## CS2105 Introduction to Computer Networks

# Lecture 10 Local Area Network

29 October 2018

## PREVIOUS LECTURE

# Link Layer

- Link layer transmits data over a single link
- Possible services by link layer protocols:
  - Framing (basic)
    - IP datagrams are encapsulated in link layer frames for transmission
  - Error Checking (optional)
    - Parity check (single, 2D); CRC
    - May also optionally perform Error Correction
  - Multiple Access Control (optional)
    - Determine in which manner multiple nodes share a broadcast channel
    - Be Nice!

## PREVIOUS LECTURE

## Multiple Access Protocols

- Channel partitioning protocols:
  - Divide channel into slots by time, frequency, or code
  - Channel is shared fairly and efficiently if most nodes have data to send
- Random access protocols:
  - Efficient at low load: single node can fully utilize channel
  - High collision rate at high load: wasted channel time
- Taking turns protocols:
  - Polling from master node or token passing
  - Efficient at both low and high load
  - Single point of failure

## **Learning Outcomes**

#### After this class, you are expected to understand:

- the role MAC address.
- the role a switch in interconnecting subnets in a LAN.
- how switching table is built and how it is used to forward link-layer frames.
- how ARP allows a host to discover the MAC address of other nodes in the same subnet.
- the link properties of a wireless link.
- how CSMA/CA works and how it addresses the hidden node problem

Application

Transport

Network

Link

Physical

You are still here

# Chapter 5

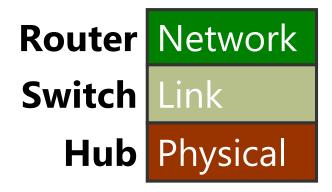
- 5.1 Introduction to the Link Layer
- 5.2 Error detecting and correction
- 5.3 Multiple access protocols
  - 5.3.1 Channel Partitioning Protocols
  - 5.3.2 Random Access Protocols
  - 5.3.3 Taking-Turns Protocols
- 5.4 Switched Local Area Networks
  - 5.4.1 Link layer addressing and ARP
  - 5.4.2 Ethernet
  - 5.4.3 Link layer switches

Discussed Last week

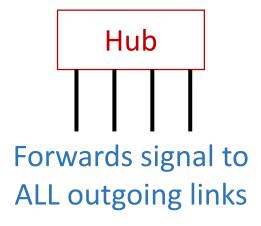
Kurose Textbook, Chapter 5 (Some slides are taken from the book)

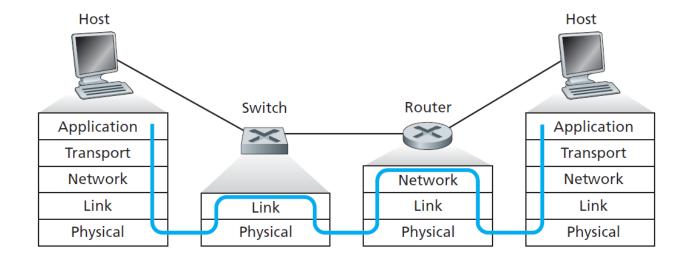
#### Local Area Network

How to inter-connect a large number of hosts in a subnet?



Router	Switches
Check IP address	Check MAC address
Store-and-forward	Store-and-forward
Compute routes to destination	Forward frame to outgoing link or broadcast





## MAC Address (1/2)

Every adapter (NIC) has a MAC address

- Used to send & receive link layer frames
- On receive of frame, checks destination MAC address
  - Match self MAC address: extracts datagram and passes up the protocol stack
  - Does not match: discard

