## **1-Error Analysis**

## 2-Solution of non-linear equation

**Example 2** Show that  $x^5 - 2x^3 + 3x^2 - 1 = 0$  has a solution in the interval [0, 1].

**Solution** Consider the function defined by  $f(x) = x^5 - 2x^3 + 3x^2 - 1$ . The function f is continuous on [0, 1]. In addition,

$$f(0) = -1 < 0$$
 and  $0 < 1 = f(1)$ .

The Intermediate Value Theorem implies that a number x exists, with 0 < x < 1, for which  $x^5 - 2x^3 + 3x^2 - 1 = 0$ .

#### **Decimal Machine Numbers**

normalized decimal floating-point form

$$\pm 0.d_1d_2...d_k \times 10^n$$
,  $1 \le d_1 \le 9$ , and  $0 \le d_i \le 9$ ,

for each i = 2, ..., k. Numbers of this form are called k-digit decimal machine numbers.

$$y = 0.d_1d_2...d_kd_{k+1}d_{k+2}...\times 10^n$$
.

$$fl(y) = 0.d_1d_2...d_k \times 10^n$$
. Chopping

#### Rounding

For rounding, when  $d_{k+1} \ge 5$ , we add 1 to  $d_k$  to obtain fl(y);

 $d_{k+1} < 5$ , we simply chop off all but the first k digits; so we round down.

## **Significant digits** are those digits that can be used with confidence.

Non zero numbers are always significant

```
1.23 45.6 6,7263
```

**❖** In between zeros are always significant

```
1.005 70206
```

**!** Leading zeros are never significant

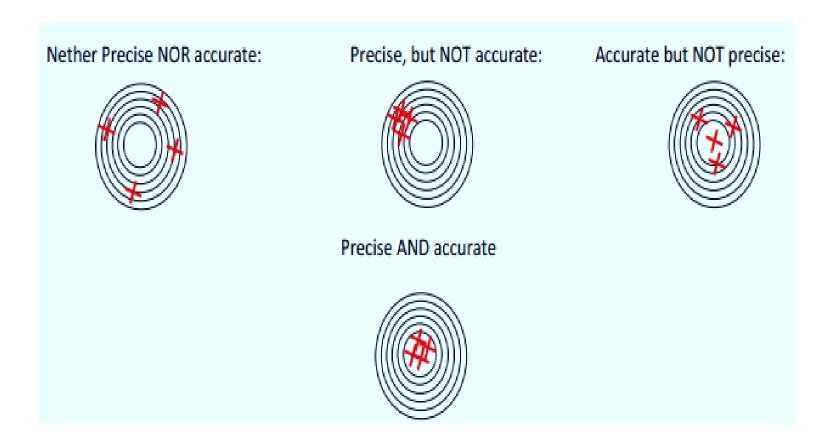
```
0.0055 0.0302
```

**❖** Trailing zeros are some time significant

70,000 70,000. 1,030 1030.0000

## **Accuracy and Precision**

- Accuracy is related to the closeness to the true value.
- Precision is related to the closeness to other estimated values.



## Rounding and Chopping

Rounding: Replace the number by the nearest machine number. OR

its impossible to represent all real numbers exactly on machine with finite

Chopping: Throw all or drop the extra digits.

Error: is difference between an approximation of number used in computation and its exact value

OR Error = True value – approximate value

```
• \sqrt{2} = 1.414213562373095048801168872
```

## **ERROR Analysis:**

Truncation Error:

are when an iterative method is terminatedOR mathematical procedure is approximated and approximate solution differs from exact solution

• Discretization Error:

are committed when a solution of discrete problem does not coincide with solution of continuous problem

## Error in CM — True Error

Can be computed if the true value is known:

Absolute Error:

$$AE = |$$
 true value – approximation |

Absolute Relative Error:

$$ARE = \frac{\text{true value - approximation}}{\text{true value}}$$

## Error in CM — Estimated Error

When the true value is not known:

Estimated Absolute Error

$$AE = |$$
 current estimate  $-$  previous estimate  $|$ 

Estimated Absolute Relative Error

$$ARE = \frac{|\text{current estimate} - \text{previous estimate}|}{\text{current estimate}}$$

Example 1 Determine the five-digit (a) chopping and (b) rounding values of the irrational number  $\pi$ .

