

ARE YOU READY?



Branch & Bound Simplex Algorithm

Pure Integer Programming models



Formulating a linear programming model

The Oakfield Corporation manufactures tables and chairs. A table requires 1 hour of labour and 9 square board metres of wood, and a chair requires 1 hour of labour and 5 square board metres of wood. Currently, 6 hours of labour and 45 square board metres of wood are available. Each table contributes R8 to profit, and each chair contributes R5 to profit.

- 1. Formulate an Integer Programming Model to maximise Oakfield's profit.
- 2. Solve the formulated Integer Programming Model using the Branch & Bound Simplex Algorithm.



Formulating a linear programming model

- First, we need to declare our decision variables. $x_i = The \ number \ of \ i \ manufactured. \ i = 1 = Tables, 2 = Chairs$
- Next, we will create the objective function. Each table contributes R8 to profit, and each chair contributes R5 to profit. $\max z = 8x_1 + 5x_2$
- Next, we will create the constraints. We will start with the labour constraint: A table requires 1 hour of labour, a chair requires 1 hour of labour. Currently, 6 hours of labour are available.

$$s.t.x_1 + x_2 \le 6$$



Formulating a linear programming model

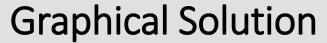
• For the next constraint we will look at the carpentry hours constraint: a table requires 9 square board metres of wood, a chair requires 5 square board metres of wood. Currently, 45 square board metres of wood are available.

$$s.t.9x_1 + 5x_2 \le 45$$

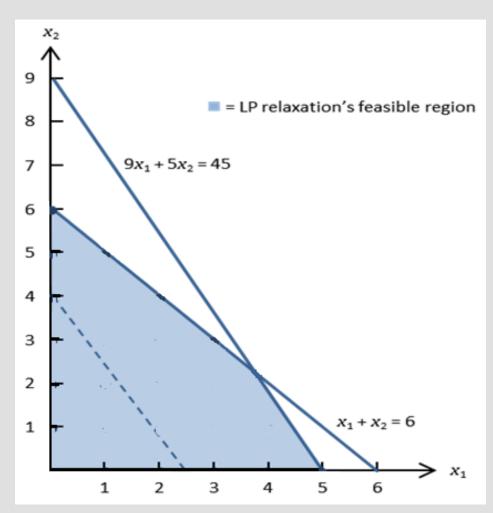
• Lastly, we need our sign restrictions:

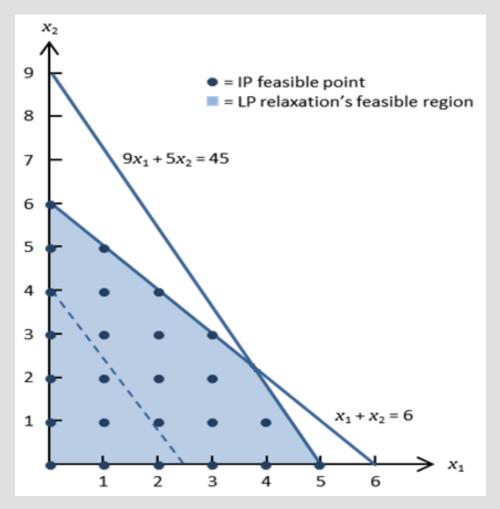
$$x_1, x_2 \ge 0$$

 $x_1, x_2 integers$



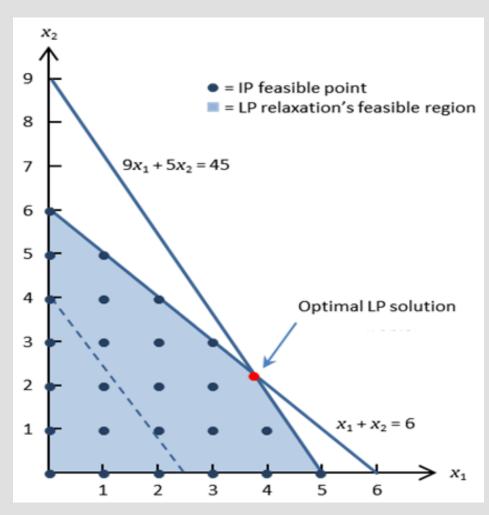


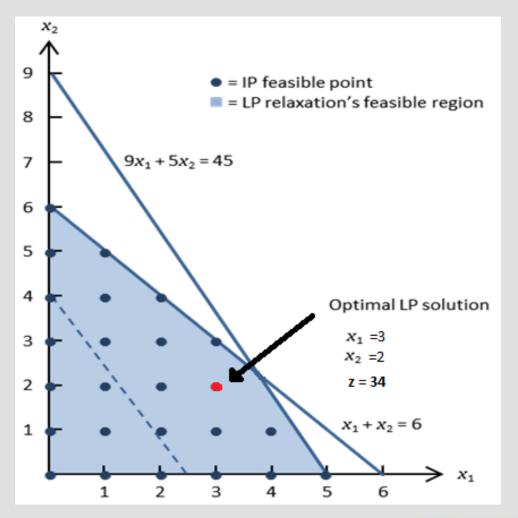










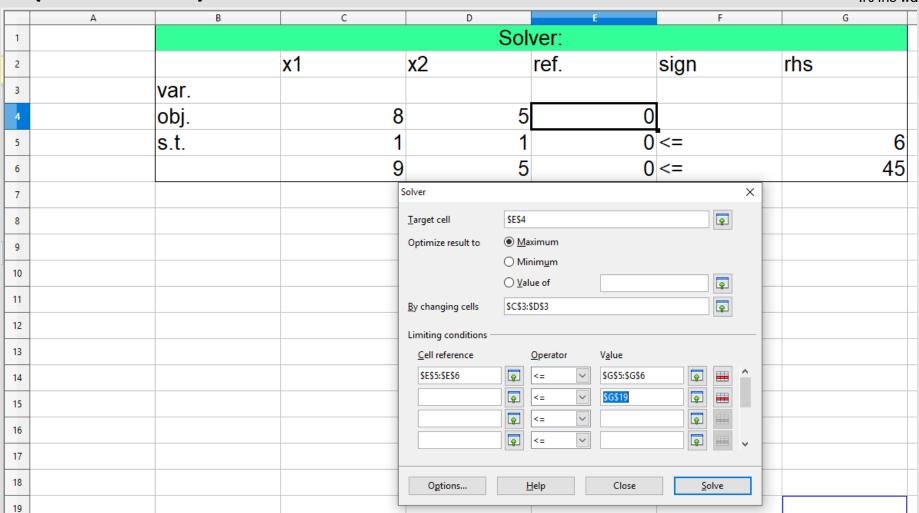






	Α	В	С	D	E	F	G	Н		
1			Solver:							
2			x1	x2	ref.	sign	rhs			
3		var.								
4		obj.	8	5	=sumproduct	(\$C\$3:\$D\$3;0	C4:D4)			
5		s.t.	1	1	1 R x 2 C	<=	6			
6			9	5		<=	45			
7										







_	Solver:								
	x1	x2	ref.	sign	rhs				
var.	3.75	2.25							
obj.	3	5	41.25						
var. obj. s.t.	1	1	6	<=	6				
	9	5	45	<= <=	45				



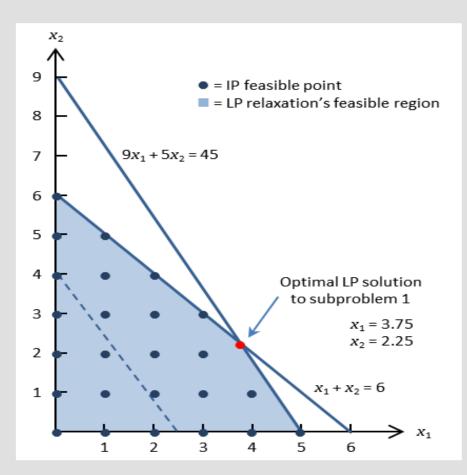
	A	В	С	D		E	F	G
1					Solv	er:		
2			x1	x2	r	ef.	sign	rhs
3		var.	3.75		2.25			
4		obj.	8		5	41.25		
5		s.t.	1		1	6	<=	6
6			9		5	45	<=	45
7			Solver				×	
8			<u>T</u> arget ce	ell \$C	\$3		•	
9			Optimize		<u>M</u> aximum			
10					Minim <u>u</u> m <u>V</u> alue of			
11			<u>B</u> y chang	_	\$3:\$D\$3			
12				conditions				
13				ference	<u>O</u> perator	V <u>a</u> lue		
14			\$E\$5:5			SG\$5:\$G\$6		
15			\$C\$3:					
16							-	
17					<u> </u>			
18			Ogt	ions	<u>H</u> elp	Close	Solve	
19								

