

Tutorial 7 Practice Problems SOLUTIONS

Discrete Programming and Formulation Practice

Here for your own benefit and practice (best to do it individually)

Recommended completion: Week 08.

Grading: 0% (Practice for assignments and tests)

Problem 1 – Bread, Math, and Beyond

You are the manager of a small bakery commonly used as a study area for first year calculus called “Bread, Math and Beyond.” You are considering the upcoming week in which you want to determine how many batches of your products to make based on some prior demand data. For the sake of argument, let’s say you can make six products: bread (BRD), cookies (CKS), Rolls (ROL), muffins (MUF), crescents (CRE), and scones (SCO). The demographic data that you anticipate (maximum demands and profit per batch) for the coming week are reported in the table below. Moreover, the mixing time (MIX), baking time (BAK) and packing time (PAK) required for each batch is reported in the same table.

Product	Max Demand (Batch)	Profit (\$/Batch)	Mixing (min/Batch)	Baking (min/Batch)	Packing (min/Batch)
BRD	20	12	25	14	6
CKS	60	16	28	26	9
ROL	25	13	16	16	7
MUF	40	19	35	10	8
CRE	75	9	16	24	4
SCO	35	20	22	21	9

There are several industrial stand mixers ovens available. You may assume that all tasks may be done **at the same time**. Based on the capacities of these units you have:

- **4800 minutes** of mixing time available per week.
- **4800 minutes** of baking time per week.
- You are the only employee working in the store and thus you have **1200 minutes of packing time** available per week.

Questions

1. Formulate this as an LP and enter it into GAMS as a starting point. We are going to introduce some binary variables as we go that will affect our solution. Run your code and **verify** that the optimal profit is \$2644 per week.

SETS

$i \triangleq$ products (BRD, CKS, ROL, MUF, CRE, SCO)

$j \triangleq$ processes (MIX, BAKE, PACK)

PARAMETERS

$P_i \triangleq$ profit earned from selling product i

$D_i \triangleq$ maximum demand for product i

$T_{i,j} \triangleq$ time required to complete process j for product i

VARIABLES

$\phi \triangleq$ objective function variable: maximize profit

$x_i \triangleq$ number of batches of product i to make

$$\begin{aligned} \max_x \phi &= \sum_i P_i \cdot x_i \\ \text{s.t.} \\ x_i &\leq D_i & (\forall i) \\ \sum_i T_{i,j} \cdot x_i &\leq A_j & (\forall j) \\ x_i &\geq 0 & (\forall i) \end{aligned}$$

The GAMS code and solution report for this LP is given below.

```
SETS
i  products / BRD, CKS, ROL, MUF, CRE, SCO /
j  processes / MIX, BAKE, PACK / ;

PARAMETERS
D(i) demand (in batches)
/ BRD 20
  CKS 60
  ROL 25
  MUF 40
  CRE 75
  SCO 35 /

P(i) profit (in $ per batch)
/ BRD 12
  CKS 16
  ROL 13
  MUF 19
  CRE 9
  SCO 20 /

A(j) time available for process j
/ MIX 4800
  BAKE 4800
  PACK 1200 / ;
```

```

TABLE      T(i,j)  time requirements of each process j for product i
            MIX      BAKE      PACK
BRD        25       14        6
CKS        28       26        9
ROL        16       16        7
MUF        35       10        8
CRE        16       24        4
SCO        22       21        9 ;

```

VARIABLES

```

PHI         objective: maximize profit
x(i)        number of batches of product i to make
POSITIVE VARIABLES x(i) ;

```

EQUATIONS

PROFIT

DEMAND

TIME ;

```

PROFIT..    PHI =E= SUM(i, P(i)*x(i)) ;
DEMAND(i).. x(i) =L= D(i) ;
TIME(j)..   SUM(i,T(i,j)*x(i)) =L= A(j) ;

```

MODEL BAKERY /ALL/;

OPTION LP = CPLEX;

SOLVE BAKERY MAXIMIZING PHI USING LP;

Optimal solution found

Objective: 2644.285714

		LOWER	LEVEL	UPPER	MARGINAL
---- EQU PROFIT		.	.	.	1.0000
---- EQU DEMAND					
	LOWER	LEVEL	UPPER	MARGINAL	
BRD	-INF	20.0000	20.0000	0.8571	
CKS	-INF	.	60.0000	.	
ROL	-INF	20.7143	25.0000	.	
MUF	-INF	40.0000	40.0000	4.1429	
CRE	-INF	75.0000	75.0000	1.5714	
SCO	-INF	35.0000	35.0000	3.2857	
---- EQU TIME					
	LOWER	LEVEL	UPPER	MARGINAL	
MIX	-INF	4201.4286	4800.0000	.	
BAKE	-INF	3546.4286	4800.0000	.	
PACK	-INF	1200.0000	1200.0000	1.8571	
---- VAR PHI					
	LOWER	LEVEL	UPPER	MARGINAL	
PHI objective: maximize profit	-INF	2644.2857	+INF	.	
---- VAR x number of batches of product i to make					
	LOWER	LEVEL	UPPER	MARGINAL	
BRD	.	20.0000	+INF	.	
CKS	.	.	+INF	-0.7143	
ROL	.	20.7143	+INF	.	
MUF	.	40.0000	+INF	.	

CRE	.	75.0000	+INF	.
SCO	.	35.0000	+INF	.

2. OK, let's start to add some yes/no decisions in here. First, consider the case where, to avoid complications and to limit your orders, you are limited to producing only 4 different products per week. Add the constraint and variables that will make this happen. Add the binary variables and constraints to your GAMS code by declaring the binary as `BINARY VARIABLE`. You must change the solve statement to solve using the `MIP` solver (NOT the `LP` solver). The MIP default is CPLEX, so you will not need to change that option. Before running your code, ask yourself: is the profit *guaranteed* to be worse, better, or the same? Verify your result by running the code.

