

Numerical Analysis

Based on SU MA232

X

Elegant \LaTeX Program

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1 Preknowledge

1.1 Gaussian elimination and Decomposition de Matrix

The core of Gussian elimination is

$$E_n E_{n-1} \dots E_1 A = U$$

where E_i is the i -th elementary row operation and U is the upper triangular matrix. Above is the language of the linear operator, we can translate each E_i to be a matrix with a good form to memorize.

Remark Suppose E and A are $n \times n$ matrix, we call E an **elementary matrix** if it is one of the following three types:

- Row addition matrix: $L_i + aL_j \rightarrow L_i$

$$E = \begin{bmatrix} 1 & & & \\ & \ddots & & \\ & & a & \\ & & & \ddots \\ & & & & 1 \end{bmatrix}, \quad E_{i,j} = a$$

- Row swap matrix: $L_i \leftrightarrow L_j$

$$E = \begin{bmatrix} 1 & & & & \\ & \ddots & & & \\ & & 0 & & 1 \\ & & & \ddots & \\ & 1 & & & 0 \\ & & & & & \ddots & \\ & & & & & & 1 \end{bmatrix}, \quad \begin{cases} E_{i,j} = E_{j,i} = 1 \\ E_{i,i} = E_{j,j} = 0 \end{cases}$$

- Row multiplication matrix: $L_i \rightarrow aL_i$

$$E = \begin{bmatrix} 1 & & & & \\ & \ddots & & & \\ & & a & & \\ & & & \ddots & \\ & & & & 1 \end{bmatrix}, \quad E_{i,i} = a$$

It is easy to find the inverse of the above matrices, notice that in the case of row addition, if the scalar a is at (i, j) upon the diagonal, then EA denotes the column

addition $C_j + aC_i \rightarrow C_j$, usually we do not need it and we just use row operation, which is more clear and convenient.

1.2 LU decomposition

Definition 1.1 (Principal Minor) For a $n \times n$ matrix A we can define its k order principal minor by

$$\Delta_k = \det \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1k} \\ a_{21} & a_{22} & \cdots & a_{2k} \\ \vdots & \vdots & \ddots & \vdots \\ a_{k1} & a_{k2} & \cdots & a_{kk} \end{pmatrix}, \quad k = 1, \dots, n$$

From Gaussian elimination, we can get the **row echelon form** of a matrix by row operation to get a upper triangular matrix, so for any matrix A there exists the upper triangular U such that

$$A = E_1^{-1} E_2^{-1} \dots E_n^{-1} U$$

If row operation does not contain row swap, it is easy to find that the composition of all inverse of the row operations equals to a lower triangular matrix, that is the LU decomposition.

Formally, The LU decomposition of a square matrix is the product of a lower triangular matrix L and an upper triangular matrix U (triangular matrix have no zero at the diagonal by its definition). For convenience, the diagonal of the L are all 1 and we denote the diagonal of U to be u_1, \dots, u_n . Next we conclude the existence of the decomposition to be a theorem.

Theorem 1.1 (LU decomposition)

- (a) *If A is an invertible matrix, then it has at most one LU decomposition*
 (b) *A square matrix has a unique LU decomposition if and only if all principal minors are not null.*

Proof. We firstly prove (a). Suppose there exists two different LU decompositions for matrix, i.e. $A = LU, A = L'U'$. so we have $L(L')^{-1} = U'U^{-1}$. A contradiction occurs since the composition of lower (upper) triangular matrices are a lower (upper) triangular matrix, so the only possibility is that $L(L')^{-1} = U'U^{-1} = I$.

Now we prove the sufficiency of (b). If A has an LU decomposition, then we write A as the form of a block matrix

$$A = \begin{pmatrix} A_1 & A_2 \\ A_3 & A_4 \end{pmatrix}, L = \begin{pmatrix} L_1 & 0 \\ B & L_2 \end{pmatrix}, U = \begin{pmatrix} U_1 & C \\ 0 & U_2 \end{pmatrix}$$

where A_1, L_1, U_1 are the $k \times k$ matrices, A_4, L_2, U_2 are the $n - k \times n - k$ matrices, so clearly k order principal minor is

$$\Delta_k = \det A_1 = L_1 U_1 \neq 0$$

Finally we prove the necessary part by induction. If A is an one order square matrix, then $A = (a) = (1)(a)$ if $\det A = a \neq 0$. Suppose a $n - 1 \times n - 1$ matrix has a LU decomposition for $n \geq 2$ if all its principal minors are not null, then for a $n \times n$ matrix A we can write by block matrix with form

