



Computer Systems B

COMS20012

Introduction to Operating Systems and Security

bristol.ac.uk

The process abstraction

bristol.ac.uk



Operating System Abstractions

- Abstractions **simplify applications design** by:
 - Hiding undesirable properties;
 - Adding new capabilities;
 - Organizing information.
- Abstractions provide an **interface** to programmers that separates **policy** – what the interface commits to accomplish – from the **mechanism** – how the interface is implemented.

bristol.ac.uk

Abstraction example: File

- What **undesirable properties** file systems hide?
 - Disk are slow!
 - Chunk of storage are distributed all over the disk.
 - Disk storage may fail.

bristol.ac.uk

Abstraction example: File

- What **undesirable properties** file systems hide?
 - Disk are slow!
 - Chunk of storage are distributed all over the disk.
 - Disk storage may fail.
- What new **capabilities** do files add?
 - Growth and shrinking.
 - Organization into directories.

bristol.ac.uk

Abstraction example: File

- What **undesirable properties** file systems hide?
 - Disk are slow!
 - Chunk of storage are distributed all over the disk.
 - Disk storage may fail.
- What new **capabilities** do files add?
 - Growth and shrinking.
 - Organization into directories.
- What **information** files help to organize?
 - Ownership and permission.
 - Access time, modification time, type etc.

bristol.ac.uk

Abstractions to come in this unit

- Threads
 - Abstract the CPU
- Address space
 - Abstract the memory
- Files
 - Abstract the disk

bristol.ac.uk

Abstractions to come in this unit

- Threads
 - Abstract the CPU
- Address space
 - Abstract the memory
- Files
 - Abstract the disk

Operating Systems are all about abstractions!

bristol.ac.uk

The process abstraction

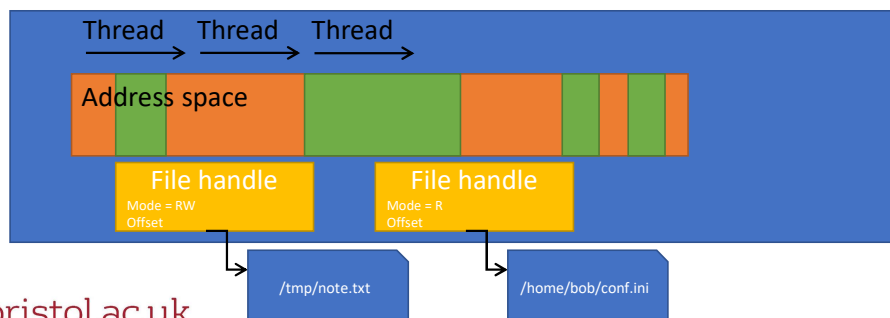
Processes are the most fundamental abstraction

- What the computer “is doing”.
- Help organize other abstractions.
- You know processes as “applications”.

bristol.ac.uk

The process abstraction

- Processes are **not tied to a hardware component**.
- They contain and organize other abstractions.



bristol.ac.uk

Processes vs Threads

- Potentially confusing due to terminology
 - both described as **running**
- Some terminology useful to remember the distinction
 - Processes require multiple resources: CPU, memory, files
 - Threads abstract the CPU
- A process contains threads, threads belong to a process
 - Except kernel threads who do not belong to a user space process
- A process is running when one or more of its threads are running
- Terminology may vary between OSes but concepts are the same

bristol.ac.uk

Process Example: Firefox

- Firefox has multiple threads. What do they do?
 - Waiting and processing interface events (e.g., mouse click)
 - Redrawing the screen as necessary (responding to user inputs)
 - Loading web pages (generally multiple elements in parallel)
- Firefox is using memory. For what?
 - The executable code itself
 - Shared library: web page parsing, TLS/SSL etc.
 - Stacks storing local variables for running threads
 - A heap storing dynamically allocated memory
- Firefox has files open. Why?
 - Fonts
 - Configuration files

bristol.ac.uk

Process as a protection boundary

- OS is responsible for **isolating processes from each others**.
 - What happened in a process **should not affect other processes**
 - ... or **crash the machine**.
- **Intra-process** communication (between threads) is **application responsibility**
 - **Shared address space**.
 - **Shared file descriptors**.
 - Use synchronization primitives to ensure consistency.

bristol.ac.uk

Process as a protection boundary

- OS is responsible for **isolating processes from each others**.
 - What happened in a process **should not affect other processes**
 - ... or **crash the machine**.
- **Intra-process** communication (between threads) is **application responsibility**
 - **Shared address space**.
 - **Shared file descriptors**.
 - Use synchronization primitives to ensure consistency.
- See Computer Systems A

bristol.ac.uk

Inter-process communication

- Allow processes to interact with each others
- A variety exists:
 - Shared files
 - Sockets
 - Signal
 - Pipes
 - Shared memory
 - etc.
- Well defined **semantics enforced by the OS**
 - Limit how process can interfere with each other
 - e.g., you cannot SIGKILL any other process on the machine

bristol.ac.uk

Process lifecycle

- `fork()`
 - Create a new process
- `exec()`
 - Replace the current process code by an executable
- `exit()`
 - Terminate the process
- `waitpid()`
 - Wait for another process to terminate

bristol.ac.uk

Process lifecycle

- `fork()`
 - Create a new process
- `exec()`
 - Replace the current process code by an executable
- `exit()`
 - Terminate the process
- `waitpid()`
 - Wait for another process to terminate

To be implemented in Lab 7

bristol.ac.uk



Thank you

bristol.ac.uk