Add the binary variable $y_i \in \{0,1\}$ to represent whether or not product i is made. The new problem formulation is therefore

$$\begin{aligned}
 \max_x \phi &= \sum_i P_i \cdot x_i \\
 \text{s.t.} \\
 x_i &\leq D_i \cdot y_i & (\forall i) \\
 \sum_i T_{i,j} \cdot x_i &\leq A_j & (\forall j) \\
 \sum_i y_i &\leq 4 \\
 x_i &\geq 0 & (\forall i) \\
 y_i &\in \{0,1\} & (\forall i)
 \end{aligned}$$

Additions and changes to the GAMS code for the base-case LP in Part 1 are given in red. All previously defines sets and parameters remain the same.

```

SCALARS
MAXPROD    maximum number of products the bakery can make    / 4 / ;

VARIABLES
PHI        objective: maximize profit
x(i)       number of batches of product i to make
y(i)       whether or not product i is made ;
POSITIVE VARIABLES x(i) ;
BINARY VARIABLES y(i) ;

EQUATIONS
PROFIT
DEMAND
TIME
PRODLIM ;

PROFIT..    PHI =E= SUM(i, P(i)*x(i))    ;
DEMAND(i).. x(i) =L= D(i)*y(i)          ;
TIME(j)..   SUM(i, T(i,j)*x(i)) =L= A(j) ;
PRODLIM..   SUM(i, y(i)) =L= MAXPROD    ;

MODEL BAKERY /ALL/;
OPTION MIP = CPLEX;

SOLVE BAKERY MAXIMIZING PHI USING MIP;

```

Proven optimal solution				
MIP Solution:		2606.111111	(4 iterations, 0 nodes)	
Final Solve:		2606.111111	(1 iterations)	
---- EQU DEMAND				
	LOWER	LEVEL	UPPER	MARGINAL
BRD	-INF	.	.	1.3333
CKS	-INF	-30.5556	.	.
ROL	-INF	.	.	0.5556
MUF	-INF	.	.	4.7778
CRE	-INF	.	.	1.8889
SCO	-INF	.	.	4.0000
---- EQU TIME				
	LOWER	LEVEL	UPPER	MARGINAL
MIX	-INF	4194.4444	4800.0000	.
BAKE	-INF	3700.5556	4800.0000	.
PACK	-INF	1200.0000	1200.0000	1.7778
---- EQU PRODLIM				
	LOWER	LEVEL	UPPER	MARGINAL
	-INF	4.0000	4.0000	.
---- VAR PHI				
	LOWER	LEVEL	UPPER	MARGINAL
PHI	-INF	2606.1111	+INF	.
PHI objective: maximize profit				
---- VAR x number of batches of product i to make				
	LOWER	LEVEL	UPPER	MARGINAL
BRD	.	.	+INF	.
CKS	.	29.4444	+INF	.
ROL	.	.	+INF	.
MUF	.	40.0000	+INF	.
CRE	.	75.0000	+INF	.
SCO	.	35.0000	+INF	.
---- VAR y whether or not product i is made				
	LOWER	LEVEL	UPPER	MARGINAL
BRD	.	.	1.0000	26.6667
CKS	.	1.0000	1.0000	EPS
ROL	.	.	1.0000	13.8889
MUF	.	1.0000	1.0000	191.1111
CRE	.	1.0000	1.0000	141.6667
SCO	.	1.0000	1.0000	140.0000

3. Repeat the above but do it for 3 maximum products and 2 maximum products. Do the active constraints change significantly as you reduce the number of products you are making? What about the combination of products you are choosing?

The only thing that changes in the GAMS codes is the value of parameter MAXPROD.

GAMS solution for 3 maximum products:

Proven optimal solution				
MIP Solution:		2420.000000	(4 iterations, 0 nodes)	
Final Solve:		2420.000000	(0 iterations)	
----- EQU DEMAND				
	LOWER	LEVEL	UPPER	MARGINAL
BRD	-INF	.	.	12.0000

CKS	-INF	.	.	16.0000	
ROL	-INF	.	.	13.0000	
MUF	-INF	.	.	19.0000	
CRE	-INF	.	.	9.0000	
SCO	-INF	.	.	20.0000	
---- EQU TIME					
	LOWER	LEVEL	UPPER	MARGINAL	
MIX	-INF	3850.0000	4800.0000	.	
BAKE	-INF	2695.0000	4800.0000	.	
PACK	-INF	1175.0000	1200.0000	.	
---- EQU PRODLIM					
	LOWER	LEVEL	UPPER	MARGINAL	
	-INF	3.0000	3.0000	.	
---- VAR PHI					
	LOWER	LEVEL	UPPER	MARGINAL	
	-INF	2420.0000	+INF	.	
PHI objective: maximize profit					
---- VAR x number of batches of product i to make					
	LOWER	LEVEL	UPPER	MARGINAL	
BRD	.	.	+INF	.	
CKS	.	60.0000	+INF	.	
ROL	.	.	+INF	.	
MUF	.	40.0000	+INF	.	
CRE	.	.	+INF	.	
SCO	.	35.0000	+INF	.	
---- VAR y whether or not product i is made					
	LOWER	LEVEL	UPPER	MARGINAL	
BRD	.	.	1.0000	240.0000	
CKS	.	1.0000	1.0000	960.0000	
ROL	.	.	1.0000	325.0000	
MUF	.	1.0000	1.0000	760.0000	
CRE	.	.	1.0000	675.0000	
SCO	.	1.0000	1.0000	700.0000	

GAMS solution for 2 maximum products:

