

Introduction to Computation

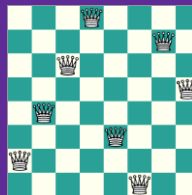
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Prof. Fan Cheng

Shanghai Jiao Tong University

chengfan85@gmail.com

<https://github.com/ichengfan/itc>





Outline

- Recursion
- Problems

Recursion



Function calls

- In python, a function can call another function defined before it
- In a program, we have a chain of function calls: $A_1 \rightarrow A_2 \rightarrow A_3 \rightarrow \dots \rightarrow A_n$
 - Call order: $A_1 \rightarrow A_2 \rightarrow A_3 \rightarrow \dots \rightarrow A_n$
 - Return order: $A_n \rightarrow A_{n-1} \rightarrow A_{n-3} \rightarrow \dots \rightarrow A_1$
 - Example: the boss of a company plans to check the progress of a task

```
def f1():  
    print("f1: begin")  
  
def f2():  
    print("f2: begin")  
    f1()  
    print("f1: finish")  
  
def f3():  
    print("f3: begin")  
    f2()  
    print("f2: finish")  
    print("f3: finish")  
  
f3()
```

调用
顺序

```
f3: begin  
f2: begin  
f1: begin  
f1: finish  
f2: finish  
f3: finish
```

返回
顺序

一个函数可以调用它自己吗？

Experiment

- Assume that we would like to implement a function $f(x)$, where $f(x)$ will call $f(y)$ inside

```
def f(x):  
    do_sth()  
  
    f(y)  
  
    return
```

- The order of function call $x \rightarrow y \rightarrow z \rightarrow w \rightarrow u \dots$ ($y \neq x$)
- If there are infinite function calls (e.g., loop) in the program, an error will occur
- In `do_sth()`, there must be a condition where the function execution will be returned, and the self call will not be invoked

```
13 def f(x):  
14     if x == 0:  
15         return 3  
16  
17     return 1 + f(x-1)  
18  
19 print(f(100))
```

103

```
1 def f(n):  
2     if n == 0:  
3         return n  
4     if n == 1:  
5         return 1  
6  
7     return f(n-1) + f(n-2)  
8  
9 print(f(10))
```

55

Recursion (递归)

It is legal for a function to call itself

- In mathematics, **recursive functions** allow a series to define itself by the other items

$$f_0 = f_1 = 1, f_{n+2} = f_{n+1} + f_n$$

- Recall: A function call is determined by the pair

(function_name, parameters)

For example, $(\text{print}, 3) \neq (\text{print}, [3])$

Thus, it does make sense to invoke itself inside the function definition

- The challenge is that recursion may be never ended (Logic error)

$$f(n) \rightarrow f(n-1) \rightarrow f(0) \rightarrow f(-1), \dots, f(-\infty)$$

There should be an stop!

- Idea: To define a function $f(n)$
 - For the simple cases like $n=0$ or 1 , we directly return its values
 - For the general cases, we try to solve $f(n)$ by reduce it to its previous solutions $f(0), f(1), \dots, f(n-1)$
 - By mathematical induction (数学归纳法), the solutions above always terminate in finite steps

Recursive function

1. A simple base case (or cases) (基本状态)
2. A set of rules which reduce (化简) all other cases toward the base case

- The formal definition: a method exhibits recursive behavior when it can be defined by **two properties**:
- The first property will make sure the function will terminate finally. 简单的情况直接计算（不递归）
- The second property will call the function itself to reach the base case. 复杂的情况转化为简单的情况（递归）

