1 SchemeFuzzyME

This scheme is only applicable to symmetric groups of prime orders.

1.1 Setup $(n, d) \rightarrow (mpk, msk)$

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\begin{split} g \leftarrow 1_{\mathbb{G}_1} \\ \text{generate } g_2, g_3 \in \mathbb{G}_1 \text{ randomly} \\ \text{generate } \vec{t} \leftarrow \{t_1, t_2, \cdots, t_{n+1}\} \in \mathbb{G}_1^{n+1} \text{ randomly} \\ \text{generate } \vec{t} \leftarrow \{t_1, t_2, \cdots, t_{n+1}\} \in \mathbb{G}_1^{n+1} \text{ randomly} \\ \text{generate } \alpha, \beta, \theta_1, \theta_2, \theta_3, \theta_4 \in \mathbb{Z}_r \text{ randomly} \\ g_1 \leftarrow g^{\alpha} \\ \eta_1 \leftarrow g^{\theta_1} \\ \eta_2 \leftarrow g^{\theta_2} \\ \eta_3 \leftarrow g^{\theta_3} \\ \eta_4 \leftarrow g^{\theta_4} \\ Y_1 \leftarrow \hat{e}(g_1, g_2)^{\theta_1 \theta_2} \\ Y_2 \leftarrow \hat{e}(g_3, g^{\beta})^{\theta_1 \theta_2} \\ Y_2 \leftarrow \hat{e}(g_3, g^{\beta})^{\theta_1 \theta_2} \\ H_1 : \{0, 1\}^* \rightarrow \mathbb{G}_1 \\ mpk \leftarrow (g_1, g_2, g_3, Y_1, Y_2, \vec{t}, \vec{l}, \eta_1, \eta_2, \eta_3, \eta_4, H_1) \\ msk \leftarrow (\alpha, \beta, \theta_1, \theta_2, \theta_3, \theta_4) \\ \mathbf{return} \ (mpk, msk) \end{split}
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1.2 $\mathbf{EKGen}(S_A) \to e\mathbf{k}_{S_A}$

$$\begin{split} g &\leftarrow 1_{\mathbb{G}_1} \\ \Delta : i, S, x &\rightarrow \prod_{j \in S, j \neq i} \frac{x - j}{i - j} \\ N &\leftarrow (1, 2, \cdots, n + 1) \\ H : x &\rightarrow g_3^{x^n} \prod_{i=1}^{n+1} l_i^{\Delta(i, N, x)} \\ \text{generate a } (d - 1) \text{ degree polynominal } q(x) \text{ s.t. } q(0) = \beta \text{ randomly } \\ \text{generate } \vec{r} &\leftarrow \{r_1, r_2, \cdots, r_n\} \in \mathbb{Z}_r^n \text{ randomly } \\ E_i &\leftarrow g_3^{q(a_i)\theta_1\theta_2} H(a_i)^{r_i}, \forall i \in \{1, 2, \cdots, n\} \\ e_i &\leftarrow g^{r_i}, \forall i \in \{1, 2, \cdots, n\} \\ e_i &\leftarrow \{E_i, e_i\}_{a_i \in S_A} \\ \mathbf{return } ek_{S_A} \end{split}$$

1.3 $\mathrm{DKGen}(id_R) o dk_{id_R}$

$$g \leftarrow 1_{\mathbb{G}_{1}}$$

$$\Delta : i, S, x \rightarrow \prod_{j \in S, j \neq i} \frac{x - j}{i - j}$$

$$N \leftarrow (1, 2, \cdots, n + 1)$$

$$T : x \rightarrow g_{2}^{x^{n}} \prod_{i=1}^{n+1} t_{i}^{\Delta(i, N, x)}$$

$$H : x \rightarrow g_{3}^{x^{n}} \prod_{i=1}^{n+1} l_{i}^{\Delta(i, N, x)}$$
generate $\gamma \in \mathbb{Z}_{r}$ randomly
generate $G_{ID} \in \mathbb{G}_{1}$ randomly