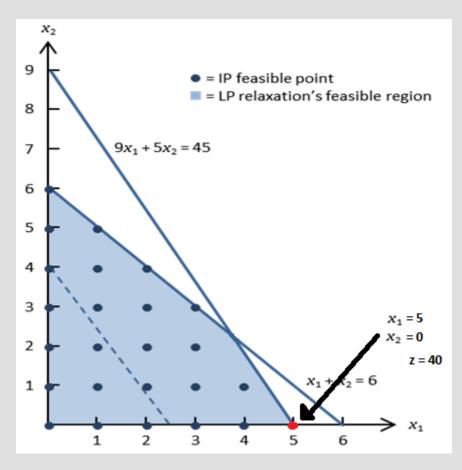


	Solver:										
	x1	x2	ref.		sign	rhs					
var.		5	0								
obj. s.t.		8	5	40							
s.t.		1	1		<=	6					
		9	5	45	<=	45					













	Integer Programming Model: Relaxed									
Max z =	8x1	+	5x 2							
s.t	X 1	+	X 2	≤	6					
	9x1	+	5 x 2	≤	45					
	X 1, X 2	≥		0						



Canonical Form

	Canonical Form: Branch & Bound Algorithm										
(z)	-	8x1	-	5x2	=		0				
	X1	+	X 2	+	S1	=	6				
	9x1	+	5x2	+	S2	=	45				





It's the way we're with

T-i	X 1	X 2	S1	S2	rhs	θ
Z	-8	-5	0	0	0	
1	1	1	1	0	6	6
2	9	5	0	1	45	5

	T-i	X1	X2	S1	S2	rhs	θ
	Z	-8	-5	0	0	0	
	1	1	1	1	0	6	6
	2	9	5	0	1	45	5
ľ							

T-2	X1	X2	S1	S2	rhs	θ
Z	0	- 5/9	0	8/9	40	
1	0	4/9	1	- 1/9	1	2 1/4
2	1	5/9	0	1/9	5	9

T-3*	X1	X2	S1	S2	rhs
Z	0	0	1 1/4	3/4	41 1/4
1	0	1	2 1/4	- 1/4	2 1/4
2	1	0	-1 1/4	1/4	3 3/4

Sub-Problem 1: x1 ≤ 3 Sub-Problem 2: x1 ≥ 4 Initial

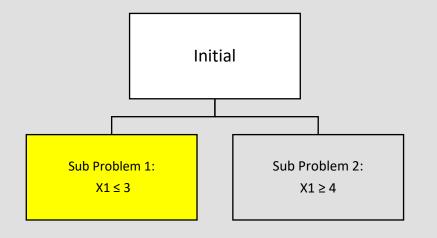
choose the one closest to 0.5

1/4 3/4

If both are the same distance, choose the lower subscript.



