

Internet Security

CS 144 Web Applications

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Outline

- Introduction to Cryptography
 - Symmetric Key
 - Asymmetric Key (Public-Key Infrastructure)
- Common Web Vulnerability
- Project 5

Some common attack techniques

- DDoS (Distributed Denial of service)
- Phishing
 - spoof web site to look like the real one
- Pharming (DNS cache poisoning)
 - e.g. wrong DNS resolution
- Packet sniffing
 - Cache theft
- SQL injection
 - `SELECT * FROM users WHERE (name='cs144') and (password='1' OR '1'='1');`

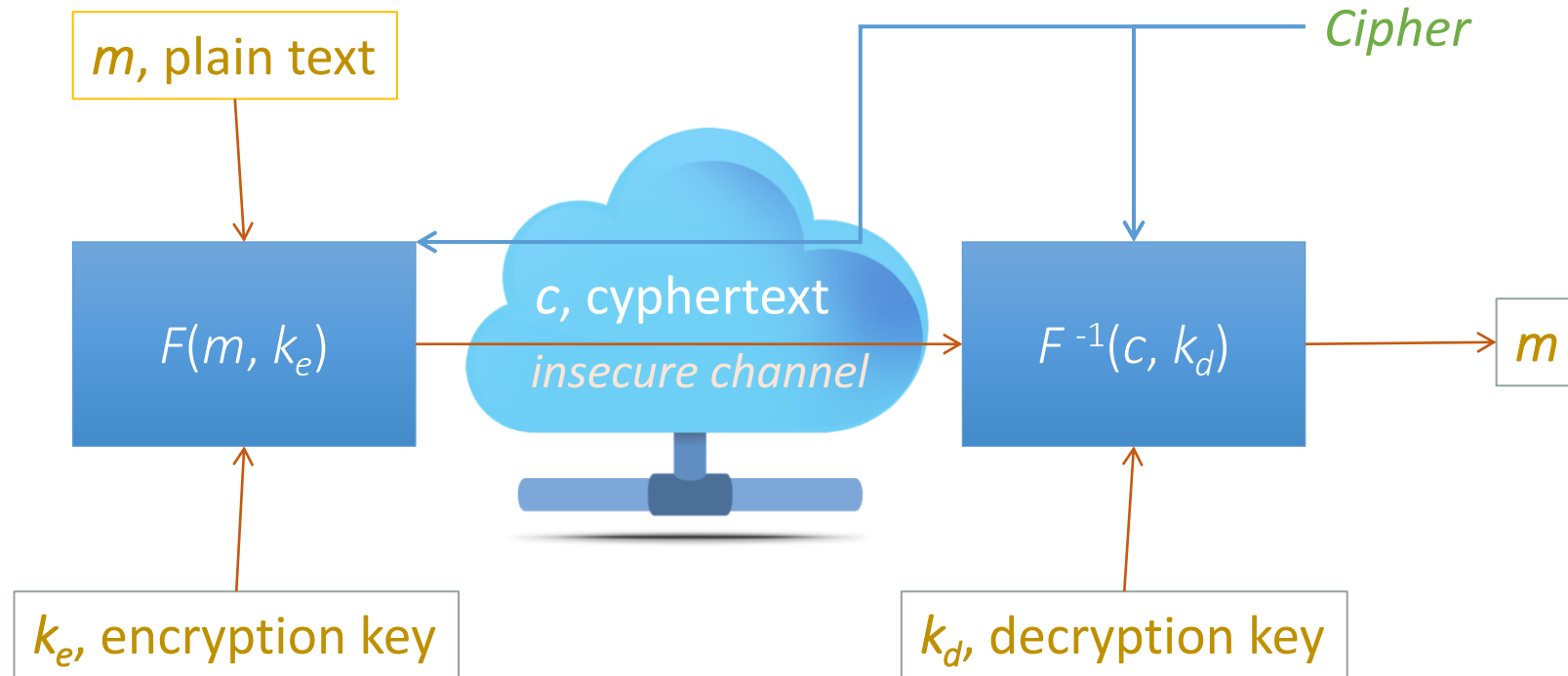


Desired Guarantees

- Confidentiality
 - Maintaining the privacy of the “conversation.”
- Message/data integrity
 - No one can modify the content of the messages.
- Authentication
 - Making sure the other party is who he/she claims to be.
- Authorization
 - Managing access to resources upon successful authentication.
 - UNIX User and Group permission mechanism