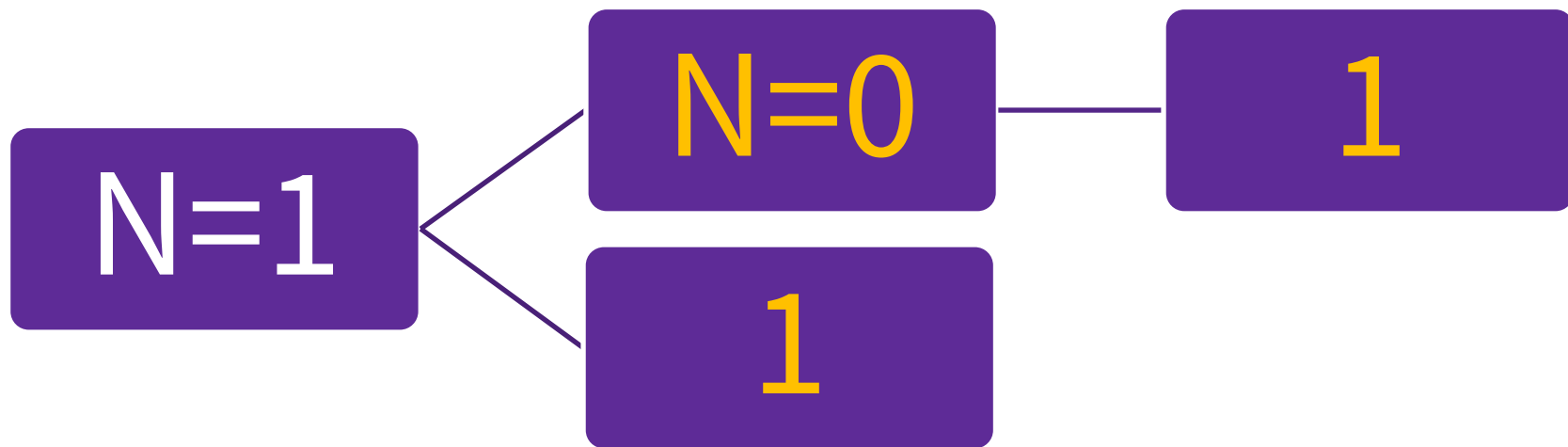
Example: compute the factorial(n): $n!$

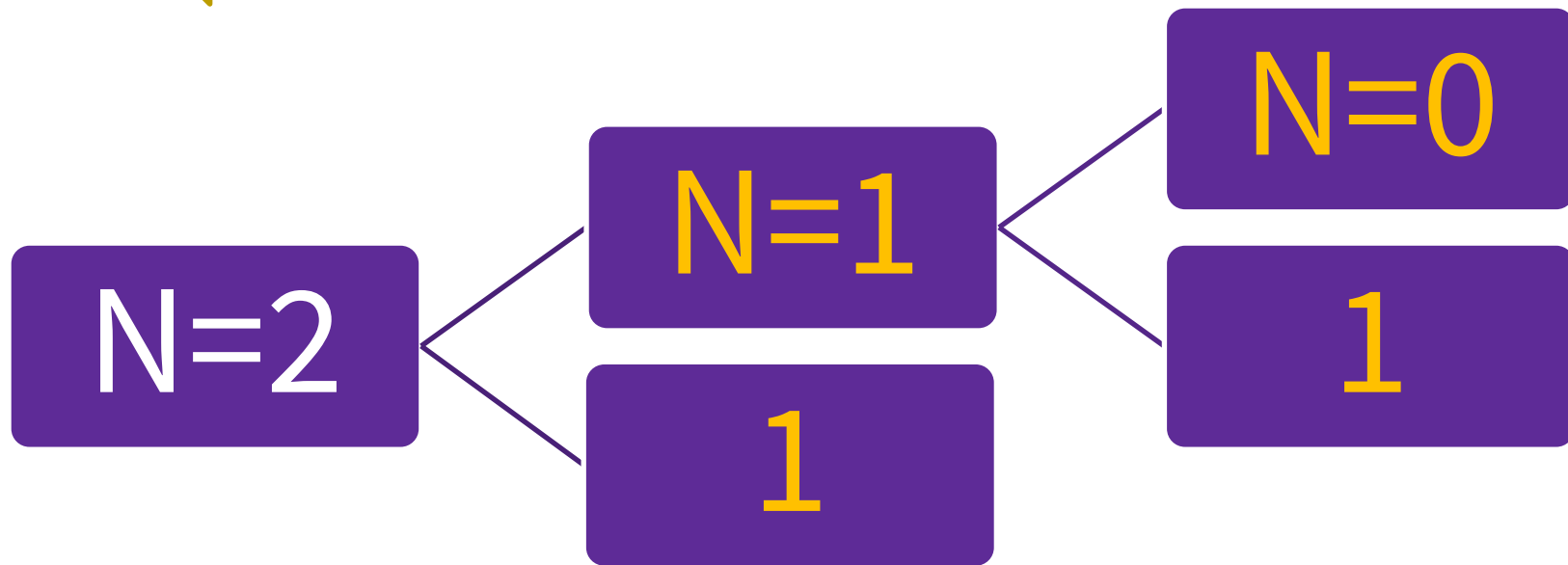
```
1 def factorial(n):
2     print(f"{n} begins")
3     if n == 0:
4         return 1
5     else:
6         f = factorial(n - 1)
7         print(f"{n-1} done")
8         return n * f
9
10
11 print(factorial(7))
```

```
7 begins
6 begins
5 begins
4 begins
3 begins
2 begins
1 begins
0 begins
0 done
1 done
2 done
3 done
4 done
5 done
6 done
5040
```

$N=0$

1





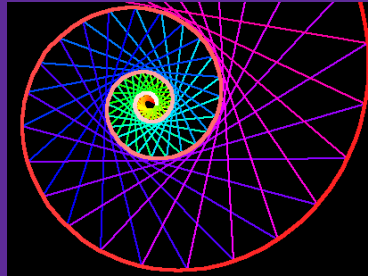
Stack (栈) diagram



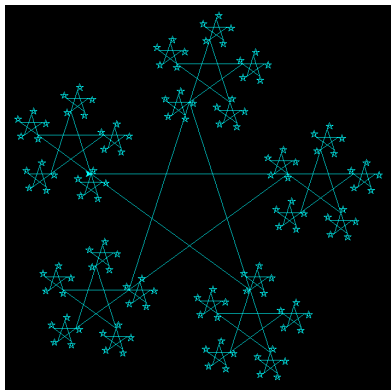
- Suppose we have a collection of books. When we pile them on a desk, we will first put a book on surface of the desktop and put the second one on the top of the first one, and so on. This process is called **push (推)**
- When we want to pick up the first book, we need to move the books from the last one to the second one. This process is called **pop (弹)**
- The whole process can be modeled by **stack**.

FIFO: First in last out (push→pop)
- In a program, we have a chain of function calls: $A_1 \rightarrow A_2 \rightarrow A_3 \rightarrow \dots \rightarrow A_n$
 - Stack diagram: when a function calls another function, the processed can also be modeled by a stack.
 - The functions A_1, A_2, \dots, A_n will be **pushed** into the stack
 - When A_n is executed, it will be **pop** up from the stack. The system will repeat popping up A_{n-1}, \dots, A_2, A_1

Problems



Recursion (递归)



$$F(n) = g(F(n-1), \dots, F(0))$$

$$F(0), F(1), F(2)$$

递归函数解决问题的特点

1. 问题具有递归的结构，可以状态转移(从一个函数值到另一个)
2. 初始状态容易处理

递归函数编写的注意事项，必须严格按照以下步骤

1. 首先判断是否达到基本状态
2. 再决定是否用递归关系处理
3. 否则，会陷入死循环

一个函数调用由 函数名+具体的参数值 决定：f(1)和f(2)是不同的函数调用
递归调用：表面上是自己调用自己，实际上参数不同，是不同的函数调用

Greatest Common Divisor (GCD)

$$(9,6) = (6,9\%6) = (6,3) = (3,6\%3) = (3,0) = 3$$

最大公约数

- 递归: $\text{gcd}(a, b) = \text{gcd}(b, a\%b)$
- 初始: $\text{gcd}(a, 0) = a$
- 证明:
 - 会在有限步后结束
 - $\log a$ 步结束

```
1 def gcd(a, b):
2     if a < b:
3         return gcd(b, a)
4     if b == 0:
5         return a
6
7     return gcd(b, a % b)
8
9
10 print(gcd(120, 88), gcd(90, 125))
```

