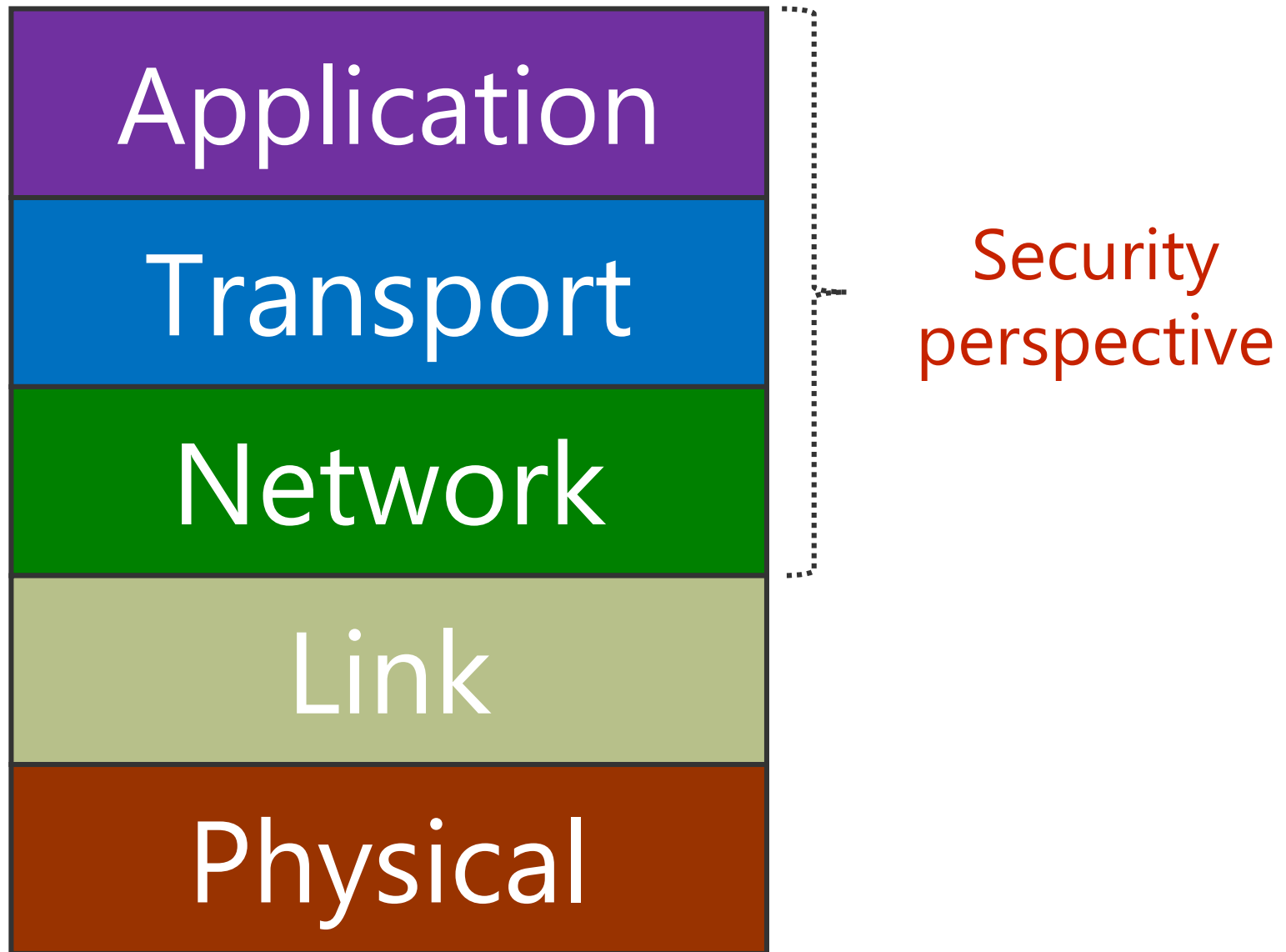


CS2105 Introduction to
Computer Network

Lecture 8

Network Security

15 October 2018



Lecture 8: Network Security

After this class, you are expected to understand:

- ❖ how *symmetric key* cryptography and *public key* cryptography can be used to ensure **message confidentiality**.
- ❖ how *message authentication code* and *digital signature* ensure **message integrity** and **authenticity**.

Lecture 8: Roadmap

8.1 What is Network Security?

8.2 Principles of Cryptography

8.3 Message Integrity and Digital Signatures

8.6 Securing TCP Connections: SSL

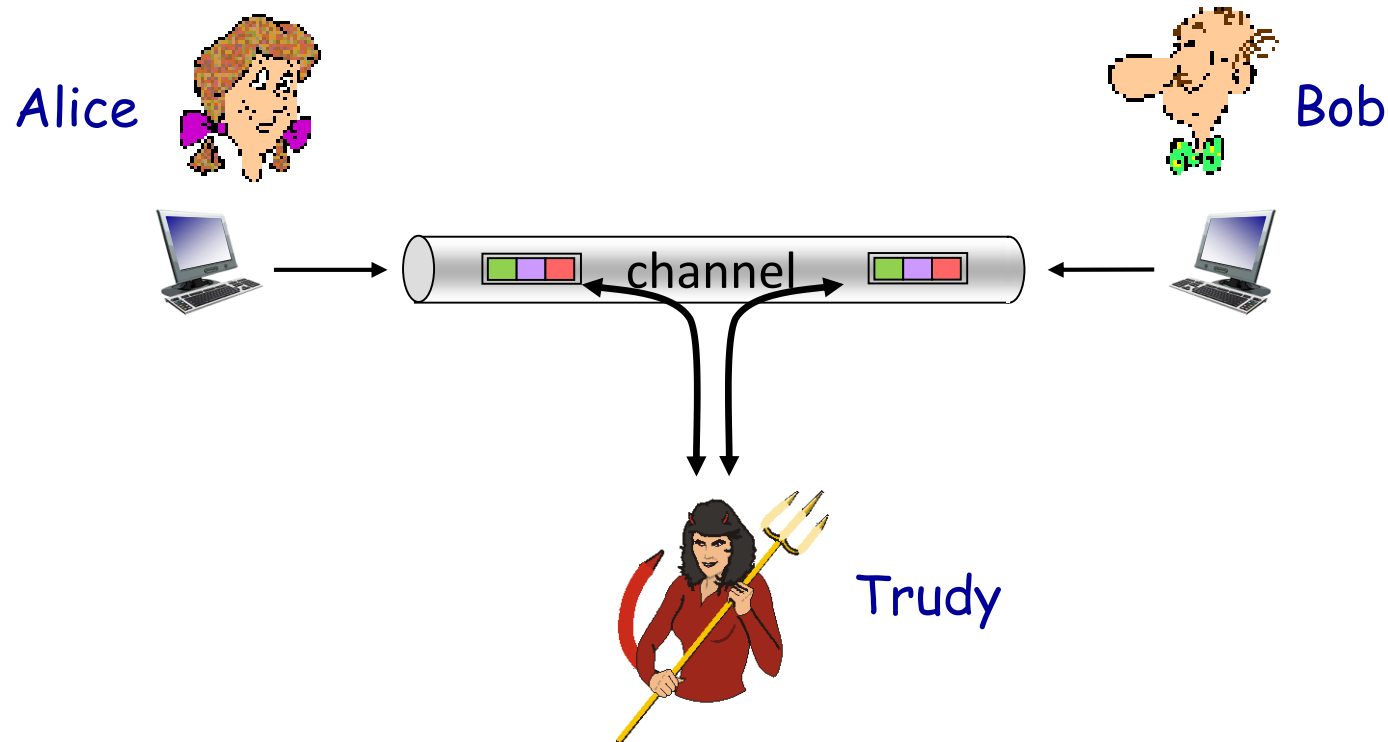
8.7 Network Layer Security: IPsec

**Non-
examinable**

Kurose Textbook, Chapter 8
(Some slides are taken from the book)

Friends and Enemies: Alice, Bob, Trudy

- ❖ Alice and Bob (lovers!) want to communicate “secretly”.
- ❖ Trudy (intruder) wants to interfere.



