
Chapter 5

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Part-1

Internal Memory

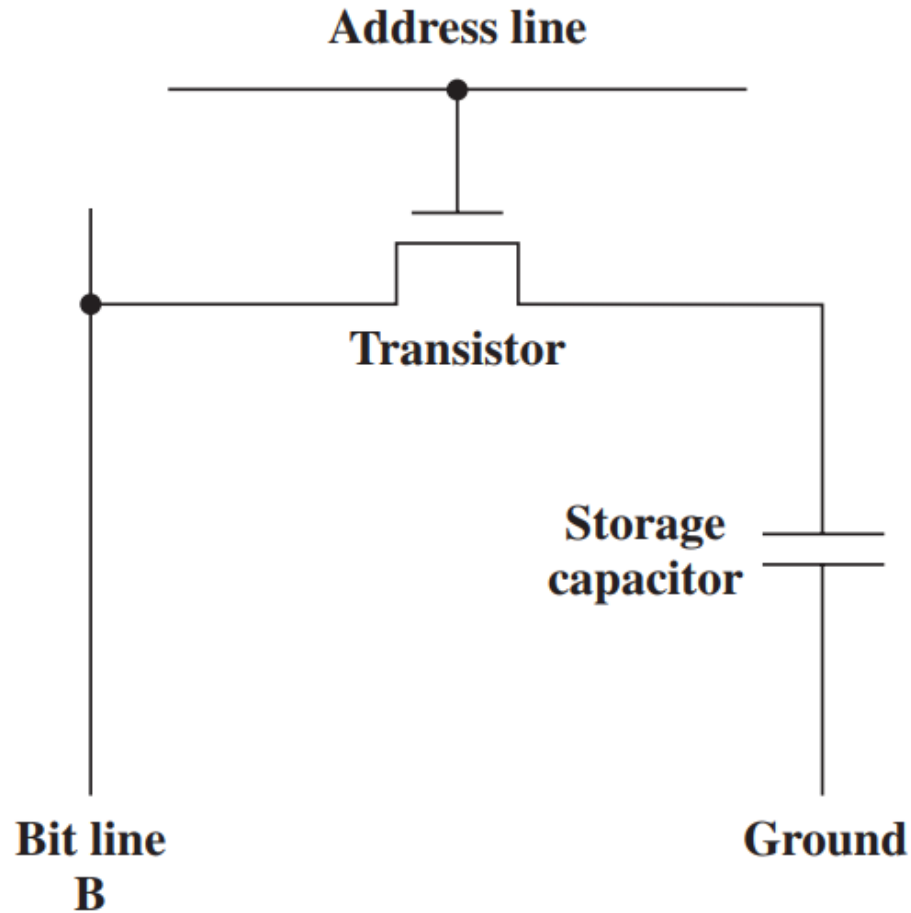
Semiconductor Memory

- RAM (Random Access Memory)
 - Misnamed as all semiconductor mem. are random access
 - Read/Write
 - Volatile
 - Temporary storage
 - Static or dynamic
- ROM (Read only memory)
 - Permanent storage
 - Read only

Dynamic RAM

- Bits stored as charge in capacitors
- Charges leak
- Need refreshing even when powered
- Simpler construction
- Smaller per bit
- Less expensive
- Need refresh circuits
- Slower
- Main memory (static RAM would be too expensive)

Dynamic RAM



Static RAM

- Bits stored as on/off switches
- No charges to leak
- No refreshing needed when powered
- More complex construction
- Larger per bit
- More expensive
- Does not need refresh circuits
- Faster
- Cache (here the faster the better)

Read Only Memory (ROM)

- Permanent storage
- Microprogramming (see later)
- Library subroutines
- Systems programs (BIOS)
- Function tables

Types of ROM

- Written during manufacture
 - Very expensive for small runs
- Programmable (once)
 - PROM
 - Needs special equipment to program
- Read “mostly”
 - Erasable Programmable (EPROM)
 - Erased by UV (it can take up to 20 minutes)
 - Electrically Erasable (EEPROM)
 - Takes much longer to write than read
 - a single byte can be erased
 - Flash memory
 - Erase memory electrically “block-at-a-time”