## MAC Address (2/2)

• 6-bytes (48-bits) burned in NIC ROM

• Example: Properties

IPv4 address: 172.26.184.138

IPv4 DNS servers: 192.168.140.2

Primary DNS suffix: comp.nus.edu.sg

Manufacturer: Intel Corporation

Description: Intel(R) Ethernet Connection (5) I219-LM

Driver version: 12.15.22.6

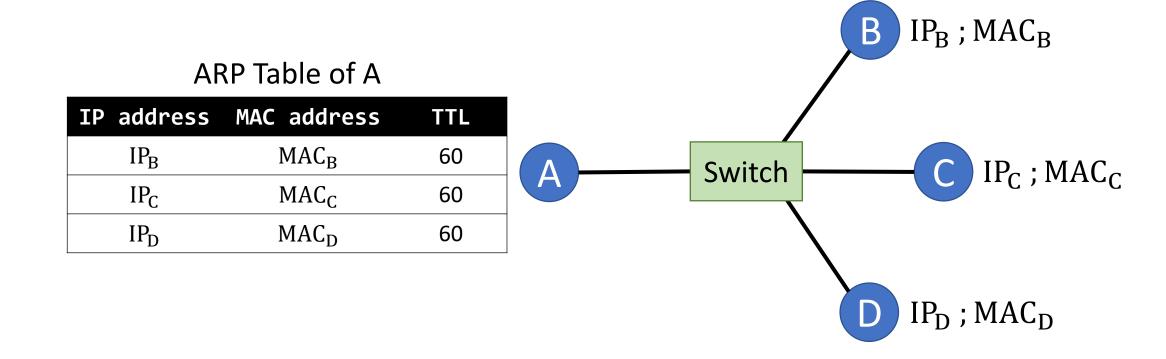
Physical address (MAC): 8C-EC-4B-50-F3-59

- Allocation administered by IEEE
  - First 3 bytes identifies the vendor
  - http://www.coffer.com/mac\_find

## MAC Address vs. IP Address

MAC Address	IP Address
Permanent	Dynamic
Hardware Assigned	Network Assigned
Node-to-node	Host-to-host
Flat	Hierarchical
Analogy: NRIC number	Analogy: Postal address

- Each IP node (e.g., host, router) has an ARP table
  - Stores mapping of IP address to MAC address of other nodes in the same subnet
  - Format: <IP address> <MAC address> <TTL>

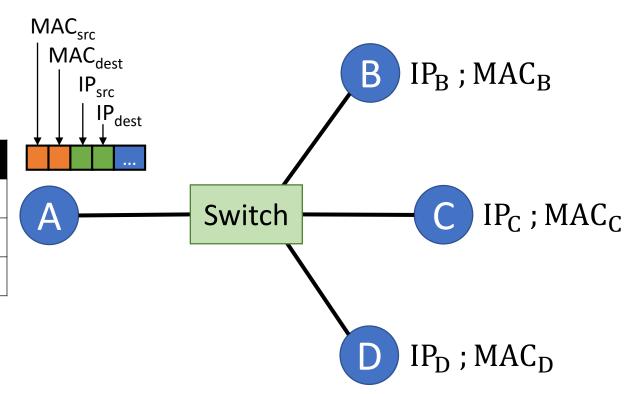


- Sending frame in the same subnet (e.g., A to B)
  - Assume A knows B's MAC address from its ARP table
  - Construct frame with B's MAC address as destination address
    - Only B will process the frame

• Others will ignore

ARP Table of A

IP address	MAC address	TTL
$IP_{B}$	$MAC_B$	60
IP <sub>C</sub>	$MAC_{C}$	60
IP <sub>D</sub>	$MAC_D$	60



Sending frame in the same subnet (e.g., A to B)

What if A does not know B's MAC address?

Broadcast ARP query packet with B's IP address

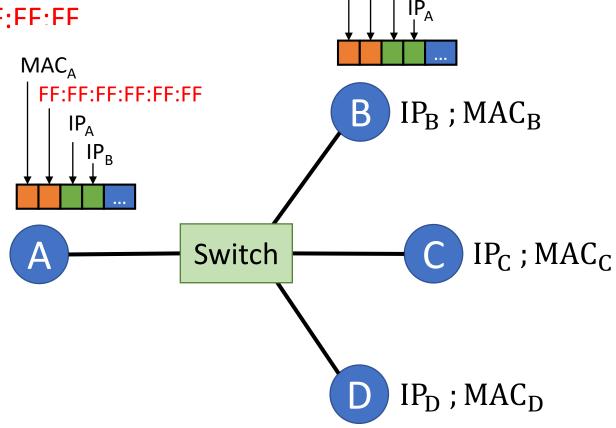
Broadcast Address: FF:FF:FF:FF:FF

B replies with MAC address

• A caches it in ARP table

ARP Table of A

IP address	MAC address	TTL
IP <sub>C</sub>	$MAC_{C}$	60
IP <sub>D</sub>	$MAC_D$	60