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generate a (d-1) degree polynominal f(x) s.t. f(0) = \alpha randomly
generate a (d-1) degree polynominal h(x) s.t. h(0) = \gamma randomly
generate a (d-1) degree polynominal q'(x) s.t. q'(0) = \beta randomly
generate \vec{k}_1 \leftarrow \{k_{1,1}, k_{1,2}, \cdots, k_{1,n}\} \in \mathbb{Z}_r^n randomly
generate \vec{k}_2 \leftarrow \{k_{2,1}, k_{2,2}, \cdots, k_{2,n}\} \in \mathbb{Z}_r^n randomly generate \vec{r}_1' \leftarrow \{r_{1,1}', r_{1,2}', \cdots, r_{1,n}'\} \in \mathbb{Z}_r^n randomly generate \vec{r}_2' \leftarrow \{r_{2,1}', r_{2,2}', \cdots, r_{2,n}'\} \in \mathbb{Z}_r^n randomly dk_{S_{B_{0,i}}} \leftarrow g^{k_{1,i}\theta_1\theta_2 + k_{2,i}\theta_3\theta_4}, \forall i \in \{1, 2, \cdots, n\}
dk_{S_{B_1}} \leftarrow g_2^{-f(b_i)\theta_2}(G_{ID})^{-h(b_i)\theta_2}[T(b_i)]^{-k_{1,i}\theta_2}, \forall i \in \{1, 2, \dots, n\}
dk_{S_{B_2}} \leftarrow g_2^{-f(b_i)\theta_1}(G_{ID})^{-h(b_i)\theta_1}[T(b_i)]^{-k_{1,i}\theta_1}, \forall i \in \{1, 2, \dots, n\}
dk_{S_{B_3,i}} \leftarrow [T(b_i)]^{-k_{2,i}\theta_4}, \forall i \in \{1, 2, \cdots, n\}
dk_{S_{B_{4,i}}} \leftarrow [T(b_i)]^{-k_{2,i}\theta_3}, \forall i \in \{1, 2, \cdots, n\}
dk_{S_{B}} \leftarrow (dk_{S_{B_{0}}}, dk_{S_{B_{1}}}, dk_{S_{B_{2}}}, dk_{S_{B_{3}}}, dk_{S_{B_{4}}})dk_{P_{A_{0,i}}} \leftarrow g^{r'_{i,1}\theta_{1}\theta_{2} + r'_{i,2}\theta_{3}\theta_{4}}, \forall i \in \{1, 2, \cdots, n\}
dk_{P_{A_{1,i}}} \leftarrow g_2^{-2q'(a_i)\theta_2}(G_{ID})^{h(a_i\theta_2)}H(a_i)^{-r'_{1,i}\theta_2}, \forall i \in \{1, 2, \cdots, n\}
dk_{P_{A_2,i}} \leftarrow g_2^{-2q'(a_i)\theta_1}(G_{ID})^{h(a_i\theta_1)}H(a_i)^{-r'_{1,i}\theta_1}, \forall i \in \{1, 2, \cdots, n\}
dk_{P_{A_2}} \leftarrow [H(a_i)]^{-r'_{2,i}\theta_4}, \forall i \in \{1, 2, \cdots, n\}
dk_{P_{A_{3,i}}} \leftarrow [H(a_i)]^{-r'_{2,i}\theta_3}, \forall i \in \{1, 2, \cdots, n\}
 dk_{P_A} \leftarrow (dk_{P_{A_0}}, dk_{P_{A_1}}, dk_{P_{A_2}}, dk_{P_{A_3}}, dk_{P_{A_4}})
 dk_{S_B,P_A} \leftarrow (d\vec{k}_{S_B}, d\vec{k}_{P_A})
 return dk_{S_B,P_A}
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1.4 Encryption $(ek_{S_A}, M) \rightarrow CT$

