

211410641

Ans:- to task 3

As we know

$V = \text{Vertex}$ & $E = \text{Edge}$

So for the problem 1 & problem 2 inside the dijkstra function, we see that there is one 'for loop' and another one is 'while loop'. If we look we can see that inside the "while loop" there are two "for loops" according to my code.

Now,

for "while loop" time complexity is $O(V)$
again, inside the "while loop" first 'for loop' time complexity is $O \rightarrow (E)$ & second loop time complexity is $O(V)$ so

$$\cancel{O(V+E)}$$

$$\Rightarrow \cancel{O(E) \cdot O(V) + O(V)}$$

$$\Rightarrow O(V) \cdot O(V+E) + O(V)$$

$$\Rightarrow O(V^2 + VE) + O(V)$$

$$\Rightarrow O(V^2)$$

$$\Rightarrow O(N^2)$$

As $N = \text{number of places}$

Similarly for ~~for~~ problem 2

$$\Rightarrow O(V) \cdot O(V+E) + O(V)$$

$$\Rightarrow O(V^2 + VE) + O(V)$$

$$\Rightarrow O(V^2) \Rightarrow O(N^2)$$

Now, if the number of fences in each road is exactly 1, BFS algorithm can solve this problem with $O(N+M)$ time complexity.

$N \Rightarrow$ places

$M \Rightarrow$ roads

We know BFS will search for the lowest number of roads needed to reach the final goal. For this, BFS algo. will need 3 arguments: The graph, starting place and final destination.

As we have n places and m roads the time complexity will be $O(N+M)$

As we know time complexity depends
of no of edges & vertices. The more
edges V we have more time complexity
will be needed to reach

As we are adding new edges, & vertices

So this is proportional to $(V+E)$

$$\text{Time complexity} = O(V+E)$$