Branch & Bound Simplex Algorithm Sub-Problem 1



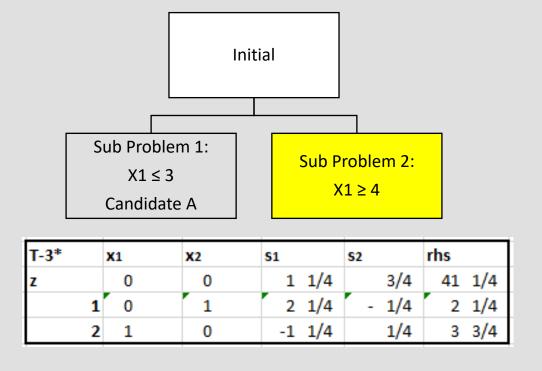
T-3*	X1	X2	S1	S2	rhs
Z	0	0	1 1/4	3/4	41 1/4
1	0	1	2 1/4	- 1/4	2 1/4
2	1	0	-1 1/4	1/4	3 3/4



	_	_	_	_	_	
T-3	X1	x2	S1	52	53	rhs
2	1	0	-1 1/4	1/4	0	3 3/4
3	1	0	0	0	1	3
2-3	0	0	-1 1/4	1/4	-1	3/4
3 x -1	0	0	1 1/4	- 1/4	1	- 3/4
T-3	x 1	x2	s 1	s2	53	rhs
Z	0	0	1 1/4	3/4	0	41 1/4
1	0	1	2 1/4	- 1/4	0	2 1/4
2	1	0	-1 1/4	1/4	0	3 3/4
3	0	0	1 1/4	- 1/4	1	- 3/4
θ				3		
T-4	x1	x2	s 1	52	53	rhs
Z	0	0	5	0	3	39
1	0	1	1	0	-1	3
2	1	0	0	0	1	3
3	0	0	-5	1	-4	3

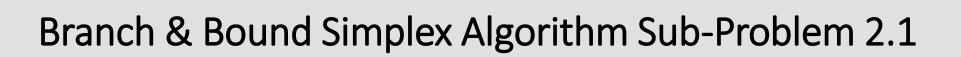


Branch & Bound Simplex Algorithm Sub-Problem 2

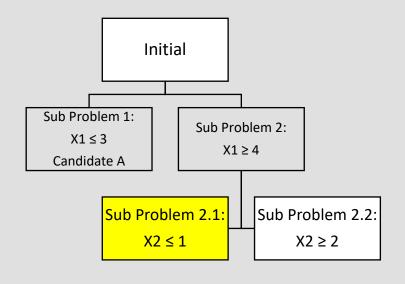


Sub-Problem 2.1: $x2 \le 1$ Sub-Problem 2.2: $x2 \ge 2$

T-3		x1		x2		51		52		е3		rhs	
	2		1		0		1/4		1/4		0	3	3/4
	3		1		0		0		0		-1		4
2-3			0		0	-1	1/4		1/4		1	-	1/4
T-3		x1		x2		51		52		е3		rhs	
Z		0		0		1	1/4		3/4	0		41	1/4
	1	0		1		2	1/4	-	1/4	0		2	1/4
	2	1		0		-1	1/4		1/4	0		3	3/4
	3		0		0	-1	1/4		1/4		1	-	1/4
θ							1						
T-4		x1		x2		51		52		e 3		rhs	
Z		0		0		0		1		1		41	
	1	0		1		0			1/5	1	4/5	1	4/5
	2	1		0		0		0		-1		4	
	3	0		0		1		_	1/5	-	4/5		1/5







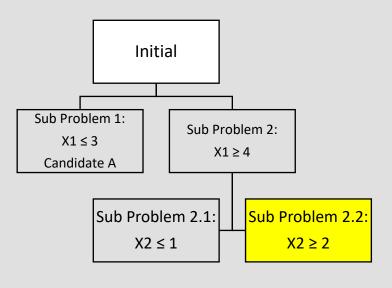
T-4	X1	X2	S1	52	e3	rhs
Z	0	0	0	1	1	41
1	0	1	0	1/5	1 4/5	1 4/5
2	1	0	0	0	-1	4
3	0	0	1	- 1/5	- 4/5	1/5