$$\begin{split} g &\leftarrow 1_{\mathbb{G}_1} \\ \Delta &: i, S, x \rightarrow \prod_{j \in S, j \neq i} \frac{x-j}{i-j} \\ N &\leftarrow (1, 2, \cdots, n+1) \\ T &: x \rightarrow g_2^{x^n} \prod_{i=1}^{n+1} t_i^{\Delta(i,N,x)} \\ H &: x \rightarrow g_3^{x^n} \prod_{i=1}^{n+1} l_i^{\Delta(i,N,x)} \\ \text{generate } s, s_1, s_2, \tau \in \mathbb{Z}_r \text{ randomly} \\ K_s &\leftarrow Y_1^s \\ K_l &\leftarrow Y_2^s \cdot \hat{e}(g_3, g^{-\tau}) \\ C_0 &\leftarrow M \cdot K_s \cdot K_l \\ C_1 &\leftarrow \eta_1^{s-s_1} \\ C_2 &\leftarrow \eta_2^{s_1} \\ C_3 &\leftarrow \eta_3^{s-s_2} \\ C_4 &\leftarrow \eta_4^{s_2} \\ C_{1,i} &\leftarrow T(b_i)^s, \forall b_i \in P_B \\ C_{2,i} &\leftarrow H(a_i)^s, \forall a_i \in S_A \\ \text{generate a } (d-1) \text{ degree polynominal } l(x) \text{ s.t. } l(0) = \tau \text{ randomly} \\ \text{generate } \vec{\xi} &\leftarrow \{\xi_1, \xi_2, \cdots, \xi_n\} \in \mathbb{Z}_r^n \text{ randomly} \\ \text{generate } \vec{\chi} &\leftarrow \{\chi_1, \chi_2, \cdots, \chi_n\} \in \mathbb{Z}_r^n \text{ randomly} \\ C_{3,i} &\leftarrow e_i \cdot g^{\xi_i}, \forall i \in \{1, 2, \cdots, n\} \end{split}$$

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\begin{split} &C_{4,i} \leftarrow g^{\chi_i}, \forall i \in \{1,2,\cdots,n\} \\ &C_{5,i} \leftarrow E_i^s \cdot g_3^{l(a_i)} H(a_i)^{s \cdot \xi_i} \cdot H_1(C_0||C_1||C_2||C_3||C_4||C_{1,i}||C_{2,i}||C_{3,i}||C_{4,i})^{\chi_i} \\ &CT \leftarrow (C_0,C_1,C_2,C_3,C_4,\vec{C}_1,\vec{C}_2,\vec{C}_3,\vec{C}_4,\vec{C}_5) \\ &\mathbf{return} \ \ CT \end{split}
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1.5 Decryption($dk_{S_B,P_A}, S_B, P_A, CT$) $\rightarrow M$

$$\begin{split} W_{A}' &\leftarrow S_{A} \cap P_{A} \\ W_{B}' &\leftarrow S_{B} \cap P_{B} \\ \text{if } |W_{A}'| &\leq d \wedge |W_{B}'| \leq d \text{ then} \\ \text{generate } W_{A} \subset W_{A}' \text{ s.t. } |W_{A}| = d \text{ randomly} \\ \text{generate } W_{B} \subset W_{B}' \text{ s.t. } |W_{B}| = d \text{ randomly} \\ g \leftarrow 1_{\mathbb{G}_{1}} \\ \Delta &: i, S, x \rightarrow \prod_{j \in S, j \neq i} \frac{x - j}{i - j} \\ K_{S}' &\leftarrow \prod_{b_{i} \in W_{B}} (\hat{e}(C_{1,i}, dk_{S_{B_{0,i}}}) \hat{e}(C_{1}, dk_{S_{B_{1,i}}}) \hat{e}(C_{2}, dk_{S_{B_{2,i}}}) \hat{e}(C_{3}, dk_{S_{B_{3,i}}}) \hat{e}(C_{4}, dk_{S_{B_{4,i}}}))^{\Delta(b_{i}, W_{B}, 0)} \\ CT_{i} &\leftarrow C_{0} ||C_{1}||C_{2}||C_{3}||C_{4}||C_{1,i}||C_{2,i}||C_{3,i}||C_{4,i}, \forall i \in \{1, 2, \cdots, n\} \\ K_{I}' &\leftarrow \prod_{a_{i} \in W_{A}} \left(\frac{\hat{e}(C_{1,i}, dk_{P_{A_{0,i}}}) \hat{e}(C_{1}, dk_{P_{A_{1,i}}}) \hat{e}(C_{2}, dk_{P_{A_{i,2}}})}{\hat{e}(H_{1}(CT_{i}), C_{4,i}) \cdot \hat{e}(C_{3,i}, C_{2,i})} \cdot \hat{e}(C_{3}, dk_{P_{A_{i,3}}}) \hat{e}(C_{4}, dk_{P_{A_{i,4}}}) \hat{e}(C_{5,i}, g)\right)^{\Delta(a_{i}, W_{A}, 0)} \\ M &\leftarrow C_{0} \cdot K_{S}' \cdot K_{I}' \\ \text{else} \\ M &\leftarrow \bot \\ \text{end if} \\ \text{return } M \end{split}$$