

MI-GLR Homework 5

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1 Problem 1: Nash Equilibrium

3,2	1,1
0,0	2,3

Nash Equilibria are in **Bold**.

Mixed Nash Equilibria:

Player 1 plays first action with probability p_1 :

Reward for Player 2:

$$2 * p_1 + 0 * (1 - p_1) = p_1 + 3 * (1 - p_1)$$

$$2 * p_1 = -2 * p_1 + 3$$

$$4 * p_1 = 3$$

$$p_1 = \frac{3}{4}$$

\implies Strategy for Player 1: $(\frac{3}{4}, \frac{1}{4})$.

Player 2 plays first action with probability p_2 :

Reward for Player 1:

$$3 * p_2 + 1 * (1 - p_2) = 0 * p_2 + 2 * (1 - p_2)$$

$$2 * p_2 + 1 = 2 - 2 * p_2$$

$$4 * p_2 = 1$$

$$p_2 = \frac{1}{4}$$

\implies Strategy for Player 2: $(\frac{1}{4}, \frac{3}{4})$.

2 Problem 2: Maxmin and minmax

1	2
0	x

Table 1: Matrix A of a zero-sum game.

If some entry a_{ij} of the matrix A has the property that:

1. a_{ij} is the minimum of the i th row and
2. a_{ij} is the maximum of the j th column

then we say a_{ij} is a saddle point. If a_{ij} is a saddle point, then Player I can then win at least a_{ij} by choosing row i , and Player II can keep her loss to at most a_{ij} by choosing column j . Hence a_{ij} is the **value of the game**. Optimal strategy for Player I is to always play action i and for Player II it is to always play action j .

Solution:

Saddle point is a_{11} , therefore value of the game is 1 and optimal strategy for both players is the following pure strategy: $(1, 0)$.

References

- [1] Ferguson, S. Thomas "GAME THEORY" UCLA Department of Mathematics. Accessible from: https://www.math.ucla.edu/~tom/Game_Theory/mat.pdf