Sub-Problem 2.1.1: $x1 \le 4$ Sub-Problem 2.1.2: $x1 \ge 5$

T-4		x1	x2		51		52		e3		54		rhs	
	1	0	1		0			1/5	1	4/5	0		1	4/5
	4	0			1	0		0	4	0		1		7
1-4	i	0	0		0			1/5	1		-1	_		4/5
4 x -1		0	0		0		-	1/5		4/5	1			4/5
		•	-		-			1/3	_	7/3	_			7/3
T-4		X1	X2		51		52		e3		54		rhs	
Z		0		()	0		1		1		0		41
	1	0		1	1	0		1/5	1	4/5		0	1	4/5
	2	1		()	0		0		-1		0		4
	3	0		()	1	-	1/5	-	4/5		0		1/5
	4	0	0	1	0		-	1/5	-1	4/5	1			4/5
θ								5		5/9				
T-5		x1	x2		51		52		e 3		54		rhs	
z		0	0)	0			8/9	0			5/9	40	5/9
	1	0	1		0		0		0		1		1	
	2	1	0		0			1/9	0		- !	5/9	4	4/9
	3	0	0		1		-	1/9	0			4/9		5/9
	4	0	0		0			1/9	1			5/9		4/9







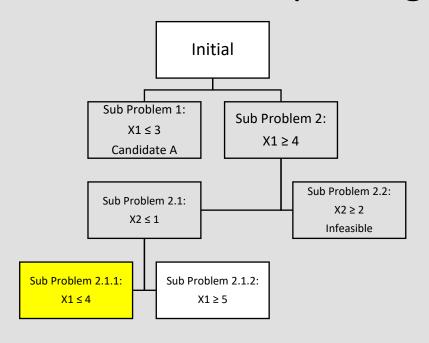
T-4	Γ-4 x1		S1	52	e3	rhs
Z	0	0	0	1	1	41
1	0	1	0	1/5	1 4/5	1 4/5
2	1	0	0	0	-1	4
3	0	0	1	- 1/5	- 4/5	1/5

T-4		x 1		x2	51		52	е3		e4		rhs	
	1	0		1		0	1/5	1	4/5	0		1	4/5
	4		0		1	0	0		0		-1		2
1-4		0		0		0	1/5	1	4/5	1			1/5
T-4		x 1		x2	51		52	е3		e4		rhs	
Z			0		0	0	1		1		0		41
	1		0		1	0	1/5	1	4/5		0	1	4/5
	2		1		0	0	0		-1		0		4
	3		1 0		0 0	0 1	- 1/5		-1 4/5		0		4 1/5









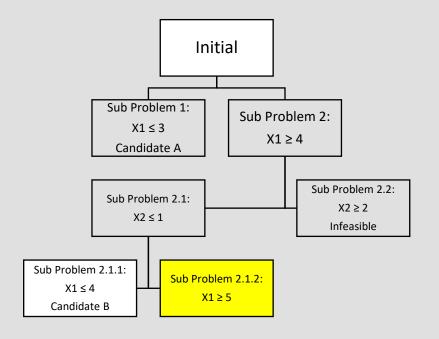
T-5	x1	x2 s1		52	e3	54	rhs
Z	0	0	0	8/9	0	5/9	40 5/9
1	0	1	0	0	0	1	1
2	1	0	0	1/9	0	- 5/9	4 4/9
3	0	0	1	- 1/9	0	- 4/9	5/9
4	0	0	0	1/9	1	- 5/9	4/9

	_	_		_	_		_	
T-5	X1	X2	51		e3		5 5	rhs
2		0	0	1/9	0	- 5/9	0	4 4/9
5	1	0	0	0	0	0	1	4
2-5	0	0	0	1/9	0	- 5/9	-1	4/9
5 x -1	0	0	0	- 1/9	0	5/9	1	- 4/9
T-5	x1	x2	s 1	s2	e 3	s4	s 5	rhs
Z	0	0	0	8/9	0	5/9	0	40 5/9
1	0	1	0	0	0	1	0	1
2	1	0	0	1/9	0	- 5/9	0	4 4/9
3	0	0	1	- 1/9	0	- 4/9	0	5/9
4	0	0	0	1/9	1	- 5/9	0	4/9
5	0	0	0	- 1/9	0	5/9	1	- 4/9
θ				8				
T-6	x1	x2	51	52	e3	54	s 5	rhs
Z	0	0	0	0	0	5	8	37
1	. 0	1	0	0	0	1	0	1
2	1	0	0	0	0	0	1	4
3	0	0	1	0	0	-1	-1	1
4	0	0	0	0	1	0	1	0
5	0	0	0	1	0	-5	-9	4