Proven optimal solution				
MIP Solution:	1720.000000	(3 iterations, 0 nodes)		
Final Solve:	1720.000000	(0 iterations)		
---- EQU DEMAND				
	LOWER	LEVEL	UPPER	MARGINAL
BRD	-INF	.	.	12.0000
CKS	-INF	.	.	16.0000
ROL	-INF	.	.	13.0000
MUF	-INF	.	.	19.0000
CRE	-INF	.	.	9.0000
SCO	-INF	.	.	20.0000
---- EQU TIME				
	LOWER	LEVEL	UPPER	MARGINAL
MIX	-INF	3080.0000	4800.0000	.
BAKE	-INF	1960.0000	4800.0000	.
PACK	-INF	860.0000	1200.0000	.
	LOWER	LEVEL	UPPER	MARGINAL

----	EQU	PRODLIM	-INF	2.0000	2.0000	.
----	VAR	PHI	LOWER	LEVEL	UPPER	MARGINAL
		PHI	-INF	1720.0000	+INF	.
		objective: maximize profit				
----	VAR	x	number of batches of product i to make			
			LOWER	LEVEL	UPPER	MARGINAL
BRD			.	.	+INF	.
CKS			.	60.0000	+INF	.
ROL			.	.	+INF	.
MUF			.	40.0000	+INF	.
CRE			.	.	+INF	.
SCO			.	.	+INF	.
----	VAR	y	whether or not product i is made			
			LOWER	LEVEL	UPPER	MARGINAL
BRD			.	.	1.0000	240.0000
CKS			.	1.0000	1.0000	960.0000
ROL			.	.	1.0000	325.0000
MUF			.	1.0000	1.0000	760.0000
CRE			.	.	1.0000	675.0000
SCO			.	.	1.0000	700.0000

4. In this scenario, we can make anything we want, BUT we are obliged to make 20 batches of cookies we promised to a local retailer ([Café Oranje](#)). Reformulate your program to account for this new constraint and re-solve the problem.

$$\begin{aligned}
 \max_x \phi &= \sum_i P_i \cdot x_i \\
 \text{s.t.} \\
 x_i &\leq D_i \cdot y_i & (\forall i) \\
 \sum_i T_{i,j} \cdot x_i &\leq A_j & (\forall j) \\
 \sum_i y_i &\leq 4 \\
 x_{CKS} &\geq 20 \\
 x_i &\geq 0 & (\forall i) \\
 y_i &\in \{0,1\} & (\forall i)
 \end{aligned}$$

Additions and changes to the GAMS code for the base-case MILP in Part 2 are given in red. All previously defines sets and parameters remain the same.

```

SCALARS
MAXPROD    maximum number of products the bakery can make    / 6 / ;

VARIABLES
PHI        objective: maximize profit
x(i)       number of batches of product i to make
y(i)       whether or not product i is made ;
POSITIVE VARIABLES x(i) ;
BINARY VARIABLES y(i) ;

```

```

EQUATIONS
PROFIT
DEMAND
COOKIES
TIME
PRODLIM ;

PROFIT..      PHI =E= SUM(i, P(i)*x(i))      ;
DEMAND(i)..   x(i) =L= D(i)*y(i)              ;
COOKIES..     x('CKS') =G= 20                ;
TIME(j)..     SUM(i,T(i,j)*x(i)) =L= A(j)     ;
PRODLIM..     SUM(i,y(i)) =L= MAXPROD         ;

MODEL BAKERY /ALL/;
OPTION MIP = CPLEX;

SOLVE BAKERY MAXIMIZING PHI USING MIP;

```

```

Proven optimal solution
MIP Solution:      2625.000000      (1 iterations, 0 nodes)
Final Solve:       2625.000000      (1 iterations)

---- EQU DEMAND
      LOWER      LEVEL      UPPER      MARGINAL
BRD      -INF      -5.8333      .          .
CKS      -INF      -40.0000      .          .
ROL      -INF      -25.0000      .          .
MUF      -INF      .          .          3.0000
CRE      -INF      .          .          1.0000
SCO      -INF      .          .          2.0000

      LOWER      LEVEL      UPPER      MARGINAL
---- EQU COOKIES      20.0000      20.0000      +INF      -2.0000

---- EQU TIME
      LOWER      LEVEL      UPPER      MARGINAL
MIX      -INF      4284.1667      4800.0000      .
BAKE      -INF      3653.3333      4800.0000      .
PACK      -INF      1200.0000      1200.0000      2.0000

      LOWER      LEVEL      UPPER      MARGINAL
---- EQU PRODLIM      -INF      6.0000      6.0000      .

      LOWER      LEVEL      UPPER      MARGINAL
---- VAR PHI      -INF      2625.0000      +INF      .
      PHI  objective: maximize profit

---- VAR x  number of batches of product i to make
      LOWER      LEVEL      UPPER      MARGINAL
BRD      .          14.1667      +INF      .
CKS      .          20.0000      +INF      .
ROL      .          .          +INF      -1.0000
MUF      .          40.0000      +INF      .
CRE      .          75.0000      +INF      .
SCO      .          35.0000      +INF      .

---- VAR y  whether or not product i is made

```


	LOWER	LEVEL	UPPER	MARGINAL
BRD	.	1.0000	1.0000	EPS
CKS	.	1.0000	1.0000	EPS
ROL	.	1.0000	1.0000	EPS
MUF	.	1.0000	1.0000	120.0000
CRE	.	1.0000	1.0000	75.0000
SCO	.	1.0000	1.0000	70.0000

5. Now consider the case where we can make anything we want again with no cookies required. HOWEVER, there is a fixed charge for making any one product of \$300 (so \$300 per product we choose to make in ANY quantity), which is the delivery charge for the necessary ingredients (also accounts for you investing your own time to unload and store all the ingredients). Reformulate your problem again to account for this new cost. Re-solve it using GAMS to report the new optimum.

