# **Disjoint Sets**

**Shusen Wang** 

# Disjoint Sets (Math)

- Disjoint sets are those whose pairwise intersections are empty sets.
- Disjoint sets:

• 
$$S_1 = \{0, 3\},$$

• 
$$S_2 = \{1, 2, 7\},$$

• 
$$S_3 = \{4, 5, 6\}.$$

# Disjoint Sets (Math)

- Disjoint sets are those whose pairwise intersections are empty sets.
- Not disjoint sets:
  - $S_1 = \{0, 3\},$
  - $S_2 = \{1, 2, 3\},$
  - $S_3 = \{4, 5, 6\}.$
  - Because the intersection is not empty:  $S_1 \cap S_2 = \{3\}$ .

#### **Data Structure**

#### **Initialization**

Element: 0 1 2 3 4 5 6 7

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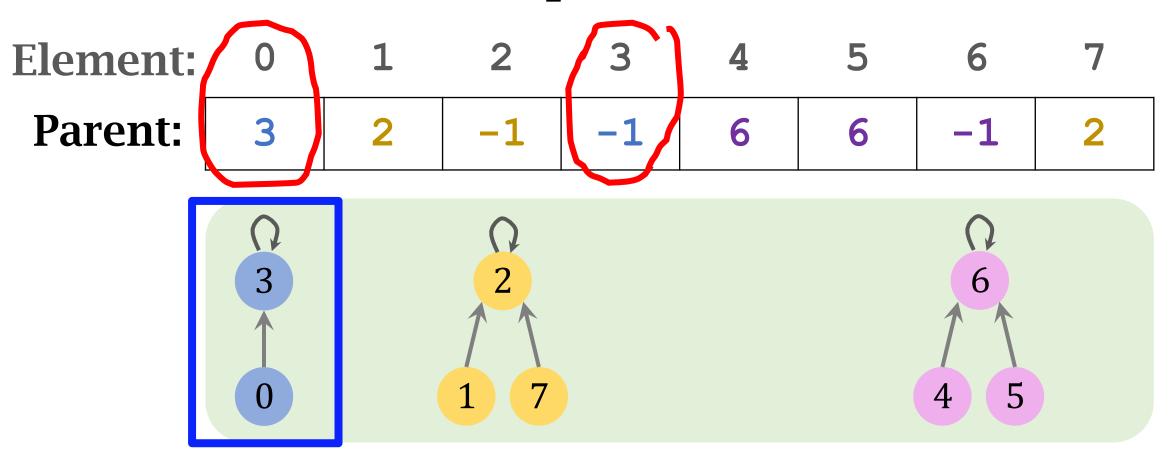
Element: 0 1 2 3 4 5 6 7

Parent: -1 -1 -1 -1 -1 -1

Element: 0 1 2 3 4 5 6 7

Parent:

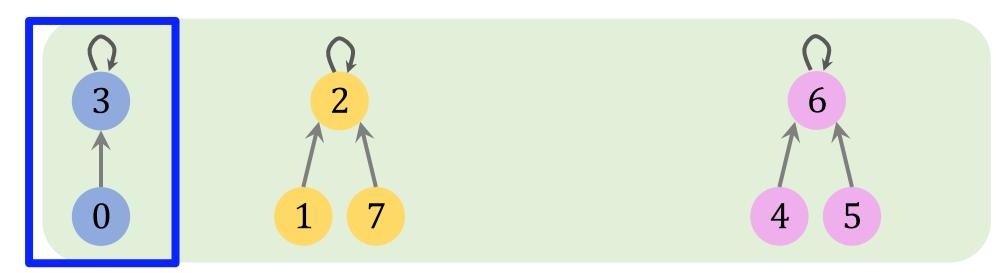
Representing sets  $S_1 = \{0, 3\}$   $S_2 = \{1, 2, 7\}$ , and  $S_3 = \{4, 5, 6\}$ .



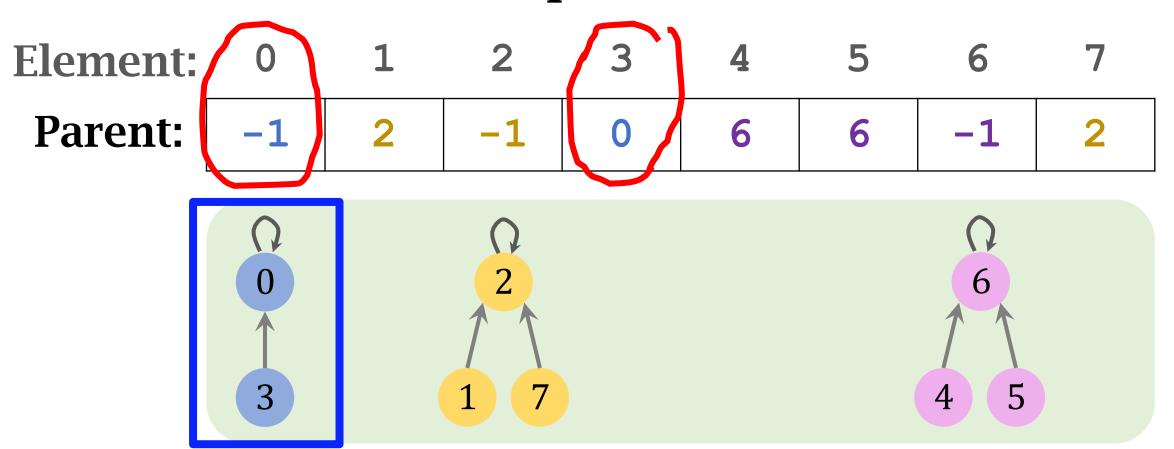
Representing sets 
$$S_1 = \{0, 3\}$$
  $S_2 = \{1, 2, 7\}$ , and  $S_3 = \{4, 5, 6\}$ .