Cadindate B







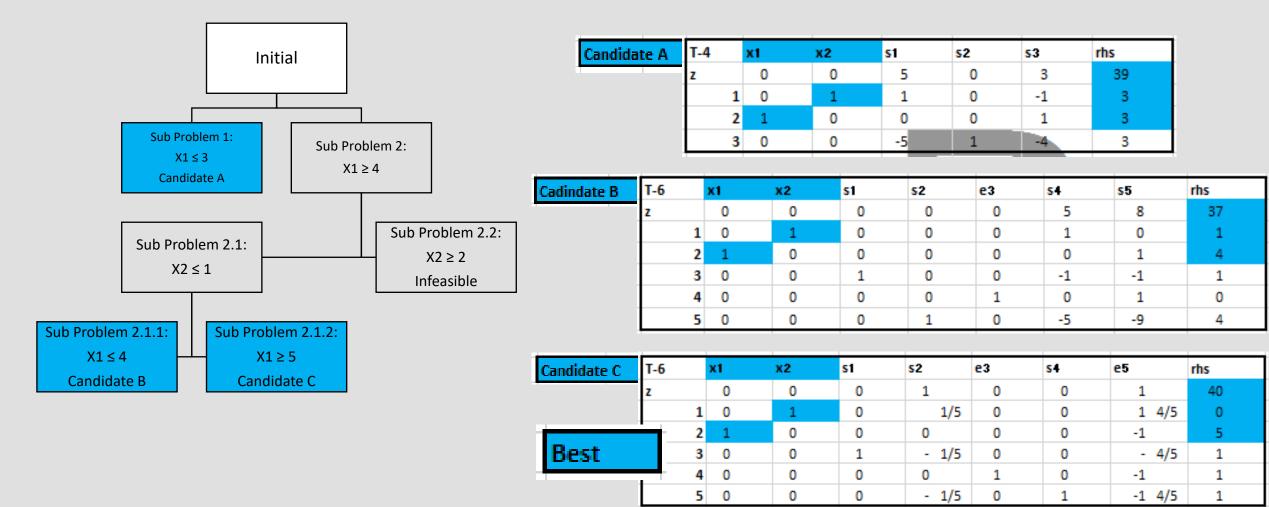
T-5	X1	x2 s1		52	e3	s4	rhs
z	0	0	0	8/9	0	5/9	40 5/9
1	0	1	0	0	0	1	1
2	1	0	0	1/9	0	- 5/9	4 4/9
3	0	0	1	- 1/9	0	- 4/9	5/9
4	0	0	0	1/9	1	- 5/9	4/9

T-5		x1	x2		s1	52		e3		54		e5		rhs	
	2	1	0		0	1	/9	0		-	5/9	0		4	4/9
	5	1		0	C)	0		0		0		-1		5
2-5		0	0		0	1	/9	0		-	5/9	1		-	5/9
T-5		X1	x2		s1	52		e3		54		e5		rhs	
z		0	0		0	8	/9	0			5/9	0		40	5/9
	1	0	1		0	0		0		1		0		1	
	2	1	0		0	1	/9	0		-	5/9	0		4	4/9
	3	0	0		1	- 1	/9	0		-	4/9	0			5/9
	4	0	0		0	1	/9	1		-	5/9	0			4/9
	5	0	0		0	1	/9	0		-	5/9	1		-	5/9
θ											1				
T-6		x1	х2		S1	52		e 3		54		e5		rhs	
z		0	0		0	1		0		0		1		40	
	1	0	1		0	1	/5	0		0		1	4/5	0	
	2	1	0		0	0		0		0		-1		5	
	3	0	0		1	- 1	/5	0		0		-	4/5	1	
	4	0	0		0	0		1		0		-1		1	
	5	0	0		0	- 1	/5	0		1		-1	4/5	1	

Candidate C



Branch & Bound Simplex Algorithm Candidate Selection







Found in the document called 'Branch & Bound Simplex Algorithm Exercises'





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