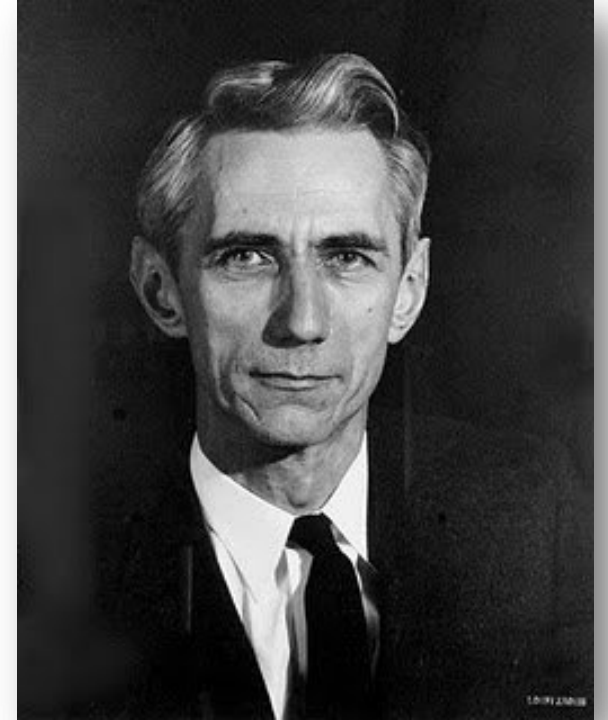
Encryption Algorithm



- Symmetric: $k_e = k_d$
- Asymmetric: $k_e \neq k_d$

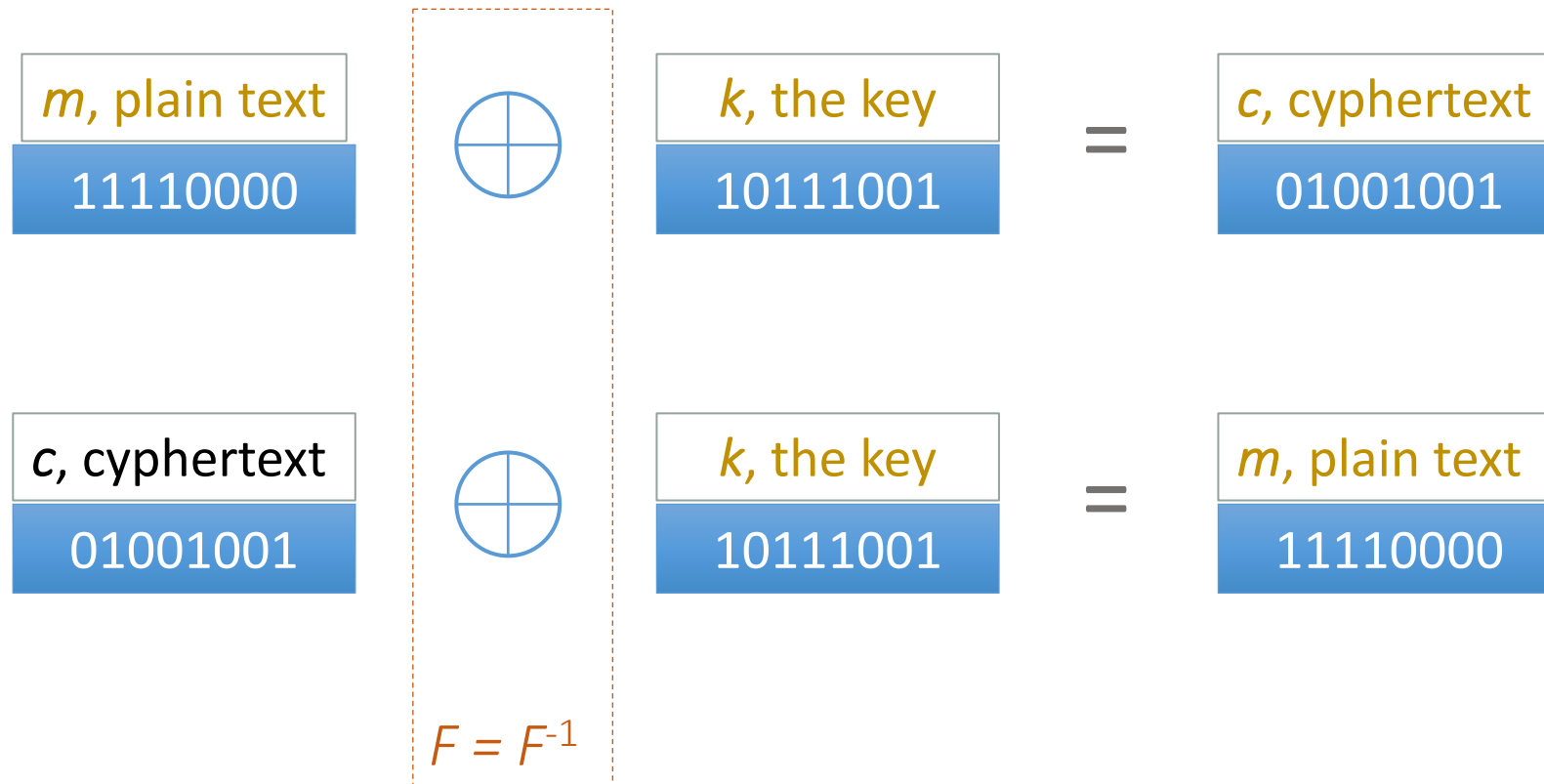
Cipher properties

- (Shannon) **Perfect Secrecy**: $P(m | c) = P(m)$
 - We should not be able to guess m by just looking at the cyphertext c .
- *One-Time Pad* (OTP)
 - Guarantees Perfect Secrecy.
 - Use a brand new key to encrypt messages every time.
 - Never use the same key again.
 - Pad value (i.e. key) must be at least as long as the message
 - Very expensive in practical terms.



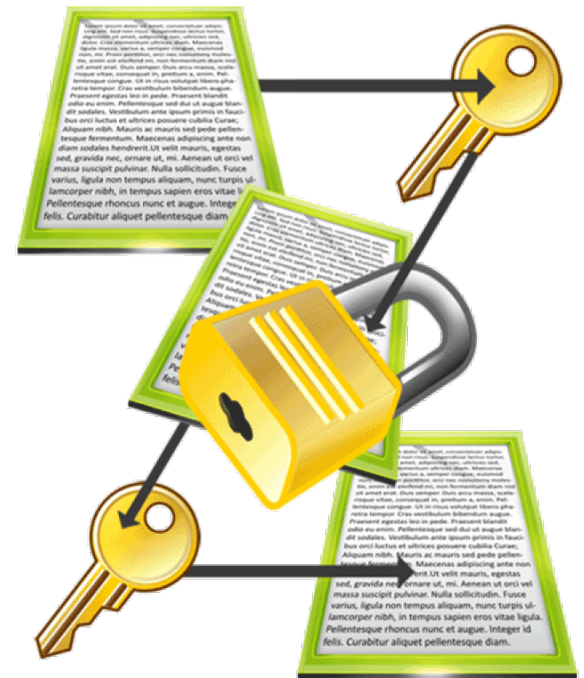
Symmetric-Key Cryptography

- Same key for encrypting and decrypting: $k_e = k_d = k$