**Definition 1.15** Suppose that  $p^*$  is an approximation to p. The absolute error is  $|p-p^*|$ , and the relative error is  $\frac{|p-p^*|}{|p|}$ , provided that  $p \neq 0$ .

**Example 2** Determine the absolute and relative errors when approximating p by  $p^*$  when

- (a)  $p = 0.3000 \times 10^1$  and  $p^* = 0.3100 \times 10^1$ ;
- **(b)**  $p = 0.3000 \times 10^{-3}$  and  $p^* = 0.3100 \times 10^{-3}$ ;
- (c)  $p = 0.3000 \times 10^4$  and  $p^* = 0.3100 \times 10^4$ .

**Example 3** Suppose that  $x = \frac{5}{7}$  and  $y = \frac{1}{3}$ . Use five-digit chopping for calculating x + y, x - y,  $x \times y$ , and  $x \div y$ .

Solution Note that

Finite-Digit Arithmetic

$$x = \frac{5}{7} = 0.\overline{714285}$$
 and  $y = \frac{1}{3} = 0.\overline{3}$ 

implies that the five-digit chopping values of x and y are

$$fl(x) = 0.71428 \times 10^0$$
 and  $fl(y) = 0.33333 \times 10^0$ .

Thus

$$x \oplus y = fl(fl(x) + fl(y)) = fl(0.71428 \times 10^{0} + 0.33333 \times 10^{0})$$
  
=  $fl(1.04761 \times 10^{0}) = 0.10476 \times 10^{1}$ .

The true value is  $x + y = \frac{5}{7} + \frac{1}{3} = \frac{22}{21}$ , so we have

Absolute Error = 
$$\left| \frac{22}{21} - 0.10476 \times 10^{1} \right| = 0.190 \times 10^{-4}$$

and

Relative Error = 
$$\left| \frac{0.190 \times 10^{-4}}{22/21} \right| = 0.182 \times 10^{-4}$$
.

$$x \oplus y = fl(fl(x) + fl(y)),$$
  
$$x \ominus y = fl(fl(x) - fl(y)),$$

Table 1.2

Operation	Result	Actual value	Absolute error	Relative error
$x \oplus y$	$0.10476 \times 10^{1}$	22/21	$0.190 \times 10^{-4}$	$0.182 \times 10^{-4}$
$x \ominus y$	$0.38095 \times 10^{0}$	8/21	$0.238 \times 10^{-5}$	$0.625 \times 10^{-5}$
$x \otimes y$	$0.23809 \times 10^{0}$	5/21	$0.524 \times 10^{-5}$	$0.220 \times 10^{-4}$
$x \oplus y$	$0.21428 \times 10^{1}$	15/7	$0.571 \times 10^{-4}$	$0.267 \times 10^{-4}$

The maximum relative error for the operations in Example 3 is  $0.267 \times 10^{-4}$ ,

Example 5 Let p = 0.54617 and q = 0.54601. Use four-digit arithmetic to approximate p - q and determine the absolute and relative errors using (a) rounding and (b) chopping.

## Loss of significance:

#### occurs in numerical calculations when too many significant digits cancel

The quadratic formula states that the roots of  $ax^2 + bx + c = 0$ , when  $a \neq 0$ , are

$$x_1 = \frac{-b + \sqrt{b^2 - 4ac}}{2a}$$
 and  $x_2 = \frac{-b - \sqrt{b^2 - 4ac}}{2a}$ . (1.1)

which simplifies to an alternate quadratic formula

$$x_1 = \frac{-2c}{b + \sqrt{b^2 - 4ac}}. (1.2)$$

The rationalization technique can also be applied to give the following alternative quadratic formula for  $x_2$ :

