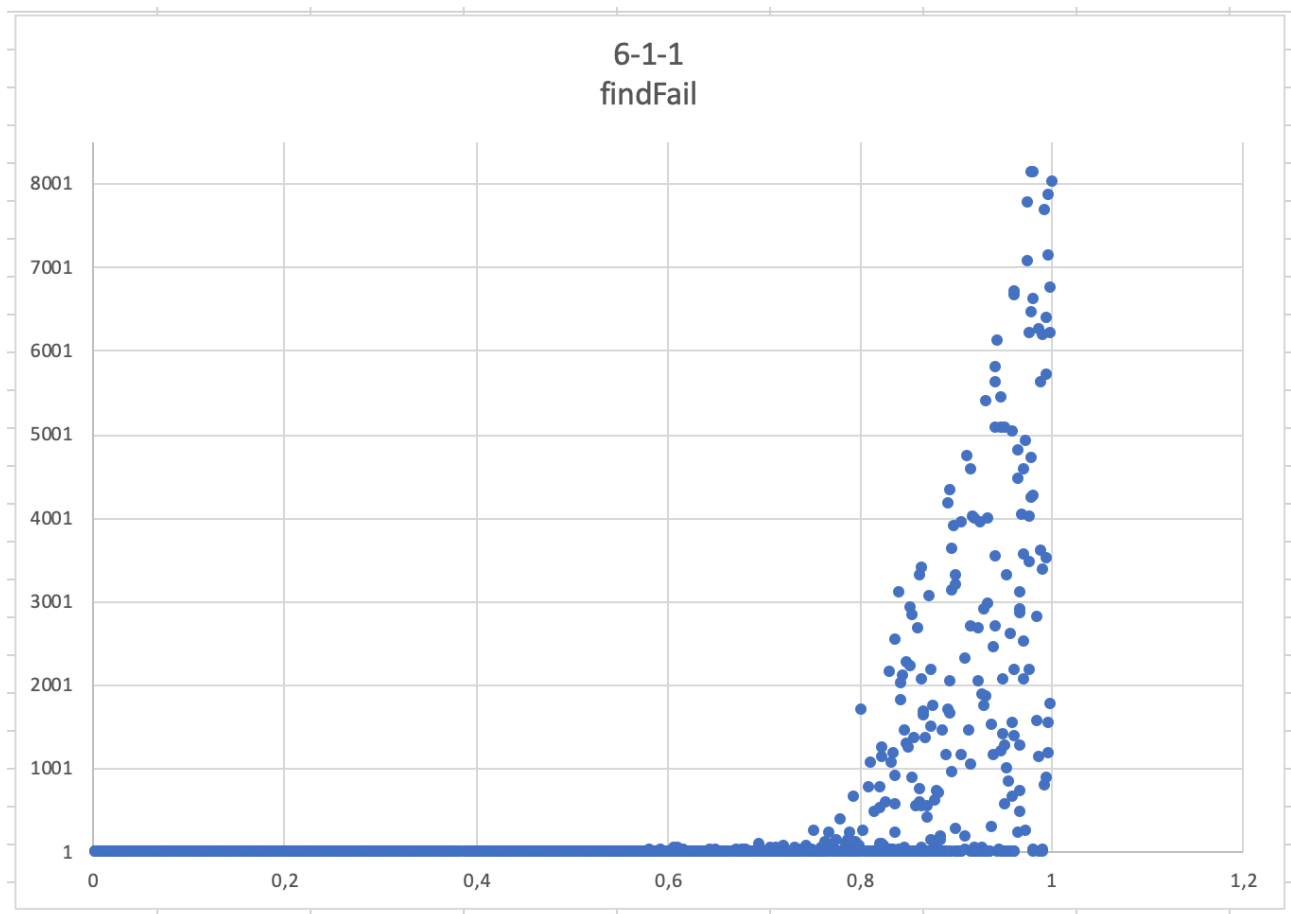
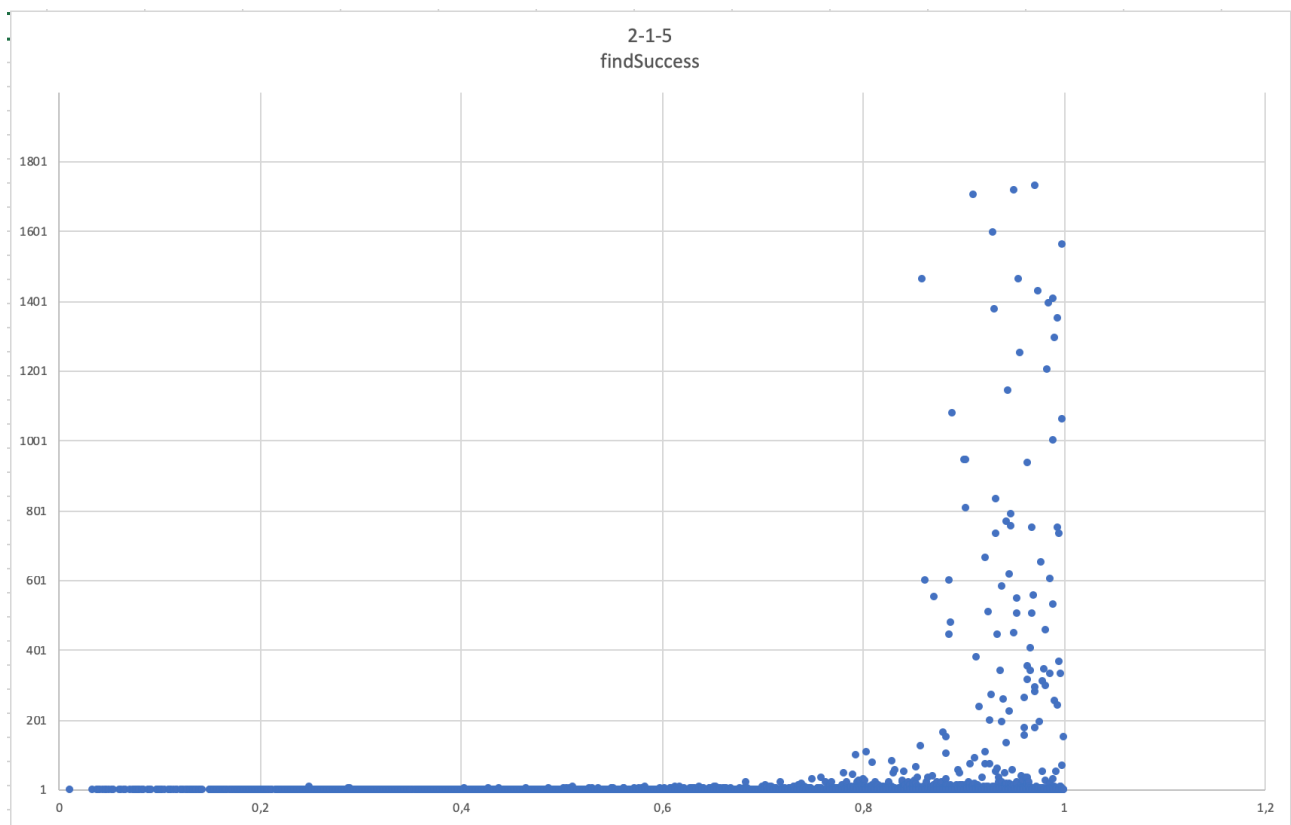
=> and my class looks like this

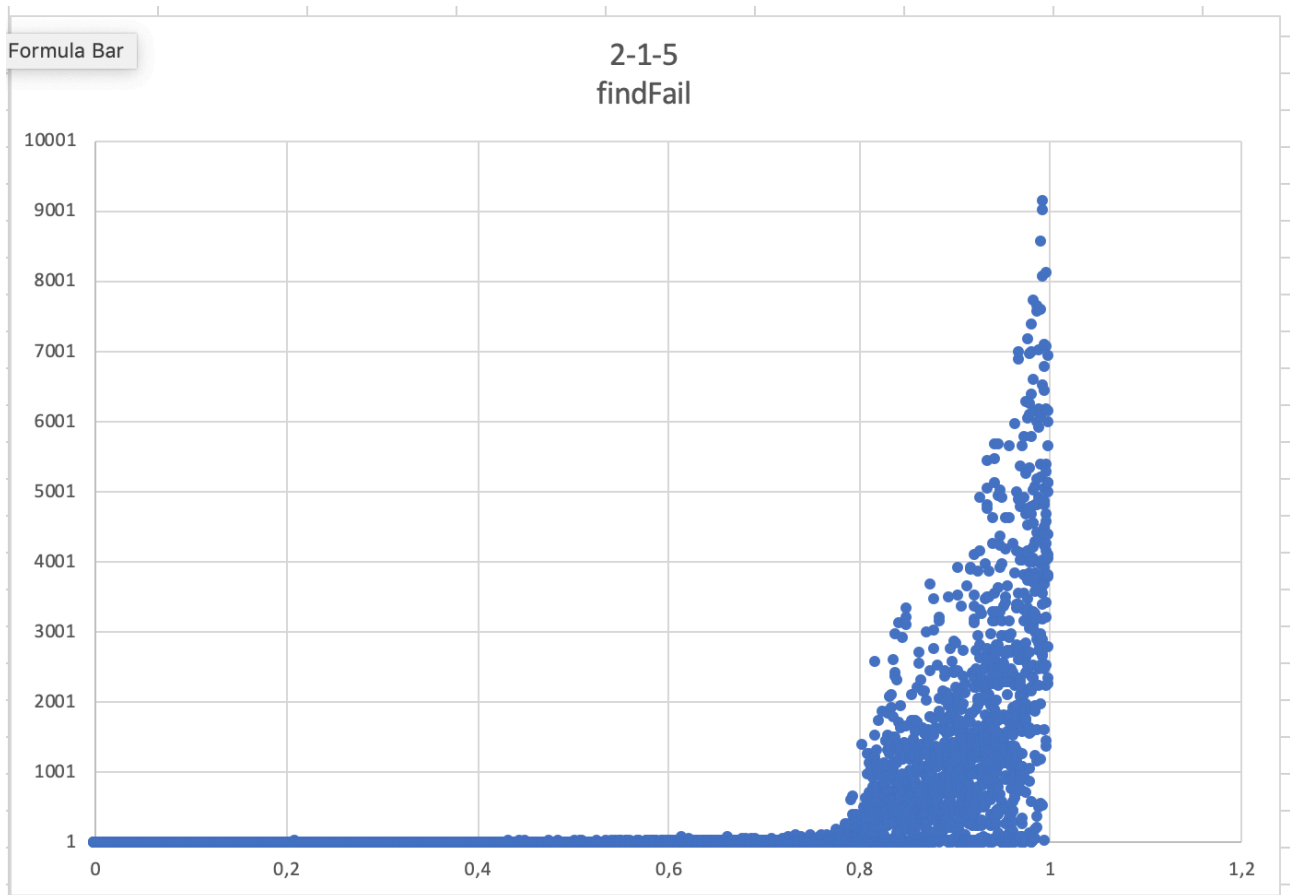