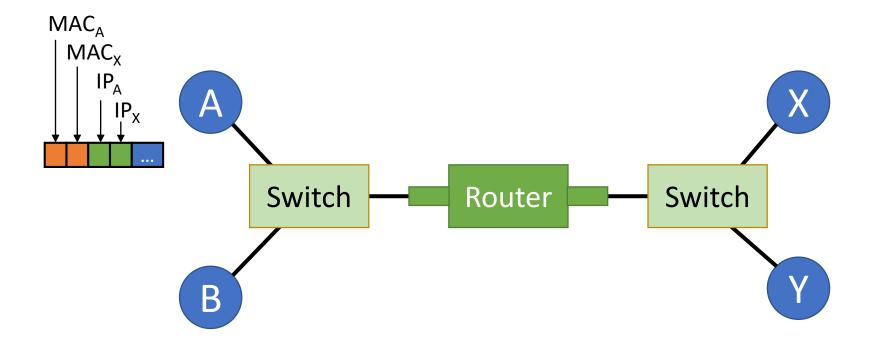


MAC

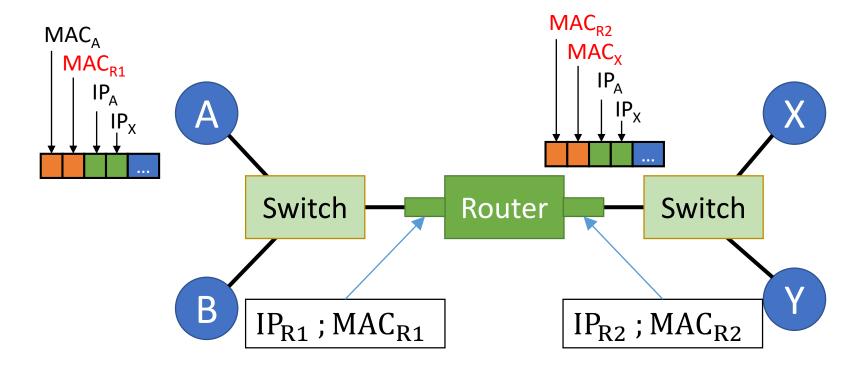
MAC<sub>A</sub>

 $IP_B$ 

- Sending frame to different subnet (e.g., A to X)
  - Attempt 1:
    - Problem?
    - Nobody has MAC<sub>X</sub> in the subnet!



- Sending frame to different subnet (e.g., A to X)
  - Attempt 2:
    - A deeper look into the router
    - Send to router! Router will forward to X!



# Chapter 5

- 5.1 Introduction to the Link Layer
- 5.2 Error detecting and correction
- 5.3 Multiple access protocols
  - 5.3.1 Channel Partitioning Protocols
  - 5.3.2 Random Access Protocols
  - 5.3.3 Taking-Turns Protocols
- 5.4 Switched Local Area Networks
  - 5.4.1 Link layer addressing and ARP
  - 5.4.2 Ethernet
  - 5.4.3 Link layer switches

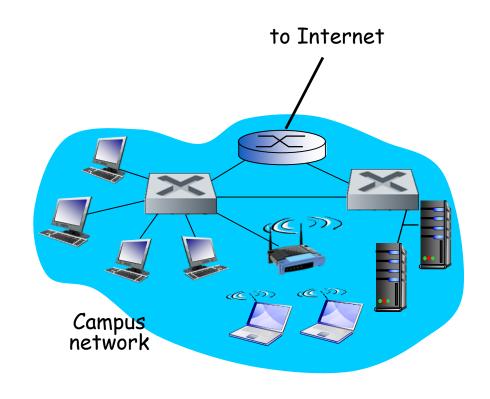
Discussed Last week

Kurose Textbook, Chapter 5 (Some slides are taken from the book)

## Local Area Network (LAN)

• LAN is a computer network that interconnects computers within a geographical area (e.g., office building or campus)

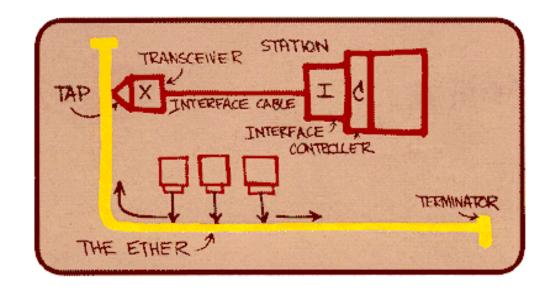
- LAN technologies:
  - IBM Token Ring: IEEE 802.5 standard
  - Ethernet: IEEE 802.3 standard
  - Wi-Fi: IEEE 802.11 standard
  - Others