What Can Bad Guy Trudy Do?



Trudy may:

- intercept messages of Alice and Bob (*eavesdrop*).
 - Need to ensure *message confidentiality*.
- modify messages between Alice and Bob or forge messages and insert into communication
 - Need to ensure *message integrity* and *message authenticity*.
- attack the communication channel between Alice and Bob (e.g. denial-of-service attack).
 - Need to ensure *service availability* (not covered).

Network Security: Algorithms

- ❖ We will not discuss any security algorithms in great details.
- ❖ Interested students may read chapter 8 of the textbook or take security courses offered by SoC, e.g.
 - **CS2107** Introduction to Information Security
 - **CS3235** Computer Security
 - **CS4236** Cryptography Theory and Practice
 - **CS5321** Network Security

Lecture 8: Roadmap

8.1 What is Network Security?

8.2 Principles of Cryptography

- 8.2.1 Symmetric Key Cryptography
- 8.2.2 Public Key Encryption

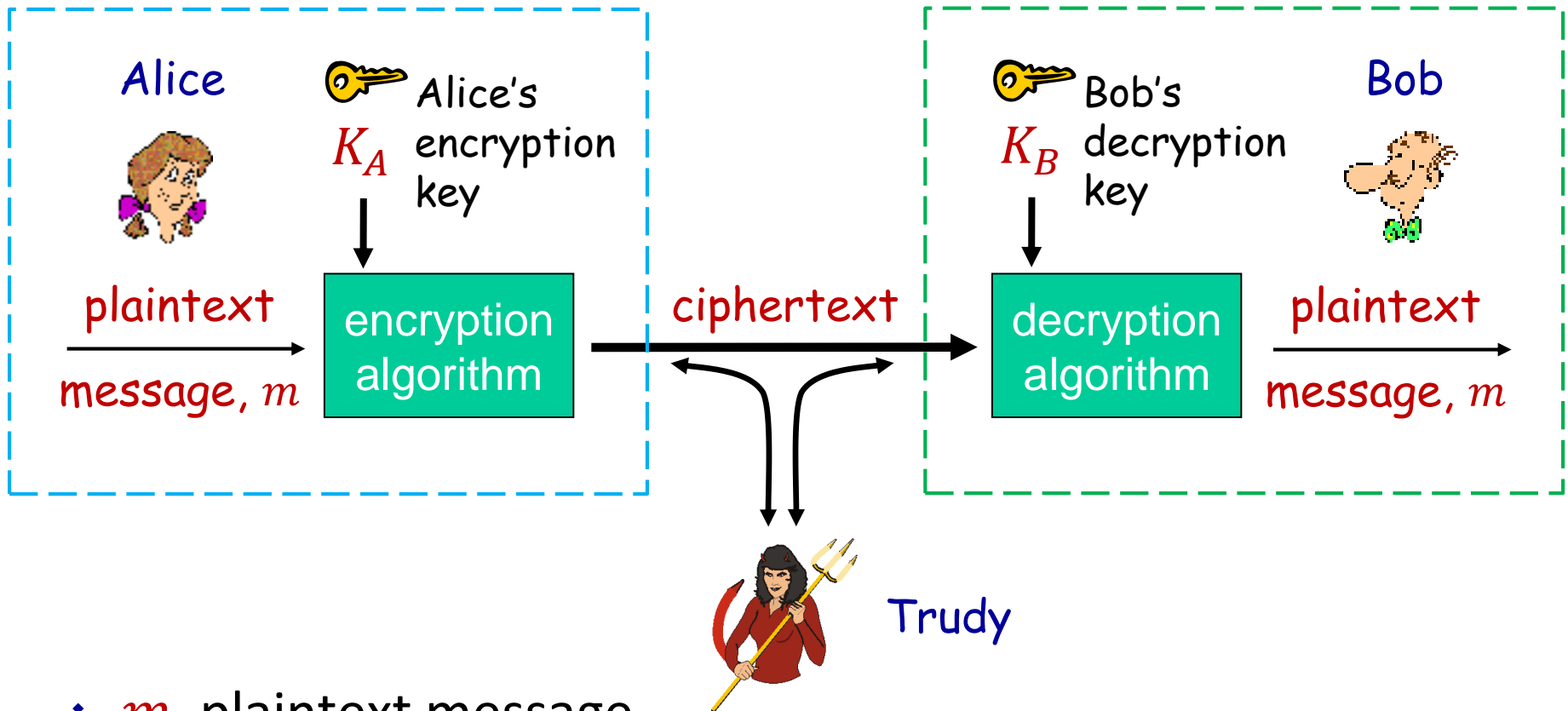
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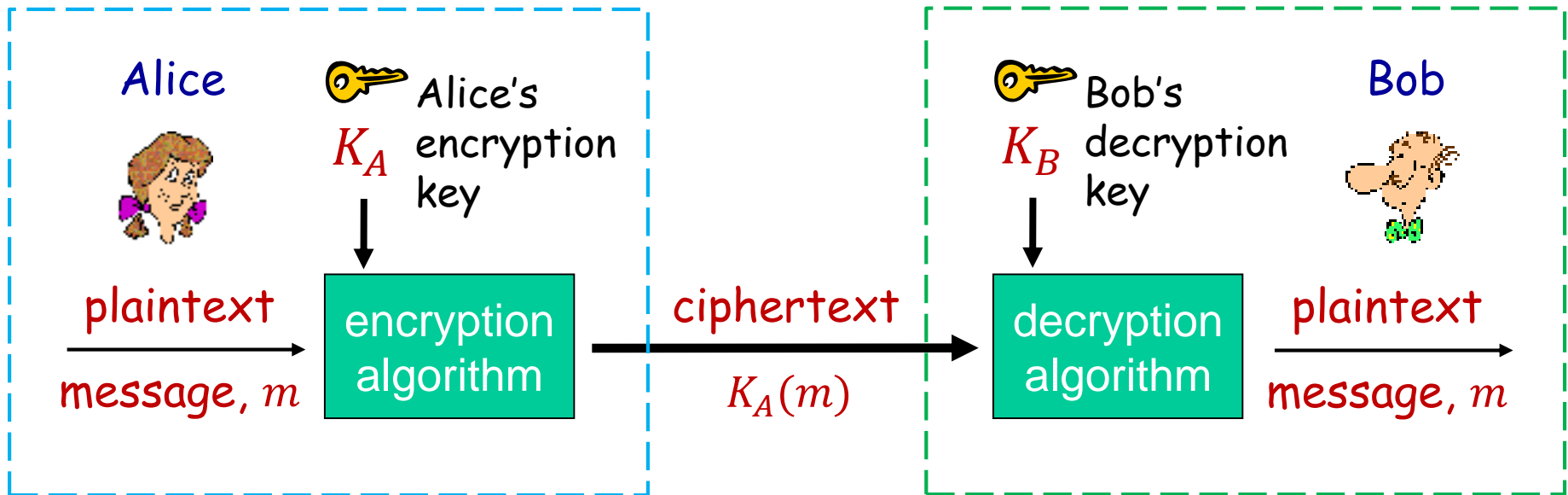
**Non-
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The Language of Cryptography