$$x_2 = \frac{-2c}{b - \sqrt{b^2 - 4ac}}. ag{1.3}$$

### **Activity**

#### (Taylor's Theorem)

$$P_n(x) = f(x_0) + f'(x_0)(x - x_0) + \frac{f''(x_0)}{2!}(x - x_0)^2 + \dots + \frac{f^{(n)}(x_0)}{n!}(x - x_0)^n$$

$$= \sum_{k=0}^n \frac{f^{(k)}(x_0)}{k!}(x - x_0)^k$$

- 7. Let  $f(x) = x^3$ .
  - a. Find the second Taylor polynomial  $P_2(x)$  about  $x_0 = 0$ .
  - **b.** Find  $R_2(0.5)$  and the actual error in using  $P_2(0.5)$  to approximate f(0.5).
  - c. Repeat part (a) using  $x_0 = 1$ .
  - **d.** Repeat part (b) using the polynomial from part (c).
- 8. Find the third Taylor polynomial  $P_3(x)$  for the function  $f(x) = \sqrt{x+1}$  about  $x_0 = 0$ . Approximate  $\sqrt{0.5}$ ,  $\sqrt{0.75}$ ,  $\sqrt{1.25}$ , and  $\sqrt{1.5}$  using  $P_3(x)$ , and find the actual errors.
- 9. Find the second Taylor polynomial  $P_2(x)$  for the function  $f(x) = e^x \cos x$  about  $x_0 = 0$ .
  - a. Use  $P_2(0.5)$  to approximate f(0.5). Find an upper bound for error  $|f(0.5) P_2(0.5)|$  using the error formula, and compare it to the actual error.

## Activity

a. 
$$p = \pi, p^* = 22/7$$

c. 
$$p = e, p^* = 2.718$$

e. 
$$p = e^{10}, p^* = 22000$$

**b.** 
$$p = \pi, p^* = 3.1416$$

d. 
$$p = \sqrt{2}, p^* = 1.414$$

0. 
$$p = \sqrt{2}, p^* = 1.414$$
  
0.  $p = 10^{\pi}, p^* = 1400$ 

c. 
$$(121 - 0.327) - 119$$

e. 
$$\frac{\frac{13}{14} - \frac{6}{7}}{2e - 5.4}$$

d. 
$$(121 - 119) - 0.327$$
  
f.  $-10\pi + 6e - \frac{3}{62}$ 

a. 
$$\frac{1}{3}x^2 - \frac{123}{4}x + \frac{1}{6} = 0$$
  
b.  $\frac{1}{3}x^2 + \frac{123}{4}x - \frac{1}{6} = 0$ 

**b.** 
$$\frac{1}{3}x^2 + \frac{123}{4}x - \frac{1}{6} = 0$$

c. 
$$1.002x^2 - 11.01x + 0.01265 = 0$$

**d.** 
$$1.002x^2 + 11.01x + 0.01265 = 0$$

Repeat Exercise 13 using four-digit chopping arithmetic.

#### **Nested Arithmetic**

Accuracy loss due to round-off error can also be reduced by rearranging calculations, as shown in the next example.

**Example 6** Evaluate  $f(x) = x^3 - 6.1x^2 + 3.2x + 1.5$  at x = 4.71 using three-digit arithmetic.

Three-digit (chopping): 
$$f(4.71) = ((104. - 134.) + 15.0) + 1.5 = -13.5$$
,

Three-digit (rounding): 
$$f(4.71) = ((105. - 135.) + 15.1) + 1.5 = -13.4$$
.

Chopping: 
$$\left| \frac{-14.263899 + 13.5}{-14.263899} \right| \approx 0.05$$
, and Rounding:  $\left| \frac{-14.263899 + 13.4}{-14.263899} \right| \approx 0.06$ .