- "Dominant" wired LAN technology:
  - Developed in mid 1970s
  - Standardized by Xerox, DEC, and Intel in 1978
  - Simpler and cheaper than token ring and ATM



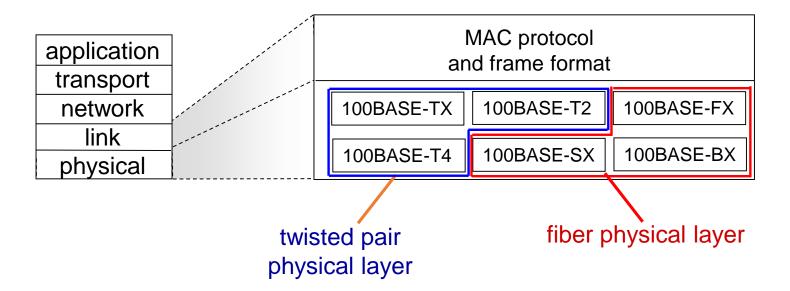
Ethernet connection (Source: Wikipedia)



Metcalfe's Ethernet sketch

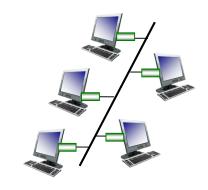
#### 802.3 Ethernet Standards

- Different speeds: 2 Mbps, 10 Mbps, 100 Mbps, 1 Gbps, 10 Gbps, 100 Gbps
- Different physical layer media: Cable, Fiber optics
- MAC protocol and frame format <u>remains unchanged</u>



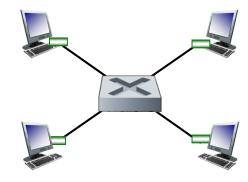
#### Physical Topology

- Bus topology: popular in mid 90s
  - All nodes can collide with each other



Ethernet with bus topology

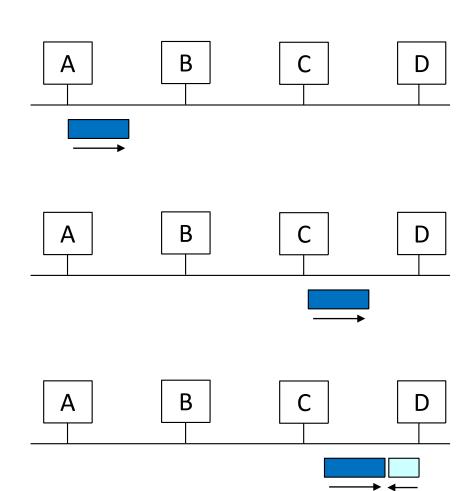
- Star topology: prevails today
  - Switch in center
  - Nodes do not collide with each other



Ethernet with star topology

#### Physical Topology

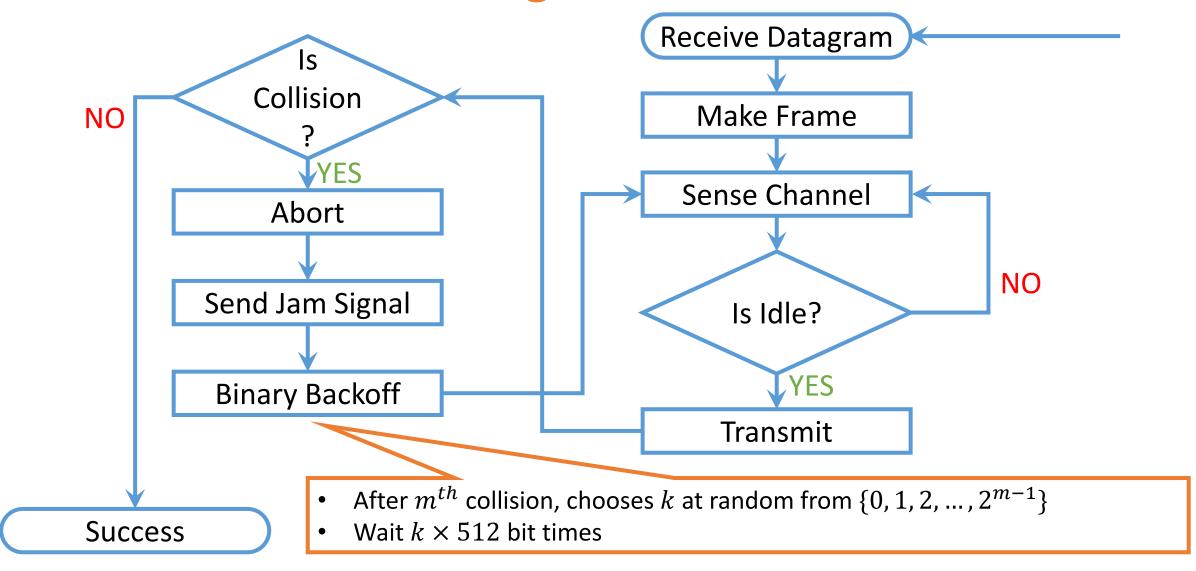
- Bus topology: popular in mid 90s
  - All nodes can collide with each other
  - Example:
    - A sends a frame at time t
    - A's frame reached D at time t+d
    - D begins transmission at time t + d 1
    - Collision!



