SQL

Part-II

SQL queries

Preview

sqltutorial.org/sql-cheat-sheet

SQL CHEAT SHEET http://www.sqltutorial.org

QUERYING DATA FROM A TABLE

SELECT c1, c2 FROM t;

Query data in columns c1, c2 from a table

SELECT * FROM t:

Ouerv all rows and columns from a table

SELECT c1, c2 FROM t

WHERE condition:

Query data and filter rows with a condition

SELECT DISTINCT c1 FROM t

WHERE condition;

Query distinct rows from a table

SELECT c1, c2 FROM t

ORDER BY c1 ASC [DESC];

Sort the result set in ascending or descending order

SELECT c1, c2 FROM t

ORDER BY cl

LIMIT n OFFSET offset:

Skip offset of rows and return the next n rows

SELECT c1, aggregate(c2)

FROM t

GROUP BY c1:

Group rows using an aggregate function

SELECT c1, aggregate(c2)

FROM t

GROUP BY cl

HAVING condition:

Filter groups using HAVING clause

OUERYING FROM MULTIPLE TABLES

SELECT c1, c2

FROM t1

INNER JOIN t2 ON condition;

Inner join t1 and t2

SELECT c1, c2 FROM t1

LEFT JOIN t2 ON condition:

Left join t1 and t1

SELECT c1, c2

FROM t1

RIGHT JOIN t2 ON condition;

Right join t1 and t2

SELECT c1, c2

FROM t1

FULL OUTER JOIN t2 ON condition:

Perform full outer join

SELECT c1, c2

FROM t1

CROSS JOIN t2:

Produce a Cartesian product of rows in tables

SELECT c1, c2

FROM t1, t2:

Another way to perform cross join

SELECT c1, c2

FROM t1 A

INNER JOIN t2 B ON condition:

Join t1 to itself using INNER JOIN clause

USING SOL OPERATORS

SELECT c1, c2 FROM t1

UNION [ALL]

SELECT c1, c2 FROM t2;

Combine rows from two queries

SELECT cl. c2 FROM tl

INTERSECT

SELECT c1, c2 FROM t2;

Return the intersection of two queries

SELECT c1, c2 FROM t1

MINUS

SELECT c1, c2 FROM t2;

Subtract a result set from another result set

SELECT c1, c2 FROM t1

WHERE cl [NOT] LIKE pattern;

Query rows using pattern matching %, _

SELECT c1, c2 FROM t

WHERE cl [NOT] IN value list;

Query rows in a list

SELECT c1, c2 FROM t

WHERE c1 BETWEEN low AND high;

Query rows between two values

SELECT c1, c2 FROM t

WHERE cl IS [NOT] NULL;

Check if values in a table is NULL or not

This Lecture

• Set operators & nested queries

Aggregation & GROUP BY

Set Operators & Nested Queries

Set algebra (reminder)

List: [1, 1, 2, 3] Set: {1, 2, 3}

Multiset: {1, 1, 2, 3}

A <u>multiset</u> is an unordered list (or: a set with multiple duplicate instances allowed)

UNIONs

```
Set: \{1, 2, 3\} \cup \{2\} = \{1, 2, 3\}
```

<u>Multiset</u>: {1, 1, 2, 3} U { 2 } = { 1, 1, 2, 2, 3 }

Cross-product

```
{1, 1, 2, 3} * { y, z } =

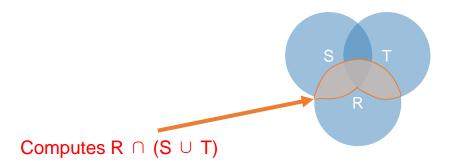
{ <1, y>, <1, y>, <2, y>, <3, y>

<1, z>, <1, z>, <2, z>, <3, z>

}
```

An Unintuitive Query

SELECT DISTINCT R.A FROM R, S, T WHERE R.A=S.A OR R.A=T.A

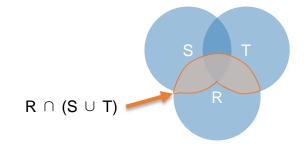


But what if $S = \Phi$?

Go back to the semantics!

What does this look like in Python?

SELECT DISTINCT R.A FROM R, S, T WHERE R.A=S.A OR R.A=T.A



- Semantics:
 - 1. Take cross-product

1. Apply <u>selections</u> / <u>conditions</u>

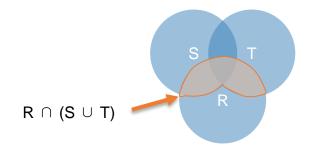
Joins / cross-products are just **nested** for loops (in simplest implementation)!