$$\begin{aligned}
 \max_x \phi &= \sum_i (P_i \cdot x_i - 300y_i) \\
 \text{s.t.} \\
 x_i &\leq D_i \cdot y_i & (\forall i) \\
 \sum_i T_{i,j} \cdot x_i &\leq A_j & (\forall j) \\
 x_i &\geq 0 & (\forall i) \\
 y_i &\in \{0,1\} & (\forall i)
 \end{aligned}$$

Additions and changes to the GAMS code for the base-case MILP in Part 2 are given in red. All previously defines sets and parameters remain the same.

```

SCALARS
MAXPROD    maximum number of products the bakery can make    / 6 /
F          fixed cost of making any one product              / 300 / ;

VARIABLES
PHI        objective: maximize profit
x(i)       number of batches of product i to make
y(i)       whether or not product i is made ;
POSITIVE VARIABLES x(i) ;
BINARY VARIABLES y(i) ;

EQUATIONS
PROFIT
DEMAND
TIME
PRODLIM ;

PROFIT..    PHI =E= SUM(i, P(i)*x(i) - F*y(i)) ;
DEMAND(i).. x(i) =L= D(i)*y(i) ;
TIME(j)..   SUM(i, T(i,j)*x(i)) =L= A(j) ;
PRODLIM..   SUM(i, y(i)) =L= MAXPROD ;

MODEL BAKERY /ALL/;
OPTION MIP = CPLEX;

```

Proven optimal solution

MIP Solution: 1520.000000 (3 iterations, 0 nodes)

Final Solve: 1520.000000 (0 iterations)

Best possible: 1520.000000

Absolute gap: 0.000000

Relative gap: 0.000000

---- EQU DEMAND

	LOWER	LEVEL	UPPER	MARGINAL
BRD	-INF	.	.	12.0000
CKS	-INF	.	.	16.0000
ROL	-INF	.	.	13.0000
MUF	-INF	.	.	19.0000
CRE	-INF	.	.	9.0000
SCO	-INF	.	.	20.0000

---- EQU TIME

	LOWER	LEVEL	UPPER	MARGINAL
MIX	-INF	3850.0000	4800.0000	.
BAKE	-INF	2695.0000	4800.0000	.
PACK	-INF	1175.0000	1200.0000	.

	LOWER	LEVEL	UPPER	MARGINAL
---- EQU PRODLIM	-INF	3.0000	6.0000	.

	LOWER	LEVEL	UPPER	MARGINAL
---- VAR PHI	-INF	1520.0000	+INF	.
PHI objective: maximize profit				

---- VAR x number of batches of product i to make

	LOWER	LEVEL	UPPER	MARGINAL
BRD	.	.	+INF	.
CKS	.	60.0000	+INF	.
ROL	.	.	+INF	.
MUF	.	40.0000	+INF	.
CRE	.	.	+INF	.
SCO	.	35.0000	+INF	.

---- VAR y whether or not product i is made

	LOWER	LEVEL	UPPER	MARGINAL
BRD	.	.	1.0000	-60.0000
CKS	.	1.0000	1.0000	660.0000
ROL	.	.	1.0000	25.0000
MUF	.	1.0000	1.0000	460.0000
CRE	.	.	1.0000	375.0000
SCO	.	1.0000	1.0000	400.0000

6. Return to the base case again. Argue convincingly that if I were to impose on you the requirement "if you make any of one product, you must make at least 10 batches of it", the optimum will either go down or stay the same in the best possible case. This would be true if, for instance, ingredients come in minimum shipments of 10 batches-worth.

If we are required to make at least 10 batches of any product we choose to make, imposing this constraint essentially reduces the size of the feasible region (tightening constraints) and thus will either make the optimal objective function value worse, or unchanged at best. In this case, the optimum will be unchanged because we can see in the base case solution that we are already making more than 10 batches for all products that are being made.

7. Re-solve the problem with this added constraint to prove your point.

$$\begin{aligned}
 \max_x \phi &= \sum_i P_i \cdot x_i \\
 \text{s.t.} \\
 x_i &\leq D_i \cdot y_i & (\forall i) \\
 x_i &\geq 10 \cdot y_i & (\forall i) \\
 \sum_i T_{i,j} \cdot x_i &\leq A_j & (\forall j) \\
 x_i &\geq 0 & (\forall i) \\
 y_i &\in \{0,1\} & (\forall i)
 \end{aligned}$$

Additions and changes to the GAMS code for the base-case MILP in Part 2 are given in red. All previously defines sets and parameters remain the same.

```

SCALARS
MAXPROD    maximum number of products the bakery can make    / 6 /
MINBATCH   minimum number of batches that must be made       / 10 / ;

VARIABLES
PHI        objective: maximize profit
x(i)       number of batches of product i to make
y(i)       whether or not product i is made ;
POSITIVE VARIABLES x(i) ;
BINARY VARIABLES y(i) ;

EQUATIONS
PROFIT
DEMAND
TIME
PRODLIM
BATCHES ;

PROFIT..    PHI =E= SUM(i, P(i)*x(i))    ;
DEMAND(i).. x(i) =L= D(i)*y(i)          ;
TIME(j)..   SUM(i,T(i,j)*x(i)) =L= A(j)  ;
PRODLIM..    SUM(i,y(i)) =L= MAXPROD     ;
BATCHES(i).. x(i) =G= MINBATCH*y(i)      ;

MODEL BAKERY /ALL/;
OPTION MIP = CPLEX;

SOLVE BAKERY MAXIMIZING PHI USING MIP;