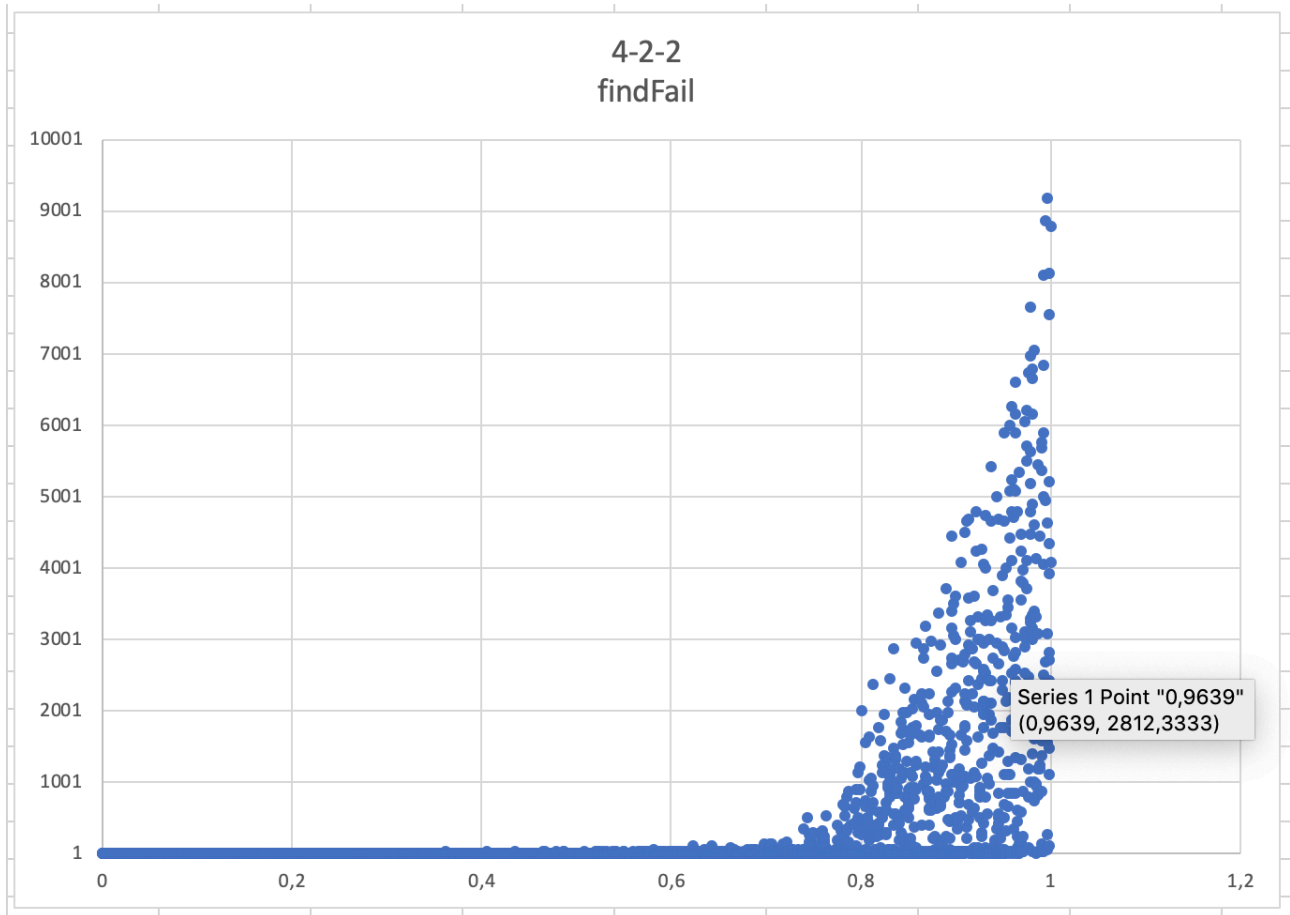
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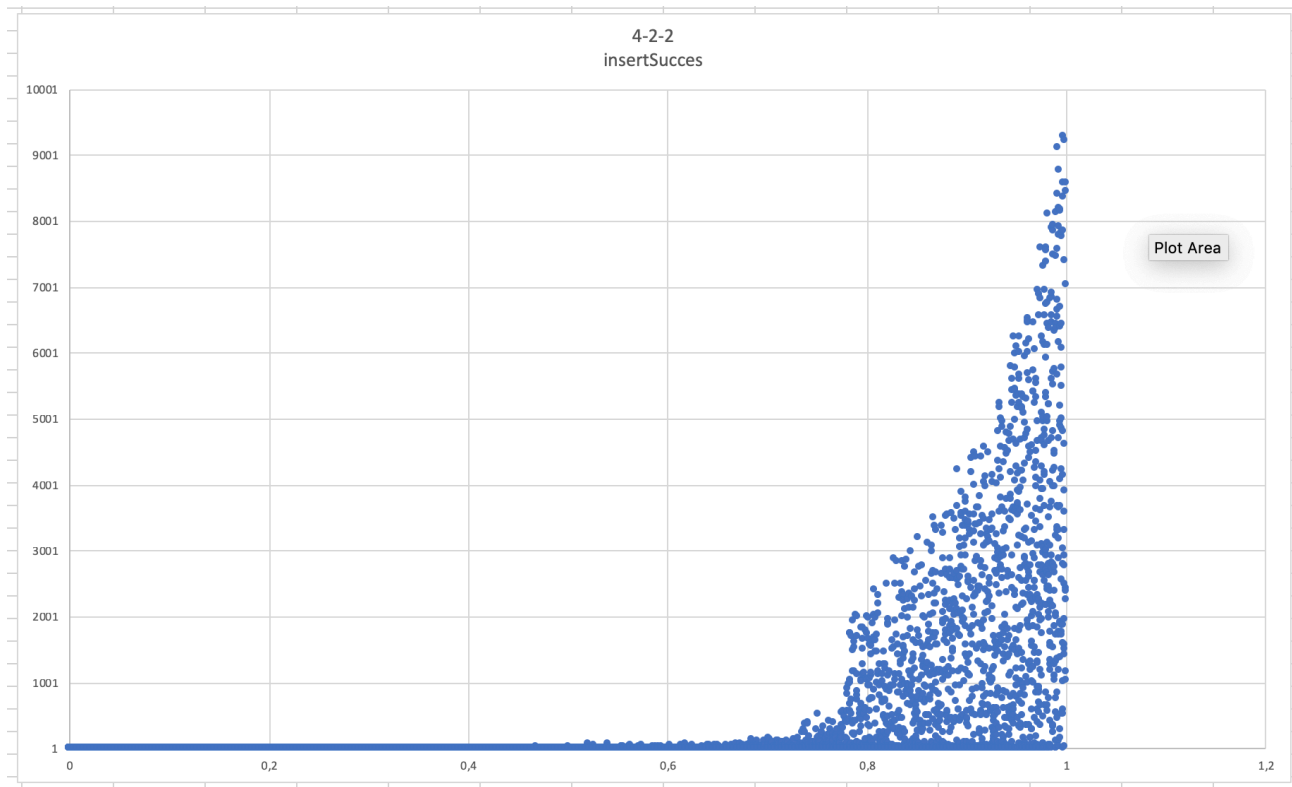