$$A = \begin{pmatrix} A_{n-1} & l \\ l' & a \end{pmatrix} = \begin{pmatrix} L_{n-1} U_{n-1} & l \\ l' & a \end{pmatrix} = \begin{pmatrix} L_{n-1} & 0 \\ l' U_{n-1}^{-1} & 1 \end{pmatrix} \begin{pmatrix} U_{n-1} & L_{n-1}^{-1} l \\ 0 & a \end{pmatrix}$$

where a is a real number and l, l' are $n - 1 \times 1$ and $1 \times n - 1$ respectively, so we prove the existence of the decomposition. The uniqueness is immediate from (a) since $\Delta_n = \det A \neq 0$.

And there is an useful technique to check the existence of the decomposition, it depends on the property of continuous two order principal minor. We notice that $\Delta_k = \prod_{i=0}^k u_k = u_k \Delta_{k-1}$ for $k \geq 2$ if the matrix has LU decomposition, according to this we have two immediate corollaries.

Corollary *for any $k = 2, \dots, n$, if $\Delta_k \neq 0$ but $\Delta_{k-1} = 0$, then the matrix does not have LU decomposition.*

Corollary *The digonal elements of the upper triangular matrix in LU decomposition is*

$$\begin{cases} u_k = \Delta_k / \Delta_{k-1}, k \geq 2 \\ u_1 = \Delta_1 \end{cases}$$

1.3 Cholesky Decomposition

LU Decomposition can be finer by add a digonal matrix D , then we can get a LDV decomposition with the following form

$$LDV = \begin{pmatrix} 1 & 0 & 0 & \cdots & 0 \\ l_{21} & 1 & 0 & \cdots & 0 \\ l_{31} & l_{32} & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ l_{n1} & l_{n2} & l_{n3} & \cdots & 1 \end{pmatrix} \begin{pmatrix} d_1 & 0 & 0 & \cdots & 0 \\ 0 & d_2 & 0 & \cdots & 0 \\ 0 & 0 & d_3 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & d_n \end{pmatrix} \begin{pmatrix} 1 & u_{12} & u_{13} & \cdots & u_{1n} \\ 0 & 1 & u_{23} & \cdots & u_{2n} \\ 0 & 0 & 1 & \cdots & u_{3n} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 1 \end{pmatrix}$$

If we distribute weights in average from digonal matrix to two triangular matrices by $D = \Sigma^2$, and if the decomposed matrix has a good property, then we can get the

decomposition form like $A = HH^T$, which is the Cholesky decomposition, a more effective technique than LU decomposition when solving linear equations.

It is not difficult to find that one good property is symmetric since $(HH^T)^T = HH^T$. Positive is another property to make the diagonal elements in the diagonal matrix to be positive. We conclude it by following theorem.

Theorem 1.2 *A square matrix has a Cholesky decomposition if and only if it is symmetric and positive-definite.*

Proof. We only prove the necessary of the proposition since the sufficiency is easy. We first prove that the principal matrix is all symmetric and positive-definite. Suppose $A_k = (a_{i,j})_{1 \leq i,j \leq k}$ is the k order principal matrix of A . Clearly, A_k is symmetric. Since A is positive-definite, then

$$X^T A_k X = \begin{pmatrix} X & 0 \end{pmatrix} \begin{pmatrix} A_k & B \\ C & D \end{pmatrix} \begin{pmatrix} X \\ 0 \end{pmatrix} > 0$$

where X is a k order non-zero vector.

An symmetric and positive-definite matrix must be invertible, otherwise there exists $X \in \mathbb{R}^k - \{0\}$ such that $A^k X = 0$, which implies $X^T A^k X = 0$. Hence

all principal minors are not null, and the matrix has an unique LDV decomposition. Similarly, A^T also has an unique LDV decomposition, so it implise $L^T = V$, so if we let $D = \Sigma^2$, then $H = L\Sigma$ such that $A = HH^T$.

2 Numercial schéma for ODE

Schéma is a vocabulary from french, in nnumerical anaylsis it can be thought as a progam to approximate the real function or an equation. Specially, each numercial method refers to a **algorithm** of the sequence which converges to the real value, for example, the Eluer method to some Cauchy problem can be wirtbe by a recurrance like

$$y_{n+1} = y_n + hf(t_n, y_n)$$

It implies what the sequence is or what the method of approximation is. We call the recurrance equation contianing all information ”schéma”.