 - XOR is the simplest operator:



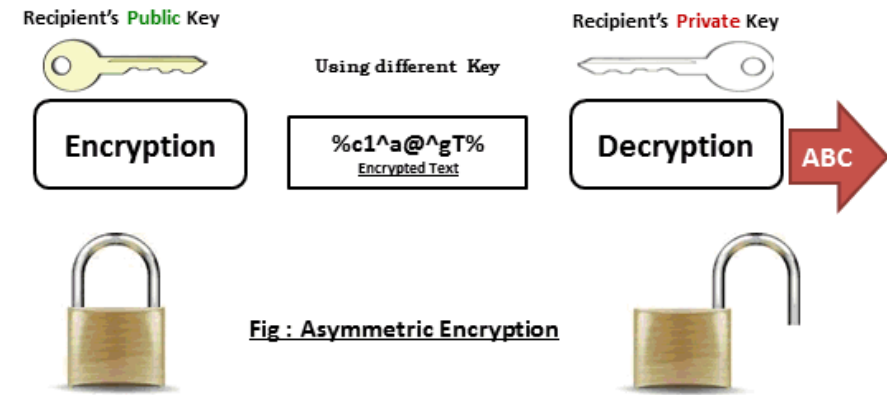
Symmetric-Key Cryptography

- **DES**, *Data Encryption Standard*, 64-bit cypher.
- **AES**, *Advanced Encryption Standard*, 128-bit cypher, 128/192/256-bit key.
- How do we ensure **authentication**?
 - Challenge message (a random number).
- How do we ensure **confidentiality**?
 - Generate individual key for each conversation.
- How do we ensure **authorization**?
 - Encrypting distinct data with distinct keys.
- How do we ensure **integrity**?
 - Encrypting a checksum.
- Any problem with symmetric-key cryptography?
 - Agreeing on key.
 - $n(n-1)/2$ keys for n parties