If-then statements!

1. Apply projection

What does this look like in Python?

```
SELECT DISTINCT R.A
FROM R, S, T
WHERE R.A=S.A OR R.A=T.A
```



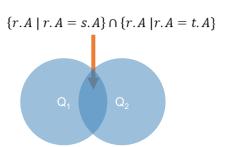
```
output = {}

for r in R:
    for s in S:
    for t in T:
        if r['A'] == s['A'] or r['A'] == t['A']:
            output.add(r['A'])
    return list(output)
```

Can you see now what happens if S = []?

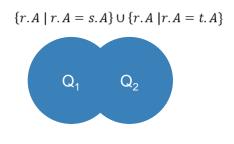
Explicit Set Operators: INTERSECT

SELECT R.A
FROM R, S
WHERE R.A=S.A
INTERSECT
SELECT R.A
FROM R, T
WHERE R.A=T.A



UNION

SELECT R.A
FROM R, S
WHERE R.A=S.A
UNION
SELECT R.A
FROM R, T
WHERE R.A=T.A



Why aren't there duplicates?

By default: SQL retains **set** semantics for UNIONs, INTERSECTs!

What if we want duplicates?

UNION ALL

SELECT R.A
FROM R, S
WHERE R.A=S.A
UNION ALL
SELECT R.A
FROM R, T
WHERE R.A=T.A

$$\{r.A \mid r.A = s.A\} \cup \{r.A \mid r.A = t.A\}$$

$$\mathsf{Q}_1 \qquad \mathsf{Q}_2$$

ALL indicates Multiset operations

Nested queries: Sub-queries Return Relations

Another example:

Company(<u>name</u>, city)
Product(<u>name</u>, maker)
Purchase(<u>id</u>, product, buyer)

66

Companies making products bought by Mickey"Location of

companies?

Nested Queries

Are these queries equivalent?

```
SELECT c.city
FROM Company c
WHERE c.name IN (
SELECT pr.maker
FROM Purchase p, Product pr
WHERE p.name = pr.product
AND p.buyer = 'Mickey')
```

```
SELECT c.city
FROM Company c,
Product pr,
Purchase p
WHERE c.name = pr.maker
AND pr.name = p.product
AND p.buyer = 'Mickey'
```

Beware of duplicates!

Nested Queries

```
FROM Company c,
Product pr,
Purchase p
WHERE c.name = pr.maker
AND pr.name = p.product
AND p.buyer = 'Mickey'
```

```
SELECT DISTINCT c.city
FROM Company c
WHERE c.name IN (
SELECT pr.maker
FROM Purchase p, Product pr
WHERE p.product = pr.name
AND p.buyer = 'Mickey')
```

Now they are equivalent (both use set semantics)

Subqueries Return Relations

You can also use operations of the form:

ANY and ALL not supported by SQLite.

- s > ALL R
- s < ANY R
- EXISTS R

Ex: Product(name, price, category, maker)

Find products that are more expensive than all those produced by "Gizmo-Works"

Subqueries Returning Relations

You can also use operations of the form:

- s > ALL R
- s < ANY R
- EXISTS R

Ex: Product(name, price, category, maker)

```
SELECT p1.name
FROM Product p1
WHERE p1.maker = 'Gizmo-Works'
AND EXISTS(
SELECT p2.name
FROM Product p2
WHERE p2.maker <> 'Gizmo-Works'
AND p1.name = p2.name)
```

Find 'copycat' products, i.e. products made by competitors with the same names as products made by "Gizmo-Works"

<> means !=

Nested queries as alternatives to INTERSECT and EXCEPT

INTERSECT and EXCEPT not in some DBMSs!

(SELECT R.A, R.B FROM R) INTERSECT (SELECT S.A, S.B FROM S)



SELECT R.A, R.B FROM R WHERE EXISTS(

SELECT * FROM S

WHERE R.A=S.A AND R.B=S.B)

(SELECT R.A, R.B FROM R) EXCEPT (SELECT S.A, S.B FROM S)



SELECT R.A, R.B FROM R WHERE NOT EXISTS(SELECT *

> FROM S WHERE R.A=S.A AND R.B=S.B)

Correlated Queries Using External Vars in Internal Subquery

Movie(<u>title</u>, <u>year</u>, director, length)

Find movies whose title appears more than once.

```
SELECT DISTINCT title
FROM Movie AS m
WHERE year <> ANY(

SELECT year

FROM Movie

WHERE title = m.title)
```

Note the scoping of the variables!

