Binary Search

Shusen Wang

Stevens Institute of Technology

$$\operatorname{arr} = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$

$$arr = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$

- Inputs: (i) an array whose elements are in the ascending order and (ii) a key.
- Goal: Search the key in the array. If found, return its index; if not found, return −1.

$$arr = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$

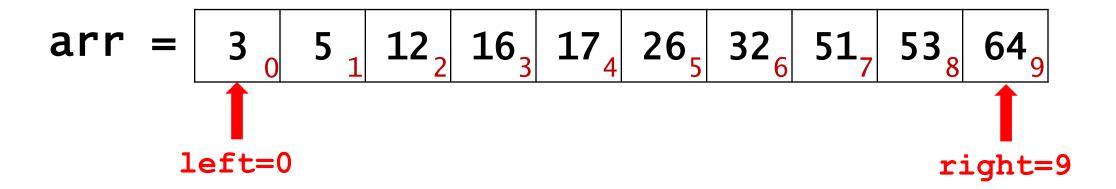
Example 1:

- Search for the element 53.
- Return 8.

$$arr = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$

Example 2:

- Search for the element 9.
- Return -1.



mid ← [(left + right)/2].
 If arr[mid] == key ==> return mid.
 If arr[mid] > key ==> right ← mid-1.
 If arr[mid] < key ==> left ← mid+1.

mid ← [(left + right)/2].
If arr[mid]==key ==> return mid.
If arr[mid]>key ==> right ← mid-1.
If arr[mid]<key ==> left ← mid+1.

```
mid ← [(left + right)/2].
If arr[mid] == key ==> return mid.
If arr[mid] > key ==> right ← mid-1.
If arr[mid] < key ==> left ← mid+1.
```

```
arr = 3 0 5 1 12 16 17 26 32 51 53 64 9

left=0 mid=4 right=9
```

```
mid ← [(left + right)/2].
If arr[mid] == key ==> return mid.
If arr[mid] > key ==> right ← mid-1.
If arr[mid] < key ==> left ← mid+1.
```

```
arr = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}
left=5
```

```
mid ← [(left + right)/2].
If arr[mid] == key ==> return mid.
If arr[mid] > key ==> right ← mid-1.
If arr[mid] < key ==> left ← mid+1.
```

arr =
$$\begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$
left=5

mid ← [(left + right)/2].
If arr[mid] == key ==> return mid.
If arr[mid] > key ==> right ← mid-1.
If arr[mid] < key ==> left ← mid+1.

mid ← [(left + right)/2].
If arr[mid] == key ==> return mid.
If arr[mid] > key ==> right ← mid-1.
If arr[mid] < key ==> left ← mid+1.

- mid $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$.
- If arr[mid] == key ==> return mid.
- If arr[mid]>key ==> right ← mid-1.
- If arr[mid] < key ==> left ← mid+1.

- mid $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$.
- If arr[mid] == key ==> return mid.
- If arr[mid]>key ==> right ← mid-1.
- If arr[mid] < key ==> left ← mid+1.

```
arr = 3 0 5 1 12 16 17 26 32 51 53 64 9 left=5 right=6
```

- mid $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$.
- If arr[mid] == key ==> return mid.
- If arr[mid]>key ==> right ← mid-1.
- If arr[mid] < key ==> left ← mid+1.

```
arr = 3 0 5 1 12 16 17 26 32 51 53 64 9 left=5 right=6
```

- mid $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$.
- If arr[mid] == key ==> return mid.
- If arr[mid]>key ==> right ← mid-1.
- If arr[mid] < key ==> left ← mid+1.

```
arr = 3 5 12 16 17 26 32 51 53 64 9

left=5

right=6
```

- mid $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$.
- If arr[mid] == key ==> return mid.
- If arr[mid]>key ==> right ← mid-1.
- If arr[mid] < key ==> left ← mid+1.

```
arr = 3 5 12 16 17 26 32 51 53 64 9

left=5

right=6
```

- mid $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$.
- If arr[mid] == key ==> return mid.
- If arr[mid]>key ==> right ← mid-1.
- If arr[mid] < key ==> left ← mid+1.