$\text{gcd}(a, b)$

0 ($b=0$)



$\text{gcd}(a, b)$

$\text{gcd}(b, a \% b)$

Fibonacci sequence

$\{f_0, f_1, f_2, \dots, f_n, \dots\}$: $f_0 = 0$, $f_1 = 1$, $f_n = f_{n-1} + f_{n-2}$

```
1 def fib(n):
2     if n == 0 or n == 1:
3         return n
4
5     return fib(n - 1) + fib(n - 2)
6
7
8 print(fib(10))
```

55

```
1 def fib(n):
2     return n if n == 0 or n == 1 else fib(n - 1) + fib(n - 2)
3
4
5 print(fib(0), fib(1), fib(10), fib(35))
```

0 1 55 9227465

Exercise:

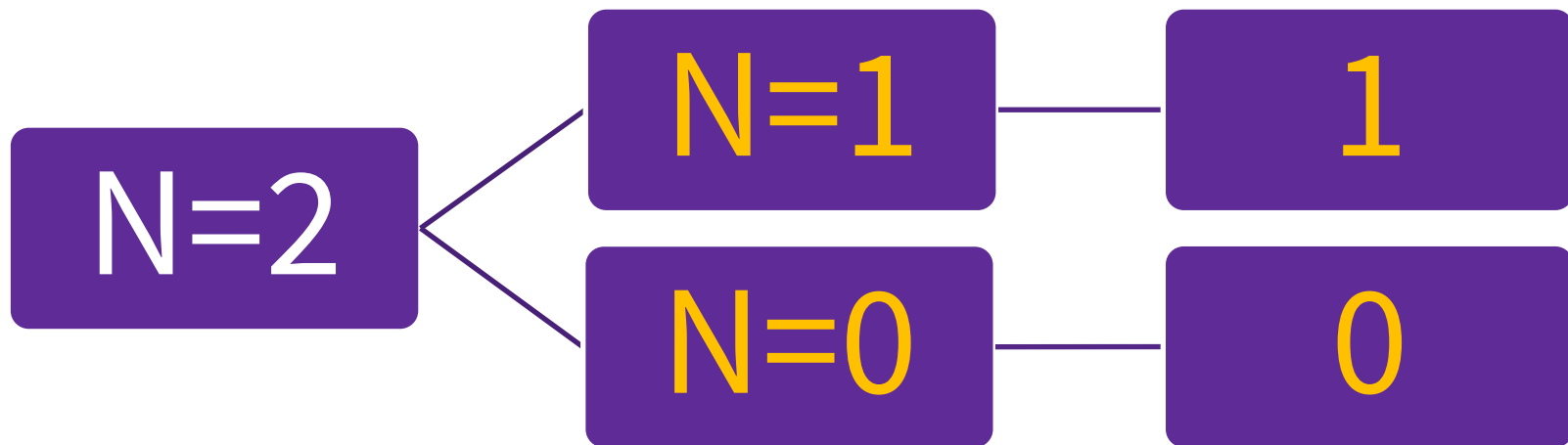
- $a_1 = 1, a_2 = 2, a_3 = 3,$
 $a_{n+3} = 3a_{n+2} - 2a_{n+1} + a_n$
- $f(0, n) = 1, f(m, 0) = m$
 $f(m, n) = f(m - 1, n) + f(m, n - 1) - f(m - 1, n - 1)$

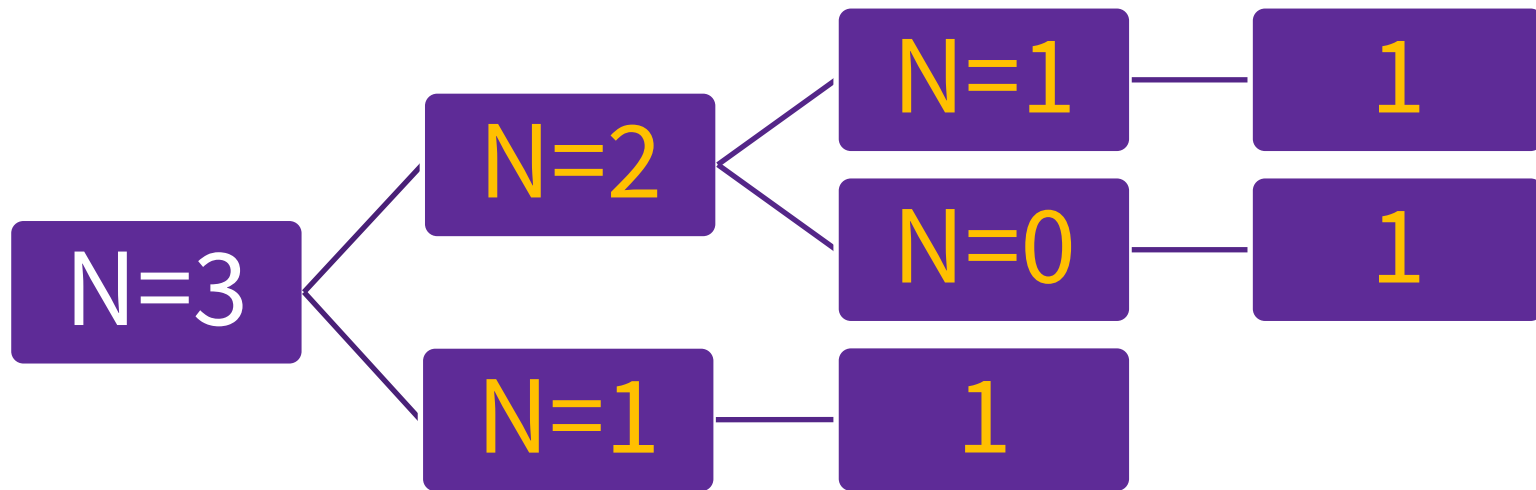
$N=0$

0

$N=1$

1





2^n 步结束

Bits Problems

- Compute the binary form of an integer
- Sum of bits: the sum of all the bits
- Length of an integer
- Inverse an integer. Don't use str

```
1 def binary_int(n):
2     return str(n) if n == 0 or n == 1 else binary_int(n // 2) + str(n % 2)
3
4
5 print(binary_int(0), binary_int(1), binary_int(2), binary_int(3), binary_int(4))
```

