

# Progressive Graph Matching: Making a Move of Graphs via Probabilistic Voting

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# Graph Matching

Given two attributed graphs  $\bar{G}^P = (\bar{V}^P, \bar{E}^P, \bar{A}^P)$  and  $\bar{G}^Q = (\bar{V}^Q, \bar{E}^Q, \bar{A}^Q)$  with  $n^P$  and  $n^Q$  nodes respectively.

A result of graph matching is a subset of possible correspondences between those graphs, which can be represented in form of assignment matrix  $X \in \{0, 1\}^{n^P \times n^Q}$ :

$$X_{ia} = \begin{cases} 1 & \text{node } v_i \in \bar{V}^P \text{ matches } v_a \in \bar{V}^Q \\ 0 & \text{otherwise} \end{cases}$$

General formulation:

$$\begin{aligned} x^* &= \arg \max S(x) \\ \text{s.t. } &\begin{cases} x \in 0, 1^{n^P n^Q} \\ \sum_{i=1}^{n^P} x_{ia} \leq 1 \\ \sum_{a=1}^{n^Q} x_{ia} \leq 1 \end{cases} \end{aligned}$$

The objective function  $S(x)$  measures the similarity between the graph attributes.

# Integer Quadratic Programming

Consider two types of similarity: node and edge similarity. Those can be combined in an affinity matrix  $W$ , where non-diagonal elements  $W_{ia,jb}$  represent edge similarity of the edges  $e_{ij} \in \bar{E}^P$  and  $e_{ab} \in \bar{E}^Q$  and diagonal elements  $W_{ia,ia}$  represent similarity between nodes  $v_i \in \bar{V}^P$  and  $v_a \in \bar{V}^Q$ .

That leads to the integer quadratic optimization problem:

$$x^* = \arg \max x^T W x$$

$$\text{s.t. } \begin{cases} x \in 0, 1^{n^P n^Q} \\ \sum_{i=1}^{n^P} x_{ia} \leq 1 \\ \sum_{a=1}^{n^Q} x_{ia} \leq 1 \end{cases}$$

# Issues

- Quadratic assignment problem is  $NP$ -complete
- Computation of  $W$  for a large graph is almost infeasible
- Reduction of the complexity often leads to worse matching results

# Progressive Graph Matching

Refer to the initial graphs  $\bar{G}^P$  and  $\bar{G}^Q$  as **maximal graphs**. For each maximal graph a subgraph induced by a reduced set of nodes is called **active graph** ( $G^P$  and  $G^Q$  respectively).

Idea: match maximal graphs by iteratively matching their active graphs ( $t = 0, 1, \dots$ )

- **Graph Matching**

match current active graphs  $G_t^P = (V_t^P, E_t^P, A_t^P)$  and  $G_t^Q = (V_t^Q, E_t^Q, A_t^Q)$ ,  $G_t^P \subset \bar{G}^P, G_t^Q \subset \bar{G}^Q$ .

- **Graph Progression**

update active graphs to improve matching score on the next matching step

# Graph Matching

To reduce the complexity of graph matching consider a given **set of candidate matches**  $C_t \subset V^P \times V^Q$ .

The corresponding active graphs are defined by the nodes appeared in  $C_t$ .

This method reduces the size of active graphs and makes a affinity matrix more sparse.

# Graph Progression I

Let  $M_t = \{m_1, \dots, m_{|M_t|}\} \subset C_t$  with  $m_i = (v_{p_i}^P, v_{q_i}^Q)$  be a **result of graph matching** of two active graphs and  $s_t$  it's **score**.

Consider conditional joint probability  $p(V^P, V^Q | M_t)$ :

$$\begin{aligned} p(V^P, V^Q | M_t) &= \sum_{m_i \in M_t} p(V^P, V^Q, M = m_i | M_t) \\ &= \sum_{m_i \in M_t} p(V^Q | V^P, M = m_i, M_t) \\ &\quad p(V^P | M = m_i, M_t) p(M = m_i | M_t) \end{aligned}$$

# Graph Progression II

$p(v_i^P, v_a^Q | M_t)$  conditional probability of match  $(v_i^P, v_a^Q)$  between two maximal graphs.

Base on the probability distribution  $p(V^P, V^Q | M_t)$  new candidate set  $C_{t+1}$  consists of  $N_c$  matches best matches.

Given new candidate set  $C_{t+1}$  we obtain new active graphs  $G_{t+1}^P$  and  $G_{t+1}^Q$ .

The condition  $M_t \subset C_{t+1}$  ensures that  $s_{t+1} \geq s_t$ , if  $M_t$  is an optimal matching.