In this part, **Truncation error** is the important concept, it often happens when replacing infinite precision or infinite process by finite number of calculations. For example, the form of series of exponential functions can be written as

$$e^x = 1 + x + \frac{e^x}{2!} + \dots$$

If we just choose part of terms to approximate e^x then error occurs, and Taylor-Lagrange theorem gives the boundary of the errors

2.1 Explicit schéma

The explicit schéma is the equation with the following equation:

$$y_{n+1} = y_n + hF(t_n, y_n, h) \quad (1)$$

where F is a continuous function defined on $[0, T] \times \mathbb{R}^m \times [0, 1]$.

Definition 2.1 (Convergence)

The schéma (??) is convergent if for any given initial condition $y(0)$, the continuous problem satisfies

$$\lim_{\substack{h \rightarrow 0 \\ y_0^N \rightarrow y(0)}} \sup_{0 \leq n \leq N} \|y_n^N - y(t_n)\| = 0$$

Definition 2.2 (Stable)

The schéma ?? est stable s'il existe une constante C indépendante de N telle que, pour toute suite de vecteurs $(\eta_n)_{0 \leq n \leq N}$, les suites $(y_n)_{0 \leq n \leq N}$ et $(z_n)_{0 \leq n \leq N}$ de \mathbb{R}^m définies respectivement par

$$\begin{cases} y_{n+1} = y_n + hF(t_n, y_n, h) & , 0 \leq n \leq N-1 \\ z_{n+1} = z_n + hF(t_n, z_n, h) + \eta_{n+1} & , 0 \leq n \leq N-1 \\ z_0 = y_0 + \eta_0 \end{cases}$$

sont telles que

$$\max_{0 \leq n \leq N} \|z_n - y_n\| \leq C \sum_{n=0}^N \|\eta_n\|$$

Remark Let us consider the motivation of the definition. For a discrete numerical method with initial value known, we start the algorithm from t_0 , then let $z_n = y(t_n)$ be the sequence of the real value at t_n , then at t_1 clearly we can deduce that

$$\begin{cases} y_1 = y_0 + hF(t_0, y_0, h) \\ z_1 = y_1 + \varepsilon_1 = z_0 + hF(t_0, z_0, h) + \varepsilon_1 \end{cases}$$

where ε_1 is the error between real value and approximation, we call it **Local Truncation Error (l'erreur de discrétisation)** but when we repeated algorithm we can get that

$$z_2 = z_1 + hF(t_1, z_1, h) + \varepsilon_2$$

where ε_2 is the LTE between real value $y(t_2)$ and apporximation of $y(t_2)$ starting from real value $y(t_1)$, and the iteration contains the previous error ε_1 in it. Hence the difference between $y(t_2)$ and y_2 relates to the two LTE. We define the **Global Truncation Error (l'erreur de schéma)** with the form

$$\max_{0 \leq n \leq N} \|z_n - y_n\|$$

So the definition of stable schéma means that **the GTE is bounded by the sum of LTE.**

Definition 2.3 (Consistent) *The sch (??) is consistent with the corresponding ODE if for any solution of the ODE $y(x)$ we have*

$$\lim_{h \rightarrow 0} \sum_{n=0}^{N-1} \|\varepsilon_n\| = 0$$

where $\varepsilon_n = y(t_{n+1}) - y(t_n) - hF(t_n, y(t_n), h)$ is LTE.

Immediately we can get the sufficant condition for the convergence of schéma by the squeezes theorem, where stable property ensures the inequality and consistent property controls the limit.

Theorem 2.1 (convergence of explicit schéma)

The explicit schéma is convergent if it is stable and consistent.

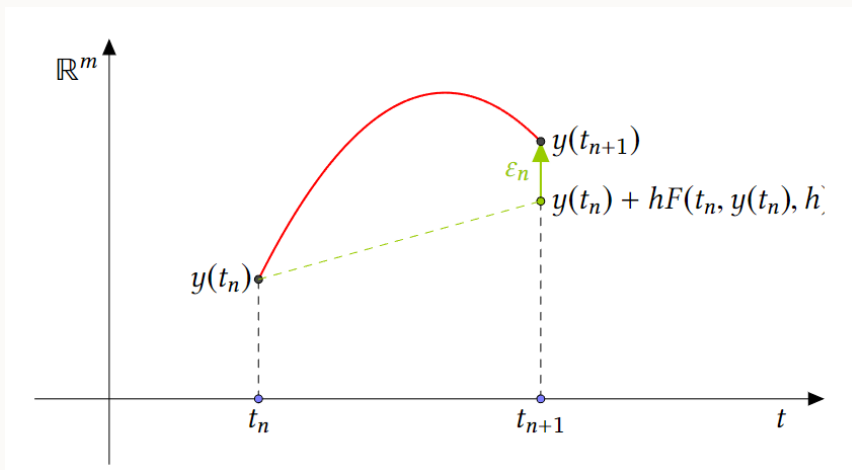


Figure 1: Local Truncation Error

We omit the proof and consider how to determine that a explicit schéma is stable and consistent, we conclude the condition below.