```

Proven optimal solution

MIP Solution: 2644.285714 (1 iterations, 0 nodes)				
Final Solve: 2644.285714 (1 iterations)				
---- EQU DEMAND				
	LOWER	LEVEL	UPPER	MARGINAL
BRD	-INF	.	.	0.8571
CKS	-INF	.	.	.
ROL	-INF	-4.2857	.	.
MUF	-INF	.	.	4.1429
CRE	-INF	.	.	1.5714
SCO	-INF	.	.	3.2857
---- EQU TIME				
	LOWER	LEVEL	UPPER	MARGINAL
MIX	-INF	4201.4286	4800.0000	.
BAKE	-INF	3546.4286	4800.0000	.
PACK	-INF	1200.0000	1200.0000	1.8571
---- EQU PRODLIM				
	LOWER	LEVEL	UPPER	MARGINAL
	-INF	5.0000	6.0000	.
---- EQU BATCHES				
	LOWER	LEVEL	UPPER	MARGINAL
BRD	.	10.0000	+INF	.
CKS	.	.	+INF	.
ROL	.	10.7143	+INF	.
MUF	.	30.0000	+INF	.
CRE	.	65.0000	+INF	.
SCO	.	25.0000	+INF	.
---- VAR PHI				
	LOWER	LEVEL	UPPER	MARGINAL
PHI	-INF	2644.2857	+INF	.
PHI objective: maximize profit				
---- VAR x number of batches of product i to make				
	LOWER	LEVEL	UPPER	MARGINAL
BRD	.	20.0000	+INF	.
CKS	.	.	+INF	-0.7143
ROL	.	20.7143	+INF	.
MUF	.	40.0000	+INF	.
CRE	.	75.0000	+INF	.
SCO	.	35.0000	+INF	.
---- VAR y whether or not product i is made				
	LOWER	LEVEL	UPPER	MARGINAL
BRD	.	1.0000	1.0000	17.1429
CKS	.	.	1.0000	EPS
ROL	.	1.0000	1.0000	EPS
MUF	.	1.0000	1.0000	165.7143
CRE	.	1.0000	1.0000	117.8571
SCO	.	1.0000	1.0000	115.0000

8. Re-formulate and re-solve the optimization so that if I want to make any one product, I am obliged to make 30 batches. You may adjust the demand data such that any product with a demand < 30 can be equal to 30 instead. Comment on your findings.

$$\begin{aligned}
 \max_x \phi &= \sum_i P_i \cdot x_i \\
 \text{s.t.} \\
 x_i &\leq D_i \cdot y_i & (\forall i) \\
 x_i &\geq 30 \cdot y_i & (\forall i) \\
 \sum_i T_{i,j} \cdot x_i &\leq A_j & (\forall j) \\
 x_i &\geq 0 & (\forall i) \\
 y_i &\in \{0,1\} & (\forall i)
 \end{aligned}$$

Additions and changes to the GAMS code for the MILP in the previous question are given in red. All previously defines sets and parameters remain the same, except the values for the maximum demand parameter D_i were increased to 30 for products $i = \text{BRD}$ and $i = \text{ROL}$.

```

SCALARS
MAXPROD    maximum number of products the bakery can make    / 6 /
MINBATCH   minimum number of batches that must be made      / 30 / ;

VARIABLES
PHI         objective: maximize profit
x(i)        number of batches of product i to make
y(i)        whether or not product i is made ;
POSITIVE VARIABLES x(i) ;
BINARY VARIABLES y(i) ;

EQUATIONS
PROFIT
DEMAND
TIME
PRODLIM
BATCHES ;

PROFIT..    PHI =E= SUM(i, P(i)*x(i)) ;
DEMAND(i).. x(i) =L= D(i)*y(i) ;
TIME(j)..   SUM(i, T(i,j)*x(i)) =L= A(j) ;
PRODLIM..   SUM(i, y(i)) =L= MAXPROD ;
BATCHES(i).. x(i) =G= MINBATCH*y(i) ;

MODEL BAKERY /ALL/;
OPTION MIP = CPLEX;

SOLVE BAKERY MAXIMIZING PHI USING MIP;

```

```

Proven optimal solution
MIP Solution:          2605.000000      (20 iterations, 0 nodes)
Final Solve:          2605.000000      (0 iterations)

---- EQU DEMAND

```

	LOWER	LEVEL	UPPER	MARGINAL
BRD	-INF	.	.	.
CKS	-INF	.	.	.
ROL	-INF	.	.	.
MUF	-INF	.	.	1.0000
CRE	-INF	-20.0000	.	.
SCO	-INF	-5.0000	.	.

---- EQU TIME					
	LOWER	LEVEL	UPPER	MARGINAL	
MIX	-INF	4170.0000	4800.0000	.	
BAKE	-INF	3250.0000	4800.0000	.	
PACK	-INF	1200.0000	1200.0000	2.2500	
---- EQU BATCHES					
	LOWER	LEVEL	UPPER	MARGINAL	
BRD	.	.	+INF	-1.5000	
CKS	.	.	+INF	.	
ROL	.	.	+INF	-2.7500	
MUF	.	10.0000	+INF	.	
CRE	.	25.0000	+INF	.	
SCO	.	.	+INF	-0.2500	
---- VAR PHI					
		LOWER	LEVEL	UPPER	MARGINAL
		-INF	2605.0000	+INF	.
PHI objective: maximize profit					
---- VAR x number of batches of product i to make					
	LOWER	LEVEL	UPPER	MARGINAL	
BRD	.	30.0000	+INF	.	
CKS	.	.	+INF	-4.2500	
ROL	.	30.0000	+INF	.	
MUF	.	40.0000	+INF	.	
CRE	.	55.0000	+INF	.	
SCO	.	30.0000	+INF	.	
---- VAR y whether or not product i is made					
	LOWER	LEVEL	UPPER	MARGINAL	
BRD	.	1.0000	1.0000	-45.0000	
CKS	.	.	1.0000	EPS	
ROL	.	1.0000	1.0000	-82.5000	
MUF	.	1.0000	1.0000	40.0000	
CRE	.	1.0000	1.0000	EPS	
SCO	.	1.0000	1.0000	-7.5000	

Problem 2 – Optimization... IN SPAAAAAAACE!