Asymmetric-Key Cryptography

- Different keys
 - For encryption: $e = k_e$ (Public Key)
 - $c = F(m, e)$
 - For decryption: $d = k_d$ (Private Key)
 - $m = F^{-1}(c, d)$
- Requirements:
 - $F^{-1}(F(m, e), d) = m$
 - $c = F(m, e) \not\Rightarrow m$ (Perfect Secrecy)
 - $e \not\Rightarrow d$



RSA Algorithm

1. Pick two large prime numbers: p and q .
2. Select e such that $1 < e < (p-1)(q-1)$
 - e doesn't have to be random.
 - e is coprime to $(p-1)(q-1)$
3. Solve for d in $de \bmod (p-1)(q-1) = 1$
4. Make $n = pq$ and throw away p and q .
5. The new keys are $k_e = (e, n)$ and $k_d = (d, n)$
 - Encryption $c = F(m, k_e) = m^e \bmod n$
 - Decryption $m = F^{-1}(c, k_d) = c^d \bmod n$



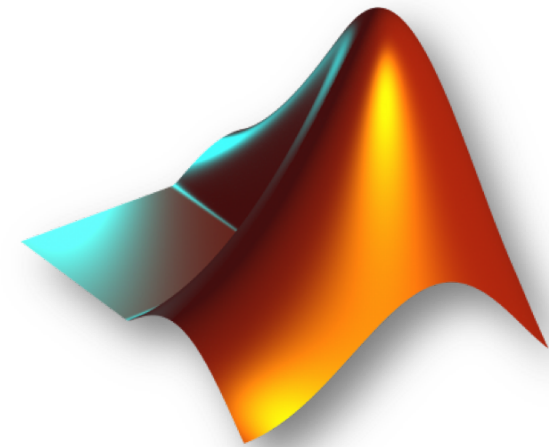
Rivest, Shamir, and Adleman

RSA Properties

- $F^{-1}(F(m, e), d) == m$?
 - $m = m^{ed} \bmod n$
- Can we obtain m from $c = m^e \bmod n$?
 - We know c , e , and n .
 - **RSA problem**
- Can we obtain d from $de \bmod (p-1)(q-1) = 1$?
 - We know e , and $n = pq$.
 - **Large-number factorization problem.**
- *RSA is one thousand times slower than symmetric-key cryptography.*

RSA Example

- Let $p = 7$ and $q = 11$
- $e < (6)(10)$
 - $e = 13$, being coprime to p and q
- Solve for d in $13d \bmod 60 = 1$
 - $d = 37$
- $n = (7)(11) = 77$, and throw away p and q
- If $m = 3$, then $c = 3^{13} \bmod 77$
 - $c = 38$
- If $c = 38$, then $m = 38^{37} \bmod 77$
 - $m = 3$



Asymmetric-Key Cryptography Applications

- Confidentiality?
 - Asymmetric-Key Cryptography is very expensive.
 - Use it to agree on a common key, and continue communicating with *symmetric-key cryptography*.
- Authentication?
 - A wants to make sure B is not someone else.
 - A generates a random number r , and sends $c = F(r, k_e^B)$ to B .
 - B decrypts c , and sends back $r' = F(c, k_d^B)$ to A .
 - If $r' = r$, then B is authentic.