0 1 10 11 100

```
1 def sum_of_bits(n):
2     return n if n == 0 or n == 1 else sum_of_bits(n // 2) + (n & 1)
3
4
5 print(sum_of_bits(7), sum_of_bits(5), sum_of_bits(4), sum_of_bits(1))
```

3 2 1 1

```
1 def length_int(n):
2     return 1 if n < 10 else length_int(n // 10) + 1
3
4
5 def inverse_int(n):
6     return n if n < 10 else (n % 10) * 10 ** (length_int(n) - 1) + inverse_int(n // 10)
7
8
9 print(inverse_int(1001), inverse_int(1000), inverse_int(123456789))
```

1001 1 987654321

Infinite recursive function

- Recursive functions without base case will lead to errors

```
94 def print_test(str1):  
95     print_test(str1)  
96  
97 print_test("Hello world")
```

```
Traceback (most recent call last):  
  File "C:\Users\fcheng\OneDrive\CS124计算导论\2018\lecture notes\2.py", line 5, in <module>  
    print_test("Hello world")  
  File "C:\Users\fcheng\OneDrive\CS124计算导论\2018\lecture notes\2.py", line 3, in print_test  
    print_test(str)  
  File "C:\Users\fcheng\OneDrive\CS124计算导论\2018\lecture notes\2.py", line 3, in print_test  
    print_test(str)  
  File "C:\Users\fcheng\OneDrive\CS124计算导论\2018\lecture notes\2.py", line 3, in print_test  
    print_test(str)  
  [Previous line repeated 995 more times]  
RecursionError: maximum recursion depth exceeded
```

RecursionError: maximum recursion depth exceeded

a^n

- How to implement a^n , where n is an integer
- $y = a^{n/2}$
- $a = y \times y$

```
1  def my_pow(a, n):
2      if n == 0:
3          return 1
4
5      t = my_pow(a, n // 2)
6
7      return t * t if n % 2 == 0 else t * t * a
8
9
10 for i in range(4):
11     print(my_pow(2, i))
```

1
2
4
8

Ackermann function

For nonnegative integers m and n , $A(m, n)$ is defined as follows:

1. $A(0, n) = n + 1$
 2. $A(m + 1, 0) = A(m, 1)$
 3. $A(m + 1, n + 1) = A(m, A(m + 1, n))$
- Its value grows very rapidly; for example, $A(4, 2)$ results in $2^{65536} - 3$, an integer of 19,729 decimal digits.

```
1 def Ackermann(m, n):
2     if m == 0:
3         return n + 1
4     if n == 0:
5         return Ackermann(m - 1, 1)
6
7     return Ackermann(m - 1, Ackermann(m, n - 1))
8
9
10 print(Ackermann(3, 6))
11 print(Ackermann(4, 0))
```

Values of $A(m, n)$

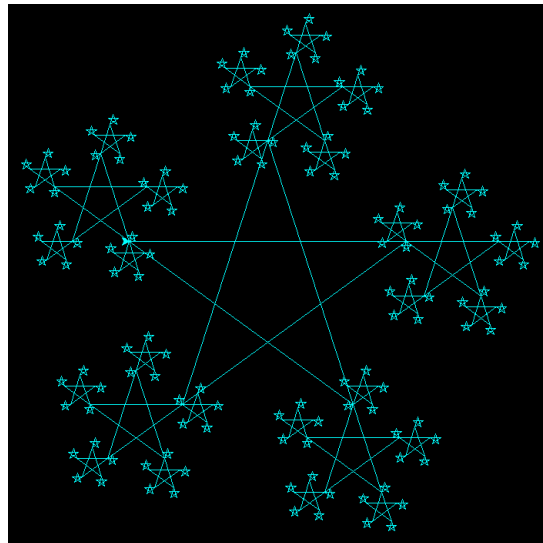
$m \backslash n$	0	1	2	3	4
0	1	2	3	4	5
1	2	3	4	5	6
2	3	5	7	9	11
3	5	13	29	61	125
4	13 $= 2^{2^2} - 3$ $= 2 \uparrow\uparrow 3 - 3$	65533 $= 2^{2^{2^2}} - 3$ $= 2 \uparrow\uparrow 4 - 3$	$2^{65536} - 3$ $= 2^{2^{2^{2^2}}} - 3$ $= 2 \uparrow\uparrow 5 - 3$	$2^{2^{65536}} - 3$ $= 2^{2^{2^{2^{2^2}}}} - 3$ $= 2 \uparrow\uparrow 6 - 3$	$2^{2^{2^{65536}}} - 3$ $= 2^{2^{2^{2^{2^{2^2}}}}} - 3$ $= 2 \uparrow\uparrow 7 - 3$

509
13

Fractal

In Python, turtle graphics provides a representation of a physical “turtle” (a little robot with a pen) that draws on a sheet of paper on the floor.

```
1  import turtle
2
3  tur = turtle.Turtle()
4  tur.speed(6)
5  tur.getscreen().bgcolor("black")
6  tur.color("cyan")
7  tur.penup()
8  tur.goto((-200, 50))
9  tur.pendown()
10
11
12 def star(turtle, size):
13     if size <= 10:
14         return
15     else:
16         for i in range(5):
17             turtle.forward(size)
18             star(turtle, size / 3)
19             turtle.left(216)
20
21
22 star(tur, 360)
23 turtle.done()
```