- ❖ m plaintext message
- ❖ $K_A(m)$ ciphertext, encrypted with key K_A
- ❖ $K_B(K_A(m)) = m$

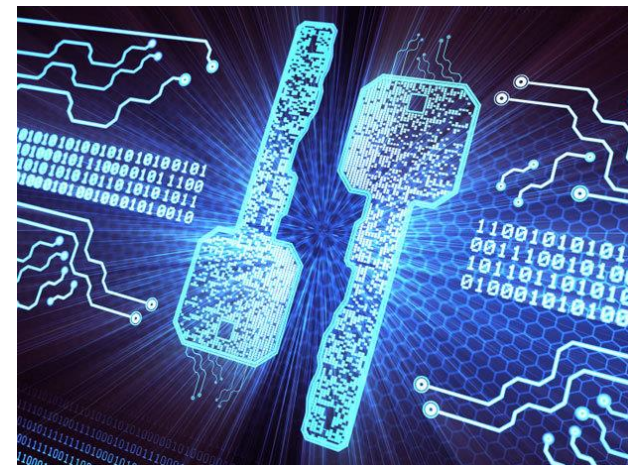
The Language of Cryptography



- ❖ Given ciphertext $K_A(m)$, it should be computationally hard to find plaintext m without knowing decryption key K_B .
- ❖ We will skip the mathematical details on how to derive K_A and K_B .

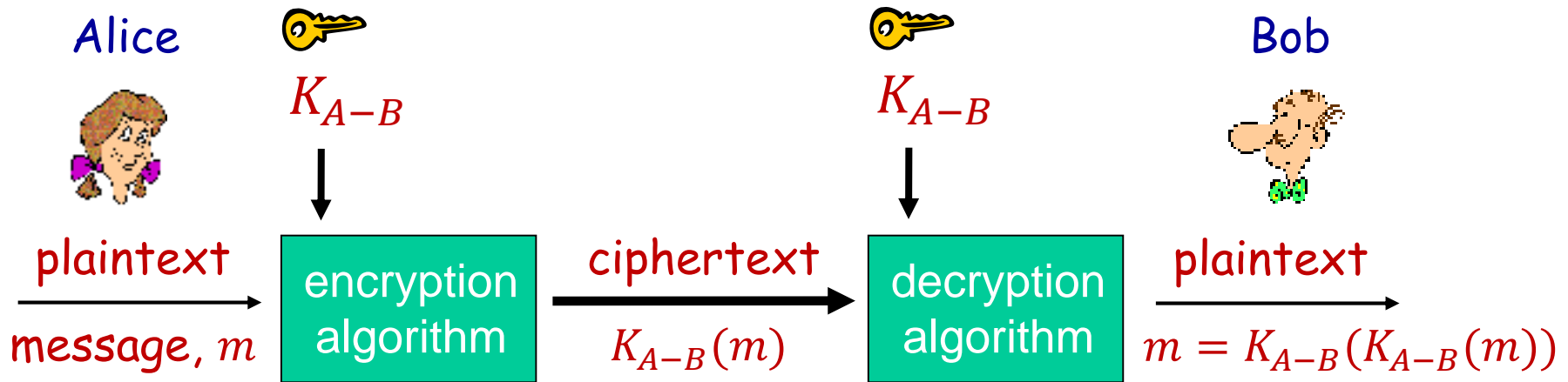
Types of Cryptography

- ❖ The purpose of cryptography is to make it difficult for an unauthorized third party to understand private communication between two parties.
- ❖ Cryptography often uses **keys**:
 - Algorithms are known to everyone
 - Only “keys” are secret
- ❖ **Symmetric key** cryptography
 - Involves the use of one key
- ❖ **Public key** cryptography
 - Involves the use of a pair of keys



Source: IEEE Spectrum

Symmetric Key Cryptography



- ❖ **Symmetric key crypto**: Bob and Alice share and use the same (symmetric) key: K_{A-B}
 - Popular algorithms: **DES** (Data Encryption Standard), **AES** (Advanced Encryption Standard)

Example Encryption Scheme

- ❖ **Mono-alphabetic cipher**: substituting one letter for another.

plaintext	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
ciphertext	m	n	b	v	c	x	z	a	s	d	f	g	h	j	k	l	p	o	i	u	y	t	r	e	w	q



Alice

Plaintext: bob, i love you. alice
 ciphertext: nkn, s gktc wky. mgsbc



Bob

- 🔑 **Encryption key**: mapping from a set of 26 letters to another set of 26 letters

Public Key Cryptography

❖ Symmetric key crypto issues:

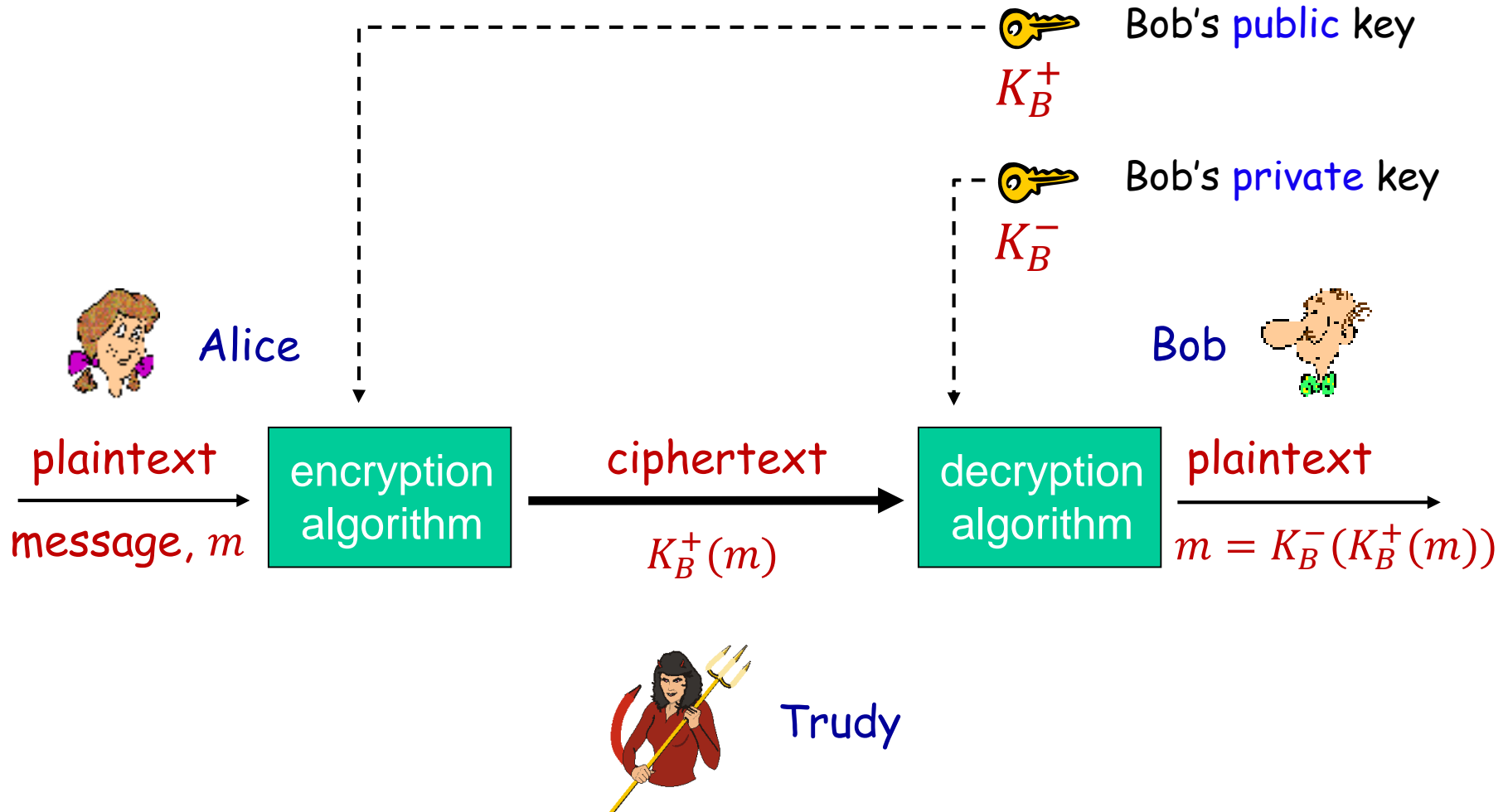
- Require sender and receiver to **share a secret key**.
- Use the same secret key to encrypt and decrypt data.
- **Question:** how to agree on a key in the first place?