Proposition 2.2 *Consider the schéma (??)*

- (a) *It is stable if $F(t, y, z)$ is Lipschitz with respect to variable y*
- (b) *It is consistent if and only if $F(t, y, 0) = f(t, y)$ for any $(t, y) \in [0, T] \times \mathbb{R}^m$, where f is function in the corresponding Cauchy problem.*

2.2 Implicit schéma

The implicit schéma is the equation with the following equation:

$$y_{n+1} = y_n + hF(t_n, y_n, y_{n+1}, h) \quad (2)$$

where F is a continuous function defined on $[0, T] \times \mathbb{R}^m \times \mathbb{R}^m \times [0, 1]$.

Remark Suppose $\phi(Z) = y_n + hF(t_n, y_n, z, h)$, then the recurrence of the sequence is equivalent to the equation $z = \phi(z)$, so we should verify that the sequence given by schéma is well-defined.

To make ϕ is a contraction function, we add a condition that $F(t, y, z, h)$ is L-Lipschitz with respect to z , then we have inequality

$$\|\phi(z) - \phi(z')\| \leq |hL| \|z - z'\|$$

Which means if $h < 1/L$, then by **fixed point theorem** the sequence is well-defined.

Similarly, we can define stable and consistent like what we do in the previous section (just turn $F(t_n, y_n, y_{n+1}, h)$ into $F(t_n, y_n, h)$), then the Local Truncation Error is

$$\varepsilon_n = y(t_{n+1}) - y(t_n) - hF(t_n, y(t_n), y(t_{n+1}), h)$$

We conclude the same result about convergent of the schéma .

Proposition 2.3 *For the implicit schéma (??)*

(a) *If $F(t, y, z, h)$ is L -Lipschitz with respect to variable z and y , then for any $h < 1/L$, the schéma is stable*

(b) *If $F(t, y, y, h) = f(t, y)$ for any $(t, y) \in [0, T] \times \mathbb{R}^m$, then the schéma is consistent.*

Theorem 2.4 *The implicit schéma is convergent if it is stable and consistent.*

3 Newton's Method

3.1 One-dimension method

We suppose $y = f(x)$ is a class C^1 real function with a root x^* , the tangent line at $(x, f(x))$ can be described by

$$Y = f'(x)(x - x_0) + f(x_0)$$

Let $Y = 0$ we can find the point of intersection of the tangent line at x-axis, the formula is that

$$x' = x^* - \frac{f(x^*)}{f'(x^*)}$$

According to this, we can define a function $h_f(x) = x - \frac{f(x)}{f'(x)}$, which maps x to the intersection of the tangent line at x with x-axis, so we can give a recurrence for some class C^1 function defined on $[a, b]$

$$\begin{cases} x_0 \text{ near } x^* \\ x_{n+1} = h_f(x_n) \end{cases}$$

we call it **Newton sequence**. Then we will consider several questions: Is it the sequence well-defined? Does the sequence converge to the root? If it converges, What's the rate of the convergence? We response it by the following theorem:

Theorem 3.1 (Quadratic convergence of Newton Method)

Suppose that f is a class C^2 function with a root at c , if f defined on $I = [c - r, c + r]$ and $f'(x) \neq 0$ on I , then there exist positive value a, K such that Newton sequence

$$\begin{cases} x_0 \in (c - a, c + a) \\ x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)} \end{cases}$$

converges to c with quadratic estimation

$$|x_n - c| \leq \frac{1}{K}(K|x_0 - c|)^{2^n}$$

where $K = \frac{\max_{x \in I} |f''(x)|}{2 \min_{x \in I} |f'(x)|}$, $a = \min(r, 1/K)$.

