# Chapter 1

# Particle orbit pusher algorithms

This chapter is dedicated to the implemented algorithms for finding the first intersection of particle orbits with the tetrahedral cell boundaries in the grids that were previously introduced. Here, a particle can either start at an arbitrary position inside a given tetrahedron or directly at a face of a tetrahedron through which it enters. These routines efficiently compute the next exiting position of the particle through the tetrahedron and the associated flight time of the trajectory. Since this procedure can be thought of as a pushing of the particle orbit through the tetrahedron, the implemented routines are denoted pusher-routines. On a sidenote, the fact that both position and time are obtained directly by the approaches used in the pusher routines, a box counting scheme can easily be implemented for future applications, allowing for a very efficient approximation of particle distribution functions, which in turn are a necessary part for possible future computations of kinetic plasma equilibria. The focus of the pusher routines lies, however, not only on the computation of the trajectory and the calculation of the next intersection but rather on finding a numerically inexpensive scheme that allows to save computational cost while reliably yielding accurate results for the exiting position. In the diploma thesis of M. Eder [?], a prior version of the presented pusher routine was discussed in great detail, this routine was named pusher\_tetra\_orb. Due to new insights and structural limitations of the previous code, this code was refactored and extended in cooperation with M. Eder. The resulting code was named pusher\_tetra\_orbit, an overview of the code is given below, however, due to large similarities with the previous approach discussed in [?], the new route will be presented in less detail. Apart from this routine, a second routine named pusher\_tetra\_analytic was implemented based on the previously derived polynomial expansion of the particle orbit. While the results are in theory equivalent for both pushing routines, the approaches are completely independent and thus may vary in both computational efficiency and numerical accuracy, depending on up to which order the analytical expansion of the orbit is computed. Furthermore, for starting a particle at a given position without knowing to which tetrahedron it belongs, an additional routine find\_tetra was constructed to find the corresponding tetrahedron index to start a calculation.

# 1.1 Pusher routine pusher\_tetra\_orbit

As discussed, the pusher subroutine pusher\_tetra\_orbit computes the position and time where the particle trajectory first exits a given current tetrahedron. In reality, however, the occurring problem is not only to directly compute the orbit of a single tetrahedron passing, but rather to let a particle start at a position in space and trace its orbit for a defined time. For such a problem one can construct a wrapping routine orbit\_timestep\_3dgeoint which is given the initial conditions of the particle and iteratively calls the pusher subroutine pusher\_tetra\_orbit until the set time is reached. Since generally the set flight time of the particle will lead to an orbit position inside the final tetrahedron, the remaining time of the trajectory must also be given to the pusher routine. The pusher routine then computes the time it takes until the particle exits the current tetrahedron and compares this value to the remaining time of the orbit integration step which was given to the wrapper routine. In case the time it takes to leave the tetrahedron is smaller than the remaining time, the pushing is computed, then the remaining time is reduced by this value and the next pushing through the adjacent tetrahedron is started. In case there is not sufficient time to complete the pushing, the orbit is instead integrated up to the value of the orbit parameter tau corresponding to the remaining time, leading to an arbitrary final position inside the tetrahedron. The code structure of the wrapping routine orbit\_timestep\_3dgeoint and the components of the module pusher\_tetra\_orbit\_mod is given in figure 1.1.

Due to this wrapping routine, one can directly start the computation of a single particle orbit for a given flight time by calling subroutine orbit\_timestep\_3dgeoint with arguments (x,vpar,vperp,t\_step,boole\_initialized,ind\_tetr,iface). This list of parameters is explained in tab. 1.1.

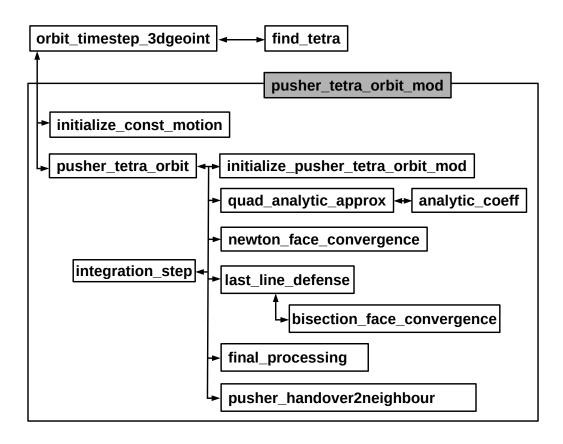


Figure 1.1: Code structure of pusher\_tetra\_orbit\_mod and associated subroutine

Table 1.1: Parameter overview for wrapper subroutine orbit\_timestep\_3dgeoint, defines initial conditions and duration of particle motion, orientations of velocities are taken with respect to the orientation of the magnetic field  $\vec{B}$ .

