Answer True or False to the following questions and briefly justify your answer:

- (a) With the Selective Repeat protocol, it is possible for the sender to receive an ACK for a packet that falls outside of its current window.
- (b) With Go-Back-N, it is possible for the sender to receive an ACK for a packet that falls outside of its current window.
- (c) The Stop&Wait protocol is the same as the SR protocol with a sender and receiver window size of 1.
- (d) Selective Repeat can buffer out-of-order-delivered packets, while GBN cannot. Therefore, SR saves network communication cost (by transmitting less) at the cost of additional memory.

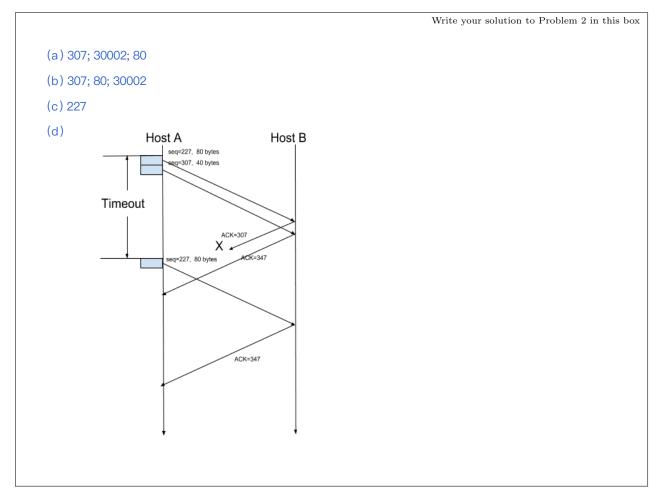
(a) True. It could happen when the delay is high. Suppose the sender sent a suquence of packests, and then

Write your solution to Problem 1 in this box

- (a) True. It could happen when the delay is high. Suppose the sender sent a suquence of packests, and then the receiver sent the corresponding ACKs back. If the network is slow, the sender times out and resend the packets before it gets the ACKs. Afte that, the sender receives the ACKs and moves on. So, when the sender receives the smae ACKs for the second time, those ACKs are outside of its current window.
- (b) True. Same principle in (a)
- (c) True. SR protocol with window size 1 is identical to Stop&Wait protocol
- (d) True. SR protocol can buffer out-of-order-delivered packets within the window. When the loss rate is high, resending the whole window of packets is expensive. So, SR protocol saves network communication cost.

Host A and B are communicating over a TCP connection, and Host B has already received from A all bytes up through byte 226. Suppose Host A then sends two segments to Host B back-to-back. The first and second segments contain 80 and 40 bytes of data, respectively. In the first segment, the sequence number is 227, the source port number is 30002, and the destination port number is 80. Host B sends an acknowledgment whenever it receives a segment from Host A. Fill in the blanks for questions (a) - (c) directly; work out the diagram in the box for question (d).

- (a) In the second segment sent from Host A to B, the sequence number is _____, source port number is _____, and destination port number is _____.
- (b) If the first segment arrives before the second segment, in the acknowledgment of the first arriving segment, the ACK number is _____, the source port number is _____, and the destination port number is _____.
- (c) If the second segment arrives before the first segment, in the acknowledgment of the first arriving segment, the ACK number is _____.
- (d) Suppose the two segments sent by A arrive in order at B. The first acknowledgment is lost and the second acknowledgment arrives after the first timeout interval. Draw a timing diagram in the box below, showing these segments and all other segments and acknowledgment sent. Assume no additional packet loss. For each segment in your diagram, provide the sequence number and the number of bytes of data; for each acknowledgment that you add, provide the ACK number.



In Fast Retransmit algorithm, we saw TCP waits until it has received three duplicate ACKs before performing a fast retransmit. Why do you think the TCP designers chose not to perform a fast retransmit after the first or second duplicate ACKs for a segment received?

Write	your solution to Problem 3 in this box
Because TCP designers want to minimize the retransmission rate. If the TCP designers decided to perform a fast retransmit after the first a segment received, a packet is retransmitted as long as it is not delive much more unnecessary traffic which makes the TCP less efficient.	or second duplicate ACKsfor ered in order. This will cause

Suppose that three measured SampleRTT values are 106 ms, 120 ms, and 140 ms. Compute the EstimatedRTT after each of these SampleRTT values is obtained, assuming that the value of EstimatedRTT was 100 ms just before the first of these three samples were obtained. Compute also the DevRTT after each sample is obtained, assuming the value of DevRTT was 5 ms just before the first of these three samples was obtained. Last, compute the TCP TimeoutInterval after each of these samples is obtained.

Write your solution to Problem 4 in this box Since alpha and beta are not given, I will assume the typical values, 0.125 and 0.25 respectively. EstimatedRTT = alpha *SampleRTT+ (1 -alpha) *EstimatedRTT DevRTT = beta * | SampleRTT - EstimatedRTT | + (1 - beta) * DevRTT TimeoutInterval = EstimatedRTT + 4 * DevRTT (1) Sample RTT = 106ms Estimated RTT = 0.125 * 106 + (1 - 0.125) * 100 = 100.75ms DevRTT = 0.25 * | 106 - 100.75 | + (1 - 0.25) * 5 = 5.06ms TimeoutInterval = 100.75 + 4 * 5.06 = 120.99ms (2) Sample RTT = 120ms EstimatedRTT= 0.125 *120 + (1 -0.125) *100.75 = 103.15ms DevRTT= 0.25 * | 120 -103.15| + (1 -0.25) *5.06 = 8ms TimeoutInterval = 103.15 + 4 *8 = 135.15ms (3) Sample RTT = 140 Estimated RTT = 0.125 * 140 + (1 - 0.125) * 103.15 = 107.75ms DevRTT = 0.25 * | 140 -107.75| + (1 -0.25) * 8 = 14.06ms TimeoutInterval = 107.75 + 4 *14.06 = 164ms

Compare Go-Back-N, Selective Repeat, and TCP (no delayed ACK). Assume that timeout values for all three protocols are sufficiently long, such that 5 consecutive data segments and their corresponding ACKs can be received (if not lost in the channel) by the receiving host (Host B) and the sending host (Host A), respectively. Suppose Host A sends 5 data segments to Host B, and the 2nd segment (sent from A) is lost. In the end, all 5 data segments have been correctly received by Host B.

- (a) How many segments has Host A sent in total and how many ACKs has Host B sent in total? What are their sequence numbers? Answer this question for all three protocols.
- (b) If the timeout values for all three protocols are much longer than 5RTT, then which protocol successfully delivers all five data segments in shortest time interval?

Write your solution to Problem 5 in this box (a) Go-Back-N: 9 segments sent by A and 8 ACKs sent by B. Since 2 is lost, A wil retransmit 2,3,4,5 after it sent 1,2,3 4, 5. Bwill first sent 4 ACKs1 indicating 2 is lost and then sent ACK 2, 3, 4, 5. Selective Repeat: 6 segments sent by A and 5 ACKs sent by B. Since 2 is lost, A will retransmit 2 while the rest remained in buffer. Bwill send ACKs in the order it receives the segments 1, 3, 4, 5, 2 6 segments sent by A and 5 ACKs sent by B. Since 2 is lost, A will retransimit 2 while the rest remained in the buffer. Bwill sent 4 ACK2, and one ACK6 (if there is segment 6). (b) TCP. because TCP use fast retransmit algorithm.