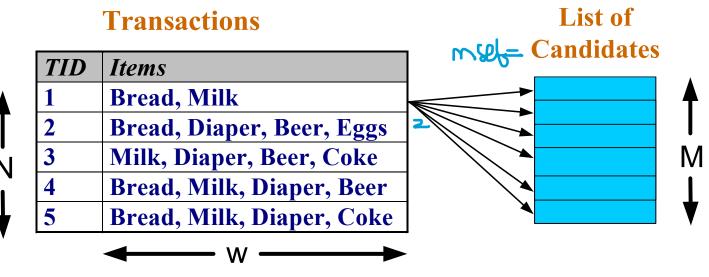
# ASSOCIATION RULE MINING

**BEIYU LIN** 

# FREQUENT ITEMSET GENERATION

- Brute-force approach:
  - Each itemset in the lattice is a candidate frequent itemset
  - Count the support of each candidate by scanning the database



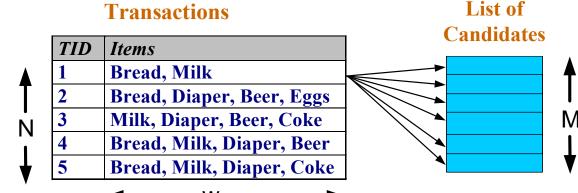
- Match each transaction against every candidate
- Complexity ~ O(NMw) => Expensive since M = 2<sup>d</sup> !!!

## FREQUENT ITEMSET GENERATION STRATEGIES

- Reduce the number of candidates (M)
  - Complete search: M=2<sup>d</sup>
  - Use pruning techniques to reduce M
- Reduce the number of transactions (N)
  - Reduce size of N as the size of itemset increases
  - Used by DHP and vertical-based mining algorithms

Given a transaction {B, M, D, C}, find all possible subset with size 3 from this transaction.

- Reduce the number of comparisons (NM) with size 3 from this transaction.
  - Use efficient data structures to store the candidates or transactions
  - No need to match every candidate against every transaction



#### REDUCING NUMBER OF CANDIDATES

- Apriori principle:
  - If an itemset is frequent, then all of its subsets must also be frequent
- Apriori principle holds due to the following property of the support measure:

$$\forall X, Y : (X \subseteq Y) \Rightarrow s(X) \ge s(Y)$$

- Support of an itemset never exceeds the support of its subsets
- This is known as the anti-monotone property of support

$$5 = \frac{6}{\text{totalT}} \times = \{M, B\}$$

$$Y = \{D\}$$

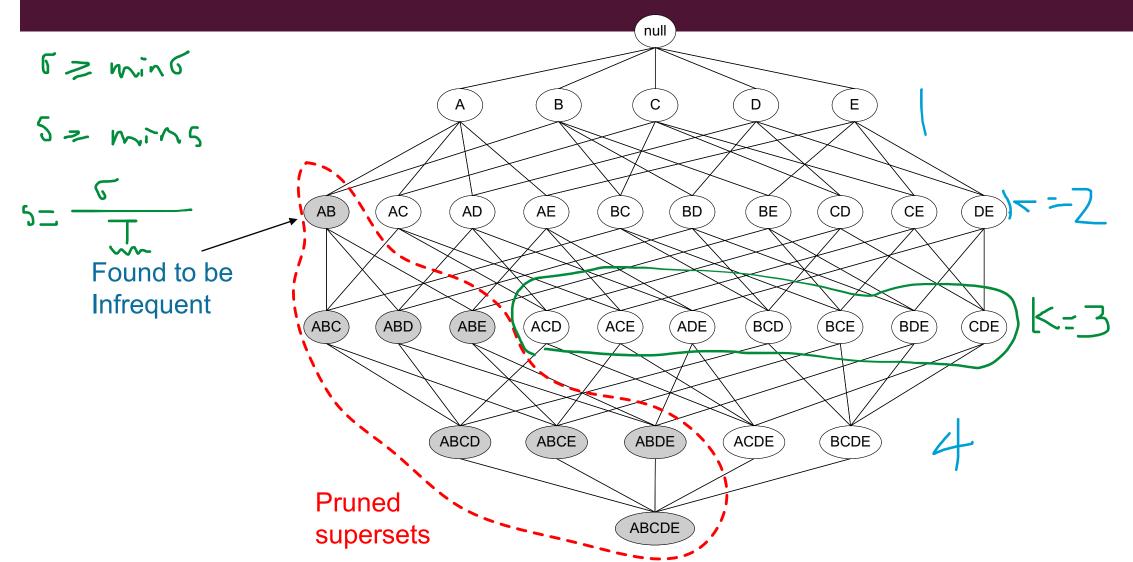
$$\times \cup Y = \{M, B, D\}$$

| TID | Items                     |
|-----|---------------------------|
| 1   | Bread, Milk               |
| 2   | Bread, Diaper, Beer, Eggs |
| 3   | Milk, Diaper, Beer, Coke  |
| 4   | Bread, Mill, Diaper, Beer |
| 5   | Bread, Milk, Diaper, Coke |

Support count: # of the itemsets
that show in the transaction 
$$= 2$$

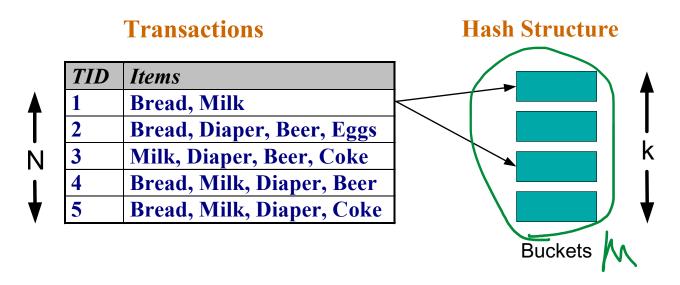
Confidence  $= \frac{\# \times \vee Y}{\#} = \frac{2}{3}$ 