#### **Nested Arithmetic**

As an alternative approach, the polynomial f(x) in Example 6 can be written in a nested manner as

$$f(x) = x^3 - 6.1x^2 + 3.2x + 1.5 = ((x - 6.1)x + 3.2)x + 1.5.$$

Using three-digit chopping arithmetic now produces

$$f(4.71) = ((4.71 - 6.1)4.71 + 3.2)4.71 + 1.5 = ((-1.39)(4.71) + 3.2)4.71 + 1.5$$
$$= (-6.54 + 3.2)4.71 + 1.5 = (-3.34)4.71 + 1.5 = -15.7 + 1.5 = -14.2.$$

Three-digit (chopping): 
$$\left| \frac{-14.263899 + 14.2}{-14.263899} \right| \approx 0.0045;$$

Three-digit (rounding): 
$$\left| \frac{-14.263899 + 14.3}{-14.263899} \right| \approx 0.0025.$$

Nesting has reduced the relative error for the chopping approximation to less than 10% of that obtained initially. For the rounding approximation the improvement has been even more dramatic; the error in this case has been reduced by more than 95%.

## 2-Solution of non linear equation in one variable

- 1-Bracketing Methods
- 2-Open Methods

## **Bracketing Methods:**

• In bracketing methods, the method starts with an <u>interval</u> that contains the root and a procedure is used to obtain a smaller interval containing the root.

- Examples of bracketing methods:
  - Bisection method
  - False position method

## Open Methods:

• In the open methods, the method starts with one or more initial guess points. In each iteration, a new guess of the root is obtained.

- Open methods are usually more efficient than bracketing methods.
- They may not converge to a root.
  - Fixed point,
  - Newton and
  - Secant are examples of open method

#### Bisection

To find a solution to f(x) = 0 given the continuous function f on the interval [a, b], where f(a) and f(b) have opposite signs:

INPUT endpoints a, b; tolerance TOL; maximum number of iterations  $N_0$ .

OUTPUT approximate solution p or message of failure.

Step 1 Set 
$$i = 1$$
;  
 $FA = f(a)$ .

Step 2 While  $i \le N_0$  do Steps 3–6.

Step 3 Set 
$$p = a + (b - a)/2$$
; (Compute  $p_i$ .)  
 $FP = f(p)$ .  
Step 4 If  $FP = 0$  or  $(b - a)/2 = TOI$  then

Step 4 If 
$$FP = 0$$
 or  $(b - a)/2 < TOL$  then  
OUTPUT  $(p)$ ; (Procedure completed successfully.)  
STOP.

Step 5 Set 
$$i = i + 1$$
.

Step 6 If 
$$FA \cdot FP > 0$$
 then set  $a = p$ ; (Compute  $a_i, b_i$ .)  
 $FA = FP$   
else set  $b = p$ . (FA is unchanged.)

Step 7 OUTPUT ('Method failed after  $N_0$  iterations,  $N_0 =$ ',  $N_0$ ); (The procedure was unsuccessful.)

STOP. Assistant Prof:Jamilusmani

## **Algorithm**

To determine a root of f(x) = 0 that is accurate within a specified tolerance value, given values  $x_1$  and  $x_2$  such that  $f(x_1) * f(x_2) < 0$ ,

### Repeat

Set 
$$x_3 = (x_1 + x_2)/2$$
.

If 
$$f(x_3) * f(x_1) < 0$$
 Then

$$Set x_2 = x_3$$

Else Set 
$$x_1 = x_3$$
 End If.

Until 
$$(|x_1 - x_2|) < 2$$
 \* tolerance value).

The final value of  $x_3$  approximates the root, and it is in error by not more than  $|x_1 - x_2|/2$ .

*Note:* The method may produce a false root if f(x) is discontinuous on  $[x_1, x_2]$ .

## Stopping criteria:

select a tolerance  $\varepsilon > 0$ , and construct  $p_1, \dots p_N$  until

$$|p_N - p_{N-1}| < \varepsilon,$$

$$\frac{|p_N-p_{N-1}|}{|p_N|}<\varepsilon,\quad p_N\neq 0,$$

$$|f(p_N)| < \varepsilon.$$

### The number of iteration required for given tolerance:

$$|p-p_n|\leq \frac{b-a}{2^n}\leq \epsilon,$$

then taking the logarithms of both sides yields

$$n \ge \frac{\log\left(\frac{b-a}{\epsilon}\right)}{\log(2)}.$$

Determine the number of iterations necessary to solve  $f(x) = x^3 + 4x^2 - 10 = 0$  with accuracy  $10^{-3}$  using  $a_1 = 1$  and  $b_1 = 2$ .

