1 SchemeIBMEMR

This scheme is only applicable to symmetric groups of prime orders.

1.1 Setup $(d) \rightarrow (mpk, msk)$

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p \leftarrow \|\mathbb{G}\|
g \leftarrow 1_{\mathbb{G}_1}
H_1: \mathbb{Z}_r \to \mathbb{G}_1
H_2: \mathbb{Z}_r \to \mathbb{G}_1
\hat{H}: \mathbb{G}_T \to \{0,1\}^{\lambda}
H_3: \{0,1\}^* \to \mathbb{Z}_r
H_4: \mathbb{G}_T \to \mathbb{Z}_r
H_5: \{0,1\}^* \to \mathbb{G}_1
generate g_0, g_1 \in \mathbb{G}_1 randomly
generate w, \alpha, \gamma, k, t_1, t_2 \in \mathbb{Z}_r randomly
\Omega \leftarrow e(g,g)^w
v_1 \leftarrow g^{t_1}
v_2 \leftarrow g^{t_2}
v_3 \leftarrow g^{\gamma}
v_4 \leftarrow g^k
mpk \leftarrow (p, g, g_0, g_1, v_1, v_2, v_3, v_4, \Omega, H_1, H_2, H_3, H_4, H_5, \hat{H})
msk \leftarrow (w, \alpha, \gamma, k, t_1, t_2)
return (mpk, msk)
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$1.2 \quad \mathrm{EKGen}(id_S) ightarrow ek_{id_S}$

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ek_{id_S} \leftarrow H_1(id_S)^{\alpha}

return ek_{id_S}
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$1.3 \quad \mathrm{DKGen}(id_R) ightarrow dk_{id_R}$

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dk_1 \leftarrow H_2(id_R)^{\alpha}
dk_2 \leftarrow g^{w/t_1}(g_0g_1^{id_R})^{\gamma/t_1}
dk_3 \leftarrow g^{w/t_2}(g_0g_1^{id_R})^{\gamma/t_2}
dk_{id_R} \leftarrow (dk_1, dk_2, dk_3)
return dk_{id_R}
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$1.4 \quad ext{TDKGen}(id_R) ightarrow td_{id_R}$

$$\begin{array}{l} td_1 \leftarrow g^{-1/t_1}(g_0g_1^{id_R})^{k/t_1} \\ td_2 \leftarrow g^{-1/t_2}(g_0g_1^{id_R})^{k/t_2} \\ td_{id_R} \leftarrow (td_1, td_2) \\ \mathbf{return} \ td_{id_R} \end{array}$$

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1.5 \operatorname{Enc}(\boldsymbol{ek_{id_S}}, \boldsymbol{id_R}, m) \rightarrow \boldsymbol{ct}
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generate S \leftarrow (id_1, id_2, \cdots, id_R, \cdots, id_d) randomly
generate s_1, s_2, \beta, \sigma, K, R \in \mathbb{Z}_r randomly
r \leftarrow H_3(\sigma||m)
ct_1 \leftarrow g^{\beta}
ct_2 \leftarrow v_1^{s_1}
ct_3 \leftarrow v_2^{s_2}
K_i \leftarrow e(H_2(id_i), ek_{id_S} \cdot ct_1), \forall i \in \{1, 2, \cdots, d\}
Compute a_0, a_1, a_2, \dots a_d that satisfy \forall x \in \mathbb{Z}_r, we have F(x) = \prod_{i=1}^d (x - H_4(K_i)) +
K = a_0 + \sum_{i=1}^{d} a_i x^is \leftarrow s_1 + s_2
R_i \leftarrow e(v_3, (g_0 g_1^{id_i})^s), \forall i \in \{1, 2, \cdots, d\}
Compute b_0, b_1, b_2, \dots, b_d that satisfy \forall x \in \mathbb{Z}_r, we have L(x) = \prod_{i=1}^d (x - H_4(R_i + R_i))
e(g,g)^{ws}) + R = b_0 + \sum_{i=1}^{d} b_i x^i
ct_4 \leftarrow \hat{H}(K) \oplus \hat{H}(R) \oplus (m||\sigma)
V_i \leftarrow e(v_4, (g_0g_1^{id_i})^s), \forall i \in \{1, 2, \cdots, d\}
Compute c_0, c_1, c_2, \dots, c_d that satisfy \forall x \in \mathbb{Z}_r, we have G(x) = \prod_{i=1}^d (x - H_4(V_i))
e(g,g)^{-s}) = c_0 + \sum_{i=1}^{d} c_i x^i
ct_6 \leftarrow H_5(ct_1||ct_2||\cdots||ct_5||a_0||a_1||\cdots||a_d||b_0||b_1||\cdots||b_d||c_0||c_1||\cdots||c_d)^r
ct \leftarrow (ct_1, ct_2, ct_3, ct_4, ct_5, ct_6)
return ct
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1.6 $\operatorname{Dec}(dk_{id_R}, id_R, id_S, ct) \rightarrow m$

return m

 $\begin{array}{l} \textbf{if } e(ct_5, H_5(ct_1||ct_2||\cdots||ct_5||a_0||a_1||\cdots||a_d||b_0||b_1||\cdots||b_d||c_0||c_1||\cdots c_d)) = e(ct_6, g) \\ \textbf{then} \\ K'' \leftarrow H_4(e(dk_1, H_1(id_S)) \cdot e(H_2(id_R), ct_1)) \\ R'' \leftarrow H_4(e(dk_2, ct_2) \cdot e(dk_3, ct_3)) \\ K' \leftarrow \sum_{i=0}^d a_i K''^i \\ R' \leftarrow \sum_{i=0}^d b_i R''^i \\ m||\sigma \leftarrow ct_4 \oplus \hat{H}(K') \oplus \hat{H}(R') \\ r \leftarrow H_3(m||\sigma) \\ \textbf{if } ct_5 \neq g^r \textbf{ then} \\ m \leftarrow \bot \\ \textbf{end if} \\ \textbf{else} \\ m \leftarrow \bot \\ \textbf{end if} \end{array}$

1.7 ReceiverVerify $(ct, td_{id_R}) \rightarrow y, y \in \{0, 1\}$

$$\begin{array}{l} \textbf{if } e(ct_5, H_5(ct_1||ct_2||\cdots||ct_5||a_0||a_1||\cdots||a_d||b_0||b_1||\cdots||b_d||c_0||c_1||\cdots c_d)) = e(ct_6, g) \\ \textbf{then} \\ V' \leftarrow H_4(e(td_1, ct_2) \cdot e(td_2, ct_3)) \\ y \leftarrow \sum_{i=0}^d c_i V'^i = 0 \\ \textbf{else} \\ y \leftarrow 0 \\ \textbf{end if} \\ \textbf{return } y \end{array}$$