 Element:
 0
 1
 2
 3
 4
 5
 6
 7

 Parent:
 3
 2
 -1
 -1
 6
 6
 -1
 2



Representing sets  $S_1 = \{0, 3\}$ ,  $S_2 = \{1, 2, 7\}$ , and  $S_3 = \{4, 5, 6\}$ .



Representing sets  $S_1 = \{0, 3\}$ ,  $S_2 = \{1, 2, 7\}$ , and  $S_3 = \{4, 5, 6\}$ .

Element: 0 1 2 3 4 5 6 7

Parent:

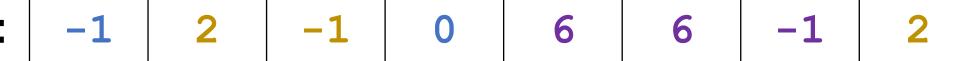




Representing sets  $S_1 = \{0, 3\}$   $S_2 = \{1, 2, 7\}$ , and  $S_3 = \{4, 5, 6\}$ .

Element: 0 1 2 3 4 5 6 7

Parent:





Representing sets  $S_1 = \{0, 3\}, S_2 = \{1, 2, 7\}, \text{ and } S_3 = \{4, 5, 6\}.$ 

## **Operations**

Element: 0 1 2 3 4 5 6 7

Parent: -1 2 -1 0 6 6 -1 2

**Operation 1:** Are elements *i* and *j* belong to the same set?

- Find the roots of *i* and *j*.
- If they have the same root, then return TRUE.
- Otherwise, return False.

## **Operations**

Element: 0 1 2 3 4 5 6 7

Parent:

 -1
 2
 -1
 0
 6
 -1
 2

**Operation 1:** Are elements *i* and *j* belong to the same set?

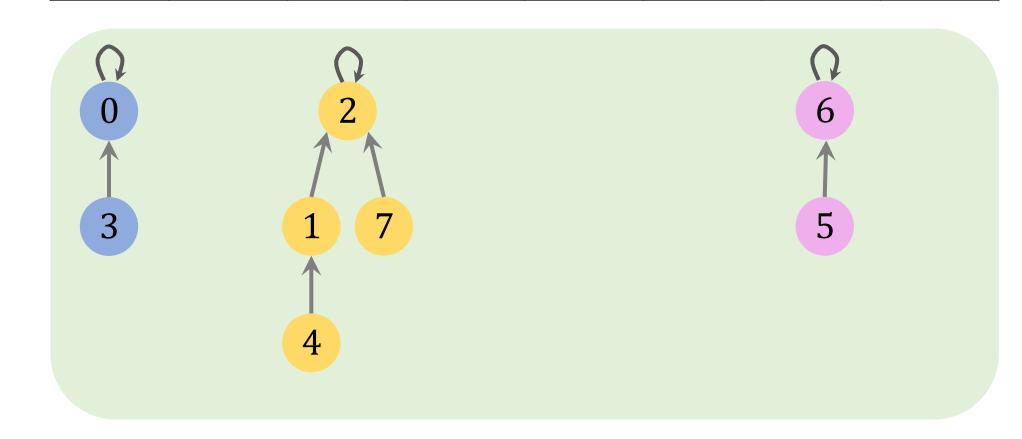
**Operation 2:** Merge (union) sets  $S_1$  and  $S_2$ .

- Find the root of  $S_1$ ; denote it by  $r_1$ .
- Find the root of  $S_2$ ; denote it by  $r_2$ .
- Make  $r_1$  a child of  $r_2$ .