Example 1 Show that  $f(x) = x^3 + 4x^2 - 10 = 0$  has a root in [1, 2], and use the Bisection method to determine an approximation to the root that is accurate to at least within  $10^{-4}$ .

	 _		
		400	- 65
400	400		-
100		-	
		-	

n	$a_n$	$b_n$	$p_n$	$f(p_n)$
1	1.0	2.0	1.5	2.375
2	1.0	1.5	1.25	-1.79687
3	1.25	1.5	1.375	0.16211
4	1.25	1.375	1.3125	-0.84839
5	1.3125	1.375	1.34375	-0.35098
6	1.34375	1.375	1.359375	-0.09641
7	1.359375	1.375	1.3671875	0.03236
8	1.359375	1.3671875	1.36328125	-0.03215
9	1.36328125	1.3671875	1.365234375	0.000072
10	1.36328125	1.365234375	1.364257813	-0.01605
11	1.364257813	1.365234375	1.364746094	-0.00799
12	1.364746094	1.365234375	1.364990235	-0.00396
13	1.364990235	1.365234375	1.365112305	-0.00194

Assistant Prof:Jamilusmani

#### **Pros**

 The method has the important property that it always converges to a solution.

#### Cons

- 1. It is relatively slow to converge.
- A good intermediate approximation might be inadvertently discarded.
- 3. If a function f(x) is such that it just **touches the x-axis** such as  $f(x) = x^2 = 0$ , it will be unable to find the lower guess a and the upper guess b such that  $f(a) \cdot f(b) < 0$ .

### Fixed-Point Iteration

**Definition 2.2** The number p is a fixed point for a given function g if g(p) = p.

**Example 1** Determine any fixed points of the function  $g(x) = x^2 - 2$ .

**Solution** A fixed point p for g has the property that

$$p = g(p) = p^2 - 2$$
 which implies that  $0 = p^2 - p - 2 = (p+1)(p-2)$ .

so g has two fixed points, one at p = -1 and the other at p = 2.

**Example 2** Show that  $g(x) = (x^2 - 1)/3$  has a unique fixed point on the interval [-1, 1].

### Example

#### Fixed Point Iteration

$$f(x) = x^2 - 2x - 3 = 0$$
 (ans:  $x = 3$  or -1)

### Case a:

$$x^{2} - 2x - 3 = 0$$

$$\Rightarrow x^{2} = 2x + 3$$

$$\Rightarrow x = \sqrt{2x + 3}$$

$$\Rightarrow g(x) = \sqrt{2x + 3}$$

### Case b:

$$x^{2} - 2x - 3 = 0$$

$$\Rightarrow x(x - 2) - 3 = 0$$

$$\Rightarrow x = \frac{3}{x - 2}$$

$$\Rightarrow g(x) = \frac{3}{x - 2}$$

### Case c:

$$\begin{vmatrix} x^2 - 2x - 3 &= 0 \\ \Rightarrow 2x = x^2 - 3 \end{vmatrix}$$
$$\Rightarrow x = \frac{x^2 - 3}{2}$$
$$\Rightarrow g(x) = \frac{x^2 - 3}{2}$$

So which one is better?

### Case a

$$x_{i+1} = \sqrt{2x_i + 3}$$

$$x_{i+1} - \sqrt{2x_i} + 3$$

1. 
$$x_0 = 4$$

2. 
$$x_1 = 3.31662$$

3. 
$$x_2 = 3.10375$$

4. 
$$x_3 = 3.03439$$

5. 
$$x_4 = 3.01144$$

6. 
$$x_5 = 3.00381$$

Converge!