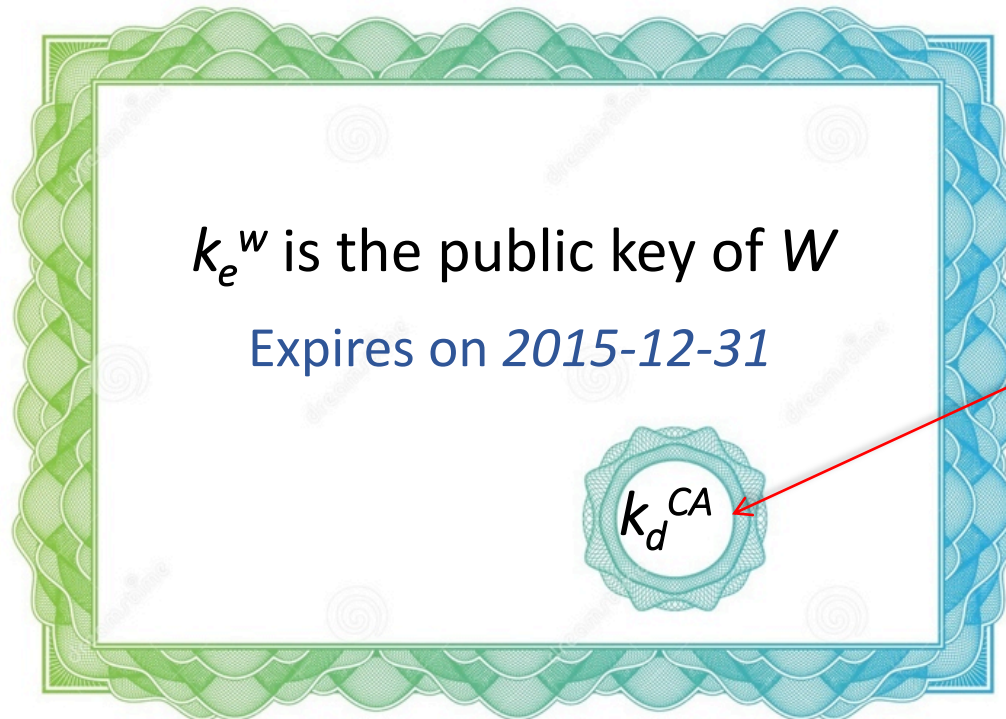
Asymmetric-Key Cryptography Applications

- Integrity?
 - Given a checksum h , A signs h with its private key k_d^A , and sends $h' = F^{-1}(h, k_d^A)$ to B .
 - B gets h' and obtains the original checksum h by using the public key of A , k_e^A . Then $h = F(h', k_e^A)$.
- How could we make sure that the public key is really the public key of the party we want to communicate in the first place?
 - Certificate Authorities

Certificate Authorities

- Allow us to trust the w website and its public key k_e^w
- The CA emits a **certificate** (with the public key k_e^w) for the w website, and **signs** it with its own private key k_d^{CA}
- Our browsers contain a list of trusted CA s with their public keys k_e^{CA}

Website W
introduces itself
with this certificate
when we start a
secure connection



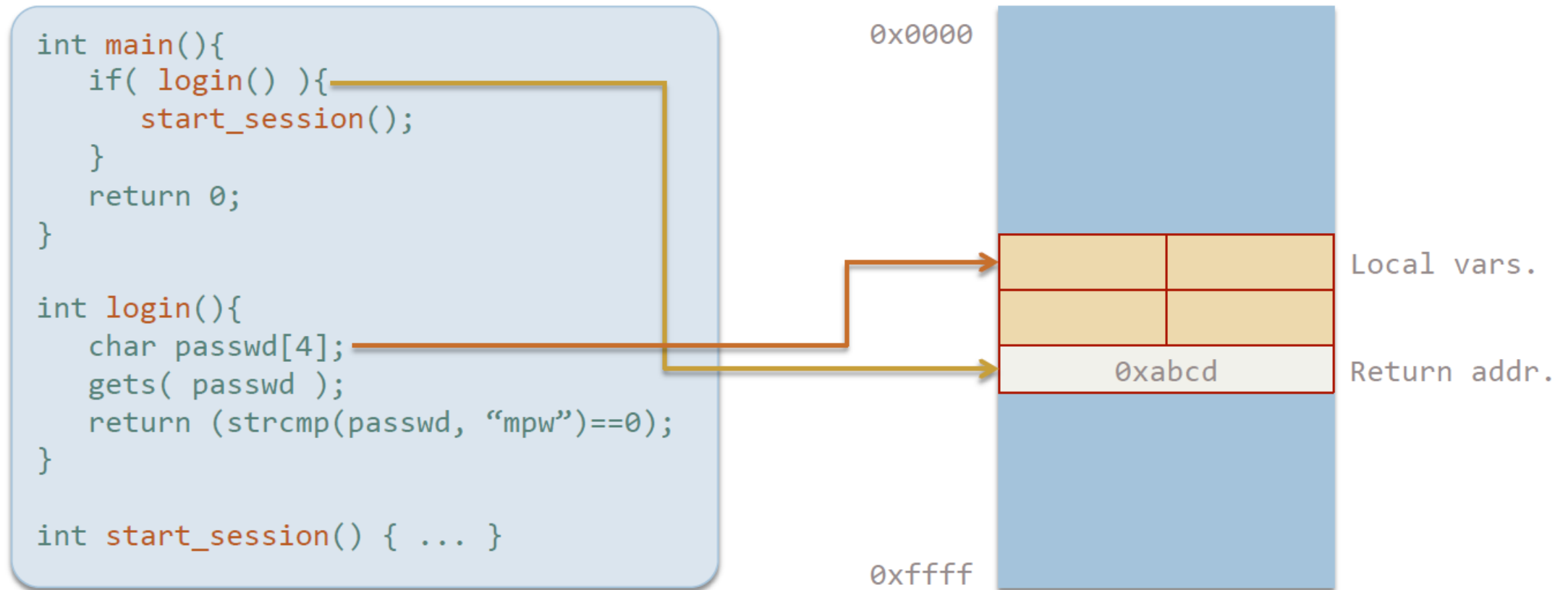
The certificate has
been signed a CA
with its private key.

