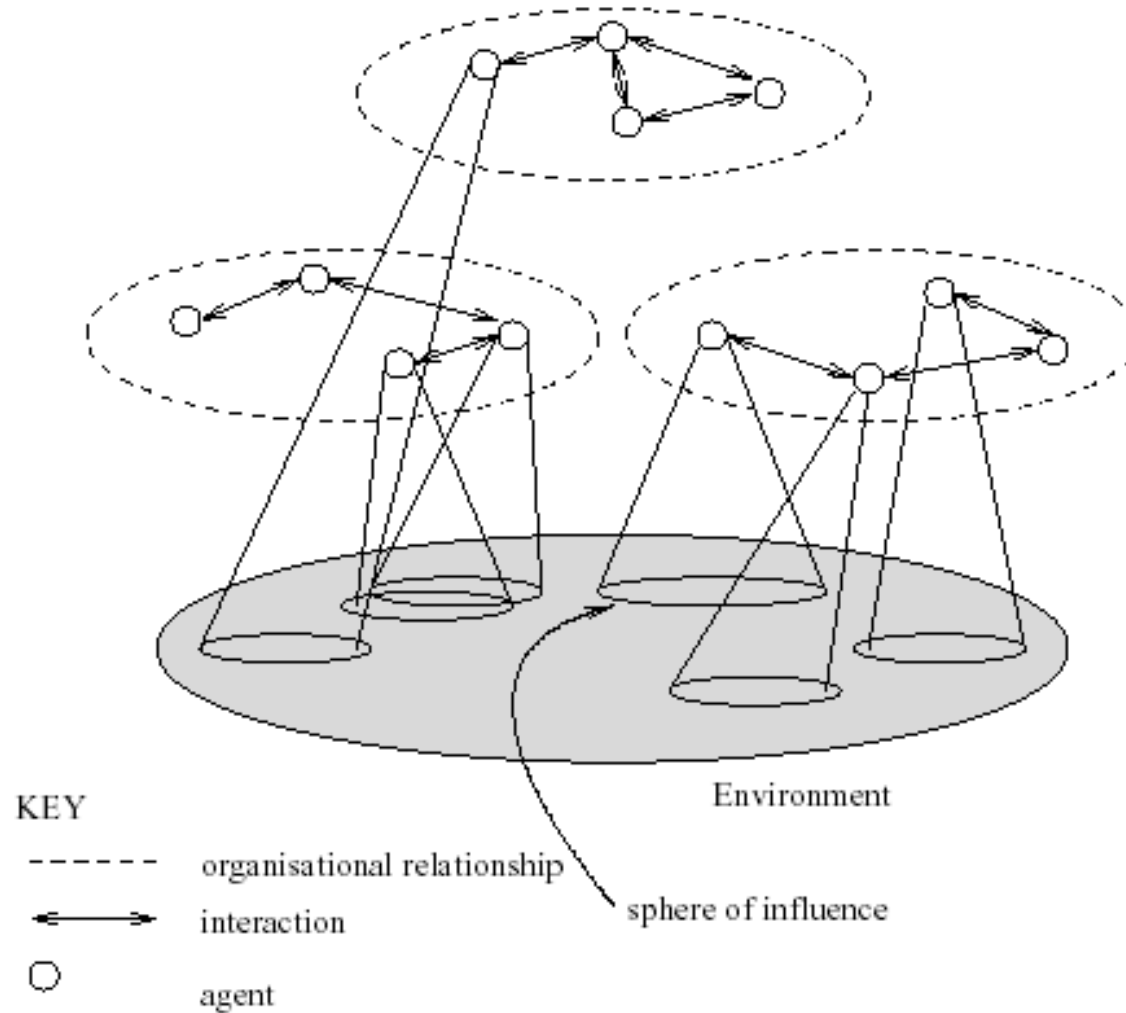


# Lecture 7: Game theory

Simon Powers  
SET10111 – Multi-Agent Systems

# Agent societies

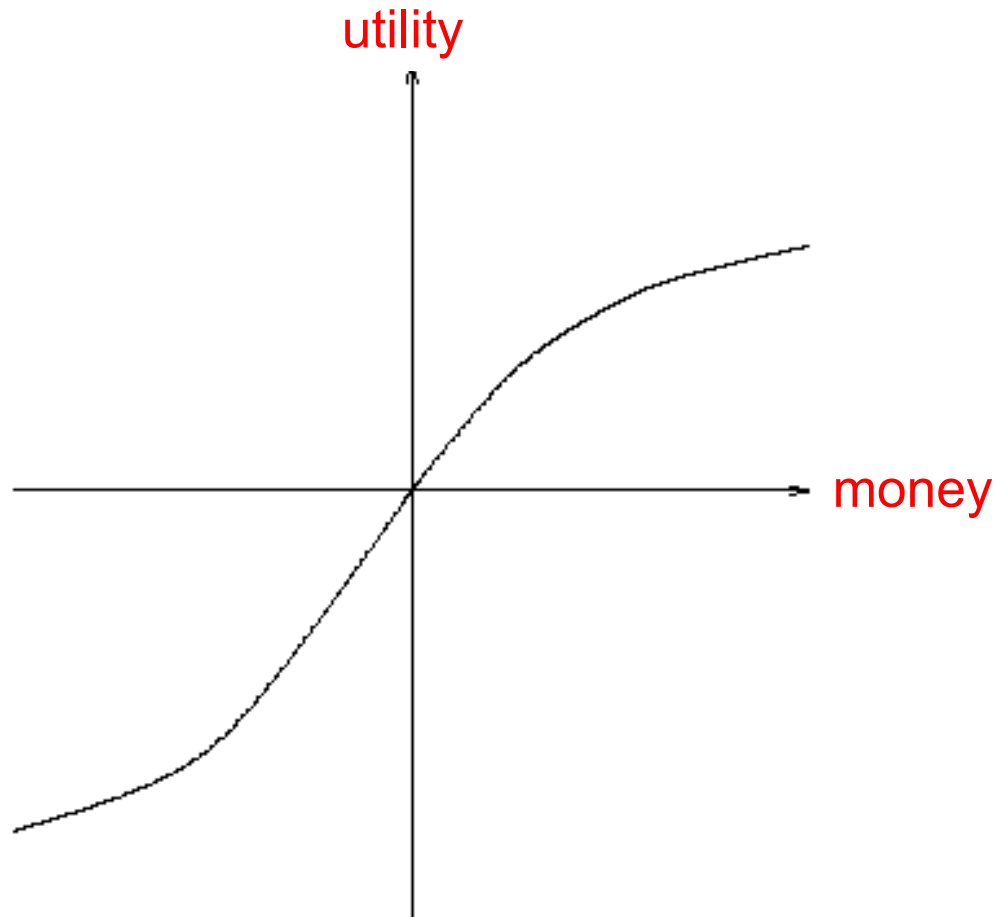


# Utilities and Preferences

- Assume two agents
- Set of **Environment states**:  $\Omega = \{\omega_1, \omega_2, \dots\}$
- **Utility functions** of Agent i and Agent j:
 
$$u_i : \Omega \rightarrow \mathbb{R}$$

$$u_j : \Omega \rightarrow \mathbb{R}$$
- **Preference ordering**: state  $\omega$  is **weakly preferred** by Ag. i over  $\omega'$  iff  
 $u_i(\omega) \geq u_i(\omega')$       Abbreviated notation:  $\omega \succeq \omega'$
- **Preference ordering**: state  $\omega$  is **strongly preferred** by Ag. i over  $\omega'$  iff  
 $u_i(\omega) > u_i(\omega')$       Abbreviated notation:  $\omega \succ \omega'$
- Relation  $\succeq$  is an ordering (reflexive, transitive, comparable)

# What is utility?



# Multi-agent encounters

- Agents simultaneously **choose action** to perform → result of the actions they select → **outcome in  $\Omega$**
- **Actual outcome** depends on the **combination of actions**
- Assume: each agent has just **two possible actions**  **$C$**  (“cooperate”) and  **$D$**  (“defect”)
- Environment behavior given by ***state transformer function***:

$$\tau : \underbrace{Ac}_{\text{agent } i\text{'s action}} \times \underbrace{Ac}_{\text{agent } j\text{'s action}} \rightarrow \Omega$$

# Examples of state transformer functions

- $\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_2 \quad \tau(C, D) = \omega_3 \quad \tau(C, C) = \omega_4$
- $\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_1 \quad \tau(C, D) = \omega_1 \quad \tau(C, C) = \omega_1$
- $\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_2 \quad \tau(C, D) = \omega_1 \quad \tau(C, C) = \omega_2$

# Examples of state transformer functions

- $\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_2 \quad \tau(C, D) = \omega_3 \quad \tau(C, C) = \omega_4$   
(environment is sensitive to actions of both agents.)
- $\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_1 \quad \tau(C, D) = \omega_1 \quad \tau(C, C) = \omega_1$   
(Neither agent has any influence in this environment.)
- $\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_2 \quad \tau(C, D) = \omega_1 \quad \tau(C, C) = \omega_2$   
(environment is controlled by  $j$ .)