Note also: this can still be expressed as single SFW query...

What you will learn about in this section

1. GROUP BY

1. GROUP BY: with HAVING, semantics

Aggregation

SELECT AVG(price)
FROM Product
WHERE maker = "Toyota"

SELECT COUNT(*)
FROM Product
WHERE year > 1995

- SQL supports several aggregation operations:
 - SUM, COUNT, MIN, MAX, AVG

Aggregation: COUNT

COUNT applies to duplicates, unless otherwise stated

SELECT COUNT(category)
FROM Product
WHERE year > 1995

Note: Same as COUNT(*). Why?

We probably want:

SELECT COUNT(DISTINCT category)
FROM Product
WHERE year > 1995

More Examples

Purchase(product, date, price, quantity)

SELECT SUM(price * quantity)
FROM Purchase

What do these mean?

SELECT SUM(price * quantity)
FROM Purchase
WHERE product = 'bagel'

Simple Aggregations

Purchase

Product	Date	Price	Quantity
bagel	10/21	1	20
banana	10/3	0.5	10
banana	10/10	1	10
bagel	10/25	1.50	20

SELECT SUM(price * quantity)
FROM Purchase
WHERE product = 'bagel'



50 (= 1*20 + 1.50*20)

Grouping and Aggregation

Purchase(product, date, price, quantity)

```
SELECT product,
```

SUM(price * quantity) AS TotalSales

FROM Purchase

WHERE date > '10/1/2005'

GROUP BY product

Find total sales after 10/1/2005 per product.

Let's see what this means...

Grouping and Aggregation

```
SELECT product,
SUM(price * quantity) AS TotalSales
FROM Purchase
WHERE date > '10/1/2005'
GROUP BY product
```

Semantics of the query:

1. Compute the FROM and WHERE clauses

2. Group by the attributes in the GROUP BY

3. Compute the SELECT clause: grouped attributes and aggregates

1. Compute the FROM and WHERE clauses

SELECT product, SUM(price*quantity) AS TotalSales

FROM Purchase

WHERE date > '10/1/2005'

GROUP BY product



Product	Date	Price	Quantity
Bagel	10/21	1	20
Bagel	10/25	1.50	20
Banana	10/3	0.5	10
Banana	10/10	1	10

2. Group by the attributes in the GROUP BY

SELECT product, SUM(price*quantity) AS TotalSales

FROM Purchase

WHERE date > '10/1/2005'

GROUP BY product

Product	Date	Price	Quantity
Bagel	10/21	1	20
Bagel	10/25	1.50	20
Banana	10/3	0.5	10
Banana	10/10	1	10





Product	Date	Price	Quantity
Bagel	10/21	1	20
	10/25	1.50	20
Banana	10/3	0.5	10
	10/10	1	10

3. Compute the SELECT clause: grouped attributes and aggregates

SELECT product, SUM(price*quantity) AS TotalSales

FROM Purchase

WHERE date > '10/1/2005'

GROUP BY product

Product	Date	Price	Quantity
Bagel	10/21	1	20
	10/25	1.50	20
Banana	10/3	0.5	10
	10/10	1	10



Product	TotalSales	
Bagel	50	
Banana	15	

GROUP BY v.s. Nested Queries

```
SELECT product, Sum(price*quantity) AS TotalSales
FROM Purchase
WHERE date > '10/1/2005'
GROUP BY product
```

```
SELECT DISTINCT x.product,

(SELECT Sum(y.price*y.quantity)

FROM Purchase y

WHERE x.product = y.product

AND y.date > '10/1/2005') AS TotalSales

FROM Purchase x

WHERE x.date > '10/1/2005'
```

HAVING Clause

SELECT product, SUM(price*quantity)
FROM Purchase
WHERE date > '10/1/2005'
GROUP BY product
HAVING SUM(quantity) > 100

Same query as before, except that we consider only products that have more than 100 buyers

HAVING clauses contains conditions on aggregates

Whereas WHERE clauses condition on individual tuples...

General form of Grouping and Aggregation

Evaluation steps:

- Evaluate FROM-WHERE: apply condition C₁ on the attributes in R₁,...,R_n
- 2. GROUP BY the attributes $a_1,...,a_k$
- 3. Apply condition C_2 to each group (may need to compute aggregates)
- 4. Compute aggregates in S and return the result

THANK YOU!