<https://docs.python.org/3/library/turtle.html>

Sort

- Given a list of integers: number= [7, 3, 4, 5, 1, 2, 3]. Rearrange the numbers in ascending order
- Divide the list into two parts in equal sizes: number1, number2
- Sort number1 and number2, respectively
- Merge the sorted number1 and number2 to recover number

```
1 def merge(number, left, right):
2     pass
3
4 def merge_sort(number, left, right):
5     if left >= right: return
6
7     merge_sort(number, left, (left+right)//2)
8     merge_sort(number, (left+right)//2+1, right)
9
10    merge(number, left, right)
11
12
13 number = [-x for x in range(10)]
14 merge_sort(number, 0, 9)
15 print(number)
```

Nested List

- Find the sum of the numbers in a nested list
- `[[1, [1], [1]], [2], [3]]`

```
1 def sum_of_nested(number):
2     if len(number) == 0:
3         return 0
4
5     if type(number[-1]) != type([]):
6         return number[-1] + sum_of_nested(number[0:-1])
7
8     return sum_of_nested(number[-1]) + sum_of_nested(number[0:-1])
9
10 number_lst = [[1, [1], [1]], [2], [3], 1, 1, 1]
11 print(sum_of_nested(number_lst))
12
13 number_lst = [[[[[1]]]]]
14 print(sum_of_nested(number_lst))
15
16 number_lst = [[[[[[[ ]]]]]]]
17 print(sum_of_nested(number_lst))
```

```
11
1
0
```

Binary_search

- Given a non-decreasing list of numbers: $S = [-3, -4, 1, 3, 5]$, find whether $x \in S$
- Halving S equally, search the left or the right part according to the comparison of x and mid

```
1 def binary_search(number, x, left, right):
2     if left > right:
3         return None
4
5     mid = (left + right) // 2
6     if number[mid] == x:
7         return mid
8
9     if number[mid] > x:
10         return binary_search(number, x, left, mid - 1)
11
12     return binary_search(number, x, mid + 1, right)
13
14
15 number = [-3, -4, 1, 3, 5]
16
17 print(binary_search(number, -6, 0, len(number) - 1))
18 print(binary_search(number, -3, 0, len(number) - 1))
19 print(binary_search(number, 1, 0, len(number) - 1))
20 print(binary_search(number, 0, 0, len(number) - 1))
21 print(binary_search(number, 5, 0, len(number) - 1))
22 print(binary_search(number, 7, 0, len(number) - 1))
```

```
None
0
2
None
4
None
```

permutation

- 问题：生成1,...,n的所有排列
- 问题具有递归的特点：1-n可以由1-(n-1)插入n得到
- 函数定义 `def perm(n):` # return list: 每个元素是1-n的一个排列 `[(), (), (), ...]`

```
def perm(n):
    if n==1:
        return [(1,)]

    lst = perm(n-1)
    ans = []

    for x in lst:
        for i in range(n):
            nx = x[:i] + (n,) + x[i:]
            ans.append(nx)

    return ans
```

```
1 import pprint as pp
2
3 pp.pprint(perm(1))
4 pp.pprint(perm(2))
5 pp.pprint(perm(3))
6 pp.pprint(perm(4))
7 pp.pprint(perm(5))
8 pp.pprint(perm(6))
```

pprint — Data pretty printer
<https://docs.python.org/3/library/pprint.html#module-pprint>

```
print(perm(1))      [(1,)]
print(perm(2))      [(2, 1), (1, 2)]
print(perm(3))      [(3, 2, 1), (2, 3, 1), (2, 1, 3), (3, 1, 2), (1, 3, 2), (1, 2, 3)]
print(perm(4))      [(4, 3, 2, 1), (3, 4, 2, 1), (3, 2, 4, 1), (3, 2, 1, 4), (4, 2, 3, 1), (2, 4, 3, 1), (2, 3, 4, 1), (2, 3, 1, 4), (4, 2, 1, 3), (2, 4, 1, 3), (2, 1, 4, 3),
(2, 1, 3, 4), (4, 3, 1, 2), (3, 4, 1, 2), (3, 1, 4, 2), (3, 1, 2, 4), (4, 1, 3, 2), (1, 4, 3, 2), (1, 3, 4, 2), (1, 3, 2, 4), (4, 1, 2, 3), (1, 4, 2, 3),
(1, 2, 4, 3), (1, 2, 3, 4)]
```

扩展问题

1. 有序排列
2. n 选 r 排列
3. n 选 r 组合

Increasing permutations

```
def perm1(n):
    if n==1:
        return [(1,)]

    lst = perm1(n-1)
    ans = []

    for i in range(1,n+1): # 把元素 i 放到首位
        for x in lst:
            lx = list(x)

            for ii in range(n-1): # 用i+1,i+2,...,n代替原来排列中的i,i+1,...,n-1
                if lx[ii]>=i:
                    lx[ii] += 1

            nx = (i,)+tuple(lx)
            ans.append(nx)

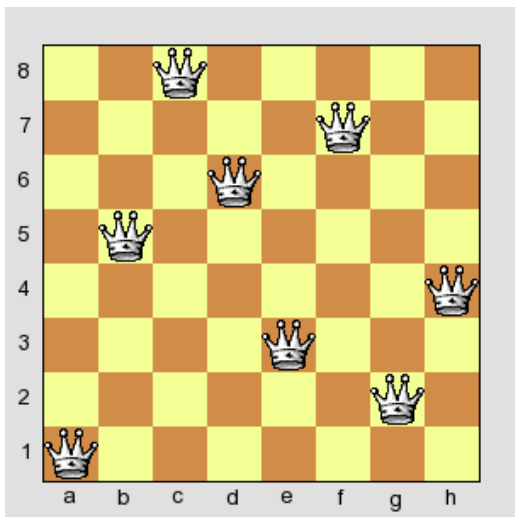
    return ans
```

```
print("Increasing permutation:")
print(perm1(1))
print(perm1(2))
print(perm1(3))
print(perm1(4))
```

```
1 import pprint as pp
2 pp.pprint("Increasing permutation:")
3 pp.pprint(perm1(1))
4 pp.pprint(perm1(2))
5 pp.pprint(perm1(3))
6 pp.pprint(perm1(4))
```