Data type	Name	Description
double precision,	х	particle position
<pre>dimension(3),intent(inout)</pre>		
double precision,	vpar	parallel velocity
intent(inout)		
double precision,	vperp	perpendicular velocity
intent(inout)		
double precision, intent(in)	t_step	defined flight time
logical, intent(inout)	boole_initialized	sets initialization of
		constants of motion
<pre>integer,intent(inout)</pre>	ind_tetr	tetrahedron index
		at position x
<pre>integer,intent(inout)</pre>	iface	index of face if x lies
		on face, 0 otherwise

## 1.1.1 Initializing constants of motion

In figure 1.1, one can see the code diagram which gives an overview of the different subroutines. When starting a calculation in orbit\_timestep\_3dgeoint for a defined step length, the subroutine that is called first is initialize\_const\_motion which sets the constants of motion for the given initial conditions. These constants of motion are E\_tot, perpinv and perinv2 which denote the total energy E, the negative perpendicular adiabatic invariant  $-J_{\perp}$  and the squared value thereof, respectively. Since these quantities are saved with attributes public, protected, the subroutine initialize\_const\_motion must be saved within the current module pusher\_tetra\_orbit\_mod, otherwise it would not be allowed to set the values. The constants of motion will retain their set values for a number of tetrahedral pushings until the next time step is executed. Usually, between time steps collision events will occur when performing Monte Carlo simulations, as a consequence the constants of motion may change and have to be defined anew.

## 1.1.2 Particle pushing algorithm

For a given time step, after initializing the constants of motion, the subroutine pusher\_tetra\_orbit is called with the initial conditions of the current pushing. At the end of the subroutine execution it returns the new starting conditions for the next pushing as well as the remaining time of the current step. An overview of the call parameters of the subroutine is given in tab. 1.2.

Table 1.2: Parameter overview for pusher\_tetra\_orbit

Data type	Name	Description
integer,	ind_tetr_inout	current tetrahedron index
intent(inout)		
integer,	iface	current face index of tetrahedron
intent(inout)		where orbit converged, 0 if not converged
double precision,	х	current particle position in global
dimension(3),		coordinates, i.e. not with respect
intent(inout)		to the first node of a tetrahedron
double precision,	vpar	parallel velocity of the
intent(inout)		particle with respect to ${f B}$
double precision,	z_final	final particle position in <i>local</i> coordinates,
dimension(3),		needed for calculation of the flux tube
intent(out)		volume used in another application
double precision,	t_remain_in	remaining time of the current integration
intent(in)		step, which consists of many pushings
double precision,	t_pass	flight time of the current pushing step
intent(out)		
logical,	boole_t_finished	boolean stating if the remaining step time
intent(out)		has been reached in the current pushing
integer,	iper_phi	+1,-1 if the particle travels through the
intent(out)		$\varphi = 0$ -plane in $-\varphi, +\varphi$ direction,
		0 otherwise

#### Initialize pusher

In the pusher\_tetra\_orbit subroutine, first an initializer subroutine initialize\_push er\_tetra\_orbit\_mod is called. Here, the initial conditions are used to compute the coefficients  $a_l^i, b^i$  for the ordinary differential equation set ??, representing the equations of motion [?]. One should emphasize here, that this ode set is solved within a shifted coordinate system, where the coordinate origin  $\mathbf{z}_0$  lies on the first vertex of a given tetrahedron. By convention in this project, when referring to a position in the global coordinate system one denotes the variable  $\mathbf{x}$ , when referring to a position inside the local shifted coordinate system one uses  $\mathbf{z}$  instead.