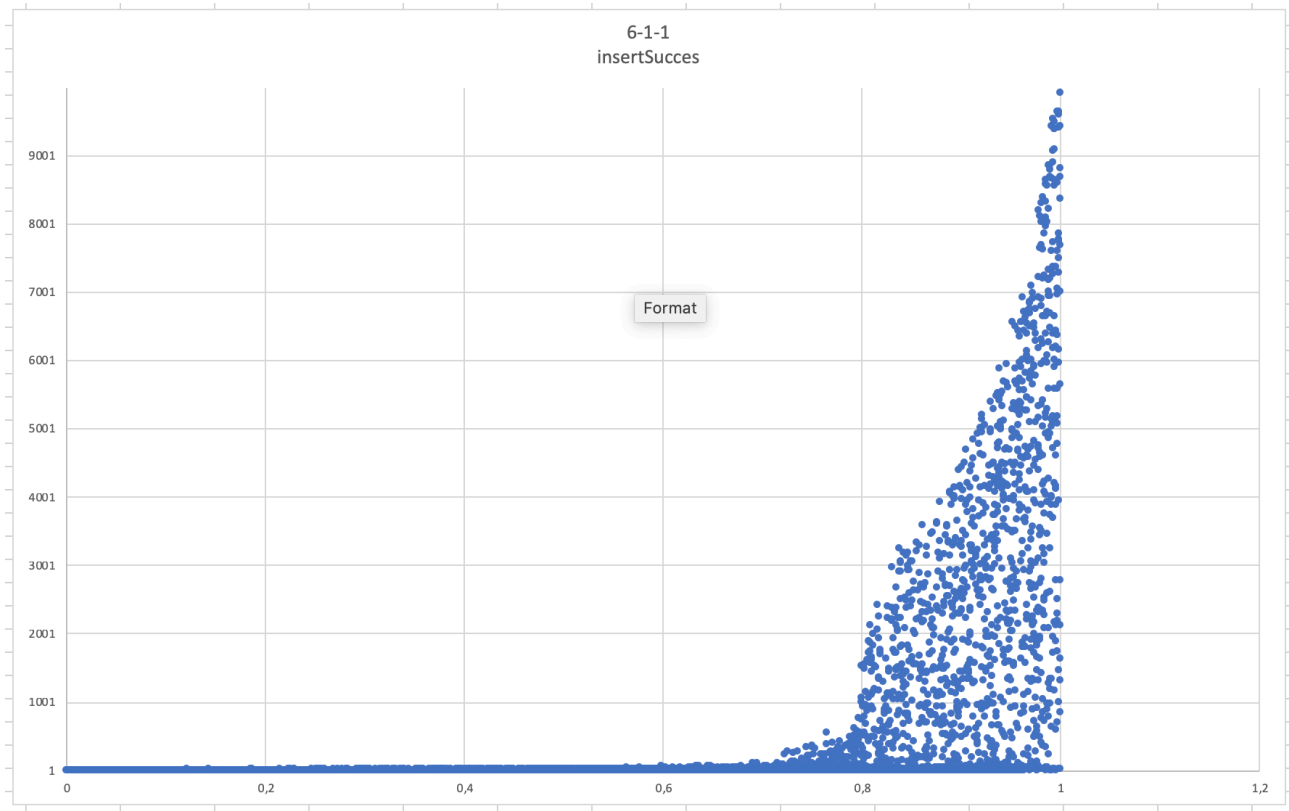
My Drawings

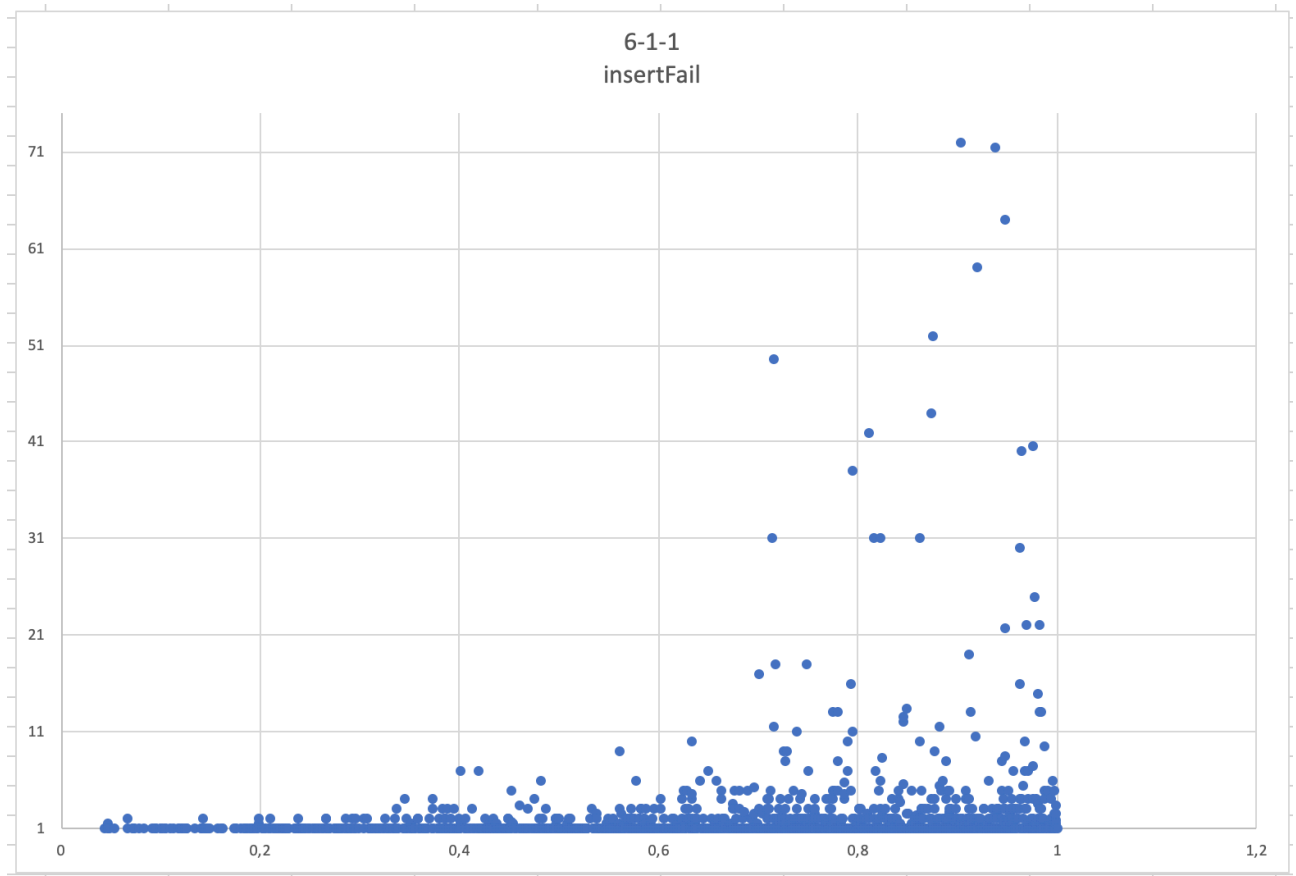
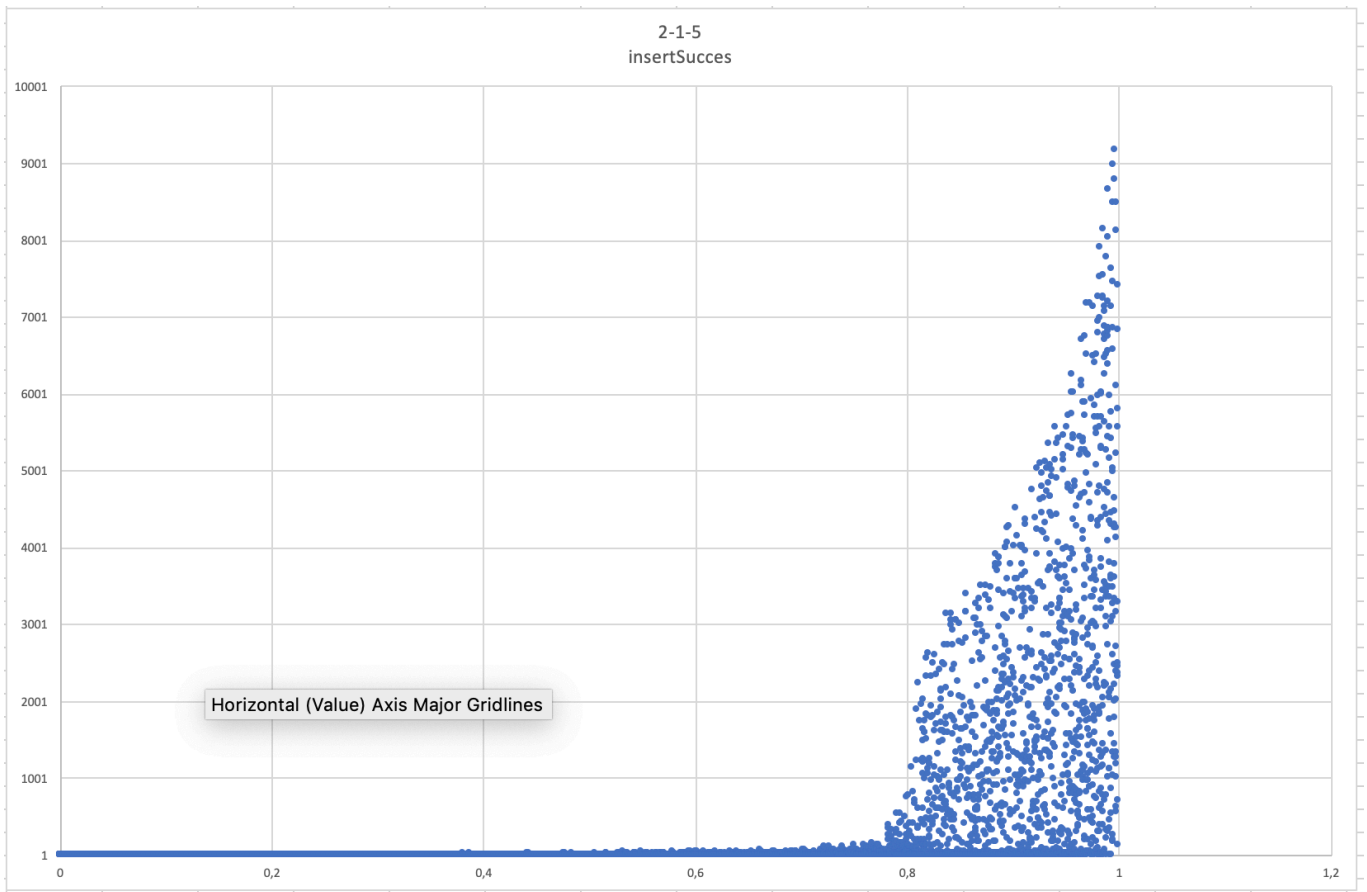