### Case b

$$x_{i+1} = \frac{3}{x_i - 2}$$

1. 
$$x_0 = 4$$

2. 
$$x_1 = 1.5$$

3. 
$$x_2 = -6$$

4. 
$$x_3 = -0.375$$

5. 
$$x_4 = -1.263158$$

6. 
$$x_5 = -0.919355$$

7. 
$$x_6 = -1.02762$$

8. 
$$x_7 = -0.990876$$

9. 
$$x_8 = -1.00305$$

### Case c

$$x_{i+1} = \frac{x_i^2 - 3}{2}$$

1. 
$$x_0 = 4$$

2. 
$$x_1 = 6.5$$

3. 
$$x_2 = 19.625$$

4. 
$$x_3 = 191.070$$

Diverge!

### Iteration Algorithm with the Form x = g(x)

To determine a root of f(x) = 0, given a value  $x_1$  reasonably close to the root,

Rearrange the equation to an equivalent form x = g(x).

Repeat

Set 
$$x_2 = x_1$$
.  
Set  $x_1 = g(x_1)$   
Until  $|x_1 - x_2|$  < tolerance value

*Note:* The method may converge to a root different from the expected one, or it may diverge. Different rearrangements will converge at different rates.

#### Fixed-Point Iteration

To find a solution to p = g(p) given an initial approximation  $p_0$ :

INPUT initial approximation  $p_0$ ; tolerance TOL; maximum number of iterations  $N_0$ .

OUTPUT approximate solution p or message of failure.

Step 1 Set 
$$i = 1$$
.

Step 2 While  $i \le N_0$  do Steps 3–6.

Step 3 Set 
$$p = g(p_0)$$
. (Compute  $p_i$ .)

Step 4 If 
$$|p - p_0| < TOL$$
 then  
OUTPUT  $(p)$ ; (The procedure was successful.)  
STOP.

Step 5 Set 
$$i = i + 1$$
.

Step 6 Set 
$$p_0 = p$$
. (Update  $p_0$ .)

#### Theorem 2.4 (Fixed-Point Theorem)

Let  $g \in C[a,b]$  be such that  $g(x) \in [a,b]$ , for all x in [a,b]. Suppose, in addition, that g' exists on (a,b) and that a constant 0 < k < 1 exists with

$$|g'(x)| \le k$$
, for all  $x \in (a, b)$ .

Then for any number  $p_0$  in [a, b], the sequence defined by

$$p_n = g(p_{n-1}), n \ge 1,$$

converges to the unique fixed point p in [a, b].

#### Newton's

To find a solution to f(x) = 0 given an initial approximation  $p_0$ :

INPUT initial approximation  $p_0$ ; tolerance TOL; maximum number of iterations  $N_0$ .

OUTPUT approximate solution p or message of failure.

```
Step 1 Set i = 1.
```

Step 2 While  $i \le N_0$  do Steps 3–6.

**Step 3** Set 
$$p = p_0 - f(p_0)/f'(p_0)$$
. (Compute  $p_i$ .)

Step 4 If 
$$|p - p_0| < TOL$$
 then OUTPUT  $(p)$ ; (The procedure was successful.) STOP.

Step 5 Set 
$$i = i + 1$$
.

Step 6 Set 
$$p_0 = p$$
. (Update  $p_0$ .)

Step 7 OUTPUT ('The method failed after  $N_0$  iterations,  $N_0 = N_0$ ); (The procedure was unsuccessful.) STOP.

#### Newton's Method

To determine a root of f(x) = 0, given  $x_0$  reasonably close to the root,

```
Compute f(x_0), f'(x_0).

If (f(x_0) \neq 0) And (f'(x_0) \neq 0) Then

Repeat

Set x_1 = x_0.

Set x_0 = x_0 - f(x_0)/f'(x_0).

Until (|x_1 - x_0| < \text{tolerance value 1}) Or |f(x_0)| < \text{tolerance value 2}).

End If.
```

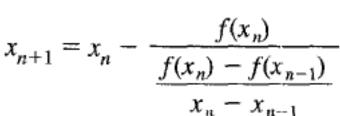
*Note:* The method may converge to a root different from the expected one or diverge if the starting value is not close enough to the root.