# Algorithm

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## Algorithmus 1 progressiveGraphMatching( $\bar{G}^P, \bar{G}^Q, N_c$ ) [1]

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1:  $C_0 = \text{Find\_Initial\_Candidates}(\bar{G}^P, \bar{G}^Q, N_c)$ 
2:  $t = 0, s_0 = 0$ 
3: while score increases do
4:    $(M_t, s_t) = \text{Graph\_Matching}(C_t)$ 
5:    $p(v_p^P, v_q^Q | M_t) = 0 \quad \forall v_p^P \in \bar{V}^P, v_q^Q \in \bar{V}^Q$ 
6:   for each  $m_i = (v_p^P, v_q^Q) \in M_t$  do
7:      $N_A = \{v^P \in \bar{V}^P | p(v^P | m_i, M_t) > \epsilon\}$ 
8:     for each  $v_j^P \in N_A$  do
9:        $N_B = \{v^Q \in \bar{V}^Q | p(v^Q | v_j^P, m_i, M_t) > \epsilon\}$ 
10:      for each  $v_b^Q \in N_B$  do
11:        
$$p(v_p^P, v_q^Q | M_t) = p(v_p^P, v_q^Q | M_t) +$$


$$p(v_b^Q | v_j^P, m_i, M_t) p(v_j^P | m_i, M_t) p(m_i | M_t)$$

12:      end for
13:    end for
14:  end for
15:   $C_{t+1} = N_c$  best matches based on  $p(V^P, V^Q | M_t)$ , which contains  $M_t$ 
16:   $t = t + 1$ 
17: end while
18: return  $M_t$ 

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# Matching of images I

Given two images.

- Detected features of the images represent nodes of the maximal graphs (and their geometric relations as edges)
- Node similarity = similarity of feature descriptors of two graphs
- Edge similarity = Symmetric Transfer Error (STE) (see [2])

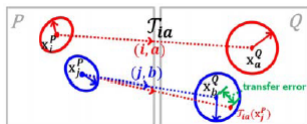
Derive affine homography transformation  $\tau_{ia}(\cdot)$  between two features

The transfer error of  $(j, b)$  with respect to  $(i, a)$

$$d_{jb|ia} = \|x_b^Q - \tau_{ia}(x_j^P)\|$$

Edge similarity  $W_{ia,jb} =$

$$\max(0, \alpha - \frac{d_{jb|ia} + d_{bj|ai} + d_{ia|jb} + d_{ai|jb}}{4})$$



# Matching of images II

- $p(M = m_i | M_t) = \text{score}(m_i) / \sum_i \text{score}(m_i)$
- Assuming  $m_i = (v_{p_i}^P, v_{q_i}^Q)$  and use function  $kNN(\cdot, k)$  to find  $k$ -nearest neighbors of a node

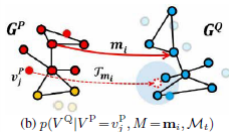
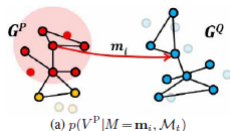
$$p(V^P = v_j^P | M = m_i, M_t) = \begin{cases} 1/k_1, & \text{if } v_j^P \in kNN(v_{p_i}^P, k_1) \\ 0, & \text{otherwise} \end{cases}$$

- Let  $NN(\cdot)$  define a nearest neighbor of a node

$$p(V^Q = v_b^Q | V^P = v_j^P, M = m_i, M_t) =$$

$$\begin{cases} 1, & \text{if } v_b^Q = NN(\tau_{m_i}(p_j^P)) \text{ and } (v_j^P, v_b^Q) \in M_t \\ \exp(-d_{jb|m_i})/Z, & \text{if } v_b^Q \in kNN(\tau_{m_i}(p_j^P), k_2) \text{ and } (v_j^P, NN(\tau_{m_i}(p_j^P))) \notin M_t \\ 0, & \text{otherwise} \end{cases}$$

$$Z = \sum_{v_b^Q \in kNN(\tau_{m_i}(p_j^P), k_2)} \exp(-d_{jb|m_i}) \text{ is a normalization constant}$$



# Evaluation

# The End

Thank you for your attention!

# References

- [1] Minsu Cho and Kyoung Mu Lee. Progressive graph matching: Making a move of graphs via probabilistic voting. *Proc. Computer Vision and Pattern Recognition (CVPR)*, 2012.
- [2] Minsu Cho, Kyoung Mu Lee, and Jungmin Lee. Reweighted random walks for graph matching. *European Conference on Computer Vision (ECCV)*, 2010.