Partern marching:

- -> Brute force approach of this is to compare each character of pattern to text and if we observe the pattern we'll return the index
- -> Le con use rabin karp algorithm.

Rabin Karp algorithmic

- -> he take a section in the text which is of parter length of partern
- -> Compare if hash code of the section is equal to hash code of pattern.
 - If they are exped
 then only we compare the string
 section and pottern.

 If they are equal then return inder.

Important thing is -

- -> To shook a good hash function i.e.,
 - --- faut to compute
 - Less Collisions Falle positives

Normal hash function in

- -> Le can add "arcii" values of all the characters and return the Sum as the hash code.
- -> In this we'll have so many collisions

Optimized hash function i-

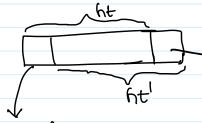
d= no of characters in the text m = length of pattern.

S = Pubstring in the middle of text.

How to compute hash function optimally i

-> If we go on for each string character and compute-the hash, it results in same time complemity as that of brute-force which is of no use.

— > We can we nolling method to compute host in O(i) time i.e., let, lit be the previous hash and lit' be the current hash



first char of ht

curr_Char we are processing

$$= \int ht' = (ht - first charofft \times d^{m-1}) \times d^{m-1}$$

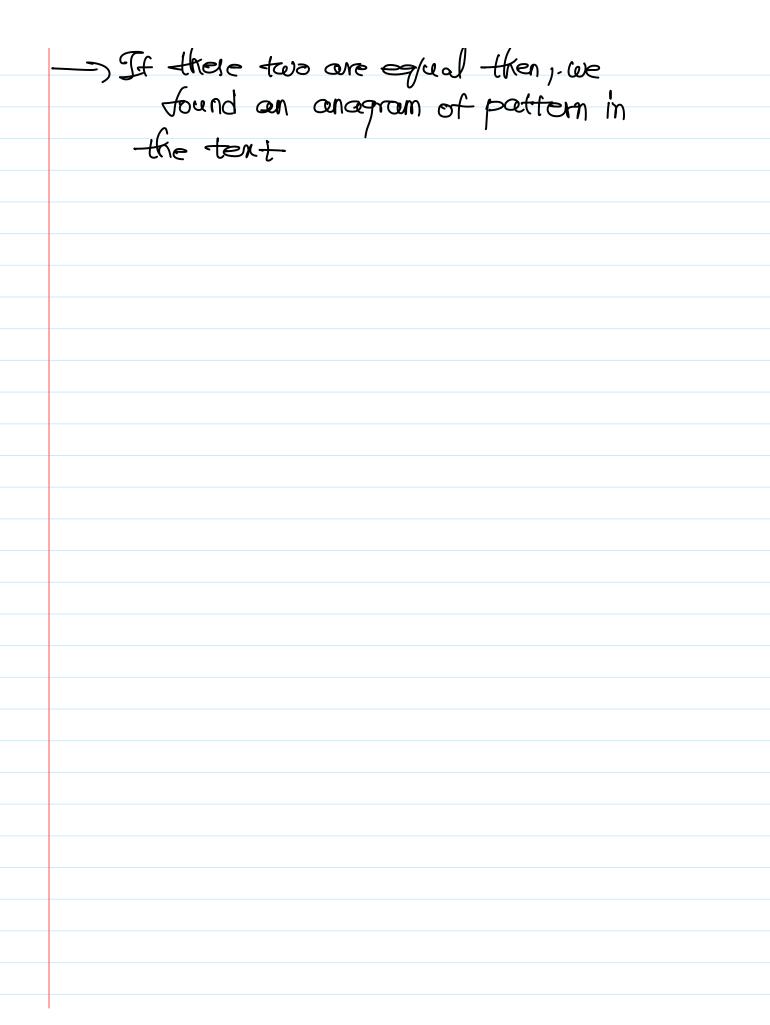
$$+ curr-char$$

Prog	ram 1
Tuesday,	19 January 2021 4:30 PM
5	
Pro	blem Statement:
	re are pair of words namely W1 and W2 with a limited of word range,
	ate a method to return a true value if W2 contains the anagram of W1.
In a	dditional, one of the anagram of first word is the substring of the second word.
You	r task is to implement the Solution class, and implement a method in it,
	- public boolean checkPalindromeSubstring(String w1, String w2){}.
Inp	ut Format:
	