# Rational behaviour

- **Assumption:** Environment is sensitive to actions of both agents:

$$\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_2 \quad \tau(C, D) = \omega_3 \quad \tau(C, C) = \omega_4$$

- **Utility functions:**

$$\begin{array}{cccc} u_i(\omega_1) = 1 & u_i(\omega_2) = 1 & u_i(\omega_3) = 4 & u_i(\omega_4) = 4 \\ u_j(\omega_1) = 1 & u_j(\omega_2) = 4 & u_j(\omega_3) = 1 & u_j(\omega_4) = 4 \end{array}$$

- **Short notation:**

$$\begin{array}{cccc} u_i(D, D) = 1 & u_i(D, C) = 1 & u_i(C, D) = 4 & u_i(C, C) = 4 \\ u_j(D, D) = 1 & u_j(D, C) = 4 & u_j(C, D) = 1 & u_j(C, C) = 4 \end{array}$$

- $\rightarrow$  agent  $i$ 's **preferences:**  
(also in short notation):

$$C, C \succeq_i C, D \quad \succ_i \quad D, C \succeq_i D, D$$



# Rational behaviour

$$\begin{array}{llll} u_i(D, D) = 1 & u_i(D, C) = 1 & u_i(C, D) = 4 & u_i(C, C) = 4 \\ u_j(D, D) = 1 & u_j(D, C) = 4 & u_j(C, D) = 1 & u_j(C, C) = 4 \end{array}$$

$$C, C \succeq_i C, D \quad \succsim_i \quad D, C \succeq_i D, D$$

$$C, C \succeq_j D, C \quad \succsim_j \quad C, D \succeq_j D, D$$

- “C” is the *rational choice* for i.  
(Because i (strongly) prefers all outcomes that arise through C over all outcomes that arise through D.)
- “C” is the *rational choice* for j.  
(Because j (strongly) prefers all outcomes that arise through C over all outcomes that arise through D.)

# Payoff matrices

- Game theory: characterize the previous scenario in a *payoff matrix*:

	Agent $j$ cooperates	Agent $j$ defects
Agent $i$ cooperates	4 4	1 4
Agent $i$ defects	4 1	1 1

$$\left[ \begin{array}{l} \text{same as:} \\ u_i(D, D) = 1 \quad u_i(D, C) = 1 \quad u_i(C, D) = 4 \quad u_i(C, C) = 4 \\ u_j(D, D) = 1 \quad u_j(D, C) = 4 \quad u_j(C, D) = 1 \quad u_j(C, C) = 4 \end{array} \right]$$

- Agent  $i$  is “*row player*” – utility given in lower-left of cell
- Agent  $j$  is “*column player*” – utility given in upper-right of cell

# Solution concepts

- How will a rational agent behave in any given scenario?
- Answered in *solution concepts*:
  - Dominant strategy
  - Nash equilibrium strategy
  - Pareto optimal strategy
  - Strategies that maximise social welfare

# Dominant strategies

- A strategy  $s_i$  is dominant for player  $i$  if no matter what strategy  $s_j$  agent  $j$  chooses,  $i$  will do at least as well playing  $s_i$  as it would playing anything else.
- Unfortunately, there isn't always a dominant strategy.

# (Pure Strategy) Nash Equilibrium

- In general, we will say that two strategies  $s_1$  and  $s_2$  are in Nash equilibrium if:
  1. under the assumption that agent  $i$  plays  $s_1$ , agent  $j$  can do no better than play  $s_2$ ; and
  2. under the assumption that agent  $j$  plays  $s_2$ , agent  $i$  can do no better than play  $s_1$ .
- Neither agent has any incentive to deviate from a Nash equilibrium.
- Unfortunately:
  1. Not every interaction has a (pure strategy) Nash equilibrium.
  2. Some interactions have more than one Nash equilibrium.

# The game of Matching Pennies

Players  $i$  and  $j$  simultaneously choose the face of a coin, either “heads” or “tails”.

If they show the same face, then  $j$  wins, while if they show different faces, then  $i$  wins.