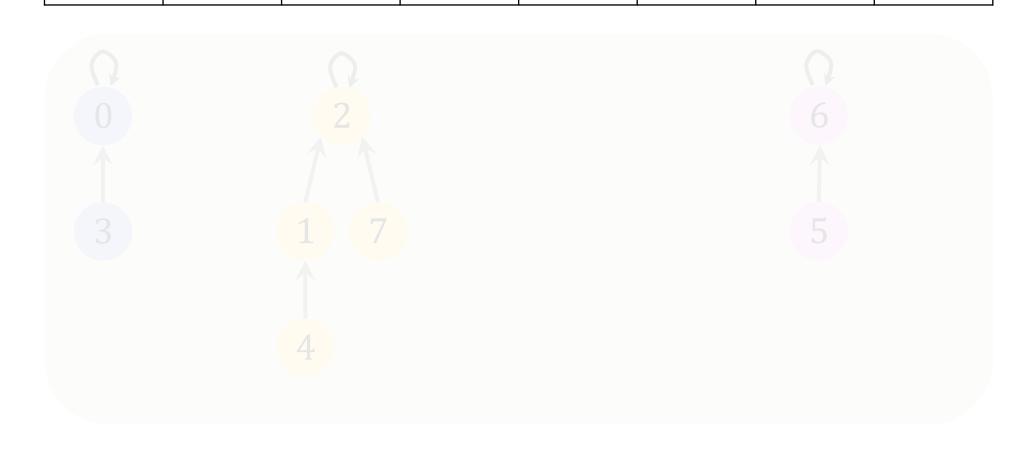
# **Find Root**

#### **Current State**

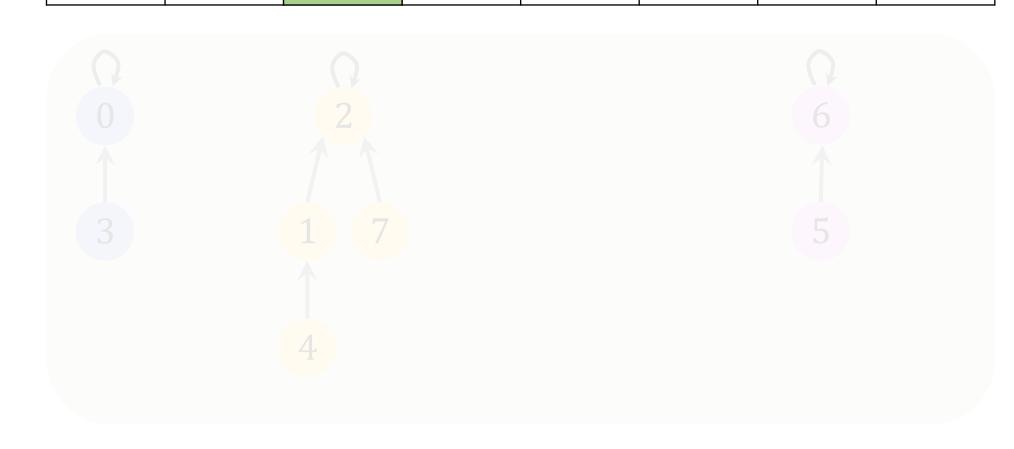
#### **Current State**



#### What is the root of 2?



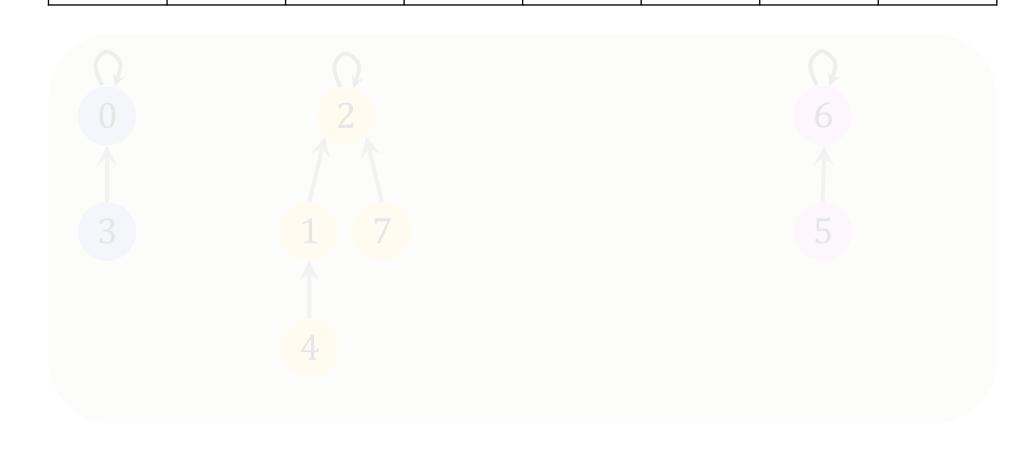
#### The root of 2 is 2 itself.



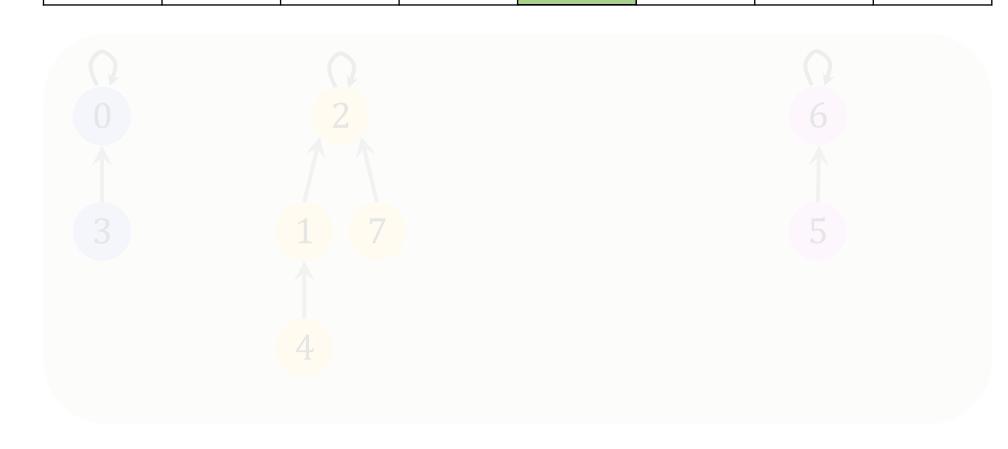
#### The root of 2 is 2 itself.

Element: 0 Parent:

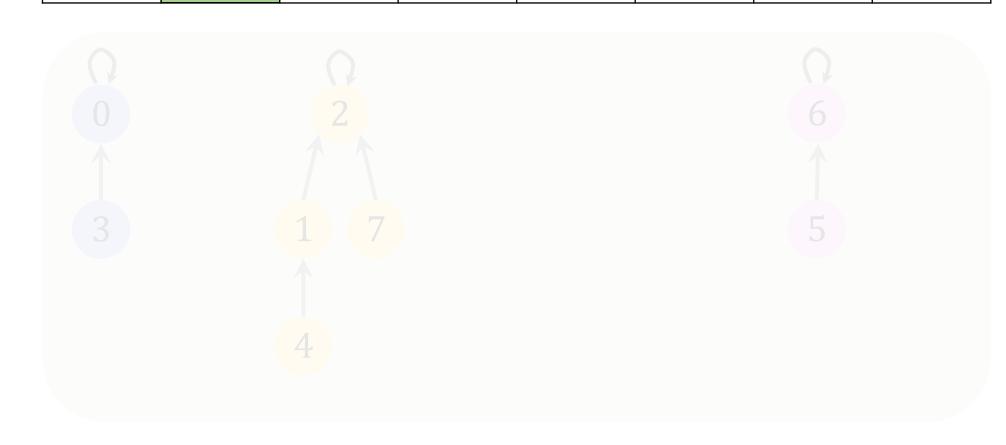
#### What is the root of 4?



