## **PROBLEMS UNIT 1**

## 3) Analyze the efficiency of the following code:

We assume the cost of Mínimo(i,j) = 1

 $T(n) \rightarrow O(n \log(n))$ 

```
T(n) = 1 + \sum_{i=1}^{n} \left(1 + \sum_{i=1}^{n} (1+i)\right) + 1 + \mathbb{E}(n/2)
= 1 + \sum_{i=1}^{n} \left(1 + \sum_{i=1}^{n} 2n + 1 + \mathbb{E}(n/2)\right)
= 1 + \sum_{i=1}^{n} \left(2 + 2n^{2} + \mathbb{E}(n/2)\right)
= 1 + 2n + 2n^{2} + \mathbb{E}(n/2)
h_{2}^{2}^{k}
T(2^{k}) = 1 + 2 \cdot 2^{k} + 2 \cdot 2^{k} + 2^{k} \cdot 2^{k} \cdot 2^{k} \cdot 2^{k}
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```

6) Program a function to determine if a number received as parameter is prime. Analyze the efficiency and complexity.

```
public static boolean isPrime(int num) {
        int temp;
        boolean isPrime = true;
        if (num == 0 || num == 1)
                                                  // Both 0 and arent prime numbers
        {
            isPrime = false;
        }
        else
             for(int i = 2; i \leftarrow num/2; i \leftrightarrow num/2; i \leftrightarrow num/2 because we are going through the possible factors,
                                                   // so we dont want to repeat them.
               temp = num % i;
               if(temp == 0)
                   isPrime = false;
                                                  // If we find a factor different from 0 or num then
                                                  // num is not prime.
                   break;
               }
             }
        return isPrime;
    }
T(n) = 1 + E 1 + 1 + 1
= 1 + 3 * n/2
```

```
= 1 + 1,5n
= O(n)
```

8) Program a recursive procedure to obtain the inverse number of a given one. Example : 627 -> 726. Analyze the efficiency and complexity.

```
public static void reverseNumber(int num, int n, int aux) {
         if (num < 10)
         {
             aux = aux * 10 + num;
             System.out.println(aux);
         }
        else
         {
             aux = aux * 10 + (num % 10);
             reverseNumber(num/10,n++,aux);
         }
    }
T(n) = 1 + 1 + t(n/10)
N = 10^k
X^k = 2 + x^k-1
X^k-1(x-1) = 0 x(H) = A
X^k-1(x-1) = 2 x(P) = Bk
X = A + Bk = A + log(n)B
= O(log(n))
```