Physical Characteristics

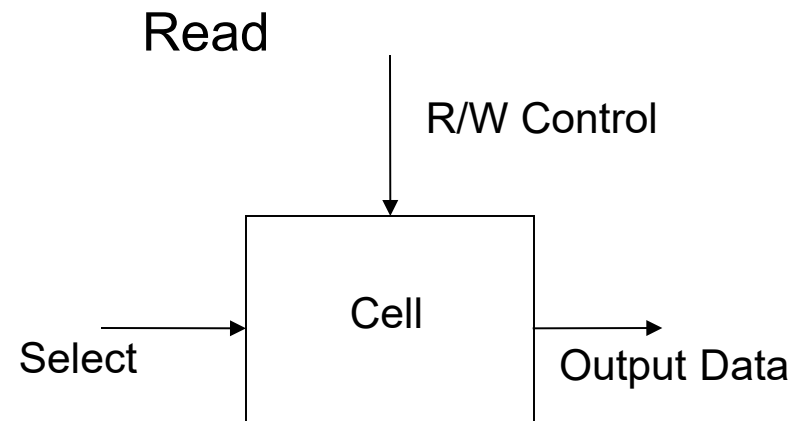
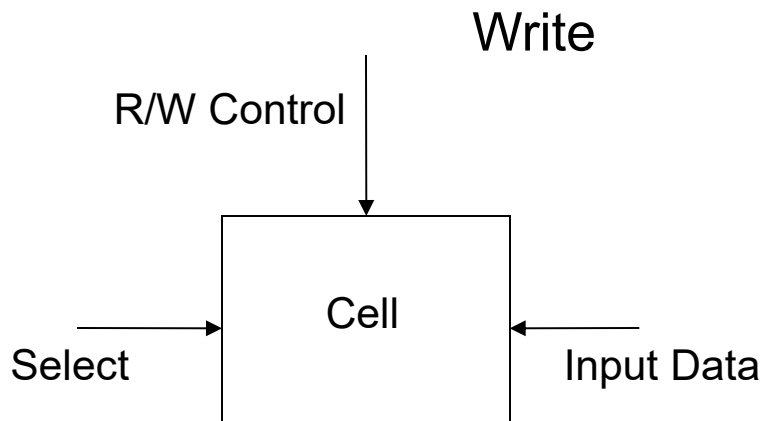
- Decay (refresh time)
- Volatility (needs power source)
- Erasable
- Power consumption

Organisation

- Physical arrangement of bits into words
- Not always obvious
 - e.g. interleaved

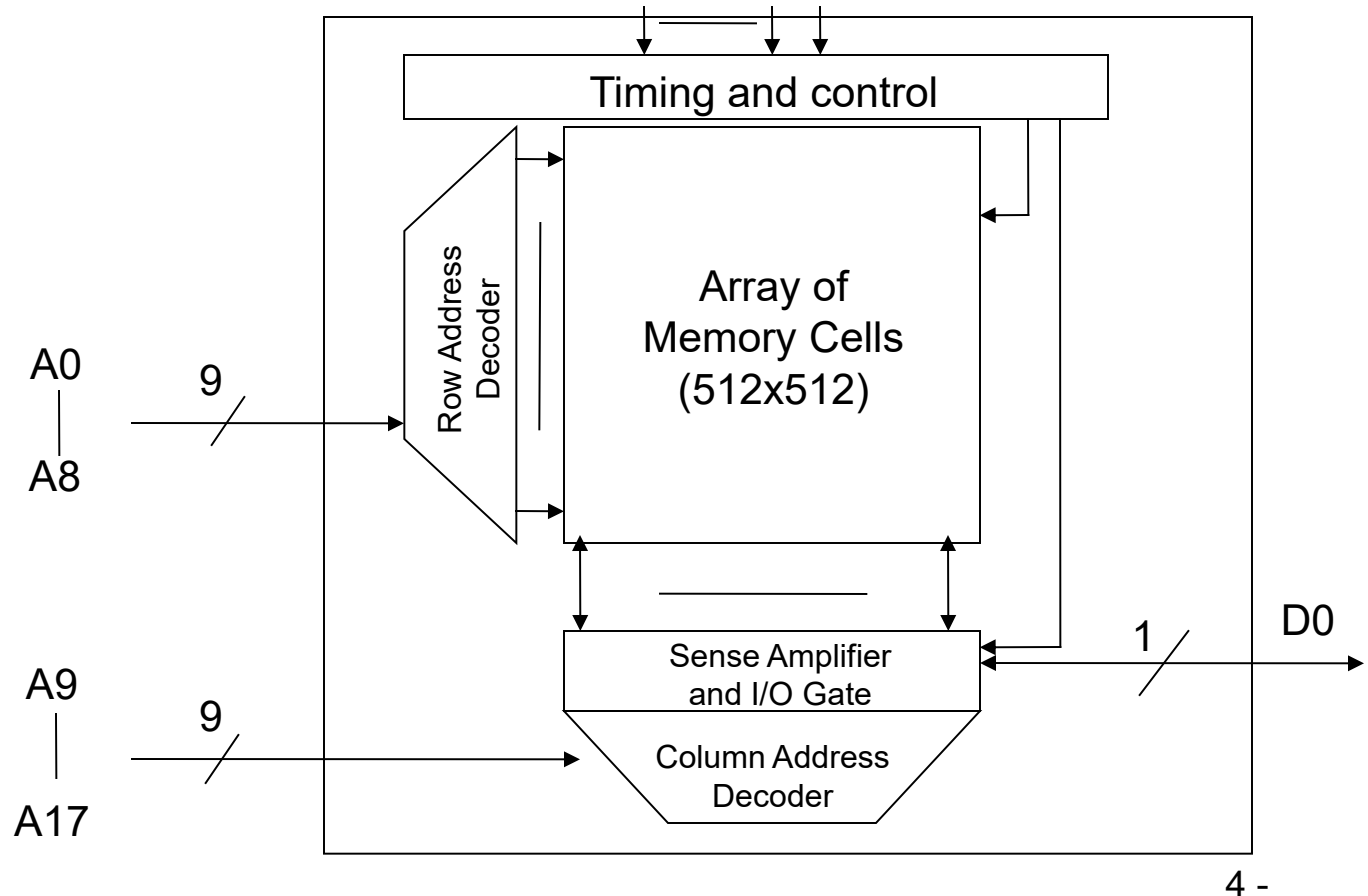
Basic Organization (1)