#### **Data Delivery Service**

- Connectionless: no handshaking between sending and receiving NICs.
- Unreliable: receiving NIC doesn't send ACK or NAK to sending NIC.
  - Data in dropped frames will be recovered only if initial sender uses higher layer RDT (e.g., TCP)
  - Otherwise, data is lost
- MAC: CSMA/CD with binary (exponential) backoff

# Ethernet CSMA/CD Algorithm



# Chapter 5

- 5.1 Introduction to the Link Layer
- 5.2 Error detecting and correction
- 5.3 Multiple access protocols
  - 5.3.1 Channel Partitioning Protocols
  - 5.3.2 Random Access Protocols
  - 5.3.3 Taking-Turns Protocols
- 5.4 Switched Local Area Networks
  - 5.4.1 Link layer addressing and ARP
  - 5.4.2 Ethernet
  - 5.4.3 Link layer switches

Discussed Last week

Kurose Textbook, Chapter 5 (Some slides are taken from the book)

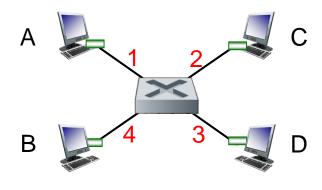
### **Ethernet Switch**

#### Properties:

- Store and forward Ethernet frames
  - Examine incoming frame's MAC address
  - Selectively forward frame to one-or-more outgoing links
- Transparent to hosts
  - No IP address
- Star topology
  - Each host has dedicated link to switch
- Full duplex
- Buffered



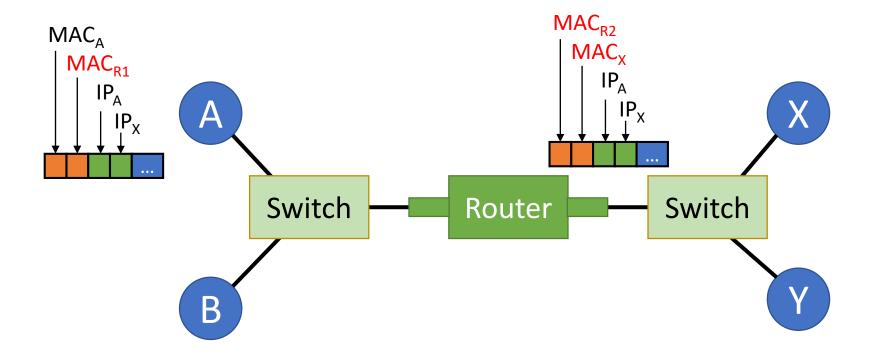
a 50-port Ethernet switch (Source: Wikipedia)



A switch with 4 interfaces (1, 2, 3, 4)

## Switch Forwarding Table

- Recall: switch does not "broadcast" like a hub
- Question: how does a switch know which interface to forward the frame to?