❖ Public key crypto:

- Sender and receiver do not share secret key.
- Use **a pair of keys**. One for encryption and the other for decryption.
- **Public encryption key:** known to the world.
- **Private decryption key:** known only to receiver.



Public Key Cryptography



Public Key Encryption Algorithms

❖ Key points of public key encryption:

- ① Need to find a pair of public/private keys such that

$$K_B^-(K_B^+(m)) = m$$

- ② Given public key K_B^+ , it should be very difficult to find private key K_B^- .

❖ Most popular algorithm: **RSA** (Rivest, Shamir, Adelson algorithm)

Public Key: RSA Algorithm

- ❖ In RSA
 - The **public key** is the product of two very large primes.
 - The **private key** is derived from these two large primes.
- ❖ The security of RSA relies on the difficulty of factoring a large composite number.
 - It would be too slow to “guess” the two large primes, given the current state of the art of number theory.
- ❖ We will skip the mathematical details.

An Important Property of RSA

- ❖ The following property of RSA will be *very* useful for our discussion later:

$$K_B^- (K_B^+ (m)) = m = K_B^+ (K_B^- (m))$$

use **public** key first,
followed by
private key

use **private** key
first, followed by
public key

Result is the same!

RSA in Practice: Session Key

- ❖ RSA (public key encryption) is computationally intensive (but doesn't require key sharing).
- ❖ DES (symmetric key encryption) is at least 100 times faster than RSA.
- ❖ Question: how to take advantage of both?
 - use public key crypto to establish secure connection, then second key – symmetric key – for encrypting data.

Session key K_S :

- ❖ Bob and Alice use RSA to exchange a symmetric key K_S .
- ❖ Once both have K_S , they use symmetric key cryptography.
- ❖ No need to remember K_S , it's valid for one session only.

Lecture 8: Roadmap

8.1 What is Network Security?

8.2 Principles of Cryptography

8.3 Message Integrity and Digital Signatures

- 8.3.1 Cryptographic Hash Functions
- 8.3.2 Message Authentication Code
- 8.3.3 Digital Signatures

8.6 Securing TCP Connections: SSL

8.7 Network Layer Security: IPsec

**Non-
examinable**

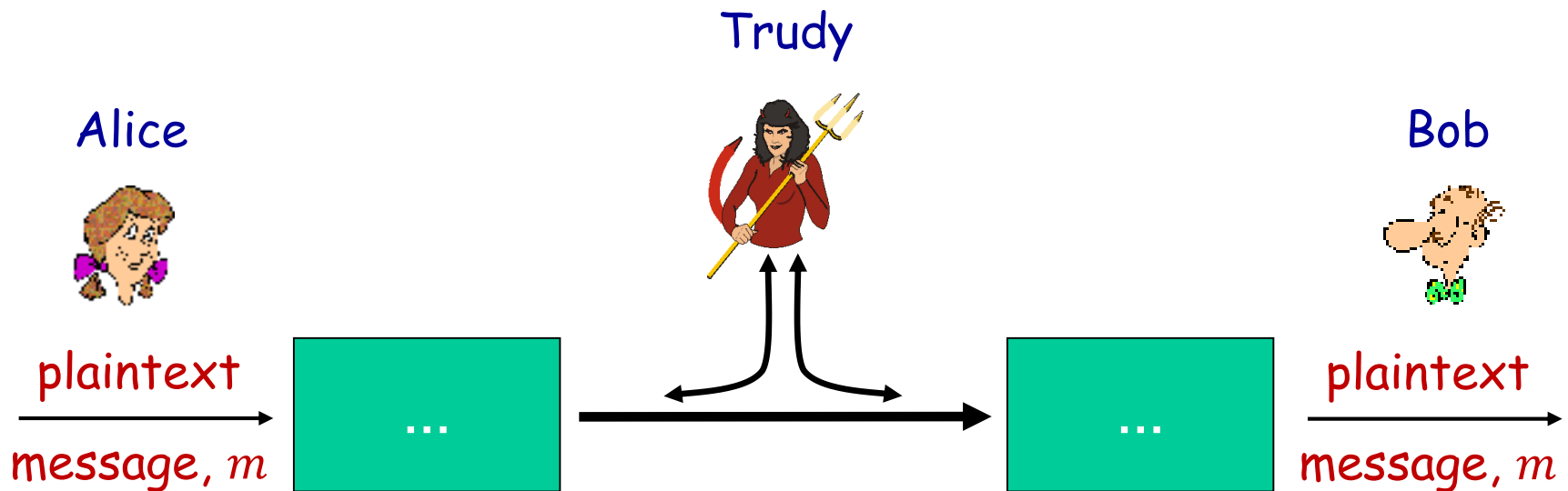
Message Integrity and Authentication

- ❖ We have seen how encryption can be used to provide confidentiality to two communicating entities.

- ❖ On the other hand, we often need to
 - ensure message has not been modified during transmission.
 - verify the creator of a message.

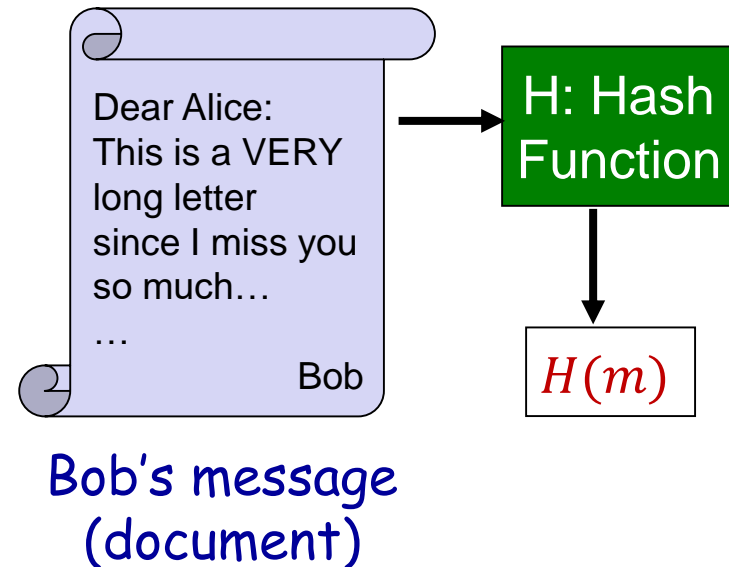
Message Integrity and Authentication

- ❖ We are going to study the following two topics:
 - ① message authentication code (MAC)
 - ② digital signature
- ❖ The basics of both is cryptographic hash function.