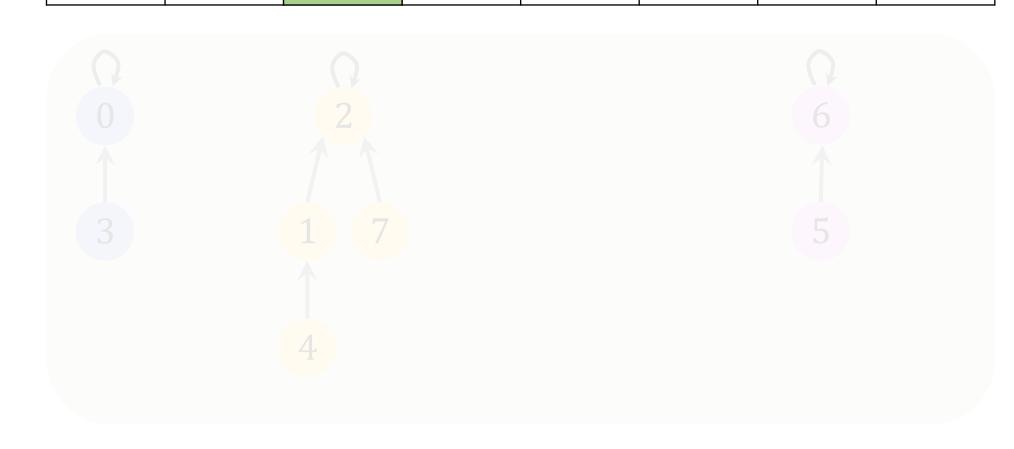
#### What is the root of 4?



#### What is the root of 4?



#### The root of 4 is 2.



#### The root of 4 is 2.

Element: 0 Parent:

#### Are two elements in the same set?

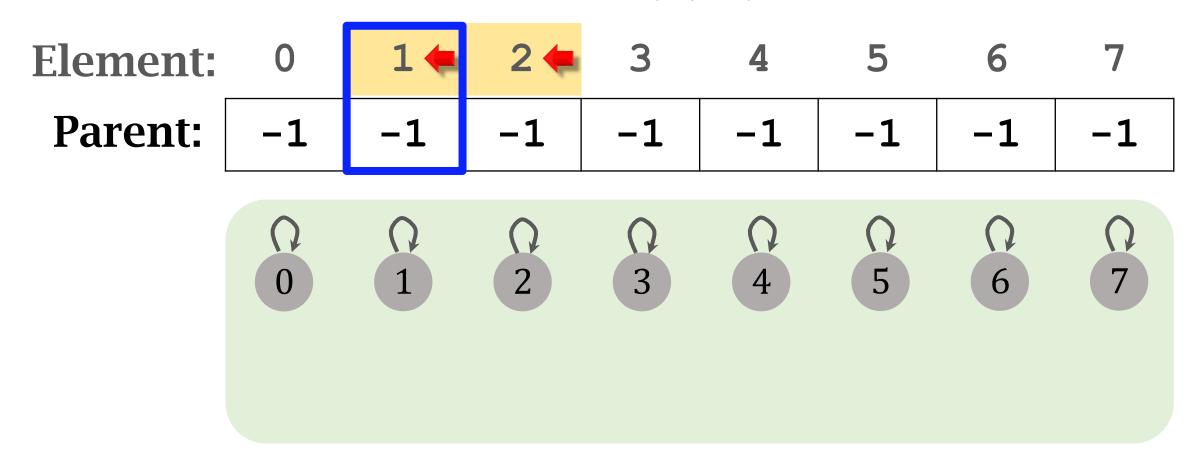
- **Inputs:** two elements (indices) *i* and *j*.
- Question: Are they in the same set?
- Solution:
  - 1.  $r_i \leftarrow$  the root of i.
  - 2.  $r_i \leftarrow \text{the root of } i$ .
  - 3. If  $r_i$  and  $r_i$  are the same, then i and j are in the same set.
  - 4. Otherwise, *i* and *j* are in different sets.

# **Union Operation**

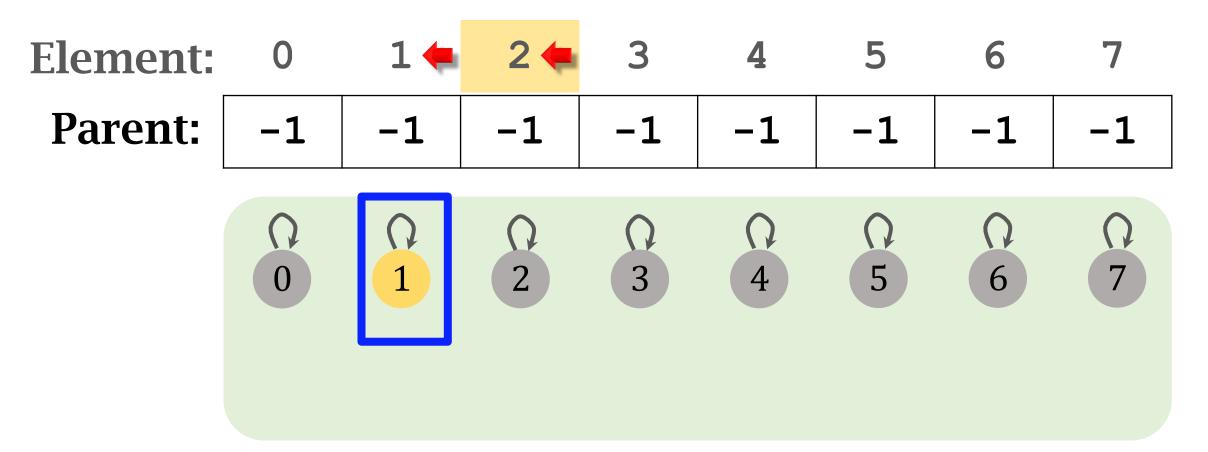
#### **Initial State**

Element: 0 1 2 3 4 5 6 7

Parent: | -1 | -1 | -1 | -1 | -1 | -1 | -1



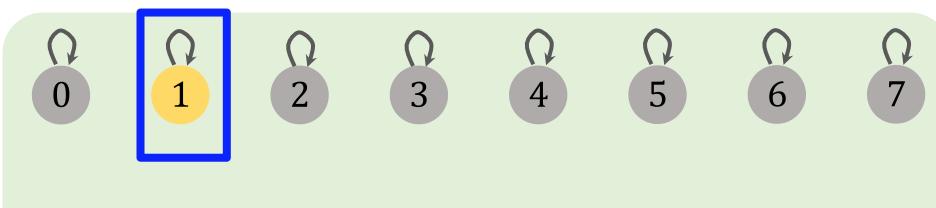
1. Find the root of 1:  $r_1 = 1$ .