Eight Queens

- Find the number of solutions for 8 Queens problem
 - How to check whether a given solution is valid
 - How to generate all the possible solutions
 - Ans: 92



如何表示一个正确的解

表示：用元组 f 表示放在各行的皇后的列号，那么 f 必须是 $1, \dots, 8$ 的一个排列。

思路：枚举1-8的所有排列，判断各个排列是不是合法的皇后放置方式（不同行、不同列，不同对角线）

最多 $8!$ 种可能

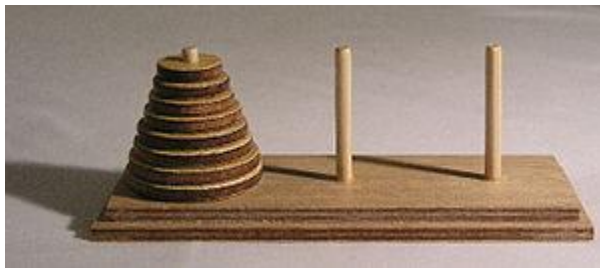
Eight Queens: Code

- Python中自带了permutation函数，可以生成各种排列组合
- 本问题中，和前面的perm函数等价（比较测试你的代码是否正确）

```
1  from itertools import permutations
2  perm = permutations(list(range(8)))
3  ans = 0
4
5  def valid(sln):
6      for i in range(8):
7          for j in range(i+1, 8):
8              if abs(sln[i]-sln[j]) == abs(i-j):
9                  return False
10     return True
11
12  for x in perm:
13      if valid(x):
14          ans += 1
15
16  print(f"The number of solutions is {ans}")
```

规范：循环的迭代深度不超过2轮，超过了用函数

Hanoi



- The Tower of Hanoi (汉诺塔) is a mathematical game or puzzle.
- It consists of three rods and a number of disks of different sizes, which can slide onto any rod.
- The puzzle starts with the disks in a neat stack in ascending order of size on one rod, the smallest at the top, thus making a conical shape.
- The objective of the puzzle is to move the entire stack to another rod, obeying the following simple rules:
 - A. Only one disk can be moved at a time.
 - B. Each move consists of taking the upper disk from one of the stacks and placing it on top of another stack or on an empty rod.
 - C. No disk may be placed on top of a smaller disk

Give a solution to this problem.

假定，任务是从圆柱1，将圆盘经过圆柱2，全部移动到圆柱3。输出移动的过程

```

1  def Hanoi(n, x, y, z):
2      if n == 1:
3          print(f"{n}: {x}->{z}")
4          return
5
6      Hanoi(n - 1, x, z, y)
7      print(f"{n}: {x}->{z}")
8      Hanoi(n - 1, y, x, z)
9
10
11 print("Hanoi: n = 1")
12 Hanoi(1, 1, 2, 3)
13
14 print("Hanoi: n = 2")
15 Hanoi(2, 1, 2, 3)
16
17 print("Hanoi: n = 3")
18 Hanoi(3, 1, 2, 3)