Note that all x-axis = load factor and y-axis =average number of probes

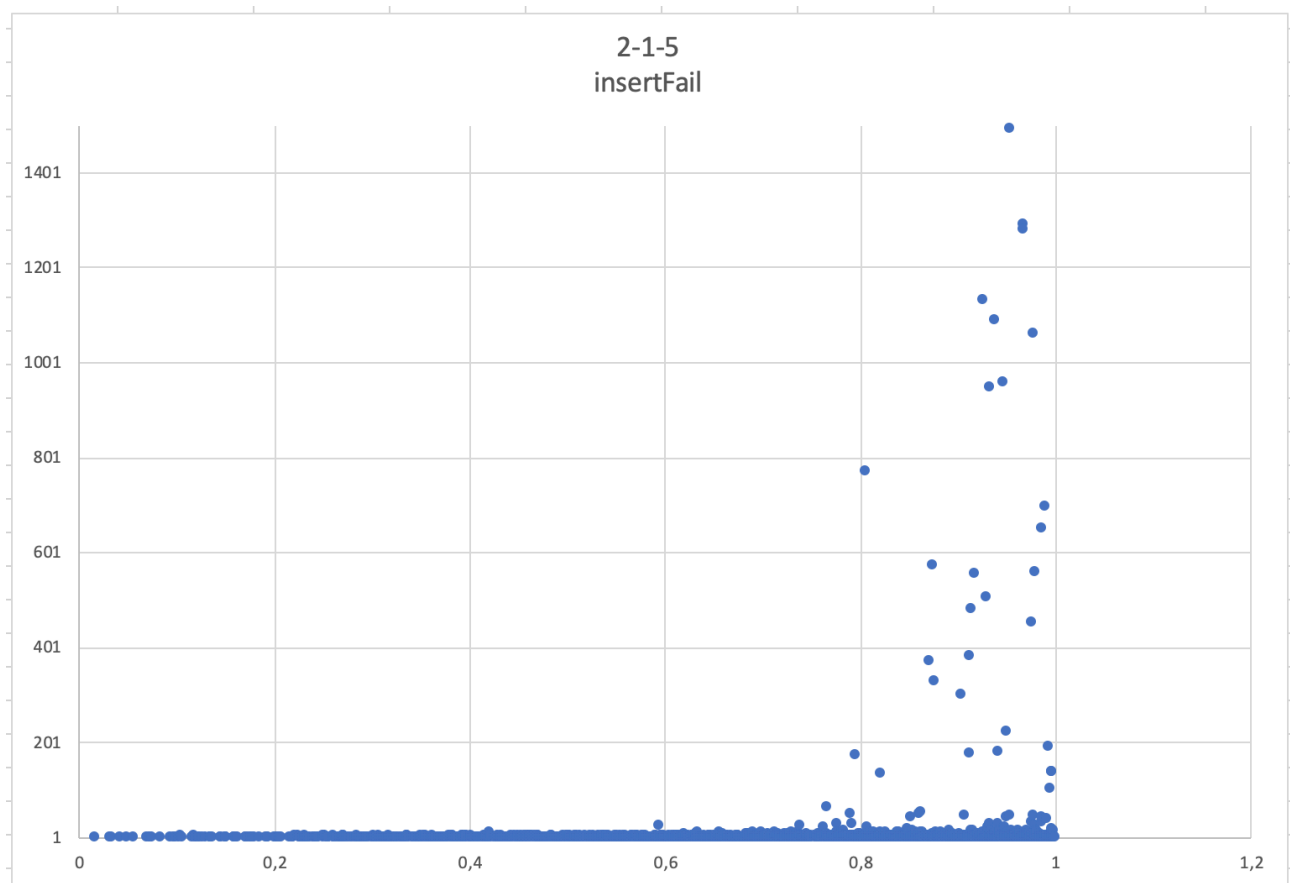
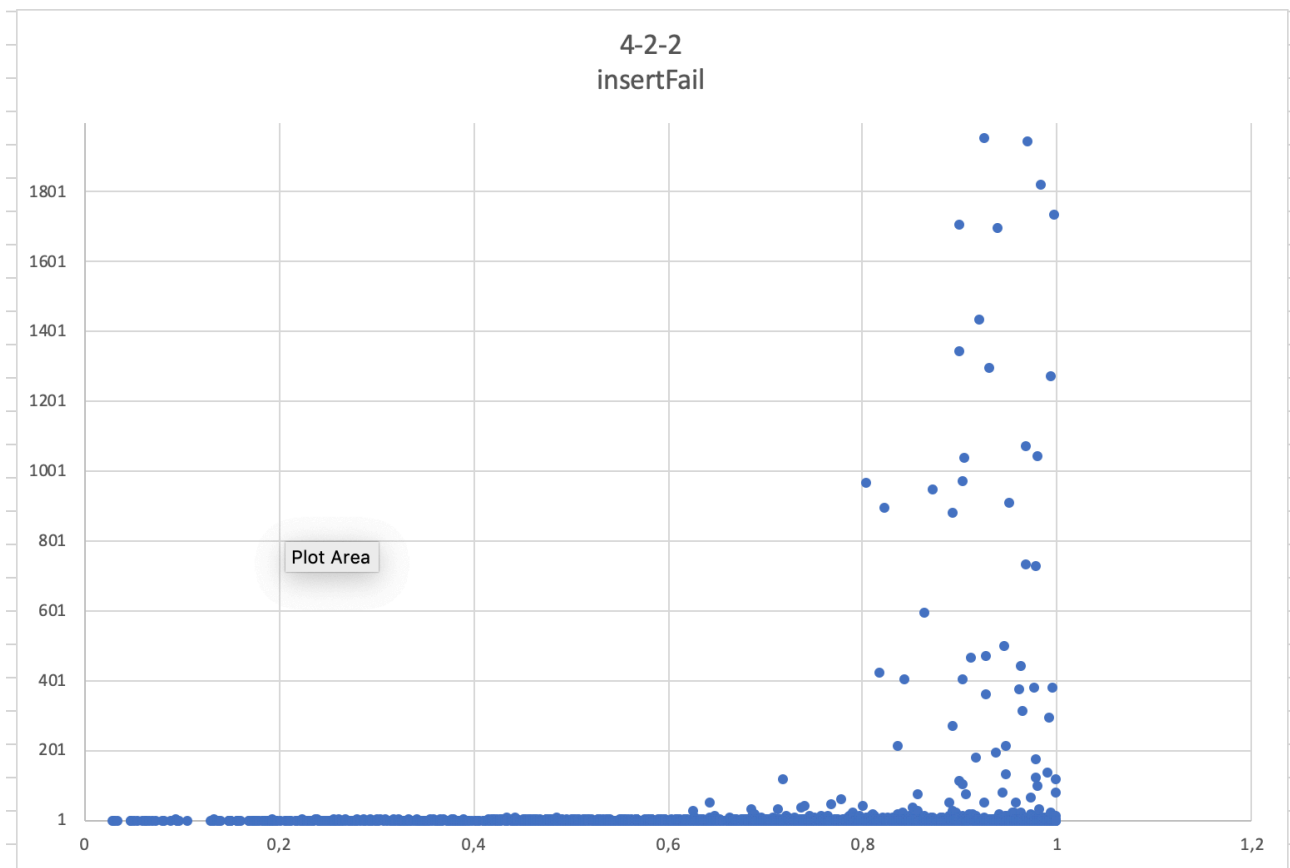