- Basic element: memory cell
 - has 2 stable states: one represent 0, the other 1
 - can be written at least once
 - can be read



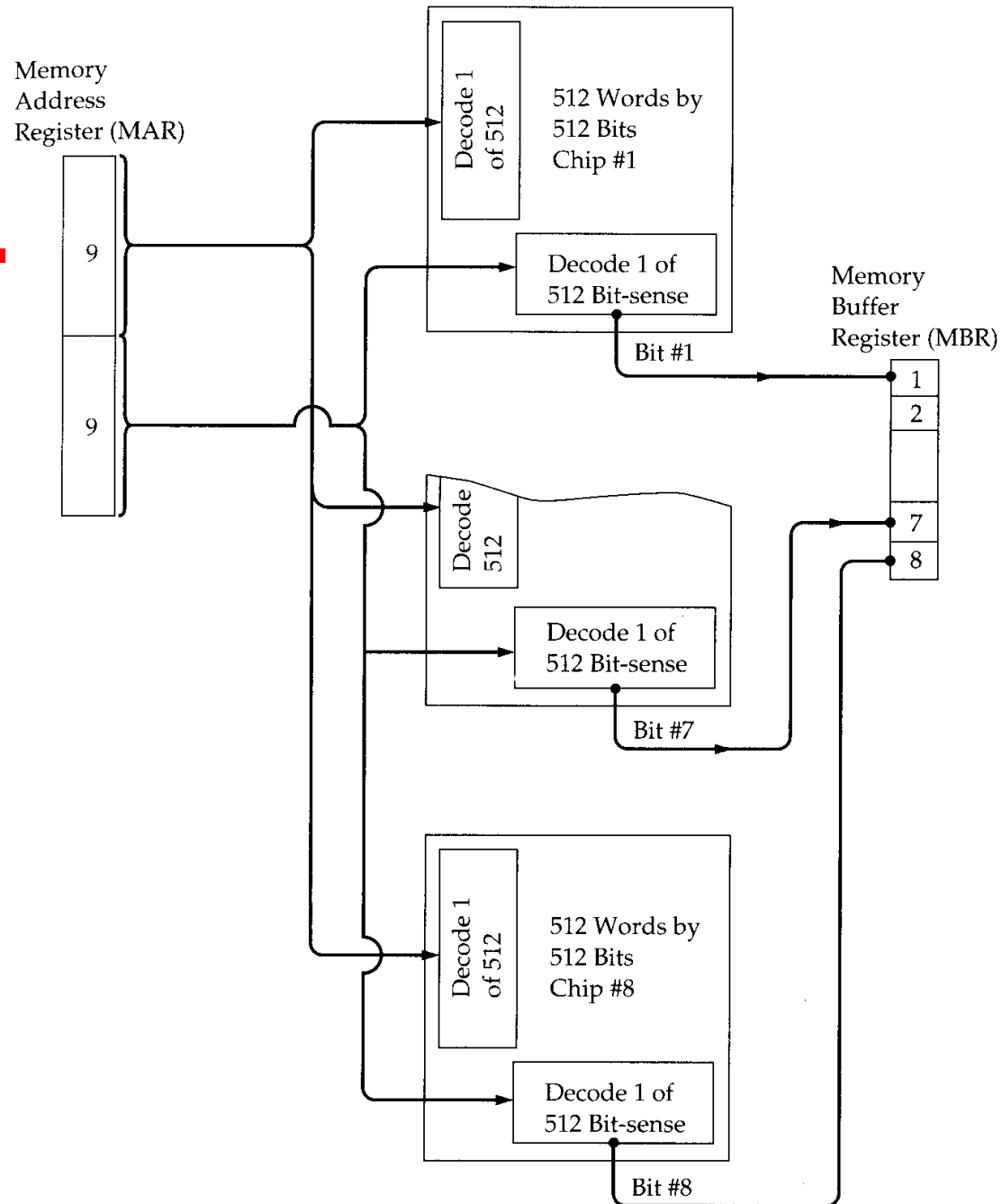
Basic Organization (2)