The US space agency known as the National Aeronautics and Space Adventuring Helionauts (NASAH) typically has to face decisions around budgeting long-term missions and projects. For example, consider the case where they have funding available for a series of projects available over a five-stage period of time. Each project requires a certain amount of money in each stage, and each stage has a well-defined budget from the government. In the **table below** is a list of prospective projects and their anticipated cost at each stage of the budgeting plan. Furthermore, in the table is the expected “benefit” (an aggregate score based on likelihood of success and information gathered as a result of the mission) of each project and a list of projects that depend on or are mutually exclusive to it. Note that if you choose a project in the table below you **must pay for all costs in each stage for that project**.

This type of problem can be defined as a “multi-dimensional knapsack” problem in which we are choosing which projects to take depending on capacity (*i.e.* budget) constraints for multiple stages. It is our objective to allocate research funds that will maximize the benefit of NASAH’s funding over the next five stages.

<i>i</i>	Mission	Budget Requirement (\$billions)					Value	Not With	Depends On
		Stage 1	Stage 2	Stage 3	Stage 4	Stage 5			
1	Comms Satellite	6	-	-	-	-	200	-	-
2	Orbital Microwave	2	3	-	-	-	3	-	-
3	IO Lander	3	5	-	-	-	20	-	-
4	Uranus Orbiter 2030	-	-	-	-	10	50	5	3
5	Uranus Orbiter 2040	-	5	8	-	-	70	4	3
6	Jupiter Probe	-	-	1	8	4	20	-	3
7	Saturn Probe	1	8	-	-	-	5	-	3
8	Infrared Imaging	-	-	-	5	-	10	11	-
9	Ground-Based SETI	4	5	-	-	-	200	13	-
10	Orbital Structures	-	8	4	-	-	150	-	-
11	Color Imaging	-	-	2	7	-	18	8	2
12	Medical Technology	5	7	-	-	-	8	-	-
13	Geosynchronous SETI	-	4	5	3	3	185	9	-
14	Jetpod Shielding	-	-	3	3	1	50	-	12
Budget		10	12	14	14	14			

Questions

1. Define the variables for this problem, clearly indicating if they are continuous or binary.
We can define the binary variable $y_i \in \{0,1\}$ to represent the choice of selecting whether or not mission i will be funded.
2. Define the objective for this problem. You may write this out in terms of hard-coded numbers or define new parameters as you see fit.
The objective for this problem is to select the projects to fund which maximize the overall value (benefit) while respecting the budget. The objective function can be defined mathematically as:

$$\phi = \sum_{i=1}^{14} V_i \cdot y_i$$

Where $i = 1 \dots 14$ is the set of available missions and V_i represents the expected value (benefit) associated with funding project i .

3. Define the constraint(s) that correspond(s) to the available budget for each stage in each for this problem. Try to be as concise as possible.

We can define the set $j = 1 \dots 5$ to be the project stages. With this, a constraint that ensures we do not exceed the maximum available budget for each project stage can be written as:

$$\sum_{i=1}^{14} C_{i,j} \cdot y_i \leq B_j \quad (\forall j)$$

Where $C_{i,j}$ represents the cost (in billions of dollars) of carrying out mission i in project stage j and B_j represents the total budget that is allocated for each stage j .

4. Code your formulation *so far* (that is, DO NOT include the mutual exclusivity or dependence constraints) in GAMS and find the solution.

Overall formulation:

$$\begin{aligned} \max_y \phi &= \sum_i V_i \cdot y_i \\ \text{s.t.} \\ \sum_i C_{i,j} y_i &\leq B_j \quad (\forall j) \\ y_i &\in \{0, 1\} \quad (\forall i) \end{aligned}$$

The GAMS code and solution are given below.

```
SETS
i   Missions to choose from / M1*M14 /
j   Stage of time period    / S1*S5 / ;

PARAMETERS
V(i)   Expected benefit of a mission
/ M1    200
  M2     3
  M3    20
  M4    50
  M5    70
  M6    20
  M7     5
  M8    10
  M9   200
  M10   150
  M11   18
  M12    8
  M13  185
  M14   50 /

B(j)   Total budget allocated for a stage($billions)
/ S1    10
  S2    12
  S3    14
  S4    14
  S5    14 / ;
```


TABLE C(i,j) Budget requirement for the missions(\$billions)

	S1	S2	S3	S4	S5
M1	6	0	0	0	0
M2	2	3	0	0	0
M3	3	5	0	0	0
M4	0	0	0	0	10
M5	0	5	8	0	0
M6	0	0	1	8	4
M7	1	8	0	0	0
M8	0	0	0	5	0
M9	4	5	0	0	0
M10	0	8	4	0	0
M11	0	0	2	7	0
M12	5	7	0	0	0
M13	0	4	5	3	3
M14	0	0	3	3	1 ;

VARIABLES

Z Objective function variable
y(i) Whether or not mission i will be funded ;
BINARY VARIABLES y(i) ;

EQUATIONS

OBJ Objective function
BUDGET Budget constraints ;

OBJ.. Z =E= SUM(i, V(i)*y(i)) ;
BUDGET(j).. SUM(i, C(i,j)*y(i)) =L= B(j) ;

Model SPACE /all/ ;
Solve SPACE MAXIMIZING Z USING MIP ;

Proven optimal solution

MIP Solution: 703.000000 (4 iterations, 0 nodes)
Final Solve: 703.000000 (0 iterations)

---- EQU BUDGET Budget constraints

	LOWER	LEVEL	UPPER	MARGINAL
S1	-INF	10.0000	10.0000	.
S2	-INF	9.0000	12.0000	.
S3	-INF	10.0000	14.0000	.
S4	-INF	13.0000	14.0000	.
S5	-INF	14.0000	14.0000	.