# Matching Pennies payoff matrix

	<i>j</i> heads	<i>j</i> tails
<i>i</i> heads	-1 1	1 -1
<i>i</i> tails	1 -1	-1 1

# Mixed strategies for matching pennies

- NO pair of strategies forms a pure strategy Nash Equilibrium: whatever pair of strategies is chosen, somebody will wish they had done something else.
- The solution is to allow **mixed strategies**:
  - Play “heads” with probability 0.5
  - Play “tails” with probability 0.5
- This is a Nash Equilibrium strategy.
- Every finite game has a Nash Equilibrium in mixed strategies.
- **But there might be more than one...**



# Pareto optimality

- An outcome is said to be *Pareto optimal* (or *Pareto efficient*) if there is no other outcome that makes one agent *better off* without making another agent *worse off*.
- If an outcome is Pareto optimal, then at least one agent will be reluctant to move away from it (because this agent will be worse off).
- If an outcome  $\omega$  is *not* Pareto optimal, then there is another outcome  $\omega'$  that makes *everyone* as happy, if not happier, than  $\omega$ .

# Social welfare

- The social welfare of an outcome  $\omega$  is the sum of the utilities that each agent gets from  $\omega$ :

$$\sum u_i(\omega)$$

- Think of it as the “total amount of money in the system”.
- As a solution concept, may be appropriate when the whole system (all agents) has a single owner (then overall benefit of the system is important, not individuals).

# The Prisoner's Dilemma

Two men are collectively charged with a crime and held in separate cells, with no way of meeting or communicating.

They are told that:

- if one confesses and the other does not, the confessor will be freed, and the other will be jailed for five years;
- if both confess, then each will be jailed for two years.

Both prisoners know that if neither confesses, then they will each be jailed for one year.

# Prisoner's Dilemma payoff matrix

	Agent <i>j</i> cooperates	Agent <i>j</i> defects
Agent <i>i</i> cooperates	3, 3	0, 5
Agent <i>i</i> defects	5, 0	2, 2

# Prisoner's Dilemma payoff matrix

	Agent $j$ cooperates	Agent $j$ defects
Agent $i$ cooperates	3, 3	0, 5
Agent $i$ defects	5, 0	2, 2

$D$  is a **dominant strategy**.

$(D, D)$  is the only **Nash equilibrium**.

All outcomes except  $(D, D)$  are **Pareto optimal**.

$(C, C)$  maximises **social welfare**.

# Rational behaviour

- suppose **I cooperate**: If j also cooperates → we both get payoff 3. If j defects → I get payoff 0. → Best **guaranteed payoff** when I cooperate is 0
- suppose **I defect**: If j cooperates → I get payoff 5. If j also defects → both get payoff 2. → Best **guaranteed payoff** when I defect is 2
- → If **I defect** I'll get a minimum guaranteed payoff of 2. If **I cooperate** I'll get a minimum guaranteed payoff of 0.
- → If prefer guaranteed payoff of 2 to guaranteed payoff of 0.  
→ I should defect

# Rational behaviour

- Only one Nash equilibrium: (D,D). (“under the assumption that the other does D, one can do no better than do D”)
- Intuition says: (C,C) is better than (D,D) so why not (C,C)? → but if agent assumes that other agent does C it is BEST to do D! → seemingly “waste of utility”
- “Shocking” truth: defect is rational, cooperate is irrational

# Examples

- Nuclear arms reduction (D: do not reduce, C: reduce)
- Paying for **public goods** like a TV license
- More generally, any form of **exchange**. Exchange is always sequential so why send the goods or be honest about their quality?
  - Buying a used car
  - Buying goods from an eBay seller



# What does all this really mean?

- “Defect more rational than cooperate” → Humans: **Machiavellism** (opposed to real altruism)
- Philosophical question: isn't even altruism ultimately some kind of optimisation towards OWN goals?!
- **Further aspect: Strict rationalism** (in case of prisoner's dilemma: defect) is usually only applied when sucker's payoff really hurts.
- What we have not yet regarded: **Multiple sequential games** between same players → “**The shadow of the future**” → What does it mean to rationalism and strategy?