- Basic organization of a 512x512 bits chip

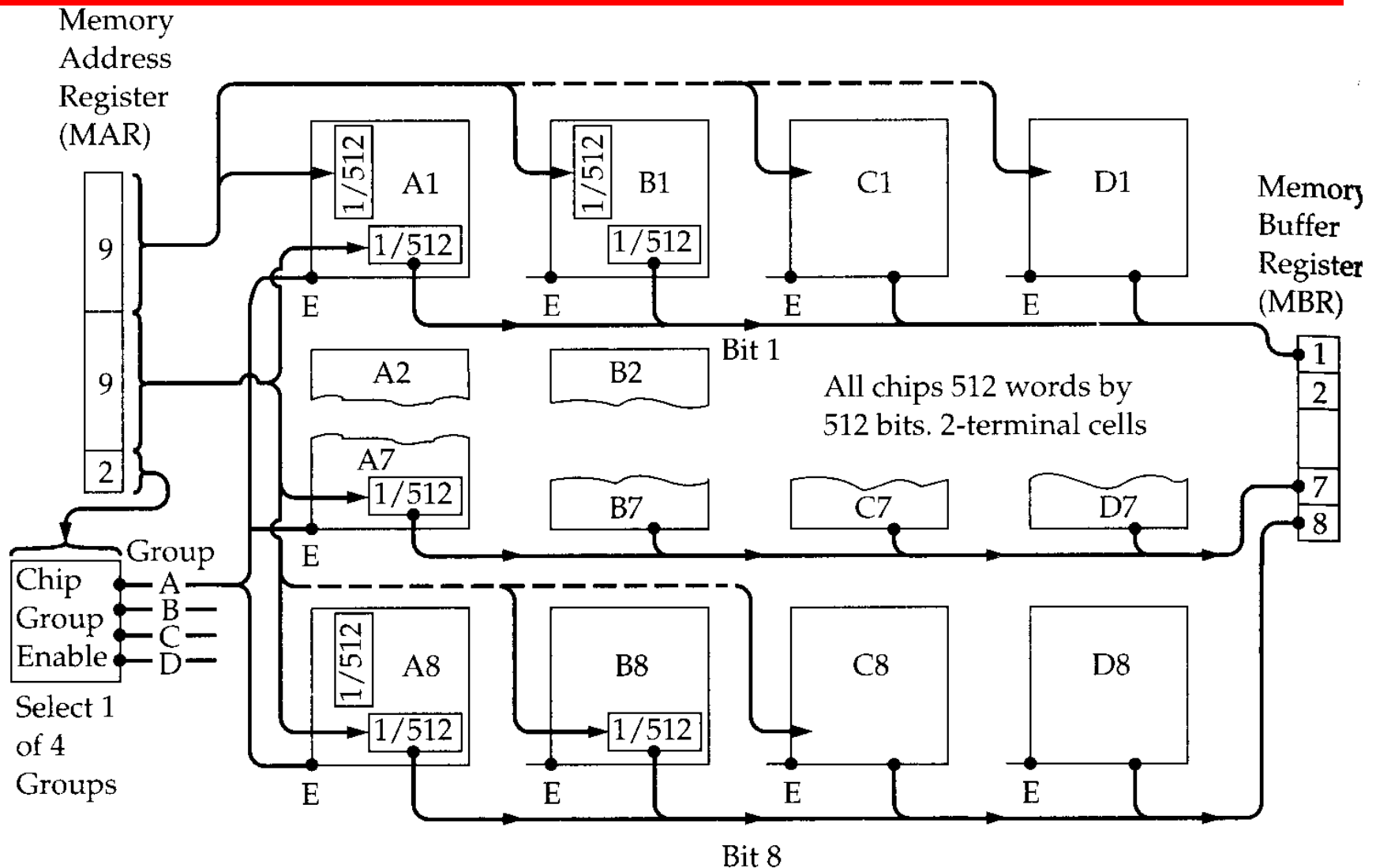


Module Organisation

- Basic organization of a 256KB chip
- 8 times a 512x512 bits chip
- ...For a 1 MB chip replicate 4 times this organization...