# ILLUSTRATING APRIORI PRINCIPLE



#### SUPPORT COUNTING OF CANDIDATE ITEMSETS

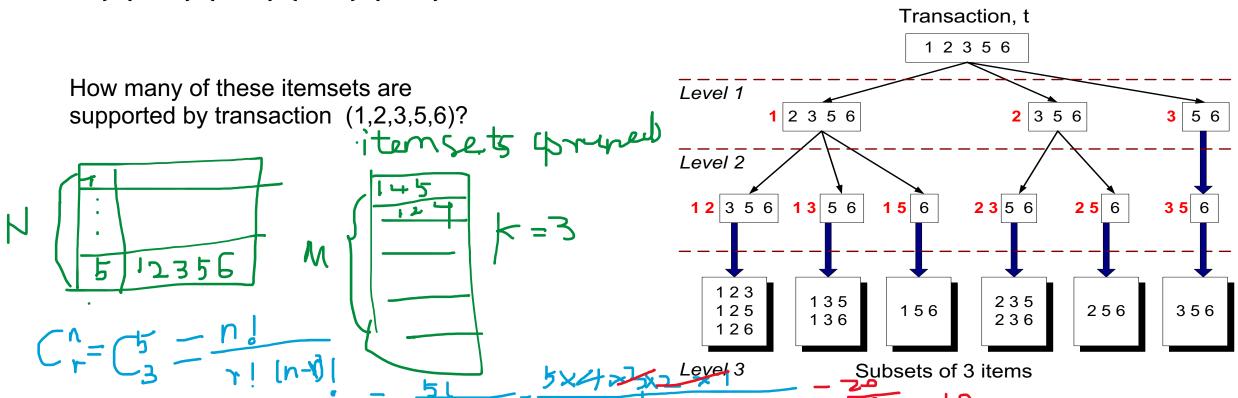
- To reduce number of comparisons, store the candidate itemsets in a hash structure / hash function
  - Instead of matching each transaction against every candidate, match it against candidates contained in the hashed buckets



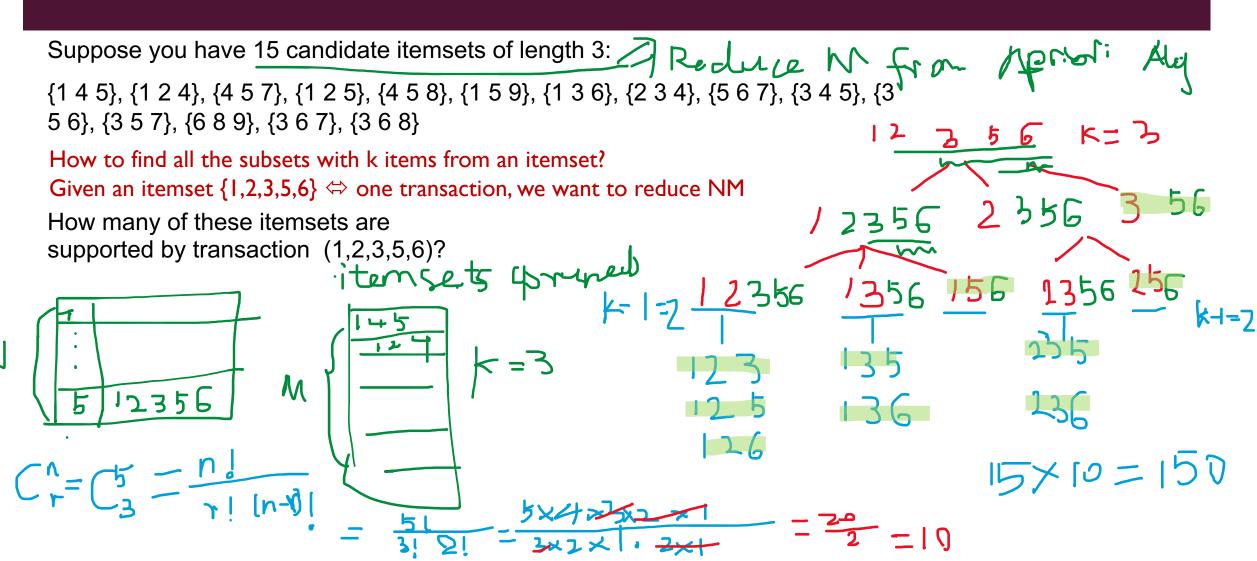


### SUPPORT COUNTING: AN EXAMPLE

Suppose you have 15 candidate itemsets of length 3: Reduce M from [1 4 5], {1 2 4}, {4 5 7}, {1 2 5}, {4 5 8}, {1 5 9}, {1 3 6}, {2 3 4}, {5 6 7}, {3 4 5}, {3 5 7}, {6 8 9}, {3 6 7}, {3 6 8}



### SUPPORT COUNTING: AN EXAMPLE



Suppose you have 15 candidate itemsets of length 3:

{1 4 5}, {1 2 4}, {4 5 7}, {1 2 5}, {4 5 8}, {1 5 9}, {1 3 6}, {2 3 4}, {5 6 7}, {3 4 5}, {3 5 6}, {3 5 7}, {6 8 9}, {3 6 7}, {3 6 8}

#### You need:

- Hash function
- Max leaf size: max number of itemsets stored in a leaf node (if number of candidate itemsets exceeds max leaf size, split the node)