Cryptographic Hash Functions

- ❖ A hash function takes an input, m , and generates a fixed size string $H(m)$ known as **message digest** (hash or **finger print**).



- ❖ Popular algorithms: **MD5** (Message Digest) and **SHA-1** (Secure Hash Algorithm)
 - Example usage: both have been widely used to ensure a file downloaded from server has arrived intact.

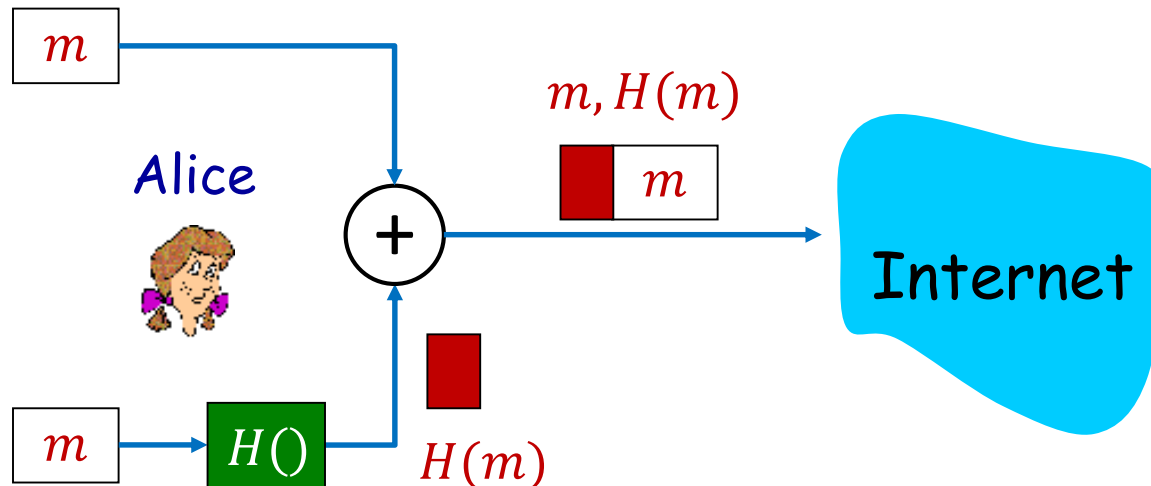
Cryptographic Hash Functions

- ❖ Cryptographic hash functions are one-way functions:
 - It is computationally infeasible to find two different messages m and m' such that $H(m) = H(m')$.
 - Therefore impossible for Trudy to forge another message m' with the same message digest as m .
- ❖ When using cryptographic hash functions,
 - A small change in the message (say, by eavesdropper) will create a significant change in the message digest.

Example Usage (1/3)

For Alice:

1. Alice creates message m and calculates the hash $H(m)$.
2. Alice then appends $H(m)$ to the message m , creating an extended message $(m, H(m))$, and sends the extended message to Bob.



Example Usage (2/3)

For Bob:

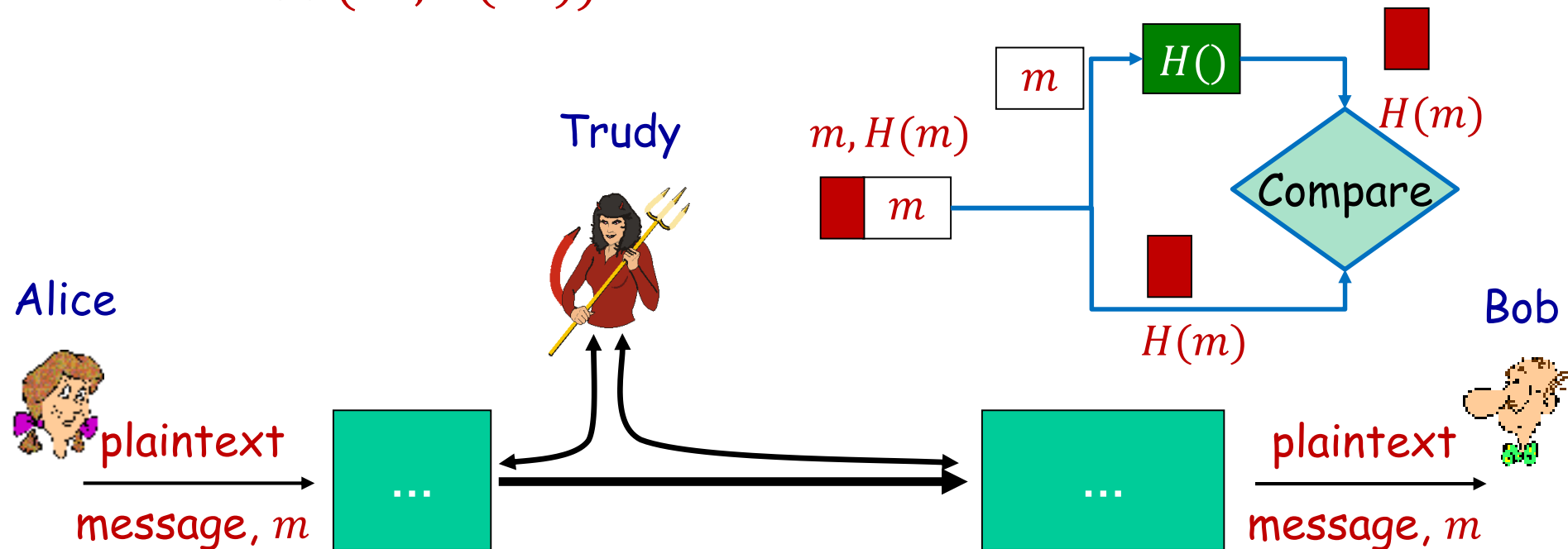
1. Bob receives an extended message (m', h') .
2. Bob calculates $H(m')$. If $H(m') = h'$, Bob concludes that everything is fine.