Module Organisation (1 MByte)

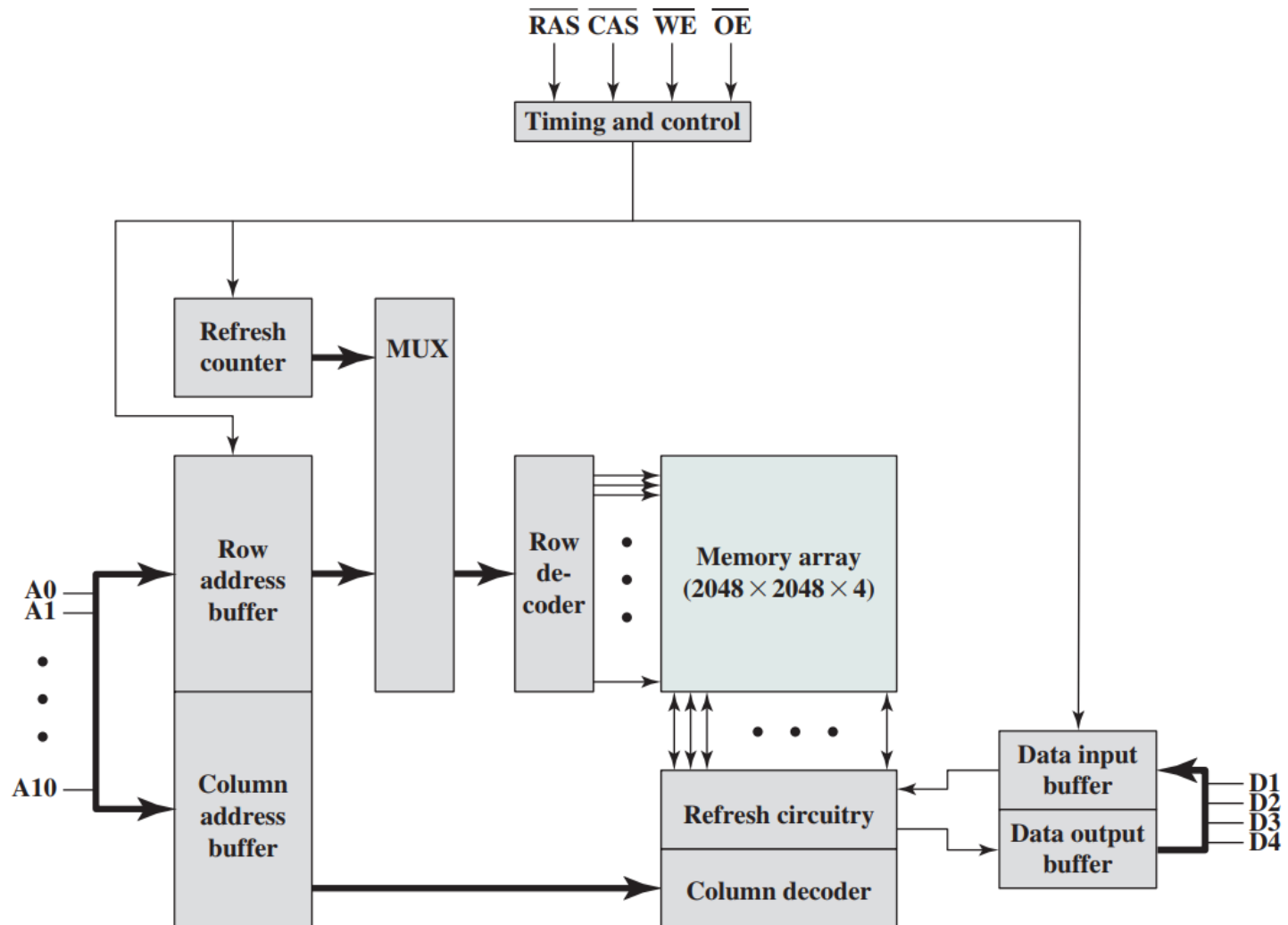


Organisation for larger sizes

- The larger the size the higher the number of address pins
- For 2^k words, k pins are needed
- A solution to reduce the number of address pins
 - Multiplex row address and column address
 - $k/2$ pins to address 2^k Bytes
 - Adding one more pin doubles range of values so x4 capacity