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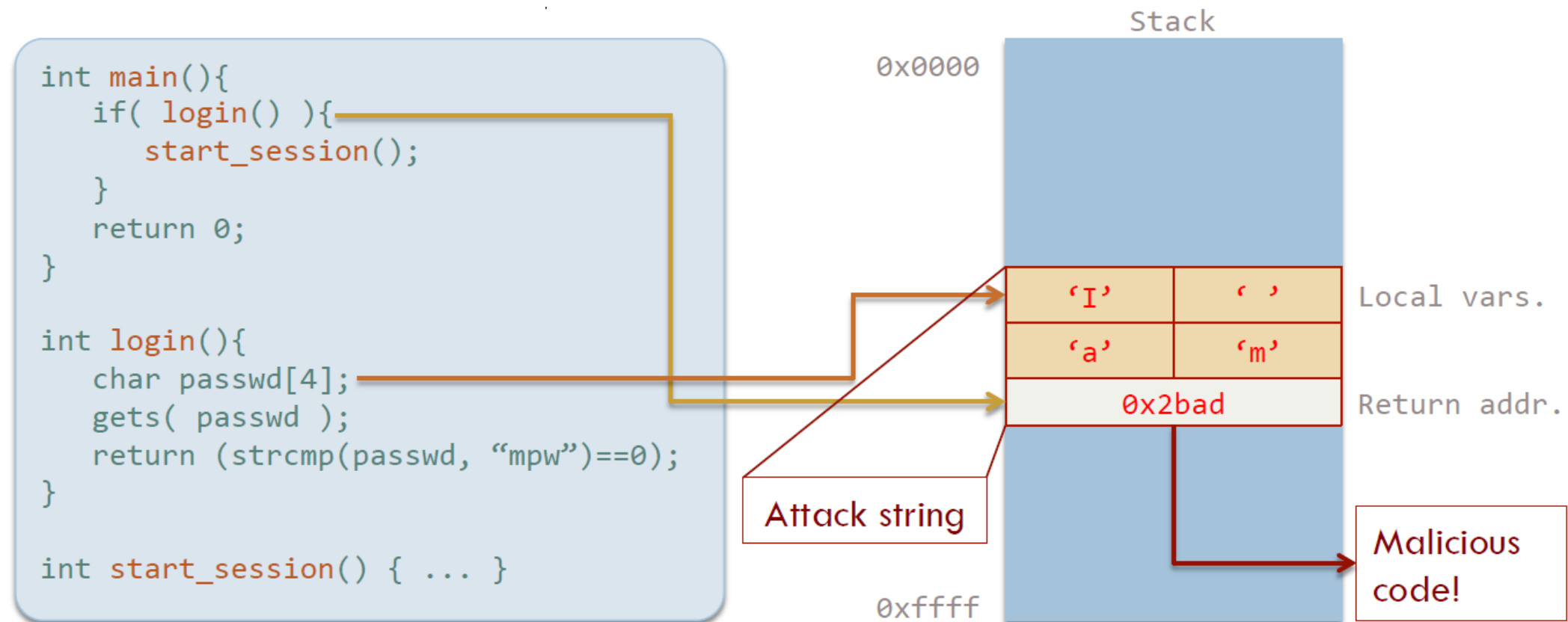
Common Vulnerabilities