```

```

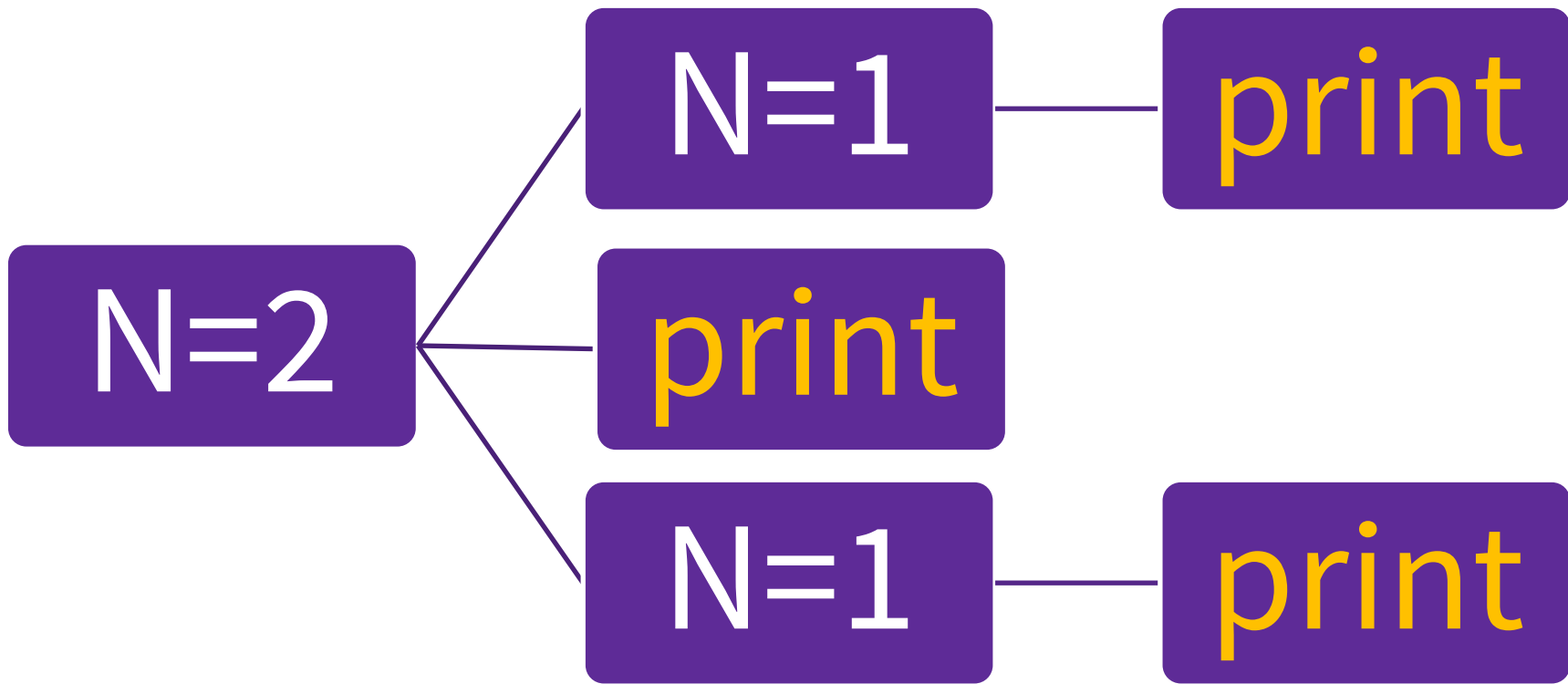
Hanoi: n = 1
1: 1->3
Hanoi: n = 2
1: 1->2
2: 1->3
1: 2->3
Hanoi: n = 3
1: 1->3
2: 1->2
1: 3->2
3: 1->3
1: 2->1
2: 2->3
1: 1->3

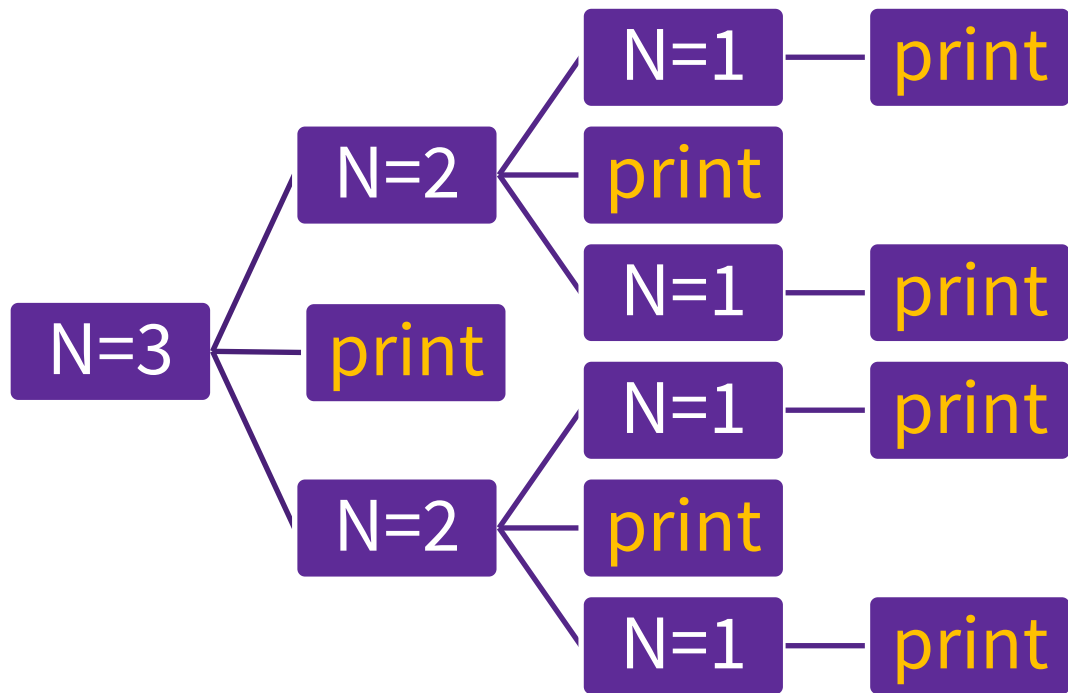
```

函数可以递归定义。但是函数调用时，即使函数名相同，参数不同也可以认为是不同的函数调用

N=1

print





递归函数

- 函数：一段具有某种功能的代码。
 - 函数执行结束，表示着已经完成了函数所定义的功能
- 功能可以是：
 - 函数值，譬如gcd（简单，好理解）
 - 抽象的一组事件。譬如Hanoi（抽象，要仔细分析）
 - 有问题？？？从定义出发思考
- 递归函数：
 - 1. 自己调用自己：从一个参数状态到另一个参数状态的转移（状态转移）
 - 2. 初始状态结束递归（保证不会死循环）
 - 3. 一个函数调用由 函数名+参数 决定，参数不同、函数名相同也是不同的调用

函数：一段具有某种功能的代码。函数执行结束，表示着已经完成了函数所定义的功能

以函数Hanoi(n, x, y, z)为例：

定义功能：输出将n个盘子从x途经y搬到z的过程（学习体会！）

由于Hanoi中每个小盘子只能放到大盘子上

首先要把盘子n，从x搬到z。必须先将1-(n-1)号盘子先搬到中间点y。

Hanoi(n-1, x, z, y) # 调用函数 n-1, x, z, y。函数结束运行后，表明已经完成了定义的功能（即输出了所有的步骤！细细体会！）

然后将n号盘子从x搬到z。

最后一步，我们需要将1-(n-1)号盘子从y搬到z，递归调用即可

Hanoi(n-1, y, x, z)#调用函数 n-1, x, z, y。函数结束运行后，表明已经完成了定义的功能（即输出了所有的步骤！细细体会！）

```
def Hanoi(n, x, y, z):  
    if n == 1: # 确保函数会结束递归  
        print(f"{n}: {x}->{z}")  
        return
```