Time Complexity

- Each iteration reduces the size of array by half.
- Let n be the size of the array.
- The total number of iterations is at most $\log_2 n$.
- Time complexity: $O(\log n)$.

```
int search(int arr[], int left, int right, int key) {
     while (left <= right) {
          int mid = (left + right) / 2;
          if (arr[mid] == key)
               return mid; // key is found
          if (arr[mid] < key)</pre>
               left = mid + 1; // search in the right half
               right = mid - 1; // search in the left half
     return -1; // key is not found
```

```
int search(int arr[], int left, int right, int key) {
     while (left <= right) {</pre>
          int mid = (left + right) / 2;
          if (arr[mid] == key)
                return mid; // key is found
          if (arr[mid] < key)</pre>
                left = mid + 1; // search in the right half
          else
                right = mid - 1; // search in the left half
```

return -1; // key is not found

```
int search(int arr[], int left, int right, int key) {
    while (left <= right) {</pre>
          int mid = (left + right) / 2;
          if (arr[mid] == key)
               return mid; // key is found
          if (arr[mid] < key)</pre>
               left = mid + 1; // search in the right half
               right = mid - 1; // search in the left half
     return -1; // key is not found
```

```
int search(int arr[], int left, int right, int key) {
     while (left <= right) {</pre>
    \longrightarrow int mid = (left + right) / 2;
          if (key == arr[mid])
                return mid; // key is found
          if (key > arr[mid])
                left = mid + 1; // search in the right half
          else
                right = mid - 1; // search in the left half
     return -1; // key is not found
```

```
int search(int arr[], int left, int right, int key) {
     while (left <= right) {</pre>
          int mid = (left + right) / 2;
    if (key == arr[mid])
               return mid; // key is found
         if (key > arr[mid])
               left = mid + 1; // search in the right half
          else
               right = mid - 1; // search in the left half
     return -1; // key is not found
```

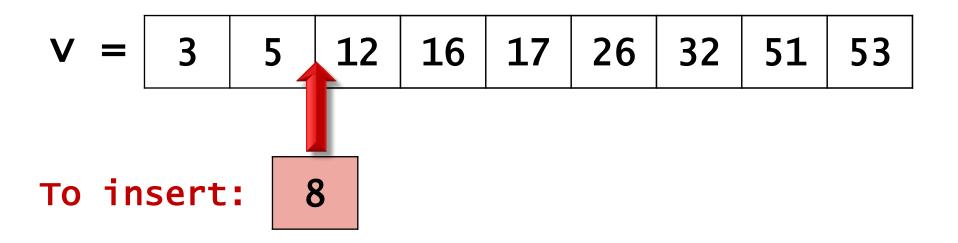
```
int search(int arr[], int left, int right, int key) {
     while (left <= right) {</pre>
          int mid = (left + right) / 2;
          if (arr[mid] == key)
               return mid; // key is found
          if (arr[mid] < key)</pre>
               left = mid + 1; // search in the right half
               right = mid - 1; // search in the left half
     return -1; // key is not found
```



Vector

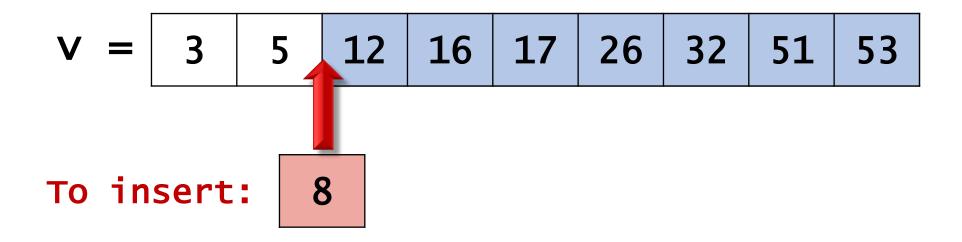
- The ascending order must be kept; otherwise, search would take O(n) time.
- Inserting an item into the middle has O(n) time complexity (on average).

Vector



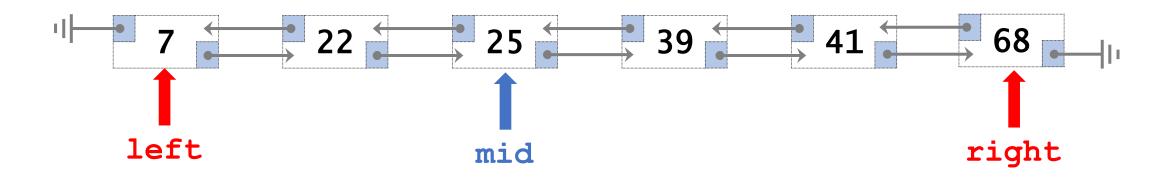
• Inserting an item into the middle has O(n) time complexity (on average).

Vector



• Inserting an item into the middle has O(n) time complexity (on average).

List



- Can we perform binary search in the list?
- No. Given left and right, we cannot get mid efficiently.

Comparisons

Vector

List

SearchInsertion $O(\log n)$ O(n)O(n)O(1)

Comparisons

	Search	Insertion
Vector	$O(\log n)$	O(n)
List	O(n)	0(1)
Skip List	$O(\log n)$	$O(\log n)$

Thank You!