1. Find the root of 1:  $r_1 = 1$ .

Element: 0 1 ← 2 ← 3 4 5 6 7

Parent: -1 -1 -1 -1 -1 -1 -1



- 1. Find the root of 1:  $r_1 = 1$ .
- 2. Find the root of 2:  $r_2 = 2$ .

Element: 0 1 **4** 2 **4** 5 6 7

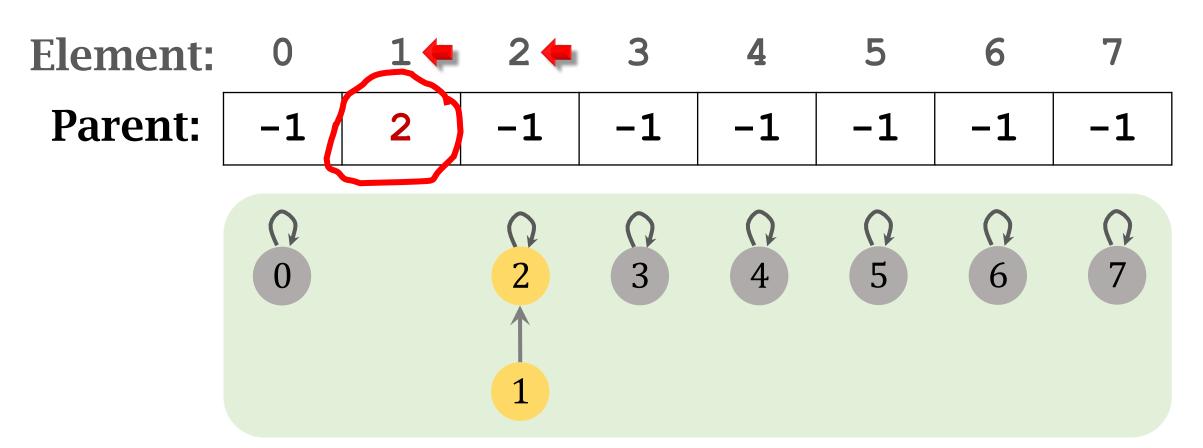
**Parent:** 

- 1. Find the root of 1:  $r_1 = 1$ .
- 2. Find the root of 2:  $r_2 = 2$ .

Element: 0 1 - 2 - 3 4 5 6 7

Parent:

- 1. Find the root of 1:  $r_1 = 1$ .
- 2. Find the root of 2:  $r_2 = 2$ .
- 3. Make  $r_2$  the parent of  $r_1$ .

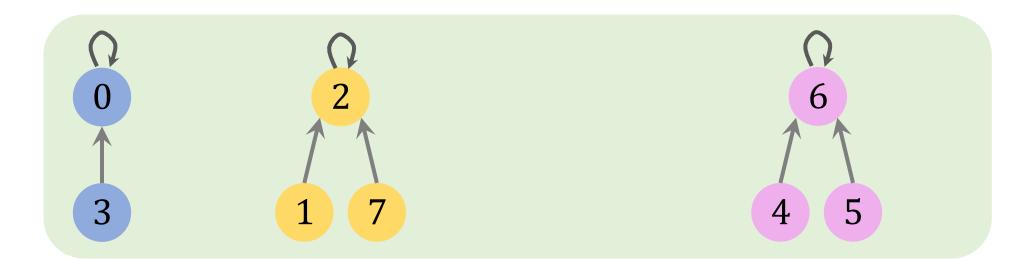


- 1. Find the root of 1:  $r_1 = 1$ .
- 2. Find the root of 2:  $r_2 = 2$ .
- 3. Make  $r_2$  the parent of  $r_1$ : parent  $[r_1] = r_2$ .

# After many unions...

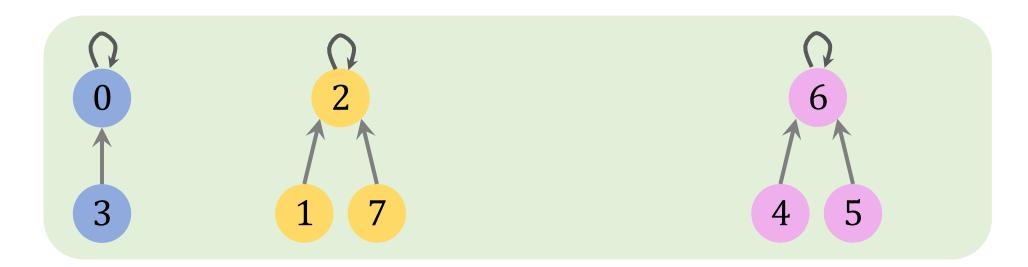
Element: 0 1 2 3 4 5 6 7

**Parent:** 



Representing sets  $S_1 = \{0,3\}$   $S_2 = \{1,2,7\}$ , and  $S_3 = \{4,5,6\}$ .

# **Union(4, 5)**



Element: 0 1 2 3 4 5 6 7

Parent:

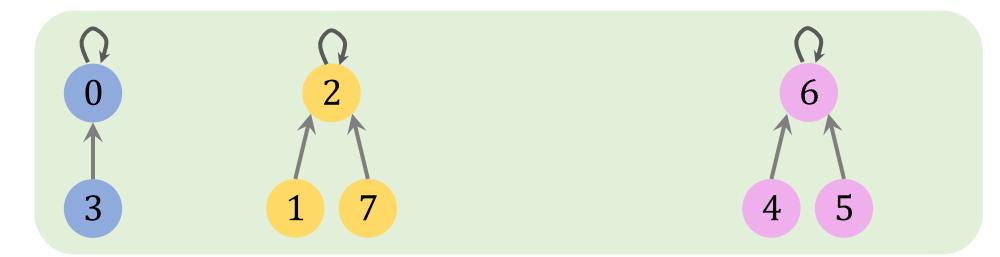
 -1
 2
 -1
 0
 6
 6
 -1
 2



1. Find the root of 4:  $r_4 = 6$ .

Element: 0 1 2 3 4 5 6 7

Parent:

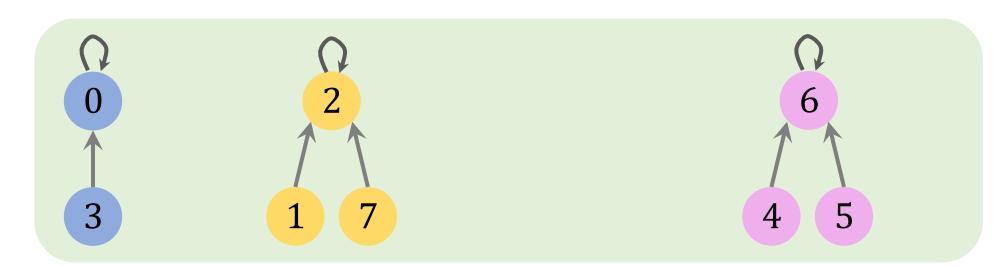


- 1. Find the root of 4:  $r_4 = 6$ .
- 2. Find the root of 5:  $r_5 = 6$ .
- 3. They have the same root; thus no operation.

Element: 0 1 2 3 4 ← 5 6 7 ← Parent: -1 2 -1 0 6 6 -1 2



Element: 0 1 2 3 4 ← 5 6 7 ← Parent: -1 2 -1 0 6 6 -1 2



1. Find the root of 4:  $r_4 = 6$ .

Element: 0 1 2 3 4 5 6

Parent:

-1	2	-1	0	6	6	-1	2



- 1. Find the root of 4:  $r_4 = 6$ .
- 2. Find the root of 7:  $r_7 = 2$ .

Element: 0 1 2 3 4 5 6 7 6

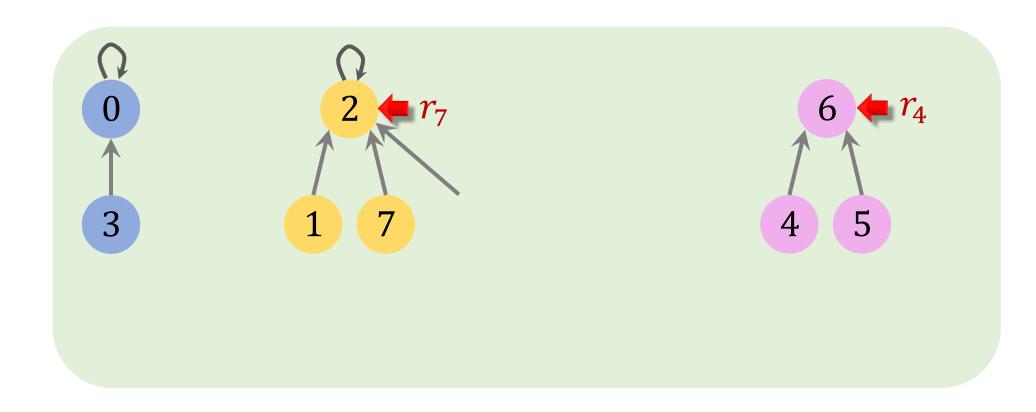
Parent:

-1	-1	2	-1	0	6	6	-1	2
----	----	---	----	---	---	---	----	---

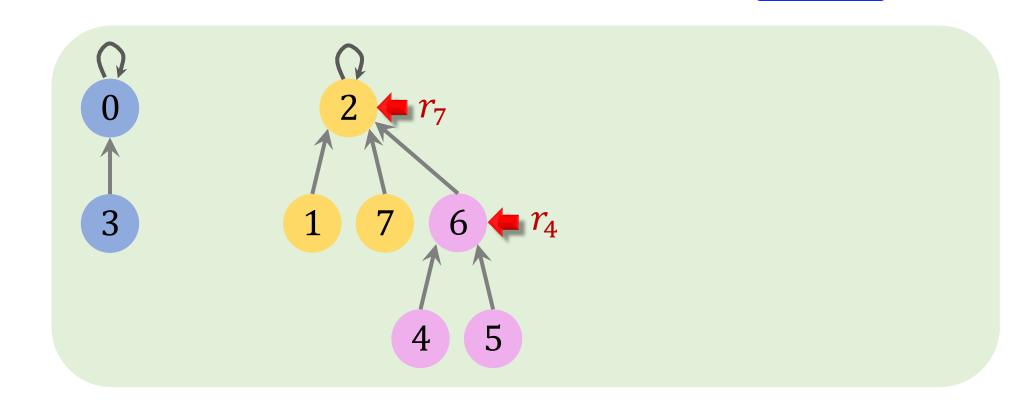


- 1. Find the root of 4:  $r_4 = 6$ .
- 2. Find the root of 7:  $r_7 = 2$ .
- 3. Make  $r_7$  the parent of  $r_4$ : parent  $[r_4] = r_7$ .

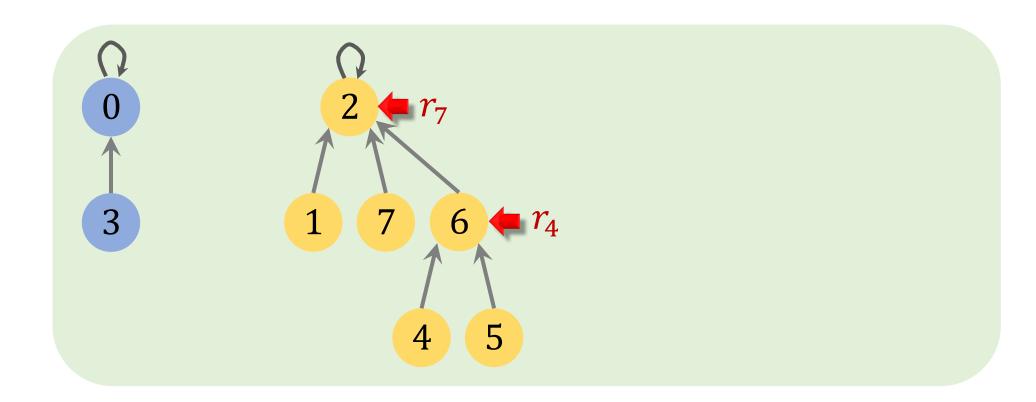
Element: 0 1 2 3 4 ← 5 6 7 ← Parent: -1 2 -1 0 6 6 -1 2