#下面三个函数调用完整地实现了Hanoi既定的功能

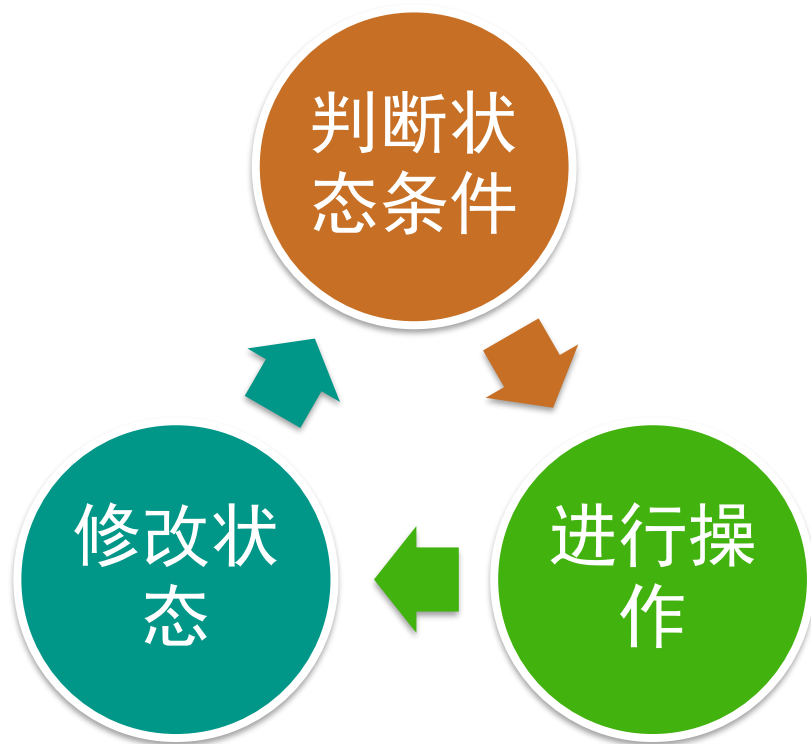
```
Hanoi(n-1, x, z, y)  
print(f"{n}: {x}->{z}")  
Hanoi(n-1, y, x, z)
```

递归与循环

- 递归调用其实是个循环过程
- 递归函数调用需要额外时间、空间开销：传递参数，保存中间值，切换函数
 - 循环比函数递归更高效
- 递归函数更容易写，更符合人的思维模式
- 递归函数是把复杂问题转化为简单问题
 - 初学者滥用递归：能用递归的地方一律递归
- 循环是从简单条件出发一步步构造复杂情况
- Life is short, use python: Python给大家很多便利，更符合人的思维
- 计算机的思维：0/1
 - 你只能用“三角形盖房子”
 - 程序员来适应计算机
- 任何程序都可以用赋值、逻辑语句、判断语句、循环语句、跳转语句实现
- 难点：循环语句
 - while是一个复杂过程
- 如何通过设计一个循环来完成一个复杂的功能

递归：从一般到特殊；循环：从特殊到一般

while: 状态更新

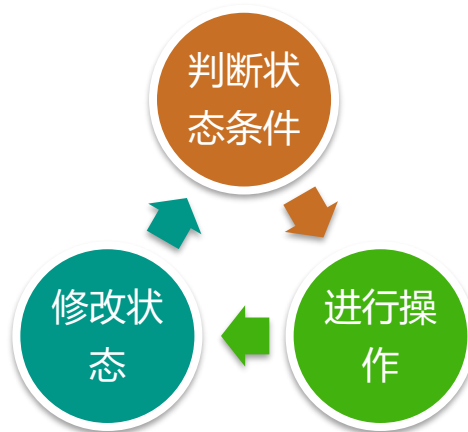


递归与循环

- 计算一个数各位数字的和

```
def digit_sum(x):  
    ans = 0  
    for i in str(x):  
        ans += int(i)  
    return ans  
  
def digit_sum_re(x):  
    if x < 10:  
        return x  
  
    return x%10 + digit_sum_re(x//10)  
  
def digit_sum_while(x):  
    ans = 0  
    while x>0:  
        ans += x%10  
        x //= 10  
  
    return ans  
  
print(digit_sum(123), digit_sum_re(123), digit_sum_while(123))
```

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while: 必须明确地想清楚整个变化过程, 从i到i+1(递归函数自动完成)
并不是所有的递归都可以很轻松地用while写

递归与循环

- 计算Fibonacci序列第n项

```
def fib(n):
    if n == 0:
        return 0
    if n == 1:
        return 1
    return fib(n-1) + fib(n-2)

def fib_loop_1(n):
    lst = [0]*(n+1)
    lst[1] = 1
    for i in range(2, n+1):
        lst[i] = lst[i-1] + lst[i-2]
    return lst[n]

def fib_loop_2(n):
    if n == 0:
        return 0
    if n == 1:
        return 1
    x1, x2 = 0, 1
    ans = 0
    for i in range(2, n+1):
        ans = x1 + x2
        x1, x2 = x2, ans
    return ans

print(fib(20), fib_loop_1(20), fib_loop_2(20))
```



- while: 必须明确地想清楚整个变化过程, 从i到i+1(递归函数自动完成)
- 并不是所有的递归都可以很轻松地用while写

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Recursive Functions → Loop

1. Write a Python program to calculate the sum of a list of numbers
2. Write a Python program to converting an integer to a string in any base.
3. Write a Python program of recursion list sum
 1. Test Data: [1, 2, [3,4], [5,6]]
 2. Expected Result: 21
4. Write a Python program to get the factorial of a non-negative integer
5. Write a Python program to get the sum of digital of a non-negative integer
 1. Test Data:
 2. sumDigits(345) -> 12
 3. sumDigits(45) -> 9
6. Write a Python program to calculate the geometric sum of n items

Note: In mathematics, a geometric series is a series with a constant ratio between successive terms

Example :

$$\sum_{i=0}^n ap^i$$
7. Write a Python program to calculate the value of 'a' to the power 'b'

Test Data : power(3,4) -> 81
8. Fibonacci, gcd, climbing steps, binary search