- Buffer overflow



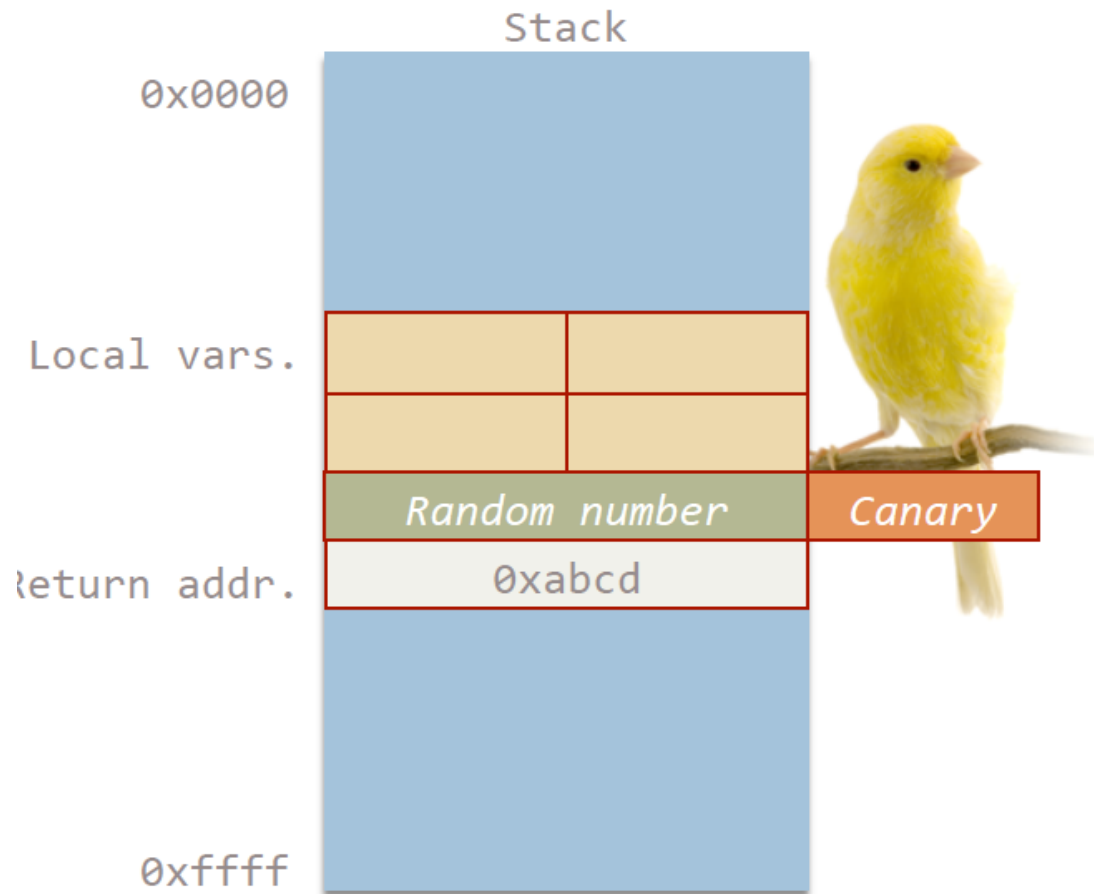
Common Vulnerabilities

- Buffer overflow



Common Vulnerabilities

- Buffer Overflow (Solution)



- Activate a “**Stack Guard**” to insert **Canary**.
- Give minimum privileges to a program
- **Never trust user input**: don't use unsafe string functions in C (like strcpy, gets, strcat, sprintf, etc.)

Common Vulnerabilities

- Client-State Manipulation



- Don't store **sensitive** information in the client (only a Session ID)
- **Encrypt a checksum** (using a **signature** prior storing in the client).
- Attach either an **expiration date** or the Session ID to client's state.

Common Vulnerabilities

- SQL Injection

```
"SELECT * FROM product  
WHERE id = " + user_input + ";"
```

123; DROP TABLE product;



- Create users with **minimum privileges**.
- Use **prepared statements**
- Encrypt sensitive data in DBMS
- Don't trust user input!