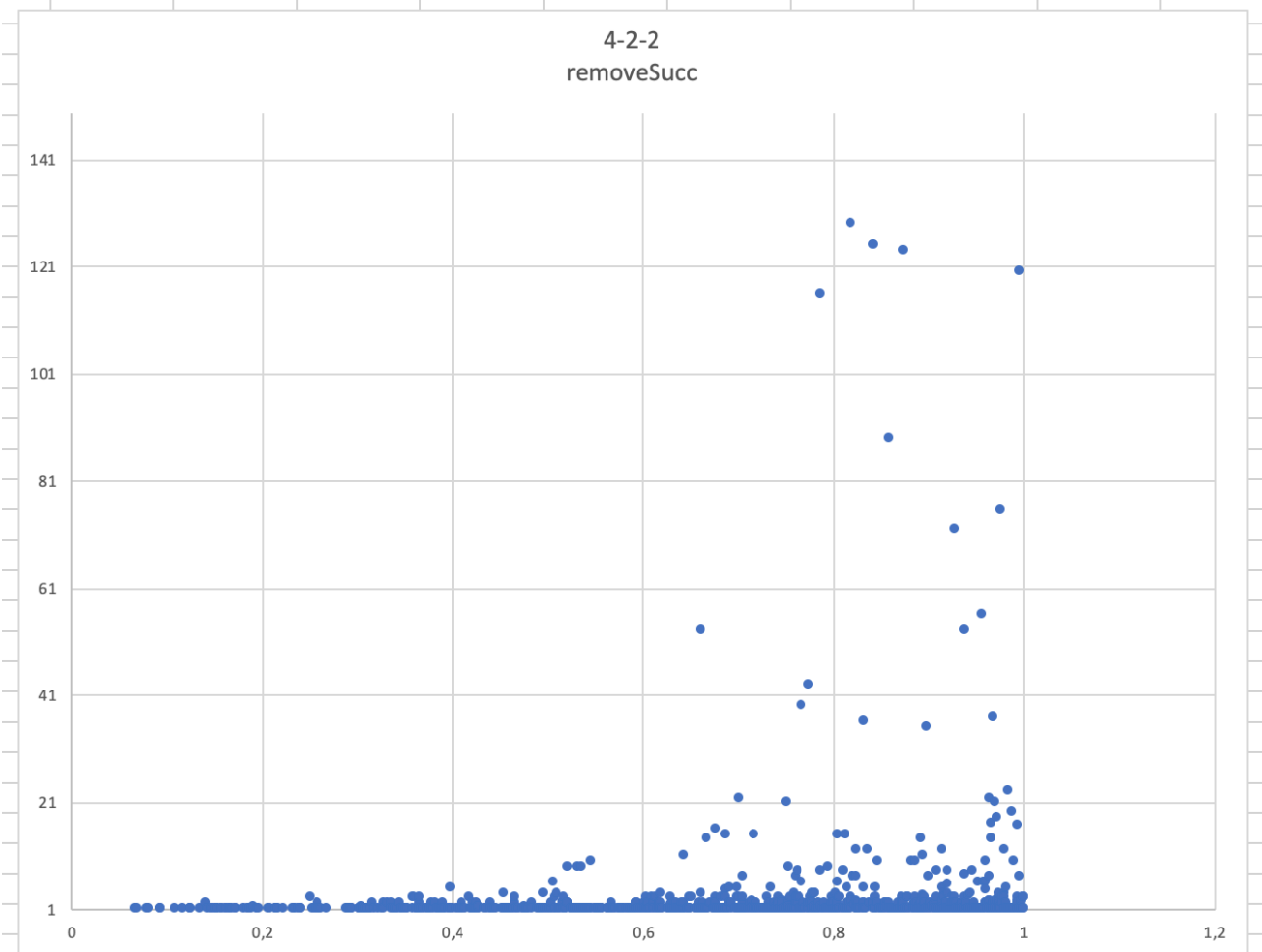
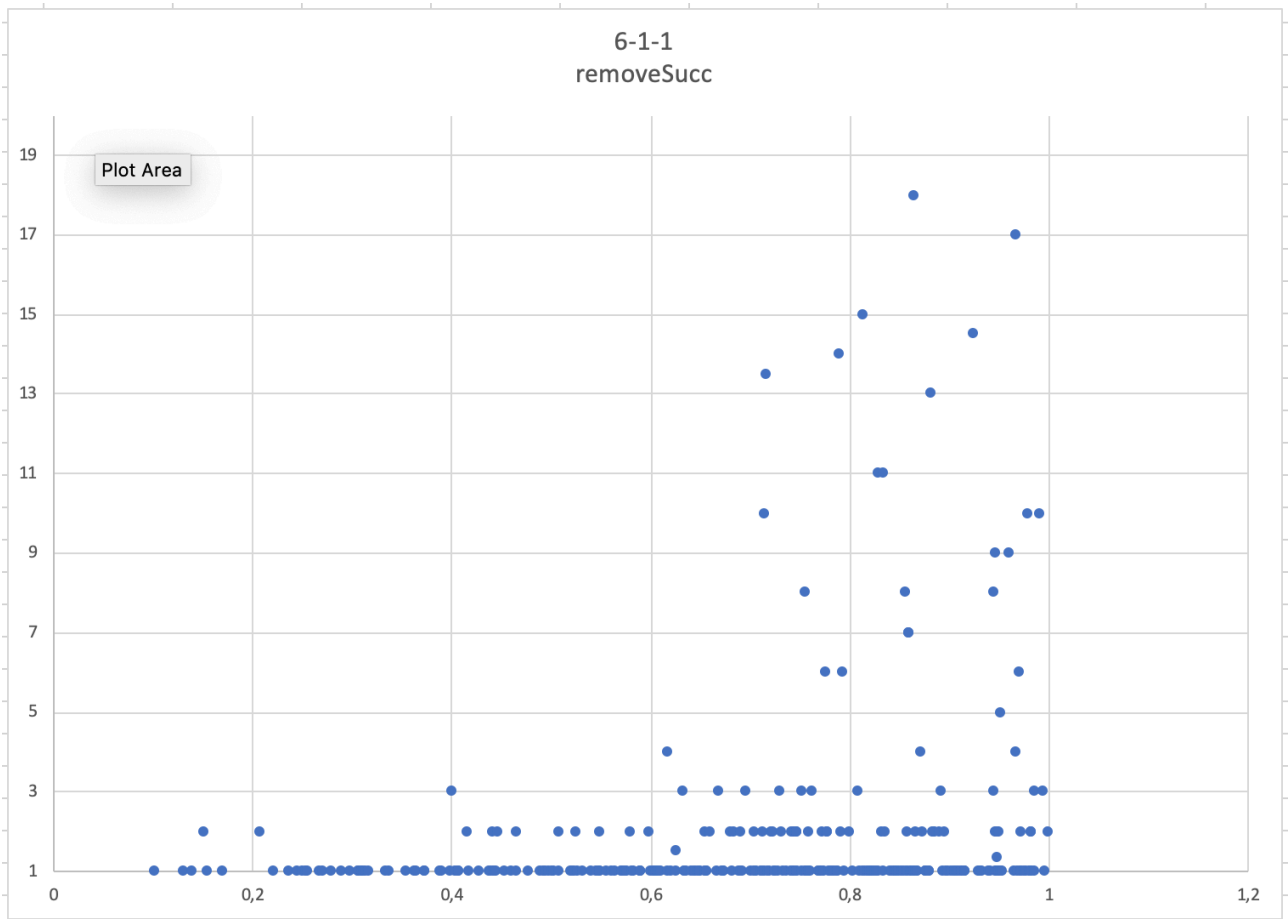


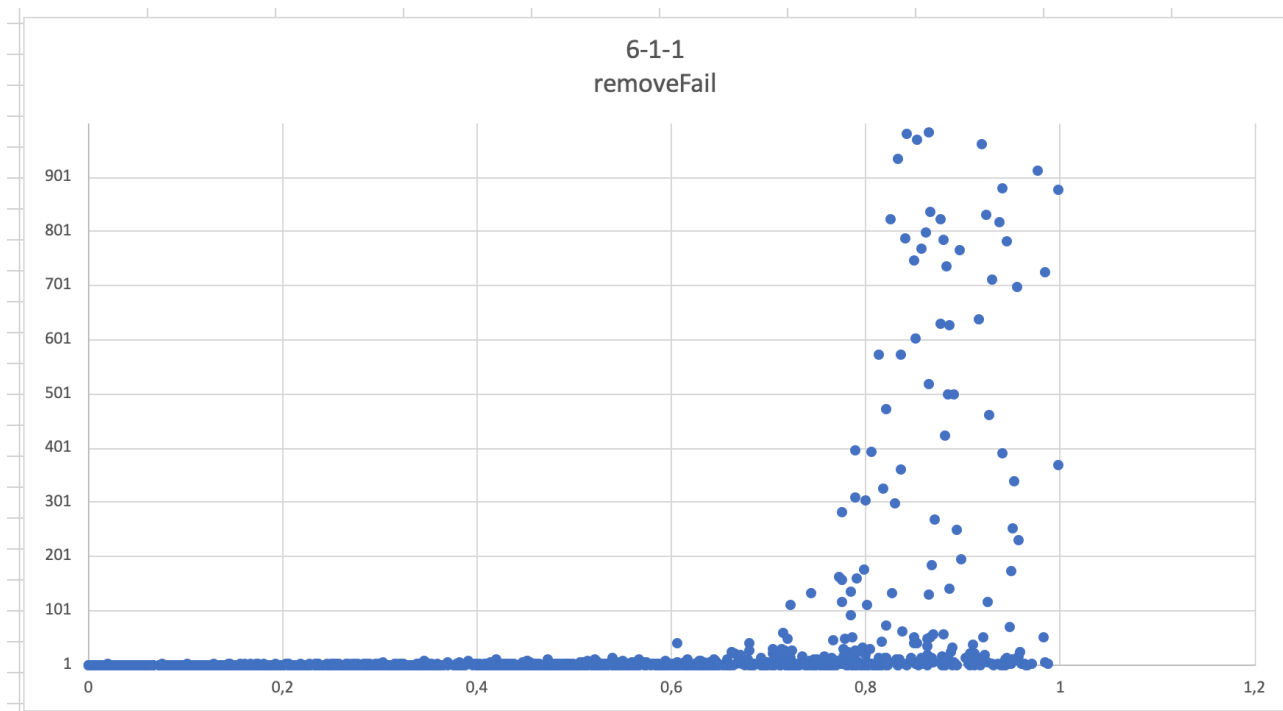
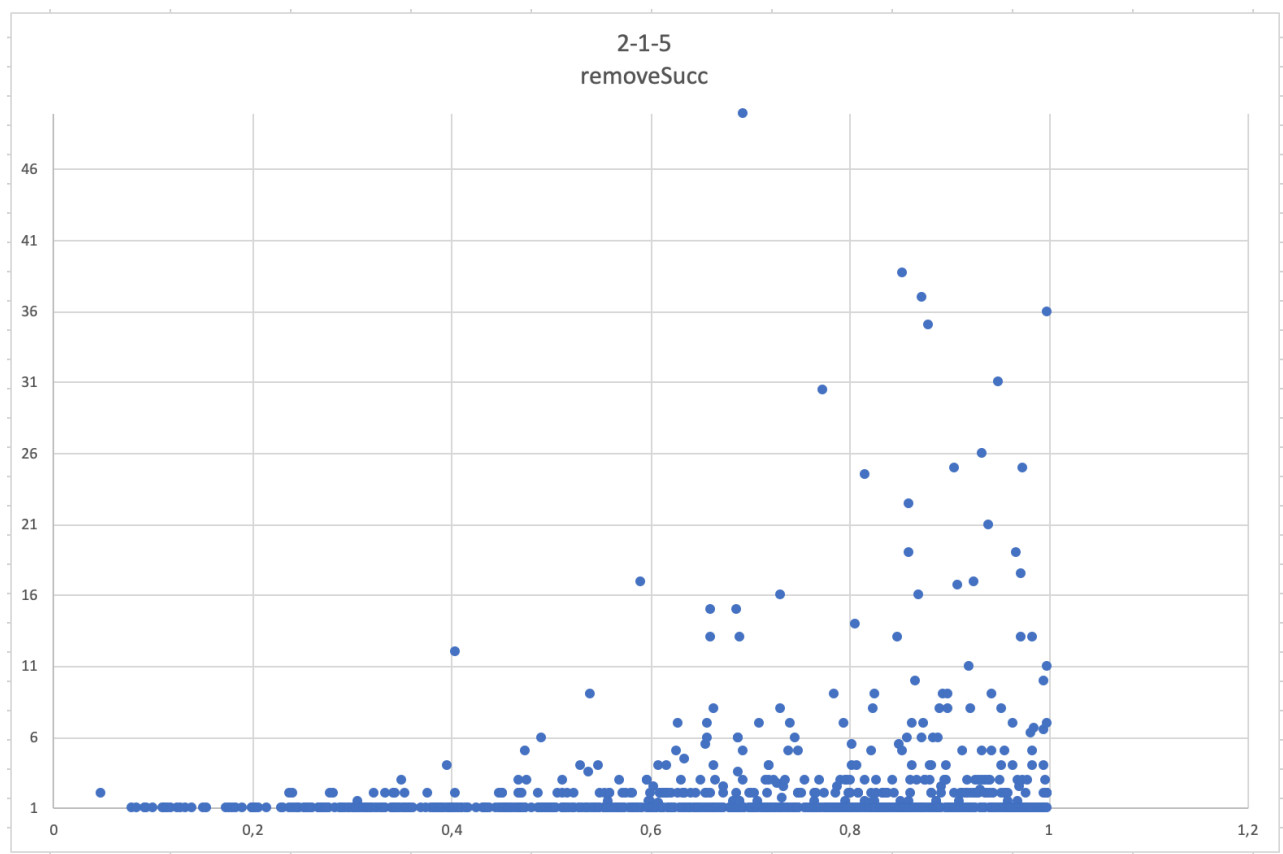