#### Initial guess of exit plane

Now, that the necessary components of the ODE set ?? have been initialized, the next occurring orbit intersection needs to be computed by the pusher routine. Since this must be done efficiently, a numerically inexpensive approximative quadratic solution is first evaluated by subroutine quad\_analytic\_approx to compute the guess for the orbit parameter tau at the first intersection of the particle trajectory with the cell boundary. Based on the result for the orbit parameter tau, an integration step is performed for the given step length using an RK4 solver, this integrator type is explained in more detail in appendix ??. The RK4 integrator subsequently returns the evaluated position for the specified value for tau. In general, due to inaccuracies in the approximation, this value does not correspond to a converged orbit position. On a sidenote, in the context of the pusher routine, converged simply means that the particle position is within a defined convergence distance to a given tetrahedral plane. This distance is given by  $10^{-10}$  times the normal distance of the first vertex within a given tetrahedron to its opposing cell boundary spanned by vertices 2,3 and 4 of the given tetrahedron. In addition, the normal velocity, which can also be computed from the output of the RK4 step, must have a negative sign in order for the convergence to be valid. The negative sign merely states that the particle is flying outwards of the tetrahedron, if there particle flew inwards, it would therefore not be accepted. Now, since the orbit position is generally not yet converged after the quadratic approximation, one next applies Newton's method for the face convergence by calling the subroutine newton\_face\_convergence. A detailed description of this approach is given by M. Eder et al [?].

## 1.1.3 Convergence and validation loop conv\_val\_loop

What one has obtained so far is a proposal for the exit plane and a converged orbit position on this plane. There may still be some problems, however, since for example Newton's method can fail if the orbit in fact turns before it intersects with the plane. Furthermore, it might happen that it does converge on the suggested plane but at the point of convergence it had already left the tetrahedron through another plane which is not allowed. Such cases need to be checked and handled appropriately. For this purpose, the convergence and validation loop conv\_val\_loop was implemented. This loop starts directly after the above mentioned quadratic approximation just before the convergence using Newton's method. Here, in each iteration of the loop, the algorithm tries to converge the particle orbit position on the currently proposed intersection face. Next, if convergence is reached the algorithm checks for the remaining planes

if the particle lies outside the tetrahedron. If this is not the case furthermore the normal velocity is checked to see if the particle flies outside. In case this is also correct, the particle is considered converged and accepted. In any case, where an error is detected, an appropriate approach is suggested. In most cases this involves using the quadratic approximation to suggest a different face, however, in some special cases this is not sufficient. Therefore, one calls inside the loop conv\_val\_loop an additional convergence routine last\_line\_defense, which no longer opts for high computational efficiency but rather for a reliable way of finding the intersection point. This subroutine is very comprehensive, but a central piece of it is a bisection scheme. For a short overview, in this scheme one computes the relative particle positions to all four faces and furthermore checks the normal velocities. Here, if the particle is inside the tetrahedron, the current step length is doubled and an integration step is performed. If the particle is now outside the tetrahedron, the last integration step is halved and integrated back in negative  $\tau$ -direction. This is done in an iterative scheme until a converged particle position has been found and albeit computationally expensive, the lastlineof defense solver remains an indispensable element of the algorithm due to its high reliability.