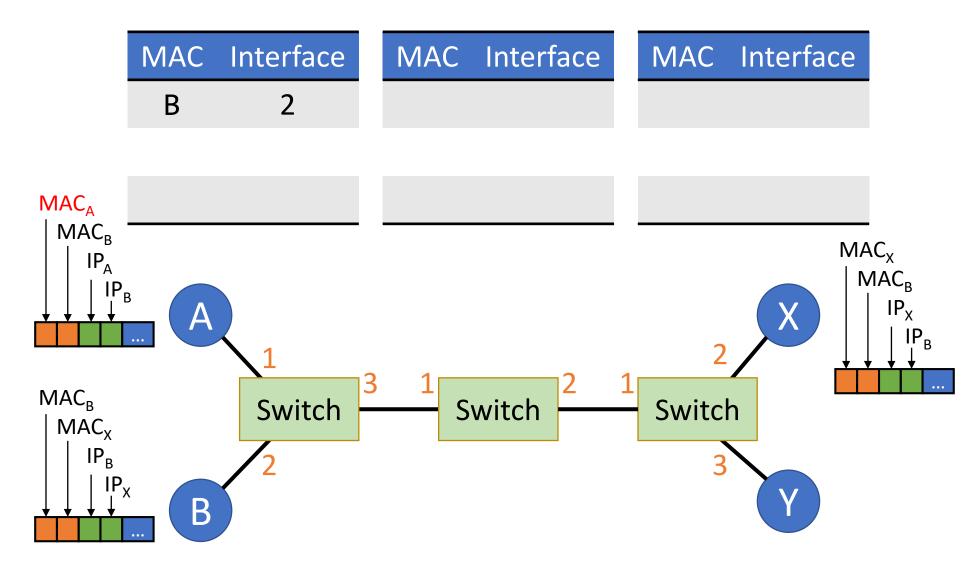
# Switch Forwarding Table

- Format of entry:
  - <MAC address> <interface> <TTL>

- Self-learning:
  - Switch learns which hosts can be reached through which interface
    - 1. Broadcast
    - 2. Forwarding
    - 3. Filtering

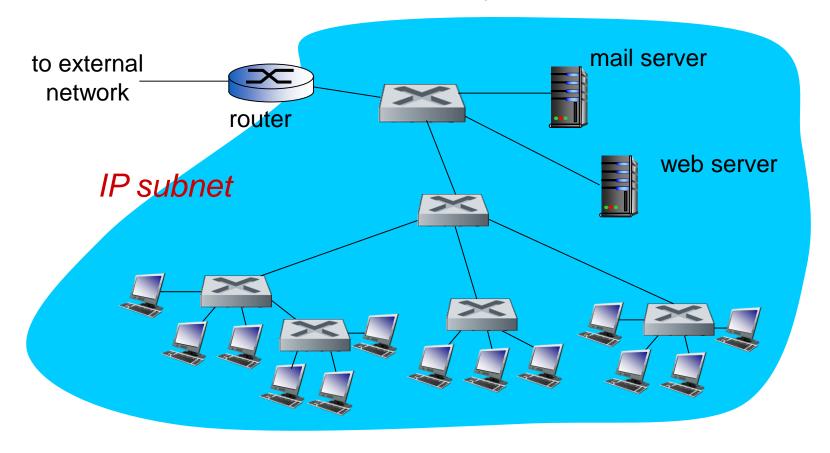
MAC Address	Interface	TTL
FA:CE:BD:0C:11:FD	1	60
66:23:C6:1D:FE:32	2	60
15:00:2A:F1:CE:A1	3	60

# **Switch Forwarding Table**



# Interconnecting Switches

• Switches can be connected in hierarchy



# Chapter 5

- 5.1 Introduction to the Link Layer
- 5.2 Error detecting and correction
- 5.3 Multiple access protocols
  - 5.3.1 Channel Partitioning Protocols
  - 5.3.2 Random Access Protocols
  - 5.3.3 Taking-Turns Protocols
- 5.4 Switched Local Area Networks
  - 5.4.1 Link layer addressing and ARP
  - 5.4.2 Ethernet
  - 5.4.3 Link layer switches

More on Wireless LAN

Kurose Textbook, Chapter 5 (Some slides are taken from the book)