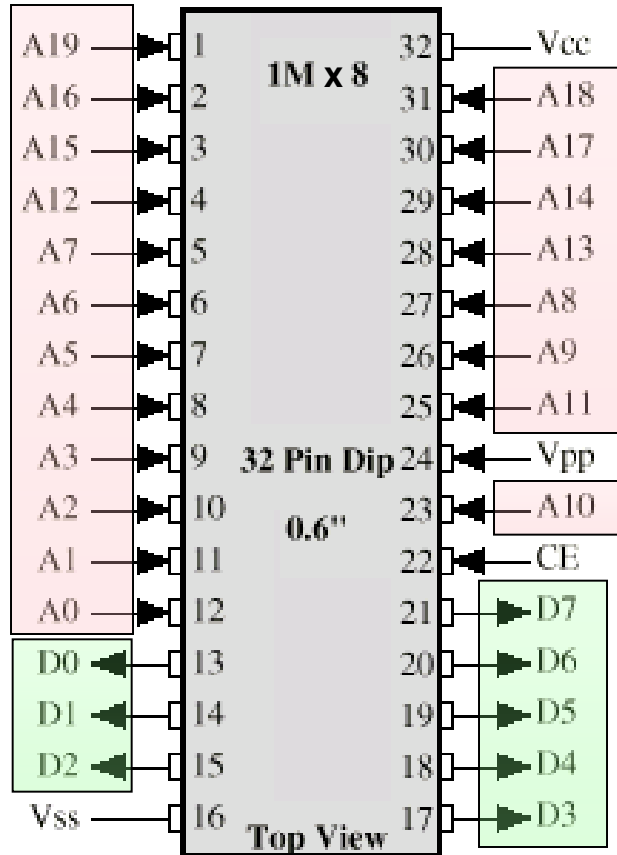
Typical 16 Mb DRAM (4M x 4)



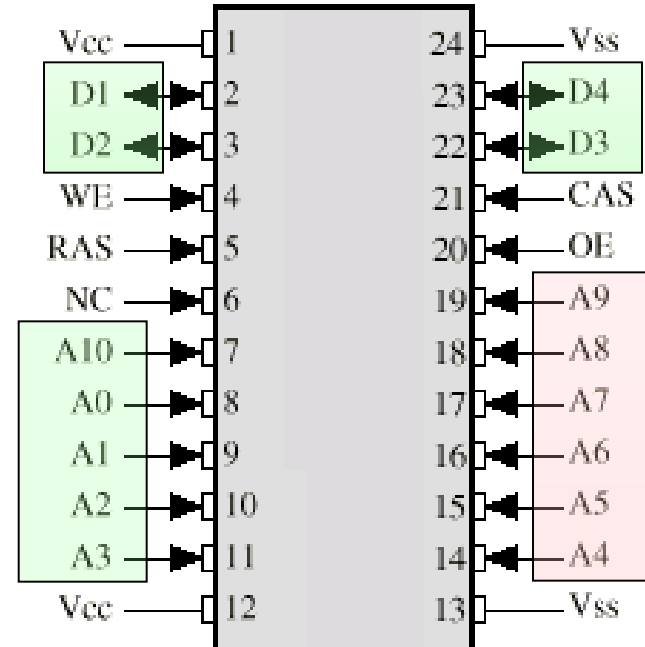
Refreshing (Dynamic RAM)

- Refresh circuit included on chip
- Disable chip
- Count through rows
- Read & Write back
- Takes time
- Slows down apparent performance