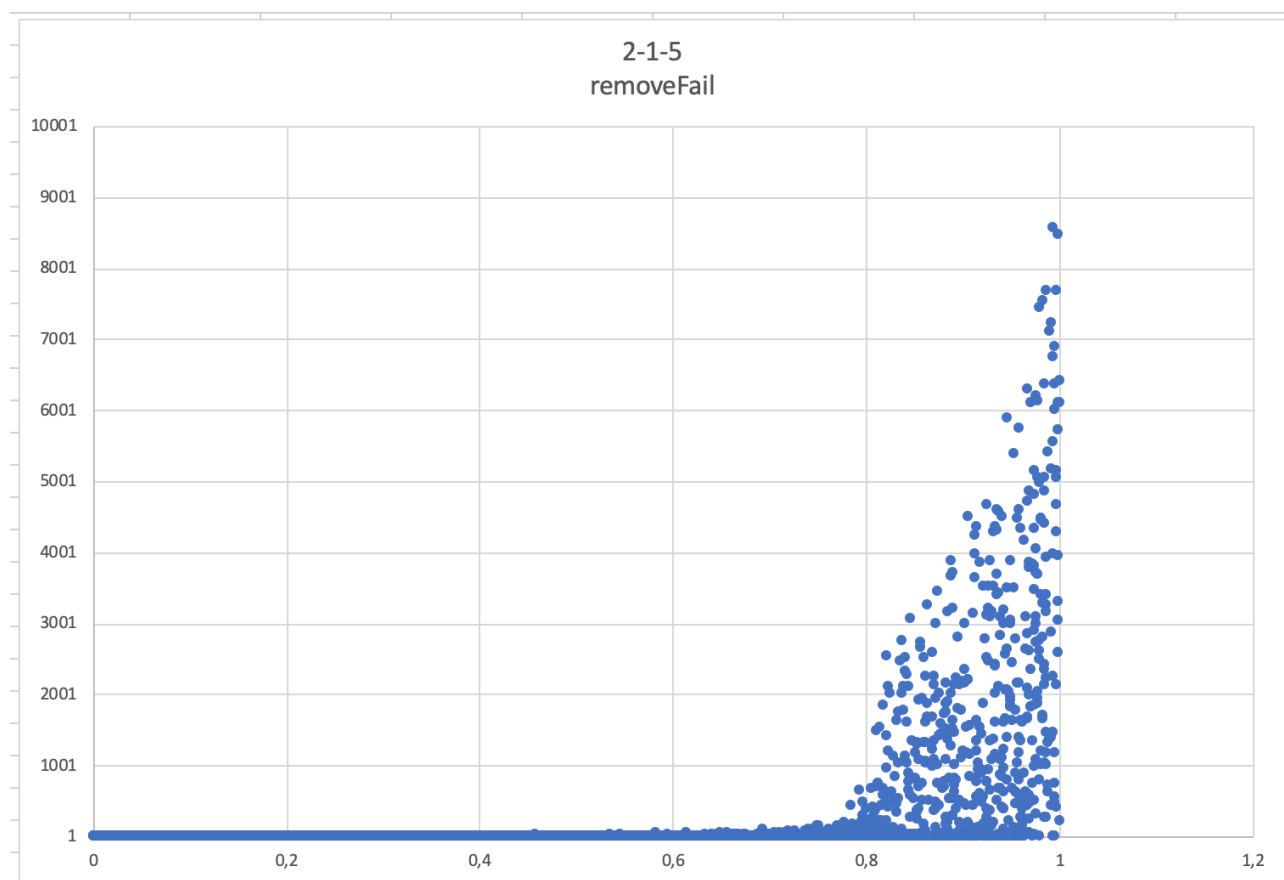
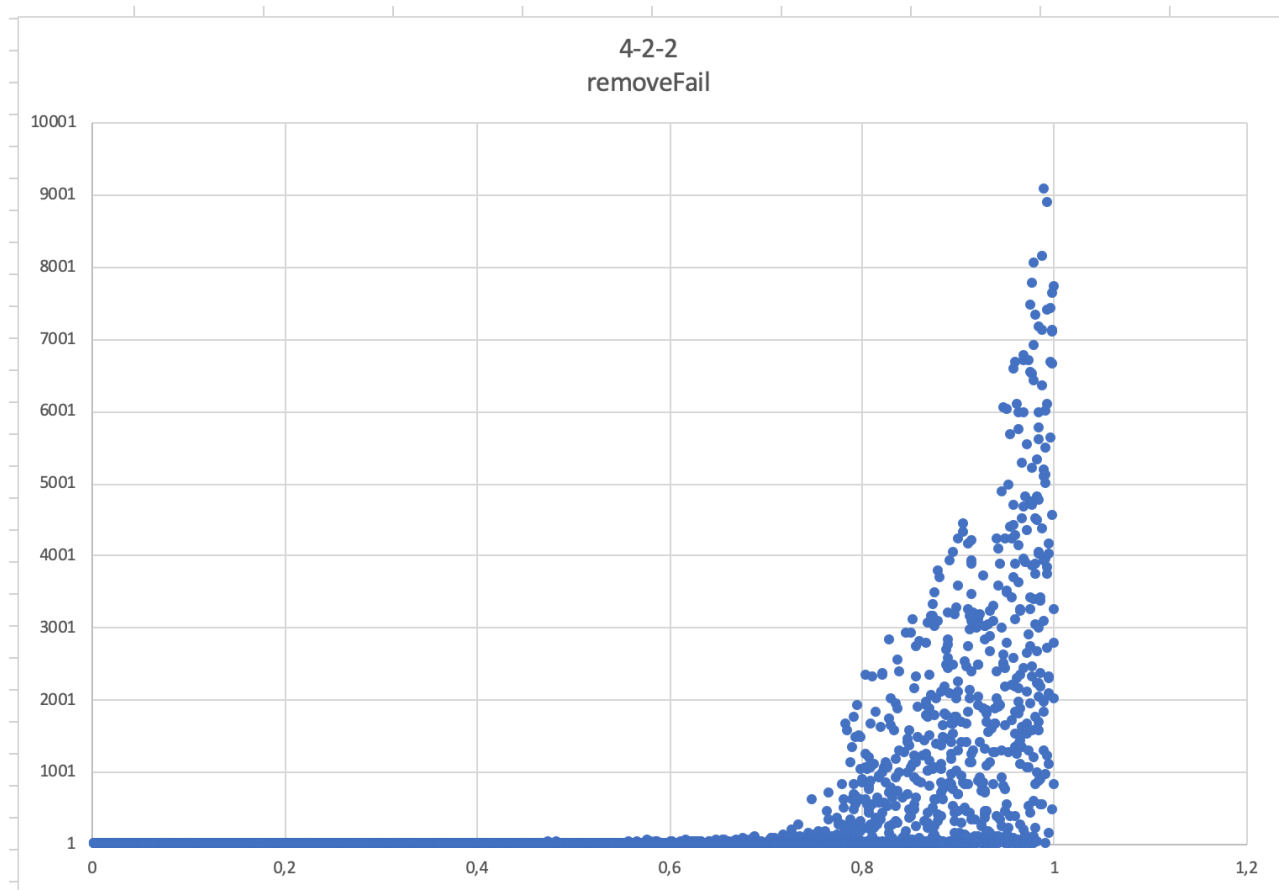












To conclude,

The graphs clearly show that, as the load factor increases, the average number of probes also increases. This means that our program runs slower when we don't resize the hash table.

Furthermore, if we compare the insertion of 6-1-1, 4-2-2 and 2-1-5 we can clearly see that as our table loads up, the traverse of probe increases.