Two	space separated words w1 and w2, consist of lowercase letters only.
Out	put Format:
	to be clean value if M/2 contains the apparam of M/1 or not
Prir	t a boolean value, if W2 contains the anagram of W1 or not.
San 	nple Input-1:
abb	cbb abbbabbcb
San	nple Output-1:
true	
_	
San	nple Input-2:
abb	cbbc abbbabbcb
1	pple Output-2:
fals	 e
Cod	e·

```
import java.util.*;
class Solution
  public int computeHash(int ht,int currchar, int firstcharofht)
    return ht-firstcharofht+currchar;
  public boolean checkSubstring(String w1,String w2)
    int alphatext[]=new int[26];
    int alphapat[]=new int[26];
    Arrays.fill(alphatext,0);
    Arrays.fill(alphapat,0);
    int lt=w2.length();
    int lp=w1.length();
    if(|t>=|p|)
       int hp=0;
       for(int i=0;i<lp;i++)
         hp+=((int)w1.charAt(i))-96;
       int ht=0;
       for(int i=0;i<lp-1;i++)
         ht+=((int)w2.charAt(i))-96;
       // System.out.println(ht);
       int firstcharofht=0;
       int currchar=0;
       int flag=0;
       int i=lp-1;
       while(i<lt)
         flag=0;
         currchar=((int)w2.charAt(i))-96;
         // System.out.println(currchar);
         // System.out.println(firstcharofht);
         ht=computeHash(ht,currchar,firstcharofht);
         // System.out.println(ht);
         if(ht==hp)
         {
           for(int j=i-lp+1;j <=i;j++)
              alphatext[(int)w2.charAt(j)-97]++;
              alphapat[(int)w1.charAt(j-i+lp-1)-97]++;
           for(int j=0;j<26;j++)
              if(alphatext[j]!=alphapat[j])
```

```
flag=1;
                break;
           }
           if(flag==0){
             return true;
           }
         }
         firstcharofht=((int)w2.charAt(i-lp+1))-96;
         Arrays.fill(alphatext,0);
         Arrays.fill(alphapat,0);
         i++;
      }
      return false;
    else
      return false;
  }
}
  Test Cases:
  case =1
  input =dinitrophenylhydrazine acetylphenylhydrazine
  output =/false/
  case =2
  input =abbcbb abbbabbcb
  output =/true/
  case =3
  input =abbcbbc abbbabbcb
  output =/false/
  case =4
  input =listentomotherinlaw wehavesilenthitlerwomantosociety
  output =/true/
  case =5
  input = mother in law keepsilent catde bit cards chool master as tronomers\\
  hit lerwoman stars peek listen act the class room credit bad no more\\
  output =/true/
  input = mother in law keepsilen tand debit card has bad credit in dirty room
  peek hit ler woman listen the class room credit bad no more\\
  output =/false/
```

case =7 input =motherinlawkeepsilentandbadcreditindirtyroom wehaveanhitlerwomanpeeklistendebitcardindormitory output =/false/ case =8 input =motherinlawkeepsilentcatdebitcardschoolmasterastronomershitlerwomanstarspeeklistenacttheclassro omcreditbadnomore there is a king lived in kosalamother in law keep silent cat debit card ten act the class room credit bad no more school and the class room credit bad noImasterastronomershitlerwomanstarspeekliswithus output =/true/ case =9 input =abcdefghijk cbcboooaaah output =/false/ case =10 input =abcd cbcb output =/false/ Logic i In this we use the normal hash function because we have to find anagrams of the pattern. > If the hour Gode is some i.e., -there is possibility that we found an anagram. So, frequency array of lubithing is alphatent frequency array of pattern is alpha pat. -) If there two one equal then, we