Element: 0 1 2 3 4 ← 5 6 7 ← Parent: -1 2 -1 0 6 6 -1 2



Element: 0 1 2 3 4 ← 5 6 7 ← Parent: -1 2 -1 0 6 6 2 2

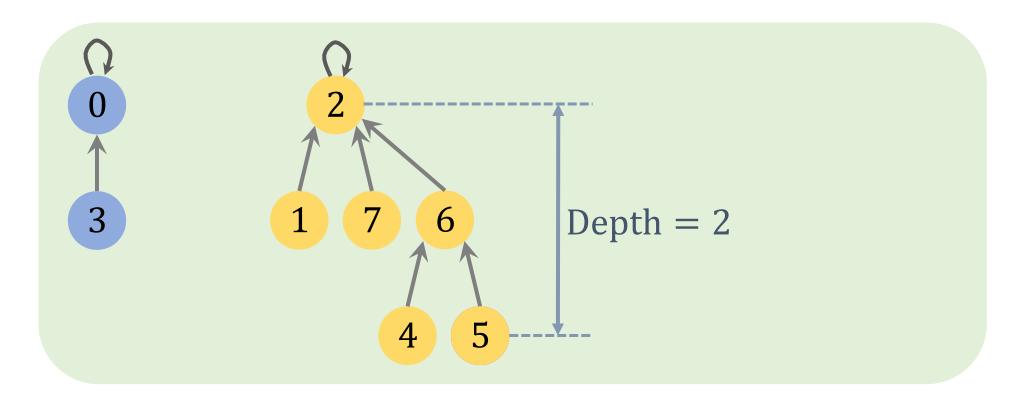


# **Time Complexities**

# Time complexity of finding root

Element: 0 1 2 3 4 5 6 7

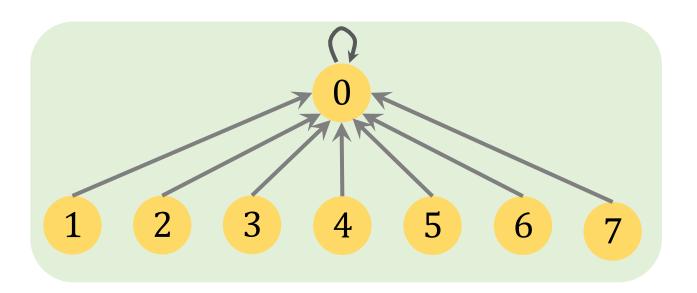
Parent: -1 2 -1 0 6 2 2



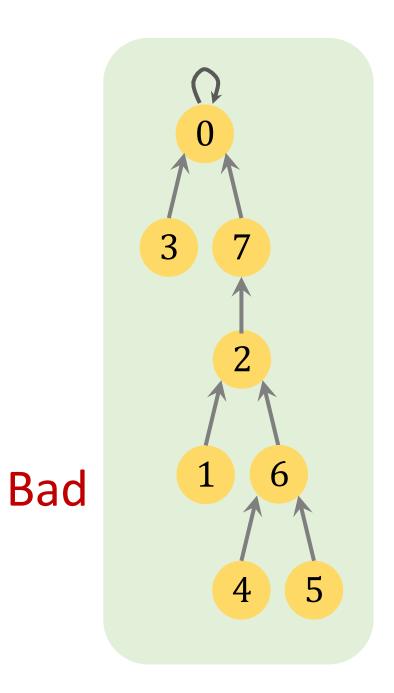
#### Time complexities

- Time complexity of finding root is the depth of the vertex.
- Time complexity of union is the depth of the tree.
- In the worst case, it can be O(n). (Here, n is the total number of elements.)

**Question:** How to reduce the time complexity?

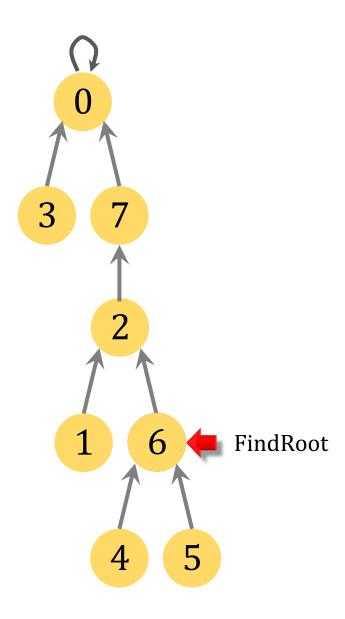


Good



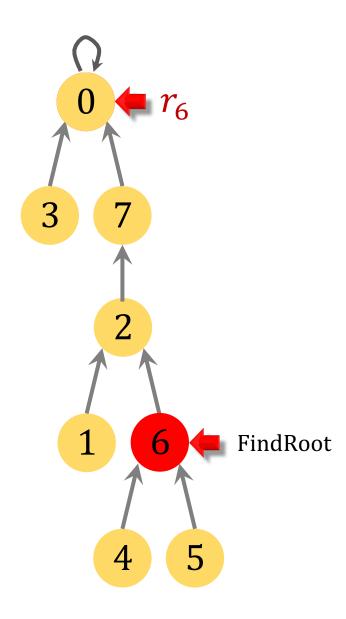
 After executing finding root, we know the root of an element:

$$r_i = FindRoot(i)$$
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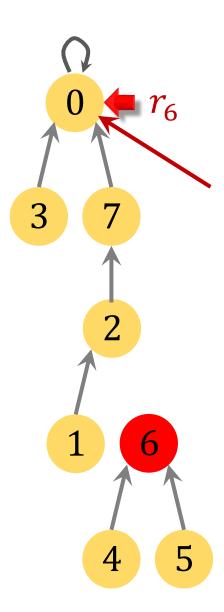
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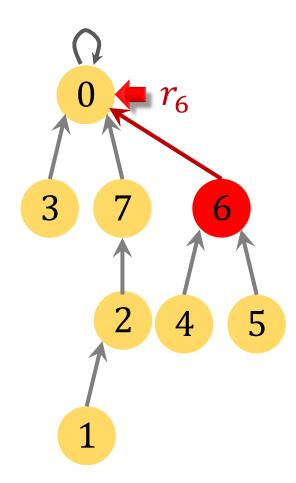
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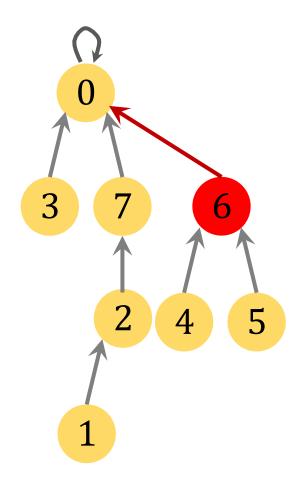
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 After executing finding root, we know the root of an element:

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.

- Let  $r_i$  be the parent of index i.
- Path compression reduces the depth without extra cost.



# **Summary**

#### Disjoint Sets Data Structure

- **Motivation:** Tracking a set of elements partitioned into disjoint (non-overlapping) subsets.
- Operation 1: Union (merge) two sets.
- Operation 2: Given two elements, tell whether they belong to the same set.
- Time complexity: Depths of the trees.
- Path compression: Reduce the depth of trees.

# Questions

#### Q1: Plot the disjoint sets as trees

Element: 0 1 2 3 4 5 6 7

Parent: | -1 | 2 | -1 | 0

### Q1: Plot the disjoint sets as trees

**Element:** Parent:

### Q1: Plot the disjoint sets as trees

Element: 0 1 2 3 4 5 6 7 8 9

Parent: 2 -1 7 -1 0 1 3 -1 1



#### Q2: Find roots

Element: 0 1 2 3 4 5 6 7 8 9

Parent: 2 -1 7 -1 0 1 3 -1 1

- 1. Find the root of element 0.
- 2. Find the root of element 4.
- 3. Find the root of element 8.

#### Q3: Are elements in the same set?

Element: 0 1 2 3 4 5 6 7 8 9

Parent: 2 -1 7 -1 0 1 3 -1 1

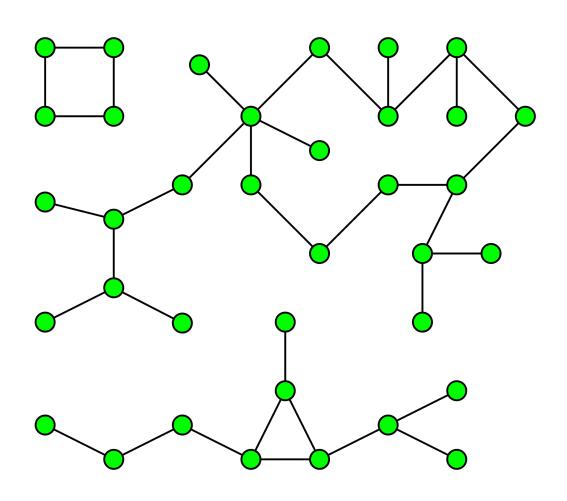
- 1. Are elements 0 and 1 in the same set?
- 2. Are elements 0 and 2 in the same set?
- 3. Are elements 5 and 6 in the same set?
- 4. Are elements 8 and 9 in the same set?

## Q4: How many disjoint sets?

Element: 0 1 2 3 4 5 6 7 8 9

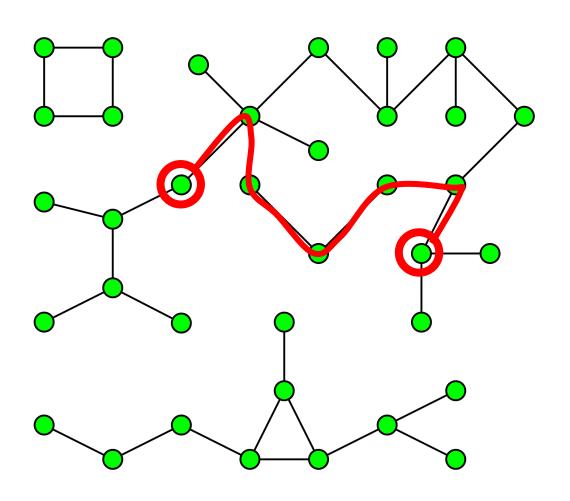
Parent: 2 -1 7 -1 0 1 3 -1 1 1

### Q5: Graph Connectivity



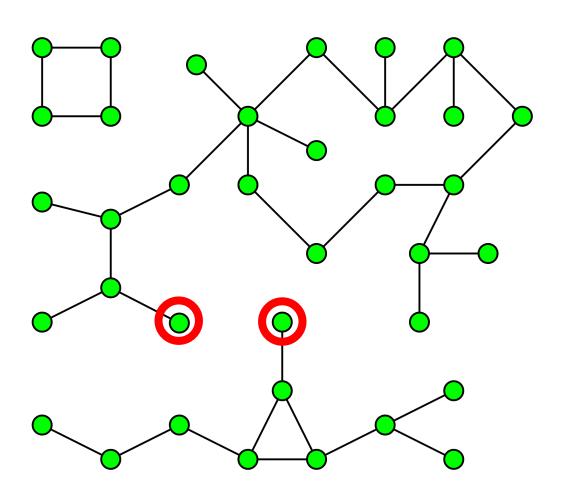
- A graph is said to be connected if there is a path between every pair of vertex.
- Given an undirected graph, decide whether it is connected.

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## Thank You!