

Lemma

Let $\mathcal{U} \in \mathcal{E}_1$

Assuming, $\exists PA \vdash \mathcal{U} \Rightarrow P^{PA} - (V)$ (in \mathcal{B}_1)

! problem solving, it should be PA .

$$\mathcal{U} := \exists y \delta(y) \quad \delta \in \Delta_0$$

Do not change
negative knowledge
in δ expression per y .

Let $\exists \delta(m)$.

Constructing $\exists PA \vdash A(x) \wedge x \leq m \Rightarrow x = 0 \vee x = 1 \vee \dots$

(the coding \downarrow)

$$\forall x = m$$

(Constructing $\exists PA$ over \mathcal{B}_1)

! δ is $\delta(m)$ for $m \leq m$, to $PA \vdash \delta(m)$

Let $\delta(m)$ to $PA \vdash \delta(m)$

$PA \vdash \delta(m)$

δ is $\delta(m)$ in $x+y = z$ ok - manipulating expression

$$\delta(x_1, \dots, x_l) := \exists z \leq y \delta(x_1, \dots, x_l, z)$$

Let $\delta(m_1, \dots, m_l)$ to $PA \vdash \delta(m_1, \dots, m_l)$

$z = p \leq m$! meaning, $\exists PA \vdash \delta(m_1, \dots, m_l, p)$

Let $\delta(m_1, \dots, m_l)$ to $PA \vdash \delta(m_1, \dots, m_l)$

$$\vdash \delta(m_1, \dots, m_l)$$

! odd/even

$$(x) \text{ do } k = m$$