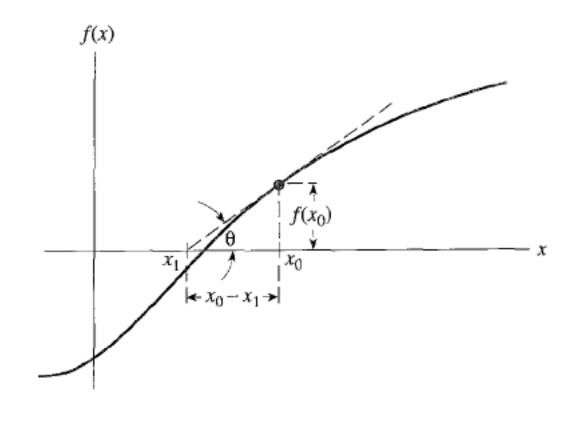
$$\tan \theta = f'(x_0) = \frac{f(x_0)}{r_0 - x_1},$$

$$x_1 = x_0 - \frac{f(x_0)}{f'(x_0)}.$$

$$x_2 = x_1 - \frac{f(x_1)}{f'(x_1)},$$

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}, \qquad n = 0, 1, 2, \dots.$$





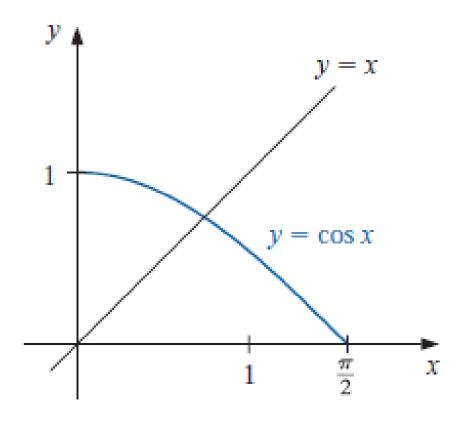
**Example 1** Consider the function  $f(x) = \cos x - x = 0$ . Approximate a root of f using (a) a fixed-point method, and (b) Newton's Method

Table 2.3

n	$p_n$
0	0.7853981635
1	0.7071067810
2	0.7602445972
3	0.7246674808
4	0.7487198858
5	0.7325608446
6	0.7434642113
7	0.7361282565

Table 2.4

	Newton's Method
n	$p_n$
0	0.7853981635
1	0.7395361337
2	0.7390851781
3	0.7390851332
4	0.7390851332



## Comparison b/w Secant and Newton

$$p_n = p_{n-1} - \frac{f(p_{n-1})}{f'(p_{n-1})}, \quad \text{for } n \ge 1.$$

$$p_n = p_{n-1} - \frac{f(p_{n-1})(p_{n-1} - p_{n-2})}{f(p_{n-1}) - f(p_{n-2})}.$$

### An Algorithm for the Secant Method

To determine a root of f(x) = 0, given two values,  $x_0$  and  $x_1$ , that are near the root,

If 
$$|f(x_0)| < |f(x_1)|$$
 Then  
Swap  $x_0$  with  $x_1$ .  
Repeat
$$Set x_2 = x_1 - f(x_1) * \frac{x_0 - x_1}{f(x_0) - f(x_1)}$$
Set  $x_0 = x_1$ .  
Set  $x_1 = x_2$ .  
Until  $|f(x_2)| <$  tolerance value.

*Note:* If f(x) is not continuous, the method may fail.

#### The Secant Method

Example 2 Use the Secant method to find a solution to  $x = \cos x$ , and compare the approximations with those given in Example 1 which applied Newton's method.

	Newton
n	$p_n$
0	0.7853981635
1	0.7395361337
2	0.7390851781
3	0.7390851332
4	0.7390851332

Secant				
n	$p_n$			
0	0.5			
1	0.7853981635			
2	0.7363841388			
3	0.7390581392			
4	0.7390851493			
5	0.7390851332			

Table 2.5

### An Algorithm for the Method of False Position (regula falsi)

To determine a root of f(x) = 0, given two values of  $x_0$  and  $x_1$  that bracket a root: that is,  $f(x_0)$  and  $f(x_1)$  are of opposite sign,

#### Repeat

Set 
$$x_2 = x_1 - f(x_1) * \frac{x_0 - x_1}{f(x_0) - f(x_1)}$$
  
If  $f(x_2)$  is of opposite sign to  $f(x_0)$  Then Set  $x_1 = x_2$   
Else Set  $x_0 = x_2$   
End If.
Until  $|f(x_2)| <$  tolerance value.