---- VAR Z	LOWER	LEVEL	UPPER	MARGINAL
	-INF	703.0000	+INF	.
Z Objective function variable				

---- VAR y Whether or not mission i will be funded

	LOWER	LEVEL	UPPER	MARGINAL
M1	.	1.0000	1.0000	200.0000
M2	.	.	1.0000	3.0000
M3	.	.	1.0000	20.0000
M4	.	1.0000	1.0000	50.0000
M5	.	.	1.0000	70.0000
M6	.	.	1.0000	20.0000
M7	.	.	1.0000	5.0000
M8	.	.	1.0000	10.0000
M9	.	1.0000	1.0000	200.0000
M10	.	.	1.0000	150.0000
M11	.	1.0000	1.0000	18.0000
M12	.	.	1.0000	8.0000

M13	.	1.0000	1.0000	185.0000
M14	.	1.0000	1.0000	50.0000

5. Define the mutual exclusivity constraint for this problem.

The mutual exclusivity constraints are based on the "Not With" column in the table. This column indicates that you cannot complete missions 4 and 5, 8 and 11, or 9 and 13 at the same time. These constraints can be written mathematically as:

$$y_4 + y_5 \leq 1$$

$$y_8 + y_{11} \leq 1$$

$$y_9 + y_{13} \leq 1$$

6. Add the mutual exclusivity constraints to GAMS and re-run the problem. Comment on the results and how adding the mutual exclusivity constraints affect your solution.

Only the EQUATIONS section of the GAMS code will change from adding the mutual exclusivity constraints. Additions to the base case are highlighted with red text.

```

EQUATIONS
OBJ          Objective function
BUDGET       Budget constraints
ME1          Mutual exclusivity constraint 1
ME2          Mutual exclusivity constraint 2
ME3          Mutual exclusivity constraint 3 ;

OBJ..        Z =E= SUM(i, V(i)*y(i)) ;
BUDGET(j)..  SUM(i, C(i,j)*y(i)) =L= B(j) ;
ME1..        y('M4') + y('M5') =L= 1 ;
ME2..        y('M8') + y('M11') =L= 1 ;
ME3..        y('M9') + y('M13') =L= 1 ;

```

```

Proven optimal solution
MIP Solution:          653.000000      (5 iterations, 0 nodes)
Final Solve:           653.000000      (0 iterations)

---- EQU BUDGET   Budget constraints
      LOWER      LEVEL      UPPER      MARGINAL
S1      -INF      6.0000      10.0000      .
S2      -INF      12.0000      12.0000      .
S3      -INF      14.0000      14.0000      .
S4      -INF      13.0000      14.0000      .
S5      -INF      14.0000      14.0000      .

      LOWER      LEVEL      UPPER      MARGINAL
---- EQU ME1      -INF      1.0000      1.0000      .
---- EQU ME2      -INF      1.0000      1.0000      .
---- EQU ME3      -INF      1.0000      1.0000      .
      ME1 Mutual exclusivity constraint 1
      ME2 Mutual exclusivity constraint 2
      ME3 Mutual exclusivity constraint 3

      LOWER      LEVEL      UPPER      MARGINAL
---- VAR Z      -INF      653.0000      +INF      .

```

Z Objective function variable				
---- VAR y Whether or not mission i will be funded				
	LOWER	LEVEL	UPPER	MARGINAL
M1	.	1.0000	1.0000	200.0000
M2	.	.	1.0000	3.0000
M3	.	.	1.0000	20.0000
M4	.	1.0000	1.0000	50.0000
M5	.	.	1.0000	70.0000
M6	.	.	1.0000	20.0000
M7	.	.	1.0000	5.0000
M8	.	.	1.0000	10.0000
M9	.	.	1.0000	200.0000
M10	.	1.0000	1.0000	150.0000
M11	.	1.0000	1.0000	18.0000
M12	.	.	1.0000	8.0000
M13	.	1.0000	1.0000	185.0000
M14	.	1.0000	1.0000	50.0000

7. Define the dependency constraints for this problem.

The dependency constraints are based on the "Depends On" column in the table, where it is indicated that: (a) missions 4, 5 and 6 depend on mission 3, (b) mission 11 depends on mission 2 and (c) mission 14 depends on completing mission 12. These dependency constraints can be written mathematically as:

$$y_4 \leq y_3$$

$$y_5 \leq y_3$$

$$y_6 \leq y_3$$

$$y_{11} \leq y_2$$

$$y_{14} \leq y_{12}$$

8. Add the dependency constraints to the **base case** in GAMS (that is, don't include the mutual exclusivity constraints) and re-run the problem. Comment on the results and how adding the dependence constraints affect your solution.

Only the EQUATIONS section of the GAMS code will change from adding the dependency constraints. Additions to the base case are highlighted with red text.