Terminology

• BSS: Basic Service Set

AP: Access Point

SSID: Service Set Identifier

• RSSI: Received Signal Strength Indicator

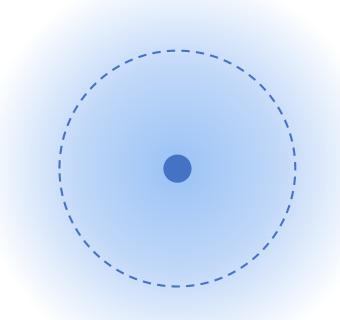






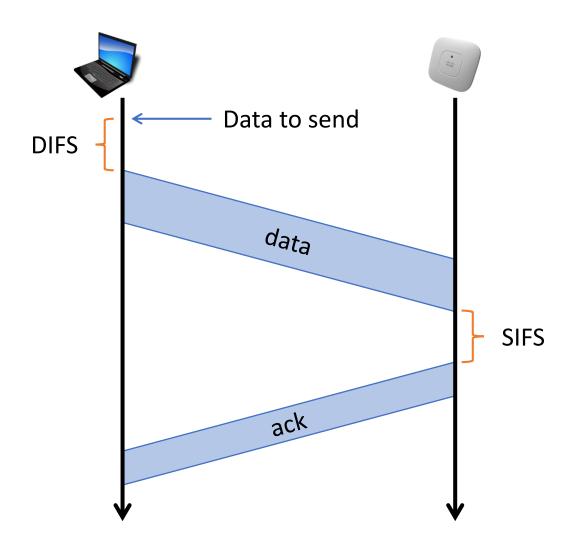


- Implications of wireless
  - Limited Range
  - Broadcast
  - High Error Rate
  - Multipath Propagation



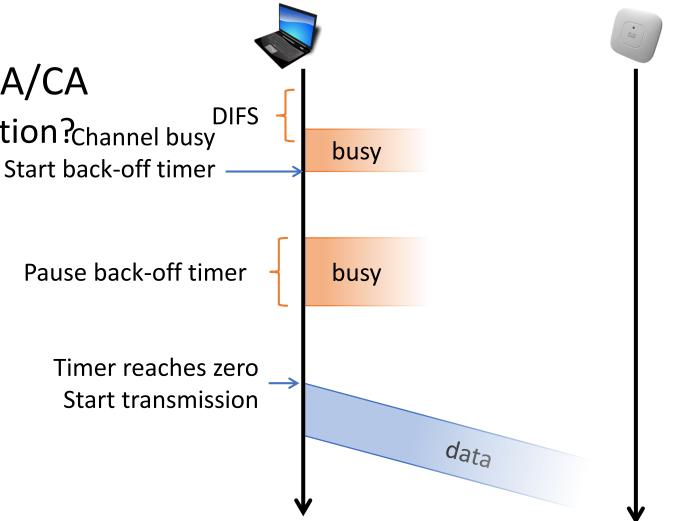
#### Recap

Use link-layer ACK to detect collision



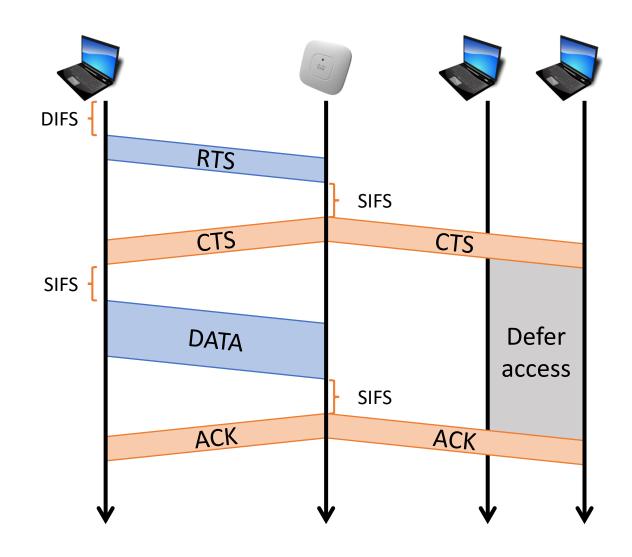
#### Recap

- Wireless LAN uses CSMA/CA
  - Why not collision detection? Channel busy
    - RSSI<sub>rcv</sub> < RSSI<sub>snd</sub>
    - Hidden Node problem



#### Recap

- Wireless LAN uses RTS/CTS
  - Hidden Node problem
  - RTS before sending
  - On CTS, send
  - Wait for ACK
  - Other nodes wait until ACK



## Summary

- ARP protocol to resolve IP ←→ MAC mapping
- Ethernet (IEEE 802.3)
  - CSMA/CD protocol with binary (exponential) back-off
  - Self-learning Ethernet switch table
    - Broadcasting
    - Forwarding
    - Filtering
- Wireless LAN (IEEE 802.11)
  - CSMA/CA with RTS/CTS ACK/NAK
  - Hidden node problem