Recap:
Lecture 5


❖ TCP/UDP Checksum

Example Usage (3/3)

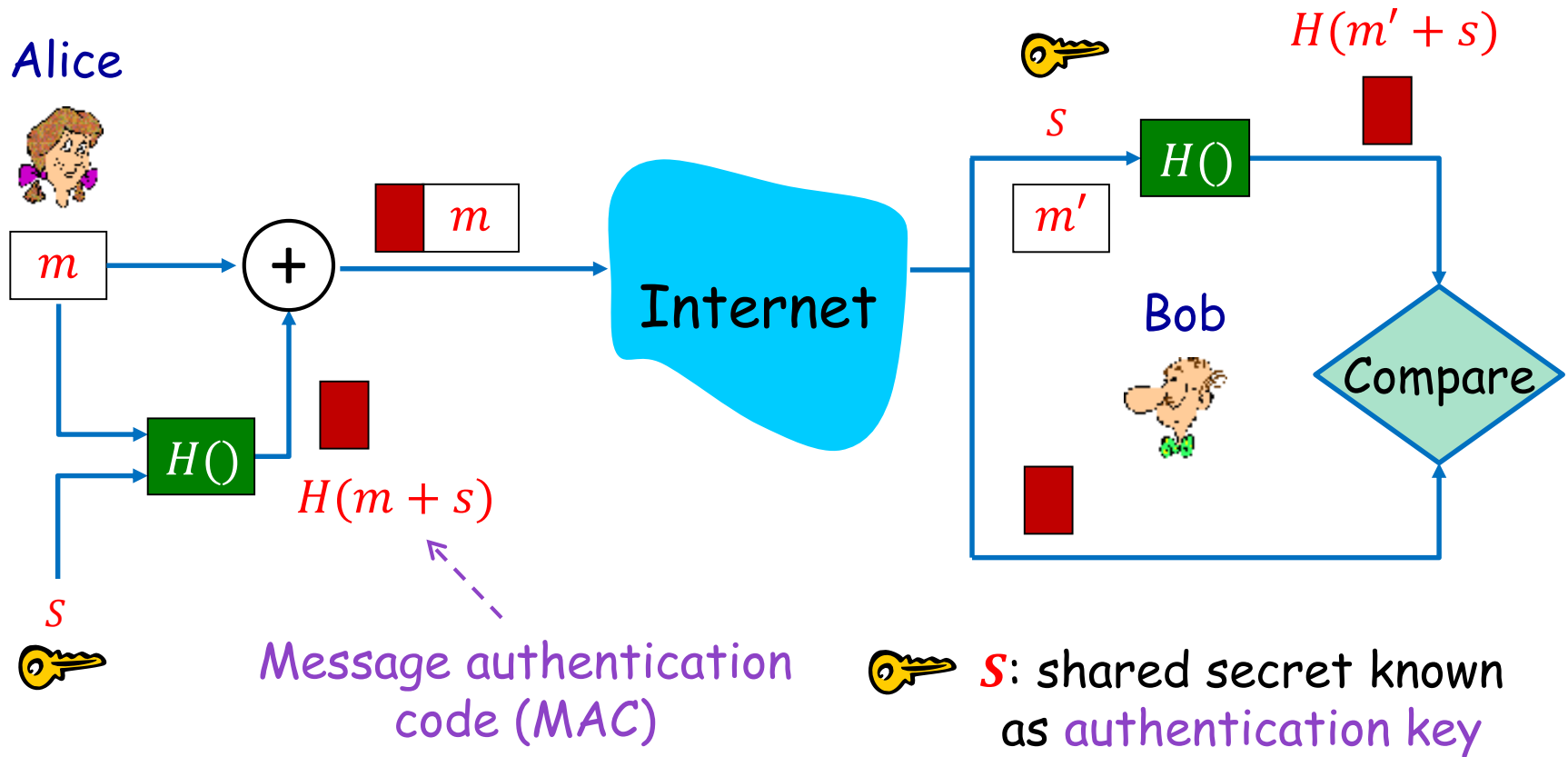
- ❖ **Q:** Can Bob be sure the source of message is Alice?
 - **No.** Because Trudy can create a bogus message m' in which she says she is Alice, calculates $H(m')$, and sends Bob $(m', H(m'))$.



Message Authentication Code

- ❖ If a **key**  is used as part of the message digest generation, such an algorithm is said to generate a **message authentication code** (MAC).
 - Can detect accidental and intentional changes to a message.
 - Can affirm to the receiver, the message's origin.
- ❖ Java supports the following standard MAC algorithms:
 - **HmacMD5, HmacSHA1, HmacSHA256**

Message Authentication Code




- ❖ MAC proves to Bob that the creator of the message is Alice and the message is not corrupted.

Digital Signature

Sender (Alice) signs a document digitally (analogous to hand-written signatures).

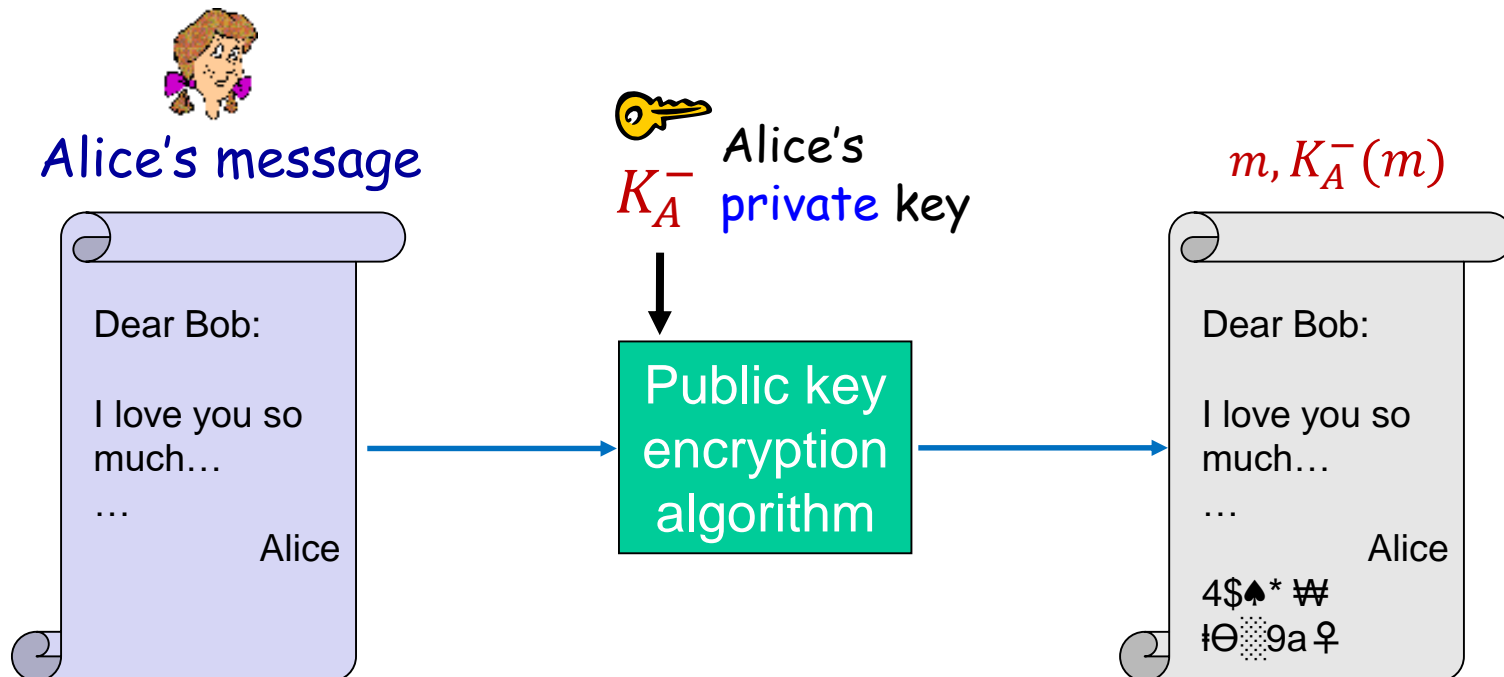
- ❖ *verifiable*: recipient (Bob) can verify that Alice, and no one else, has signed this document.
- ❖ *non-repudiation*: If Bob shows this document and digital signature to a third party (e.g. court), the third party is confident that this document is indeed signed by Alice (but no one else including Bob).