#### The Method of False Position

Example 3 Use the method of False Position to find a solution to x = cos x, and compare the approximations with those given in Example 1 which applied fixed-point iteration and Newton's method, and to those found in Example 2 which applied the Secant method.

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	False Position	Secant	Newton
n	$p_n$	$p_n$	$p_n$
0	0.5	0.5	0.7853981635
1	0.7853981635	0.7853981635	0.7395361337
2	0.7363841388	0.7363841388	0.7390851781
3	0.7390581392	0.7390581392	0.7390851332
4	0.7390848638	0.7390851493	0.7390851332
5	0.7390851305	0.7390851332	
6	0.7390851332		

## Comparison of methods, $f(x) = 3x + \sin(x) - e^x = 0$ , $x_0 = 0$ , $x_1 = 1$

	Interval halving		False position		Secant method	
Iteration	x	f(x)	x	f(x)	х	f(x)
1	0.5	0.330704	0.470990	0.265160	0.470990	0.265160
2	0.25	-0.286621	0.372277	0.029533	0.372277	0.029533
3	0.375	0.036281	0.361598	$2.94 * 10^{-3}$	0.359904	$-1.29*10^{-3}$
4	0.3125	-0.121899	0.360538	$2.90*10^{-4}$	0.360424	$5.55 * 10^{-6}$
5	0.34375	-0.041956	0.360433	$2.93 * 10^{-5}$	0.360422	$3.55 * 10^{-7}$
Error after 5 iterations	0.0	01667	-1.1	7 * 10 <sup>-5</sup>	<-	1 * 10 <sup>-7</sup>

## Features of Bisection:

- √ Type closed bracket
- ✓ No. of initial guesses 2
- ✓ Convergence linear
- ✓ Rate of convergence slow but steady
- ✓ Accuracy good
- ✓ Programming effort easy
- ✓ Approach middle point

# Features of Reguli false:

- ✓ No. of initial guesses 2
- √ Type closed bracket
- ✓ Convergence linear
- ✓ Rate of convergence slow
- ✓ Accuracy good
- ✓ Approach interpolation
- ✓ Programming effort easy

## Features of Newton Raphson:

- ✓ Type open bracket
- ✓ No. of initial guesses 1
- ✓ Convergence quadratic
- ✓ Rate of convergence faster
- ✓ Accuracy good
- ✓ Programming effort easy
- ✓ Approach Taylor's series

## **Futures of Secant Method:**

- ✓ No. of initial guesses 2
- ✓ Type open bracket
- ✓ Rate of convergence faster
- ✓ Convergence super linear
- ✓ Accuracy good
- ✓ Approach interpolation
- ✓ Programming effort tedious

# Summary

Method	Pros	Cons
Bisection	<ul> <li>Easy, Reliable, Convergent</li> <li>One function evaluation per iteration</li> <li>No knowledge of derivative is needed</li> </ul>	<ul><li>Slow</li><li>Needs an interval [a,b] containing the root, i.e., f(a)f(b)&lt;0</li></ul>
Newton	<ul><li>- Fast (if near the root)</li><li>- Two function evaluations per iteration</li></ul>	<ul> <li>May diverge</li> <li>Needs derivative and an initial guess x<sub>0</sub> such that f'(x<sub>0</sub>) is nonzero</li> </ul>
Secant	<ul> <li>Fast (slower than Newton)</li> <li>One function evaluation per iteration</li> <li>No knowledge of derivative is needed</li> </ul>	- May diverge - Needs two initial points guess x <sub>0</sub> , x <sub>1</sub> such that f(x <sub>0</sub> )- f(x <sub>1</sub> ) is nonzero