Program 2

Tuesday, January 19, 2021 9:02 PM

Problem Statement: KMIT hosting a Keshav Memorial Badminton League. They planned to conduct N number of games. Each game begin and ends in perticular time slot. You are given an array of time slots of the N games, consisting of begin and end times (b1,e1),(b2,e2),... (b < e). Your task is to determine minimum number of badminton courts required to conduct all the games smoothly. NOTE: If a game begins at time 'a' ends at time 'b', another game can start at 'b'. Input Format: Line-1: An integer N, number of games. Next N lines: Two space separated integers, begin and end time of each game. Output Format: Print an integer, minimum number of badminton courts required. Sample Input-1: 030 5 10 15 20 Sample Output-1: 2 Sample Input-2: 3 0 10 15 25 25 35 Sample Output-2: 1 Code: import java.util.*; class Badminton public static int partition(int arr[], int start, int end) int pindex=start; int pivot=arr[end];

```
int temp=0;
  for(int i=start;i<end;i++)</pre>
    if(arr[i]<=pivot)
    {
      temp=arr[i];
       arr[i]=arr[pindex];
       arr[pindex]=temp;
       pindex++;
  }
  temp=arr[end];
  arr[end]=arr[pindex];
  arr[pindex]=temp;
  return pindex;
}
public static void quickSort(int arr[],int start, int end)
{
  if(start<end)
  {
    int pindex=partition(arr,start,end);
    quickSort(arr,start,pindex-1);
    quickSort(arr,pindex+1,end);
  else
    return;
}
public static void main(String args[])
  Scanner sc=new Scanner(System.in);
  int n=sc.nextInt();
  int start[]=new int[n];
  int end[]=new int[n];
  for(int i=0;i<n;i++)
  {
     start[i]=sc.nextInt();
     end[i]=sc.nextInt();
  quickSort(start,0,n-1);
  quickSort(end,0,n-1);
  int minCourts=0;
  int maxEndTimeInSpecificCourt=0;
  for(int i=0;i<n;i++)
    if (start[i] < end [maxEndTimeInSpecificCourt]) \\
       minCourts++;
    }
    else
       maxEndTimeInSpecificCourt++;
    }
  System.out.println(minCourts);
}
```

Test Cases:
case =1
input =3
0 30
5 10 15 20
output =2
case =2
input =3
0 10 15 25
25 35
output =1
case =3
input =10 1 10
15 25
30 40 45 60
11 15
61 70
41 50 75 90
80 95
91 100 output =2
case =4 input =10
1 15
20 35 30 45
35 50
25 40 10 25
60 75
45 60 40 55
50 65
output =3
case =5
input =15
1 25 10 20
10 35
15 30 25 40
30 50
25 50 40 75
35 60
20 40 40 60
35 50
20 45 25 60
50 75
output =8

```
case =6
input =20
1 25
10 20
10 35
15 30
45 60
25 40
35 55
25 50
50 90
55 75
50 80
40 75
35 60
20 40
40 60
70 90
35 50
20 45
25 60
50 75
output =9
case =7
input =10
10 40
40 70
50 80
70 100
100 130
130 150
65 95
55 85
45 75
35 65
output =5
case =8
input =15
1 15
20 35
30 45
35 50
25 40
10 25
60 75
45 60
40 55
50 65
15 35
35 60
30 50
45 70
60 90
output =6
```

Logici

-> First we bort both start and end arrays independently. -> Now, we loop i in the Hart and initialize j=0; court =0; i is the manumumpossible end time of current court -> b, if (Start ?i] < end[i]) = another moter starts before the current court metch ends. So, we need another court to play that match. _____ court ++; > else i.e., stort [i] > and [i]. - another match starts after the current match some court: So, we increment j because the match is completed and we have to move further 19

```
6
       31
    7
       29
       15
       ΤÇ
Stort: 2 4 5 7 9
                        courts = 0
         15 19 29 31
end:
       Ţ
       U start [i] < end ?i]
          í,t
       2 4
                        Courts = 1
end: 6 15 19 29 31
       1
        | Start [i] < end [i]
 Start: 2 4 5 7 9
 end: 6 15 19 29 31
        | Hort[i] < end[j]
  Start; 2 4 5 7 9
                          Court =
  end: 6 15 19 29 31
             Stoort (i) > end (i)
     Stoent; 2 4 5 7
                                 Court = 3
     end : 6 15 19 29 31
```