Digital Signature vs. MAC

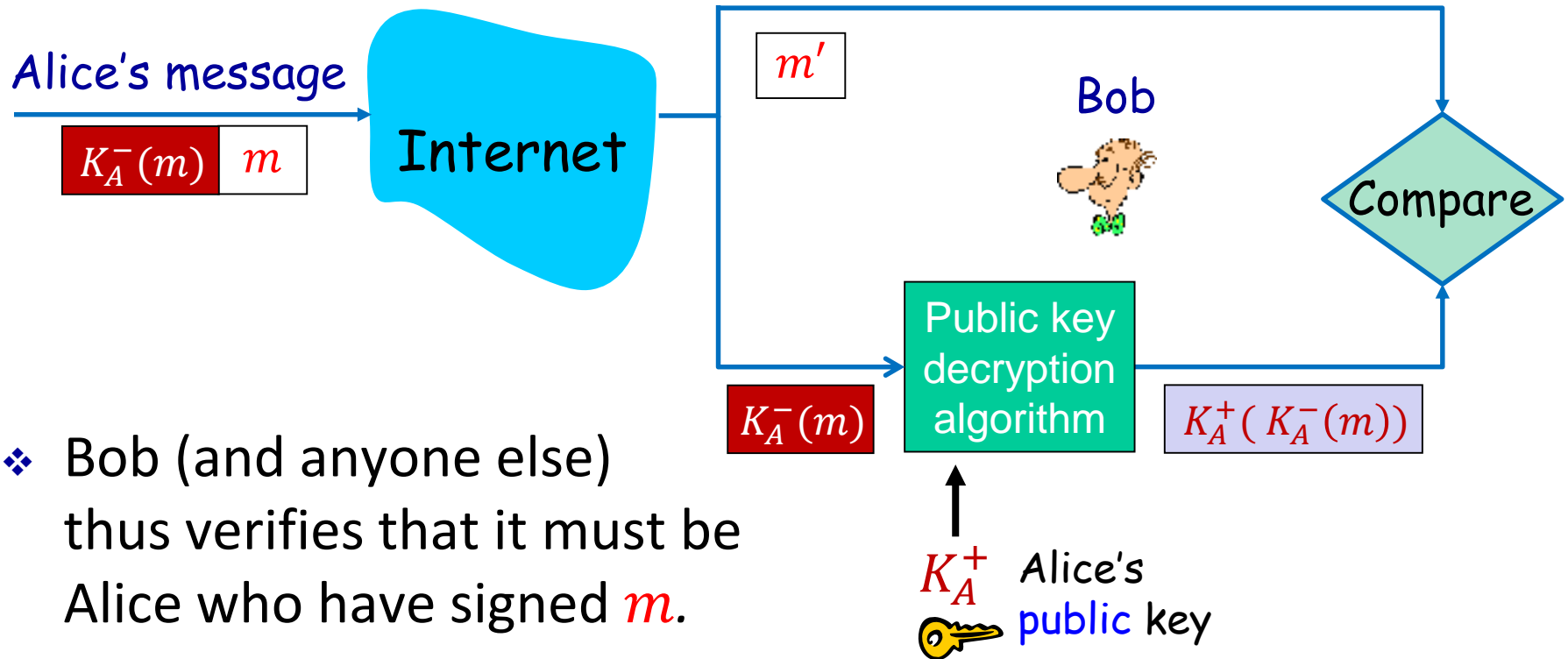
- ❖ Message authentication code (MAC) uses an authentication key **shared** between sender (Alice) and receiver (Bob).
 - Either Alice or Bob can produce the same MAC on a document, using the shared key.
 - Cannot prove to a third party MAC is produced by Alice or Bob.
- ❖ When Alice signs document digitally, she must put something on the doc that is **unique to her**.
 - her private key 

Digital Signature Example (1/2)

- ❖ Alice signs m by encrypting it with her private key K_A^- , creating a “signed” message, $K_A^-(m)$.
 - Send both m and $K_A^-(m)$ to Bob.

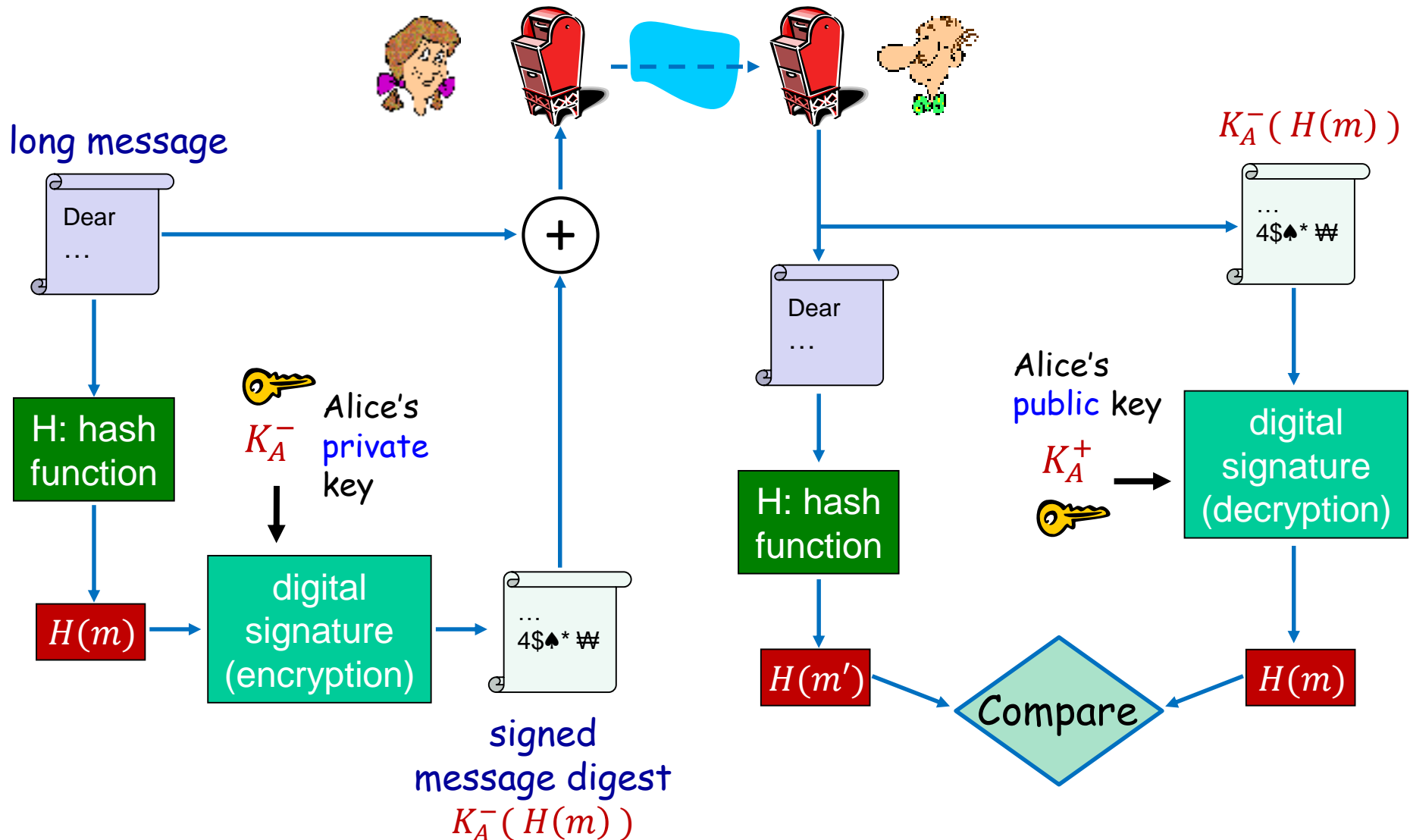


Digital Signature Example (2/2)