Common Vulnerabilities

- SQL Prepared Statement

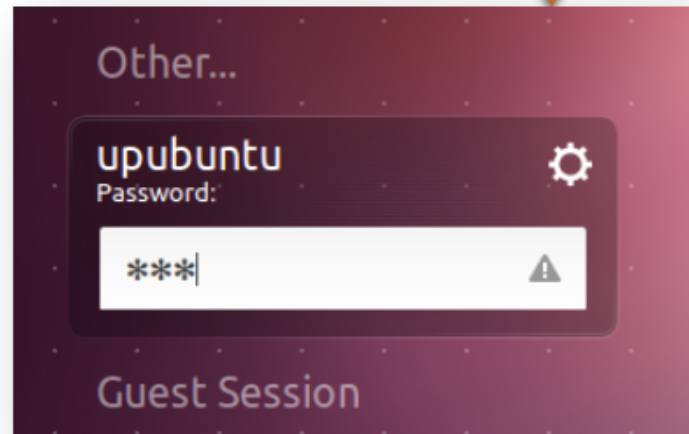
```
PreparedStatement s = db.prepareStatement("SELECT * from  
                                     Product WHERE id = ?");  
s.setInt(1, Integer.parseInt(user_input));  
ResultSet rs = s.executeQuery();
```

Common Vulnerabilities

- Command Injection

```
system( "cp f1.dat $user_input " );
```

f2.dat; rm /etc/passwd



- Don't use the System command. Use Runtime.exec() instead.
- “**Taint**” variables.
- Give **minimum privileges** to your application.

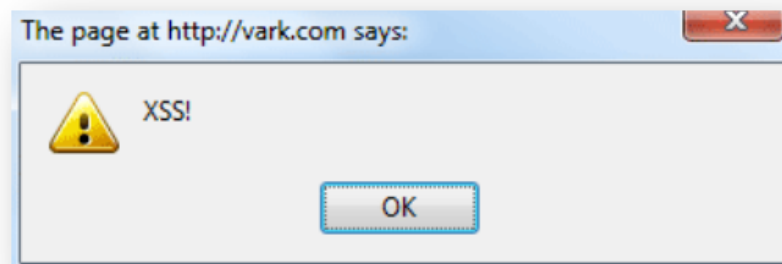
Common Vulnerabilities

- Cross-Site Scripting (XSS)

```
<body>
  Welcome to $user_input$ website
</body>
```



You're doomed `<script>hacked();</script>`

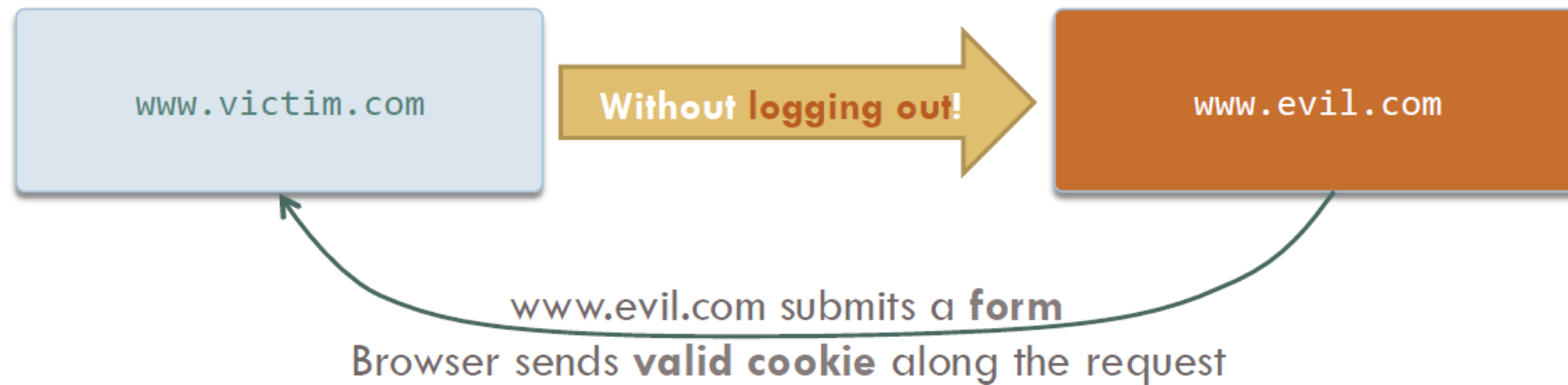


- Why is it called Cross-Site?
- It's very difficult to protect against it – we want to allow users insert HTML
- 2 options: white listing, black listing. *Which one is better?*

```
<div fu="alert('XSS!');" STYLE="background-image: url(javascript:eval(this.fu))">
```


Common Vulnerabilities

- Cross-Site Request Forgery (XSRF)



- Always logout!
- On server side, use an **action-token**.

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Project 5

- Part A - Testing Performance of Server
 - Learn to use Locust
 - Write some test cases for Locust
 - Find the max user number that the Server could handle
- Part B – Apache Spark
 - Learn to use Spark (shell)
 - Learn some basic Scala
 - Write a 10-line code in Scala to solve a problem.