# The shadow of the future: Iterated Prisoner's Dilemma

- Game is played **multiple times**. Agents can see all past actions of other agent.
- Course of reasoning:
  - If I **defect**, the other agent can **punish** me by defecting in the next run. (not a point in the one shot Prisoner's Dilemma game)
  - **Testing cooperation** (and possibly getting the sucker's payoff) is **not tragic**, because “on the long run” one (or several) sucker's payoff(s) is (are) “statistically” not important (can e.g. be equaled by gains through mutual cooperation)
- → in an iterated PD-game: cooperation is rational

# The Folk Theorem of game theory

- In a game repeated for an **indefinite number of times**, with information about the **past actions of other agents**, cooperation will be a Nash equilibrium through conditional strategies.
  - “Cooperate if my partner cooperated last time, otherwise defect”
- Technically, any strategy that gives more than the **minimax payoff** will be an equilibrium.
- Doesn't work if:
  - End time known
  - Insufficient information about the past actions of other agents.

# So is the Folk Theorem useless?

- How many times will our agents meet???
- Institutions can provide reputational information about the past actions of other agents
  - e.g. Ebay's reputation system

# Axelrod's Prisoner's Dilemma tournament

Some strategies competing:

- **ALL-D**: Always defect
- **All-C**: Always cooperate
- **RANDOM**: Choose D or C randomly
- **TIT-FOR-TAT**: On the first round cooperate; on round  $t$  do what opponent did on round  $t-1$
- **TESTER**: Intention: Exploit "nice" programs that not punish D: On first round, test opponent with D. If opponent retaliates with D  $\rightarrow$  play TIT-FOR-TAT. If opponent does C  $\rightarrow$  play CCD.
- **JOSS**: like TIT-FOR-TAT, except periodically defect.

# Axelrod's Prisoner's Dilemma tournament

- Winner was TIT-FOR-TAT (5 lines of Fortran code).
- Why? Overall Score of a strategy computed as average of performance **against all other strategies**. TIT-FOR-TAT was defeated by ALL-D.
- TIT-FOR-TAT won against “cooperative strategies” → Again (as in analysis of single PD): being not too cooperative pays out.
- → Axelrod distilled some rules from outcome of tournament:

# Axelrod's Prisoner's Dilemma tournament

1. Do not be envious.
2. Do not be first to defect.
3. Reciprocate C and D.
4. Don't be too clever.

# Snowdrift game (game of Chicken)

	Agent $j$ cooperates	Agent $j$ defects
Agent $i$ cooperates	2 2	1 3
Agent $i$ defects	3 1	0 0



# Snowdrift game (game of Chicken)

	Agent $j$ cooperates	Agent $j$ defects
Agent $i$ cooperates	2, 2	1, 3
Agent $i$ defects	3, 1	0, 0

- Two (pure) Nash Equilibria: (D,C), (C,D)  
(Assuming the other does D you can do no better than do C  
Assuming the other does C you can do no better than do D)
- There is no dominant strategy.
- All outcomes except (D, D) are Pareto optimal.
- All outcomes except (D, D) maximise social welfare.

## Pure coordination game (T-shirt game)

	Agent $j$ wears red	Agent $j$ wears blue
Agent $i$ wears red	3 3	0 0
Agent $i$ wears blue	0 0	3 3

# Pure coordination game (T-shirt game)

	Agent $j$ wears red	Agent $j$ wears blue
Agent $i$ wears red	3 3	0 0
Agent $i$ wears blue	0 0	3 3

- Two Nash equilibria: (red, red), (blue, blue)
- There is no dominant strategy.
- (red, red) and (blue, blue) are Pareto optimal.
- (red, red) and (blue, blue) maximise social welfare.

# Stag-hunt game (risky coordination game)

	Agent $j$ wears red	Agent $j$ wears blue
Agent $i$ wears red	3 3	0 1
Agent $i$ wears blue	1 0	1 1

# Stag-hunt game (risky coordination game)

	Agent $j$ wears red	Agent $j$ wears blue
Agent $i$ wears red	3, 3	0, 1
Agent $i$ wears blue	1, 0	1, 1

- Two Nash equilibria: (red, red), (blue, blue)
- (blue, blue) is **risk-dominant** -- it has the largest basin of attraction (is less risky).
- (red, red) is **payoff-dominant** – it offers to each player at least as much payoff as the other Nash equilibria.
- (red, red) is Pareto optimal.
- (red, red) maximises social welfare.

# Acknowledgements

- Much of this lecture is based on Chapter 11 of Wooldridge “An Introduction to MultiAgent Systems”, 2<sup>nd</sup> edition
  - <http://www.cs.ox.ac.uk/people/michael.wooldridge/pubs/imas/IMAS2e.html>
  - And draws upon the accompanying slides by Dr. Georg Groh, TU-München, Germany  
<http://www.cs.ox.ac.uk/people/michael.wooldridge/pubs/imas/alt-slides.zip>