end: 6 15 19 29 31

Stantfil < and 5iHaut: 2 4 5 7 9

end: 6 15 19 29 31

A

1

 \rightarrow Court = 4

Program 3

Tuesday, January 19, 2021 9:04 PM

Problem Statement: Mr. James professor of at Illinois state university, as a part of assignment he created a problem statement related to strings. He gave a String S, and asked them to design a method to return the longest substring in S, which is a palindrome. NOTE: Alphabets are case sensitive "Aa" is not considered a palindrome here. Input Format: A string S, consist of lowercase/uppercase letters or/and digits **Output Format:** Print a string, longest palindrome substring. Sample Input-1: abbbabbcbbacdb Sample Output-1: abbcbba Sample Input-2: thedivideriswide Sample Output-2: ----edivide Code: import java.util.*; class LongestPalindromicSubString public static void main(String args[])

```
Scanner sc=new Scanner(System.in);
    String s=sc.next();
    int n=s.length();
    boolean dp[][]=new boolean[n][n];
    int maxi=0;
    int maxj=0;
    dp[n-1][n-1]=true;
    for(int i=0;i<n-1;i++)
      dp[i][i]=true;
      if(s.charAt(i)==s.charAt(i+1))
         dp[i][i+1]=true;
         maxi=i;
         maxj=i+1;
      }
      else
         dp[i][i+1]=false;
    }
    int row=0,col=0;
    for(int j=2;j<n;j++)
    {
      row=0;
      col=j;
      while(row<n && col<n)
         if(s.charAt(row)==s.charAt(col))
           dp[row][col]=dp[row+1][col-1];
           if(dp[row][col]==true)
             maxi=row;
             maxj=col;
          }
         }
         else
           dp[row][col]=false;
         }
         row++;
         col++;
    }
    System.out.println(s.substring(maxi,maxj+1));
 }
}
```

Test Cases:

case -1
input =oneofthemsaidwepanicinapewinabook
output =wepanicinapew
case =2
input =forgeeksskeegfor output =geeksskeeg
output –geeksskeeg
case =3
input =abbbabbcbbacdb
output =abbcbba
case =4
input =thedivideriswide
output =edivide
case =5
input =itisneveroddorevenwhenyougivenazeroasinput
output =neveroddoreven
case =6
input =yourmaidmadeademandthatsiridemandiamamaidnamedirispleaseconsider
output =siridemandiamamaidnamediris
7
case =7 input =theyhaveprintedaphrasewontloversrevoltnowinapaperheading
output =wontloversrevoltnow
case =8
input =wepanicaspewbutdontnodsoneveroddorevensiridemandiamamaidnameiris
output =emandiamamaidname
Coaro
ADOIL :
-> This can be done using DP.
the cost be dotte divide DE.
-> Bale cales are -
Dave cover one =
\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \
a) on the substained music remata is 1
1. 1 1.75:75:7 - +200.
are palinames. — ap sois is - cide;
6) Substrings with Jength 2 -
1) they are pollind nomes if both characters
a) all the substrings whose length is I are palindromes. b) substrings with length 2— i) they are palindromes if both characters are equal.
are ey/ual,
\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \

one equal.

-> Now 1 from length 3 lets ray in the given string this substring starks from i and ends at j (i.i. 60th inclusive).

Subrtning = $\times \times \times ---- A B$ $i \quad i \in \{1, \dots, j-1, j\}$

For this substring to be a palindrome.

a) x and B mut be equal.

(If they are not equal if we reverse it doesn't become palindrome)

6) Substring - y - ... A mult be a polindrome.

1.6.1

a) s. charAt(i) = s. charAt(i)

and dp [i+i] [i-i] = true.

Fgi

a66a