Packaging



(a) 8 Mbit EPROM

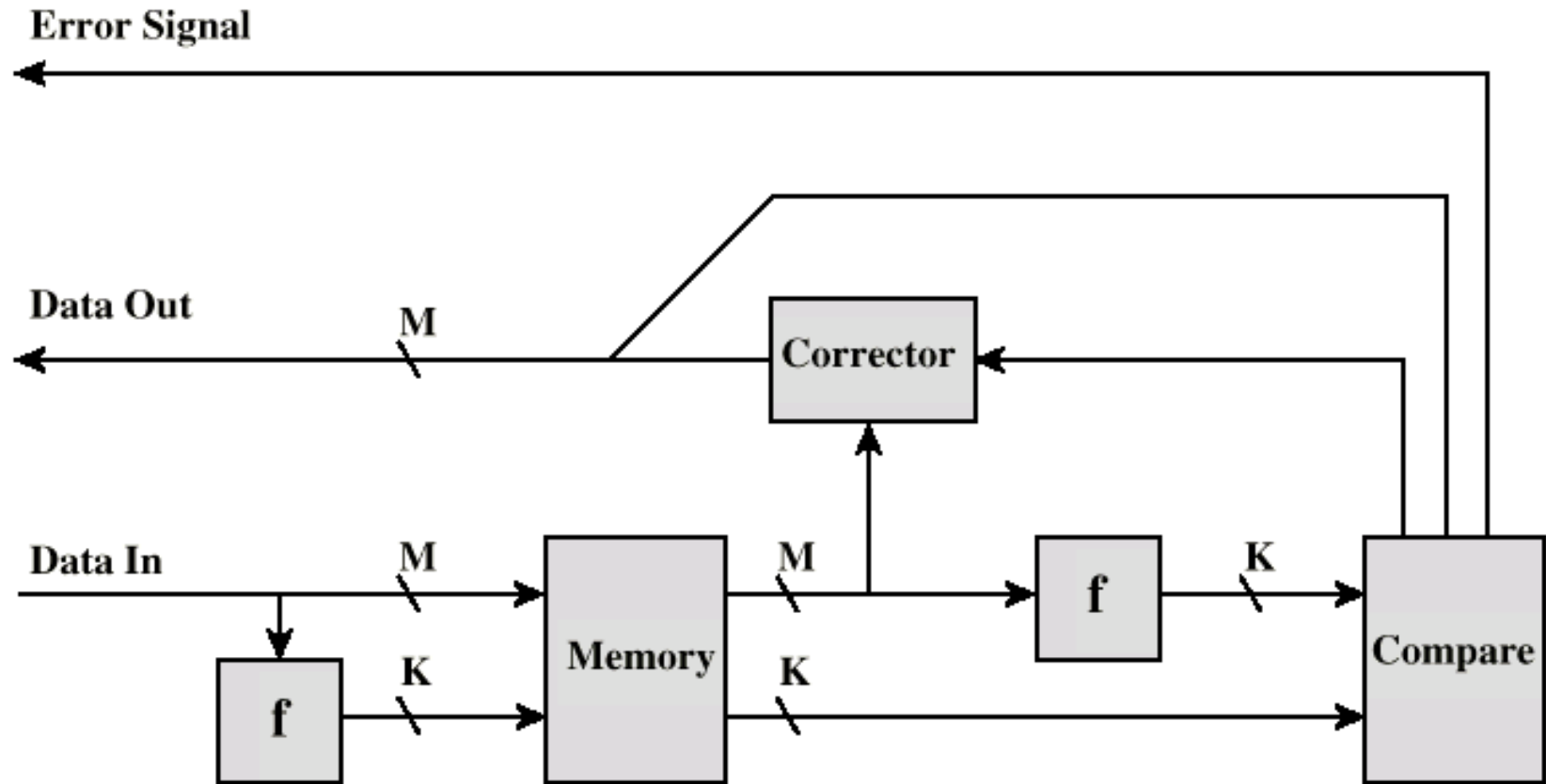


(b) 16 Mbit DRAM

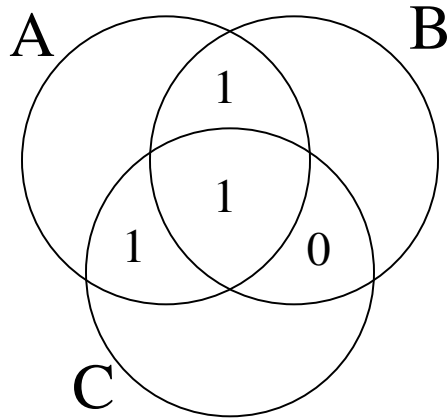
Error Correction

- Hard Failure
 - Permanent defect
- Soft Error
 - Random, non-destructive
 - No permanent damage to memory
- Detected using Hamming error correcting code
 - it is able to detect and correct 1-bit errors

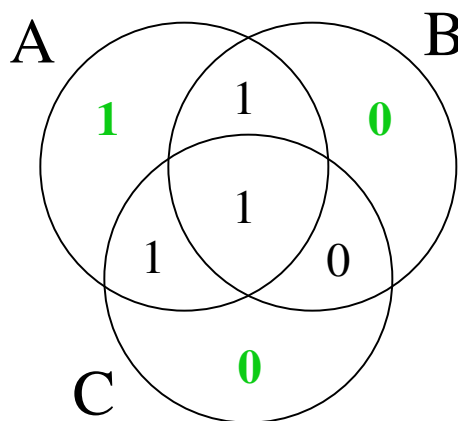
Error Correcting Code Function



A simple example of correction (1)