EQUATIONS	
OBJ	Objective function
BUDGET	Budget constraints
DC1	Dependency constraint 1
DC2	Dependency constraint 2
DC3	Dependency constraint 3
DC4	Dependency constraint 4
DC5	Dependency constraint 5 ;
OBJ..	Z =E= SUM(i, V(i)*y(i)) ;
BUDGET(j) ..	SUM(i, C(i,j)*y(i)) =L= B(j) ;
DC1..	y('M4') =L= y('M3') ;
DC2..	y('M5') =L= y('M3') ;
DC3..	y('M6') =L= y('M3') ;
DC4..	y('M11') =L= y('M2') ;
DC5..	y('M14') =L= y('M12') ;

Proven optimal solution					
MIP Solution:		595.000000	(6 iterations, 0 nodes)		
Final Solve:		595.000000	(0 iterations)		
---- EQU BUDGET Budget constraints					
	LOWER	LEVEL	UPPER	MARGINAL	
S1	-INF	10.0000	10.0000	.	
S2	-INF	9.0000	12.0000	.	
S3	-INF	5.0000	14.0000	.	
S4	-INF	8.0000	14.0000	.	
S5	-INF	3.0000	14.0000	.	
		LOWER	LEVEL	UPPER	MARGINAL
----	EQU DC1	-INF	.	.	.
----	EQU DC2	-INF	.	.	.
----	EQU DC3	-INF	.	.	.
----	EQU DC4	-INF	.	.	.
----	EQU DC5	-INF	.	.	.
DC1	Dependency constraint 1				
DC2	Dependency constraint 2				
DC3	Dependency constraint 3				
DC4	Dependency constraint 4				
DC5	Dependency constraint 5				
		LOWER	LEVEL	UPPER	MARGINAL
----	VAR Z	-INF	595.0000	+INF	.
	Z Objective function variable				
---- VAR y Whether or not mission i will be funded					
	LOWER	LEVEL	UPPER	MARGINAL	
M1	.	1.0000	1.0000	200.0000	
M2	.	.	1.0000	3.0000	
M3	.	.	1.0000	20.0000	
M4	.	.	1.0000	50.0000	
M5	.	.	1.0000	70.0000	
M6	.	.	1.0000	20.0000	
M7	.	.	1.0000	5.0000	
M8	.	1.0000	1.0000	10.0000	
M9	.	1.0000	1.0000	200.0000	
M10	.	.	1.0000	150.0000	
M11	.	.	1.0000	18.0000	
M12	.	.	1.0000	8.0000	
M13	.	1.0000	1.0000	185.0000	
M14	.	.	1.0000	50.0000	

9. Re-run your GAMS code with *both* of the mutual/dependence constraints included. Comment on *this* solution.

Only the EQUATIONS section of the GAMS code will change from adding the mutual exclusivity constraints. Additions to the base case are highlighted with red text.

EQUATIONS	
OBJ	Objective function
BUDGET	Budget constraints
ME1	Mutual exclusivity constraint 1
ME2	Mutual exclusivity constraint 2
ME3	Mutual exclusivity constraint 3
DC1	Dependency constraint 1
DC2	Dependency constraint 2
DC3	Dependency constraint 3
DC4	Dependency constraint 4

```

DC5          Dependency constraint 5 ;

OBJ..        Z =E= SUM(i, V(i)*y(i)) ;
BUDGET(j)..  SUM(i, C(i,j)*y(i)) =L= B(j) ;
ME1..        y('M4') + y('M5') =L= 1 ;
ME2..        y('M8') + y('M11') =L= 1 ;
ME3..        y('M9') + y('M13') =L= 1 ;
DC1..        y('M4') =L= y('M3') ;
DC2..        y('M5') =L= y('M3') ;
DC3..        y('M6') =L= y('M3') ;
DC4..        y('M11') =L= y('M2') ;
DC5..        y('M14') =L= y('M12') ;

```

Proven optimal solution

MIP Solution: 545.000000 (8 iterations, 0 nodes)

Final Solve: 545.000000 (0 iterations)

---- EQU BUDGET Budget constraints

	LOWER	LEVEL	UPPER	MARGINAL
S1	-INF	6.0000	10.0000	.
S2	-INF	12.0000	12.0000	.
S3	-INF	9.0000	14.0000	.
S4	-INF	8.0000	14.0000	.
S5	-INF	3.0000	14.0000	.

	LOWER	LEVEL	UPPER	MARGINAL
---- EQU ME1	-INF	.	1.0000	.
---- EQU ME2	-INF	1.0000	1.0000	.
---- EQU ME3	-INF	1.0000	1.0000	.
---- EQU DC1	-INF	.	.	.
---- EQU DC2	-INF	.	.	.
---- EQU DC3	-INF	.	.	.
---- EQU DC4	-INF	.	.	.
---- EQU DC5	-INF	.	.	.

ME1 Mutual exclusivity constraint 1
 ME2 Mutual exclusivity constraint 2
 ME3 Mutual exclusivity constraint 3
 DC1 Dependency constraint 1
 DC2 Dependency constraint 2
 DC3 Dependency constraint 3
 DC4 Dependency constraint 4
 DC5 Dependency constraint 5

	LOWER	LEVEL	UPPER	MARGINAL
---- VAR Z	-INF	545.0000	+INF	.
Z Objective function variable				

---- VAR y Whether or not mission i will be funded

	LOWER	LEVEL	UPPER	MARGINAL
M1	.	1.0000	1.0000	200.0000
M2	.	.	1.0000	3.0000
M3	.	.	1.0000	20.0000
M4	.	.	1.0000	50.0000
M5	.	.	1.0000	70.0000
M6	.	.	1.0000	20.0000
M7	.	.	1.0000	5.0000
M8	.	1.0000	1.0000	10.0000
M9	.	.	1.0000	200.0000
M10	.	1.0000	1.0000	150.0000
M11	.	.	1.0000	18.0000
M12	.	.	1.0000	8.0000
M13	.	1.0000	1.0000	185.0000

10. Feel free to play around with costs, budgets and mutual exclusivities to see how they affect the problem!