Homework 0

Problem 1: (Stable Marriage Problem)

- (i) SMP is the problem of finding a stable matching between two disjoint equally sized sets. SMP can be presented as computationally as shown as below.
 - o Input:
 - A set of men and a set of women which are sized equally.

Men =
$$\{m_1, m_2, ..., m_n\}$$
 Women = $\{w_1, w_2, ..., w_n\}$

• Ordered preference list for both sets which contains all members of the opposite sex.

Ex:
$$P_{\text{Men(i)}} = \{w_8, w_{14}, ..., w_3\}$$
 ($P_{\text{Men(i)}}$ equally sized as Men) $P_{\text{Women(i)}} = \{m_8, m_{12}, ..., m_7\}$ (equally sized as Women)

- Output:
 - A stable matching of partners Ex: $SM = \{(m_4, w_{14}), ..., (m_6, w_2)\}$
 - To have a stable match;
 - SM have to be sized as equally as Men and Women.
 - Nobody should be left out in SM.
 - \lor $\forall_{w \in Women} \exists_{m \in Men} ((m, w) \in SM)$
 - $\circ \forall_{m \in Men} \exists_{w \in Women} ((m, w) \in SM)$
 - There have to be no blocking pairs which means, at the end if m_x and w_x are not partners, there must be no case such as, m_x prefers w_x to P_{Man}(w_i) and/or w_x prefers m_x to P_{Women}(m_i).
- (ii)
- o Input:
 - A set of 3 NBA teams: $T = \{t_1, t_2, t_3\}$
 - A set of 3 players: $P = \{p_1, p_2, p_3\}$
- Output:
 - A stable matching of 3 NBA teams and players. SM_{t-p}
- Example of input and output:
 - Input:
 - $T = \{t_1, t_2, t_3\}$
 - $P = \{p_1, p_2, p_3\}$
 - $\bullet \quad P_{player(1)} = \left\{t_1, t_2, t_3\right\}, \ P_{player(2)} = \left\{t_2, t_1, t_3\right\}, \ P_{player(3)} = \left\{t_1, t_2, t_3\right\}$
 - $P_{\text{teams}(1)} = \{p_2, p_1, p_3\}, P_{\text{teams}(2)} = \{p_1, p_2, p_3\}, P_{\text{teams}(3)} = \{p_1, p_2, p_3\}$
 - Output:
 - $SM_{t-p} = \{(t_1, p_2), (t_2, p_1), (t_3, p_3)\}$

Problem 2: (Gale-Shapley Algorithm)

- (i)
- While any man is free (let's call him x),
 - he proposes to a woman he most prefers.
 - If the woman (let's call her y) is single, then they become a couple (x, y) and this pair is added to the set SM.
 - If y is paired with another man (x'), then there are two possibilities:
 - If y prefers x over x', then she breaks the pair and is coupled with x. Then, (x, y) is added to SM. x' starts proposing to his other preferences.
 - If y prefers x' over x, then x remains single and continues proposing until he is paired with a woman.
- o Return final SM as the stable matching result.
- (ii)
- o Let's say that the set of men and the set of women are sized as n.
- o If we go over the algorithm by using the same sets that explained at above, we can analyze the asymptotic time complexity of this algorithm.
- o In the worst case, the algorithm will iterate n times, and, in each iteration, the man will propose to a woman that he didn't propose before. This operation will iterate as many times as the while loop statement which iterates n times. In conclusion, in worst case iteration of the algorithm takes n² times. In big-oh notation the asymptotic time complexity of the algorithm is O(n²).

Problem 3:

```
Teams = ['Lakers','Warriors','Celtics']
Players = ['LeBron','Curry','Kyrie']
Pteams = {'Lakers':['LeBron', 'Curry','Kyrie'],
            'Warriors':['Curry','LeBron','Kyrie'],
            'Celtics':['LeBron','Curry','Kyrie']
Pplayers = {'LeBron':['Lakers','Warriors','Celtics'],
            'Curry':['Warriors', 'Lakers','Celtics'],
            'Kyrie':['Lakers', 'Warriors','Celtics']
for i in Pteams:
    Pteams[i].reverse()
for i in Pplayers:
    Pplayers[i].reverse()
Contracts = {}
teamsWithFreeContrats = list (Teams)
numOfContracts = len (Teams)
while teamsWithFreeContrats:
    T = teamsWithFreeContrats.pop()
    P = Pteams[T].pop()
    if P not in Contracts:
        Contracts[P]=T
    else:
        pickedTeam= Contracts[P]
        if Pplayers[P].index(T) > Pplayers[P].index(pickedTeam):
            Contracts[P] = T
            teamsWithFreeContrats.append(pickedTeam)
            teamsWithFreeContrats.append(T)
print (Contracts)
```

Output:

```
{'LeBron': 'Lakers', 'Curry': 'Warriors', 'Kyrie': 'Celtics'}
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