Proof. $f'(x)$ is not null dans I means that iteration will be endless, so we just need to there exist a subinterval such that h_f is invriant on it, which ensures that the sequence is well-defined. By the formula of Taylor with Reminder,

$$\begin{aligned} 0 &= f(c) = f(x + (c - x)) \\ &= f(x) + (c - x)f'(x) + \frac{(c - x)^2}{2}f''(\eta) \end{aligned}$$

with $x \in I$ and η is between x and c , rearranging it implies

$$\left(x - \frac{f(x)}{f'(x)}\right) - c = \frac{(c - x)^2}{2} \frac{f''(\eta)}{f'(x)}$$

which gives

$$|h_f(x) - c| \leq \frac{\max_{x \in I} |f''(x)|}{2 \min_{x \in I} |f'(x)|} |x - c|^2 = K|x - c|^2$$

To make h_f invriant, then $K|x - c|^2 \leq |x - c|$, so the length of the subinterval must be no larger than $a = \min(r, 1/K)$ to ensure the sequence to converges c . And the quadratic convergence is clearly.

Example 3.1 (Héon Method) *Newton method can bu taken use to approximate the $\sqrt{2}$, we define $f(x) = x^2 - 2$, then $\sqrt{2}$ is the root of the equation $f(x) = 0$, so the*

recurrence of the sequence is

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)} = \frac{x_n}{2} + \frac{1}{x_n}$$

which is a classic problem in analysis.

3.2 N-dimension method

We can generalize Newton method to \mathbb{R}^n . We consider the affine map $g(x) = Ax - b$ with $A \in M_n(\mathbb{R})$ and $b \in \mathbb{R}^n$, if A is invertible, then the equation $g(x) = 0$ has the unique solution $A^{-1}b$, but if A is not invertible and b is in the column space of the matrix, then the situation will be complex, because the null space is always convex and connected.

Let us consider the case of $n = 2$, we take an affine map $g(x, y) = (x, 0)$, it is clear that the set of solutions is y-axis. However, if we want to approximate the root $(0, 0)$, and we choose some point near it to be the first term of Newton sequence, then the sequence can not converge to $(0, 0)$ if some term x_k is at the y-axis ($h_f(x_{k+1}) = 0$). That is not avoidable since for any small open ball containing $(0, 0)$, there exists $(0, y_0)$ of the null space in the open ball, so we need a concept about the isolation

of the root.

Definition 3.1 We call c is an isolated zero of the function f defined on U if for any $\varepsilon > 0$, there exists no other zero on $B(c, \varepsilon) \cap U$.

Proposition 3.2 (Sufficient condition of an isolated zero)

Suppose $f : U \subset \mathbb{R}^n \rightarrow \mathbb{R}^n$ is a function defined on an open set, and x is a zero of the function on U . If f is differentiable and $D_f(x)$ is invertible, then x is an isolated zero of the function.

Proof. It is a weak form of **the local inversion theorem**. Considering a function $h(y) = D_f(x)^{-1}f(y)$, it is also differentiable, and $D_h(x) = I$, then by definition of the differentiable function, we have

$$\begin{aligned} h(y) &= h(x + y - x) \\ &= h(x) + I(y - x) + o(\|y - x\|) \\ &= y - x + \varepsilon(y - x)\|y - x\| \end{aligned}$$

where ε converges to 0 when $(y - x)$ converges to 0, so there exists a positive r such that $\|x - y\| < r$ implies $\|\varepsilon\| < 1/2$. Again by the triangular inequality $|x + y| \geq ||x| - |y||$, we have

$$\|h(y)\| \geq 1/2\|y - x\| > 0$$

for any $0 < \|x - y\| < r$, by the definite property of the norm we can conclude that x is an isolated zero.

The exact form of the Newton iteration can be deduced from Talyor Formula of multivariable functions. We suppose that f is a class C^1 function with an islated root at c , then

$$0 = g(c) \approx g(x) + D_f(x)(c - x)$$

for some x approching c , so we can get

$$c \approx x - D_f(x)^{-1}f(x)$$

it gives the recurrence of the Newto method on \mathbb{R}^n

$$\begin{cases} x_0 \text{ near } c \\ x_{n+1} = x_n - D_f(x_n)^{-1}f(x_n) \end{cases}$$

If the new sequence is well-defined? To prove that we need a topological property.

Lemma 3.3 $GL_n(\mathbb{R})$ is open in the norme space $M_n(\mathbb{R})$, where the euipped norm is

$$\|A\| = \max_{1 \leq i, j \leq n} |a_{i,j}|.$$

The proof of the Lemma can be found in TD 1 of the course of Topologie, according to this we can prove the convergence theorem in n-dimension.

Theorem 3.4 (General quadratic convergence of New Mehod) *If f is a class C^2 defined on U with a zero x , if $D_f(x)$ is invertible, then there exist an open set and a postive value K such that for any given initial conition in it, the Newton sequence is well-defined and converges to x with rate de convergence*

$$\|x_{k+1} - x\| \leq K\|x_k - x\|^2$$

Proof. wait...

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