The structure of the loop conv\_val\_loop is presented in pseudo-code below.

```
conv_val_loop: do i = 1,5
  boole_converged = .true.
  call newton_face_convergence(z,tau,iface_new,..)
  if (Newton's method failed) then
     allowed_faces(iface_new)=.false.
     if (all allowed_faces forbidden) then
        call last_line_defense(z,tau,iface_new,..)
        cycle conv_val_loop
  endif
  call quad_analytic_approx(z,allowed_faces,dtau,..)
  if (quadratic approximation failed) then
     call last_line_defense(z,tau,iface_new,..)
     boole_converged = .false.
```

```
cycle conv_val_loop
    endif
    call integration_step(z,dtau,..)
    tau = tau + dtau
    boole_converged = .false.
    cycle conv_val_loop
endif
three_planes_loop: do j=1,3
    k = modulo(iface_new+j-1,4)+1
    if (particle is outside face k) then
         allowed_faces(iface_new)=.false.
         if (all allowed_faces forbidden) then
              call last_line_defense(z,tau,iface_new,..)
              boole_converged = .false.
              cycle conv_val_loop
         endif
         if (face k is not forbidden in allowed_faces) then
              iface_new = k
              boole_converged = .false.
              cycle conv_val_loop
         endif
    endif
enddo three_planes_loop
    if (normal velocity at iface_new points inwards) then
         allowed_faces(iface_new)=.false.
         if (all allowed_faces forbidden) then
              call last_line_defense(z,tau,iface_new,..)
              boole_converged = .false.
              cycle conv_val_loop
         endif
         call quad_analytic_approx(z,allowed_faces,dtau,..)
```

```
if (quadratic approximation failed) then
                   call last_line_defense(z,tau,iface_new,..)
                   boole_converged = .false.
                   cycle conv_val_loop
              endif
              call integration_step(z,dtau,..)
              tau = tau + dtau
              boole_converged = .false.
              cycle conv_val_loop
         endif
         if (tau is negative) then
              allowed_faces(iface_new)=.false.
              z = z_{init}
              tau = 0.d0
              if (all allowed_faces forbidden)
                   call last_line_defense(z,tau,iface_new,..)
                   boole_converged = .false.
                   cycle conv_val_loop
              endif
              call quad_analytic_approx(z,allowed_faces,dtau,..)
              if (quadratic approximation failed) then
                   call last_line_defense(z,tau,iface_new,..)
                   boole_converged = .false.
                   cycle conv_val_loop
              endif
              call integration_step(z,dtau,..)
              tau = tau + dtau
              boole_converged = .false.
              cycle conv_val_loop
         endif
         exit conv_val_loop
enddo conv_val_loop
```

From the pseudo-code one can see, that a lot of thought has gone into efficiently computing the next intersection. However, due to the used approach of guessing the exit face with an approximation, many special cases of particle trajectories had to be taken into account such that the logic deals with them correctly.

## 1.1.4 Final processing

The last steps of the orbit integration are to first check if the computed time is in fact smaller than the remaining flight time. If this is not the case, instead of the converged orbit position, the corresponding position at the remaining time is evaluated and assumed. In this case the current tetrahedron index will be returned with iface\_new being set to 0. If the computed time is smaller than the remaining time, the remaining time is reduced by the current value and the tetrahedron index adjacent to iface\_new will be returned by calling the subroutine pusher\_handover2neighbour. Furthermore, iface\_new is changed to neighbour\_face(iface\_new) of the adjacent tetrahedron to mark the new entry face. Now, one must check if the intersection face is at a periodic boundary of the coordinate system. In this case the corresponding value of i\_per\_theta/phi times  $2\pi$  is added to the respective coordinate component. Finally, the values for the current position z in local coordinates are converted to x in global coordinates, then the output values are returned and the orbit pushing is completed.

# 1.2 Pusher routine pusher\_tetra\_orbit\_analytic

Implementation of analytic pusher from analytical solution

## 1.3 Search routine for tetrahedra with starting points

Since generally, a user defines for a particle a starting position x in global coordinates, rather than a tetrahedron index and a local position, a search routine is needed for finding the corresponding tetrahedron index. For this, a function is\_inside is implemented which allows to check if the particle position lies inside a proposed tetrahedron. This function uses the Hesse normal form to compute the distances to all tetrahedral planes. The necessary quantities are the coordinate position of the first vertex and the four normal vectors of the planes, these are accessable in module tetra\_grid\_mod. Next, due to the axisymmetry of the grid, based on the current phi position one can vastly reduce the number of possible tetrahedra by allowing only tetrahedra of the current  $\varphi$ -slice. Now, a loop over all possible tetrahedron indices is implemented to check if the particle lies inside. Once the correct tetrahedron has been found, the distances to the four planes need to be checked for random convergence on a plane. If this is the case, also the normal velocity with respect to that plane must be evaluated, since the pusher always assumes that a given particle flies inwards which could lead to errors in the logics. If the particle, however, is converged on a plane and flies outwards, instead of the current tetrahedron index the adjacent neighbor at this face index is returned with the corresponding value for iface\_new.