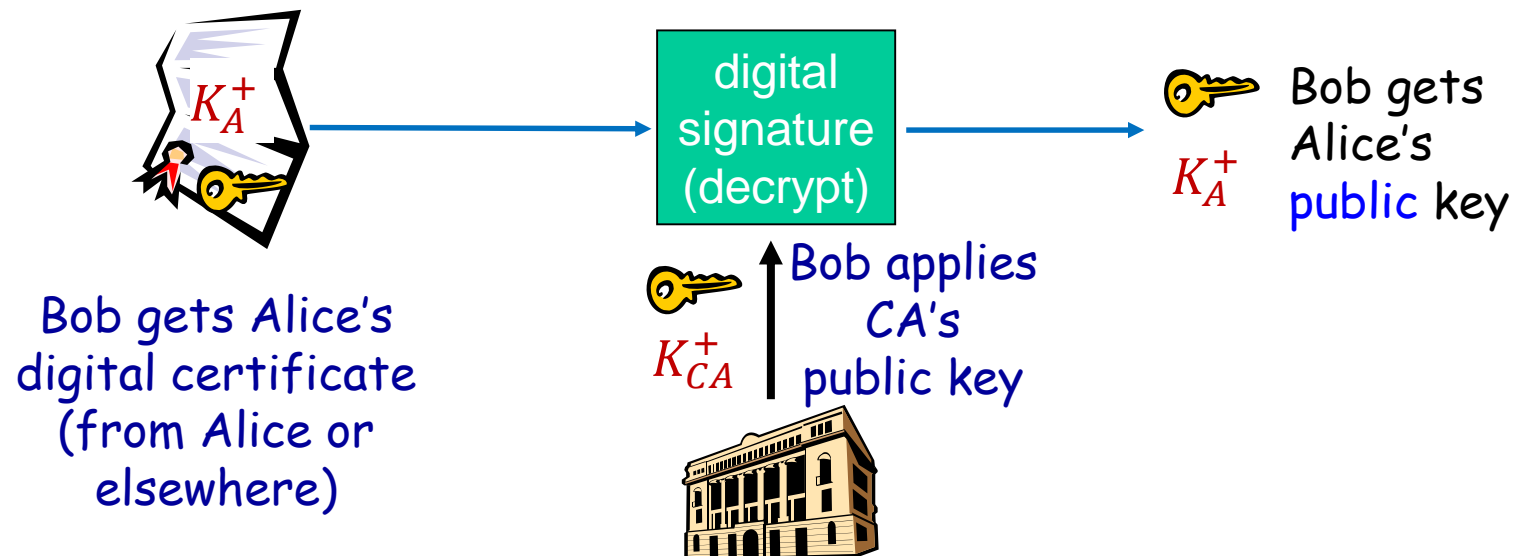
- ❖ Bob (and anyone else) thus verifies that it must be Alice who have signed m .
- ❖ Just one minor point:
 - Public key encryption is very slow.
 - Efficiency is a concern if m is long.

Digital Signature = Signed Message Digest



Digital Certificate

- ❖ Bob may wonder if the public key he uses is indeed Alice's.
- ❖ Certificate authority (CA) is an entity that issues digital certificates.
 - A digital certificate certifies the ownership of a public key by the named subject of the certificate.



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**Non-
examinable**

SSL: Secure Sockets Layer

SSL is a widely deployed security protocol.

- ❖ Applicable to TCP applications
- ❖ A variation is **TLS (Transport Layer Security)** defined in RFC 2246.
- ❖ Supported by almost all modern browsers and web servers.
- ❖ For example, **https = http + SSL/TLS**
 - adding security capabilities of SSL/TLS to standard HTTP communications.

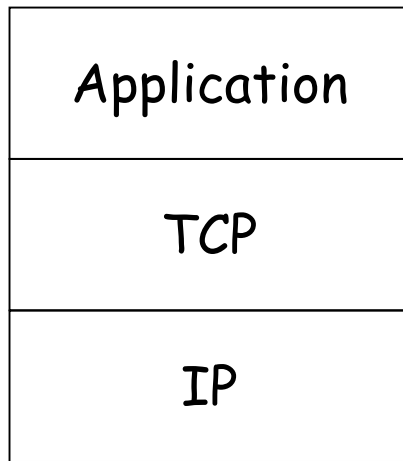
Common SSL symmetric ciphers

- DES - Data Encryption Standard: block
- 3DES - Triple strength: block
- RC2 - Rivest Cipher 2: block
- RC4 - Rivest Cipher 4: stream

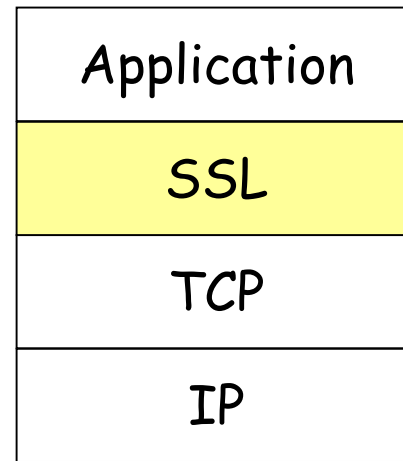
SSL public key encryption

- RSA

SSL: Secure Sockets Layer



Normal Application

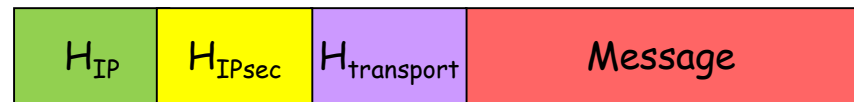
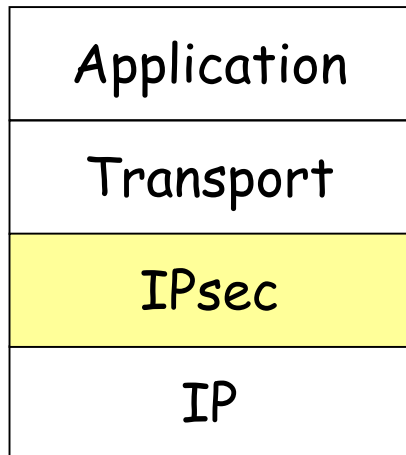


Application with SSL

- ❖ SSL provides application programming interface (API) to applications.
 - Java SSL libraries/classes readily available.

Internet Protocol Security (IPsec)

- ❖ IPsec is a suite of protocols that secure communications by authenticating and encrypting each IP packet of a communication session.



Packet structure w/ IPsec

- ❖ Both SSL and IPsec can be used to build VPN.
 - SoC and NUS WebVPN run over SSL.

Lecture 8: Summary

basic techniques

- data confidentiality (symmetric and public keys)
- message digest
- message authentication code
- digital signature

.... used in many different security scenarios

- https
- secure transport (SSL)
- IPsec
- 802.11 WEP