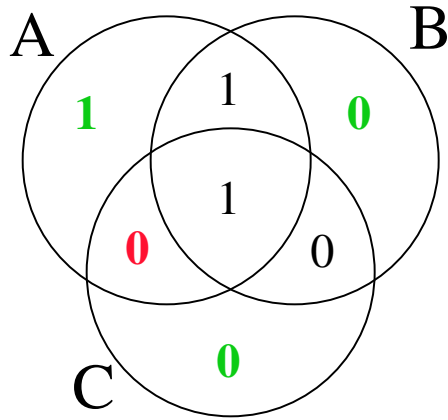


- Correcting errors in 4 bits words
- 3 control groups

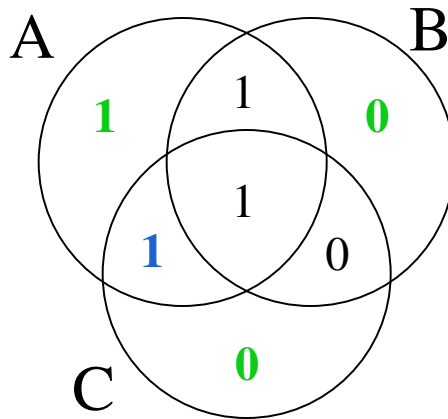


- In each control group add 1 parity bit

A simple example of correction (2)



- One of the bits change value



- Using control bit the right value is restored

Compare Circuit

- it takes two K-length binary strings X, Y as input
 - $X = X_K \dots X_1$
 - $Y = Y_K \dots Y_1$
- it returns a K-length binary string Z (syndrome)
 - $Z = Z_K \dots Z_1$
 - $Z_i = X_i \oplus Y_i$ for each $i = 1, \dots, K$
- $Z = 0 \dots 0$ means no error

Relation between M and K

- Z may assume 2^K values
- the value $Z=0\dots 0$ means no error
- the error may be in any bit among the $M+K$ bits
- it must be

$$2^K - 1 \geq M + K$$

Data bits (M)	Control Bits (K)	Additional Memory (%)
4	3	75
8	4	50
16	5	31,25
32	6	18,75
64	7	10,94
128	8	6,25
256	9	3,52

How to arrange the $M+K$ bits

- the $M+K$ bits are arranged so that
 - if Z contains a single bit equal to 1
 - error occurred in the corresponding control bit
 - if Z contains more than one bit equal to 1
 - error occurred in the i -th bit where i is the value (in binary) of Z

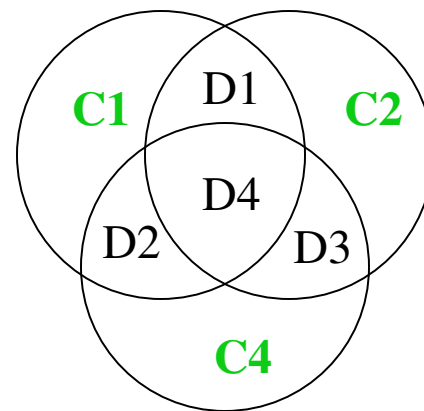
The case M=4

bit position	7	6	5	4	3	2	1
position number	111	110	101	100	011	010	001
data bits	D4	D3	D2		D1		
control bits				C4		C2	C1

$$C1 = D1 \oplus D2 \oplus D4$$

$$C2 = D1 \oplus D3 \oplus D4$$

$$C4 = D2 \oplus D3 \oplus D4$$



Exercise

- Design a Hamming error correcting code for 8-bit words
- See the textbook for the solution

Chapter 5

Part 2

External Memory

Types of External Memory

⌘ Magnetic Disk

- ☑ RAID

⌘ Magnetic Tape

⌘ Optical

- ☑ CD-ROM

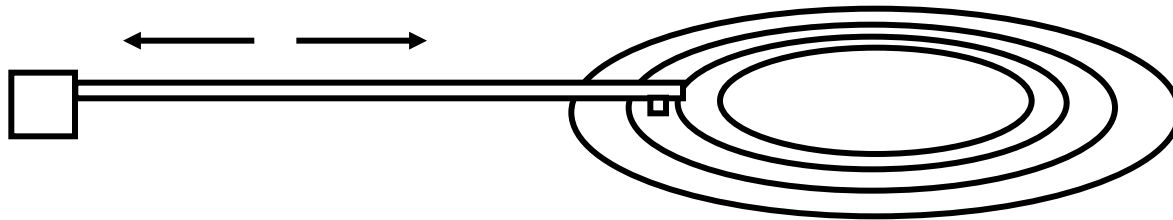
- ☑ CD-R

- ☑ CD-RW

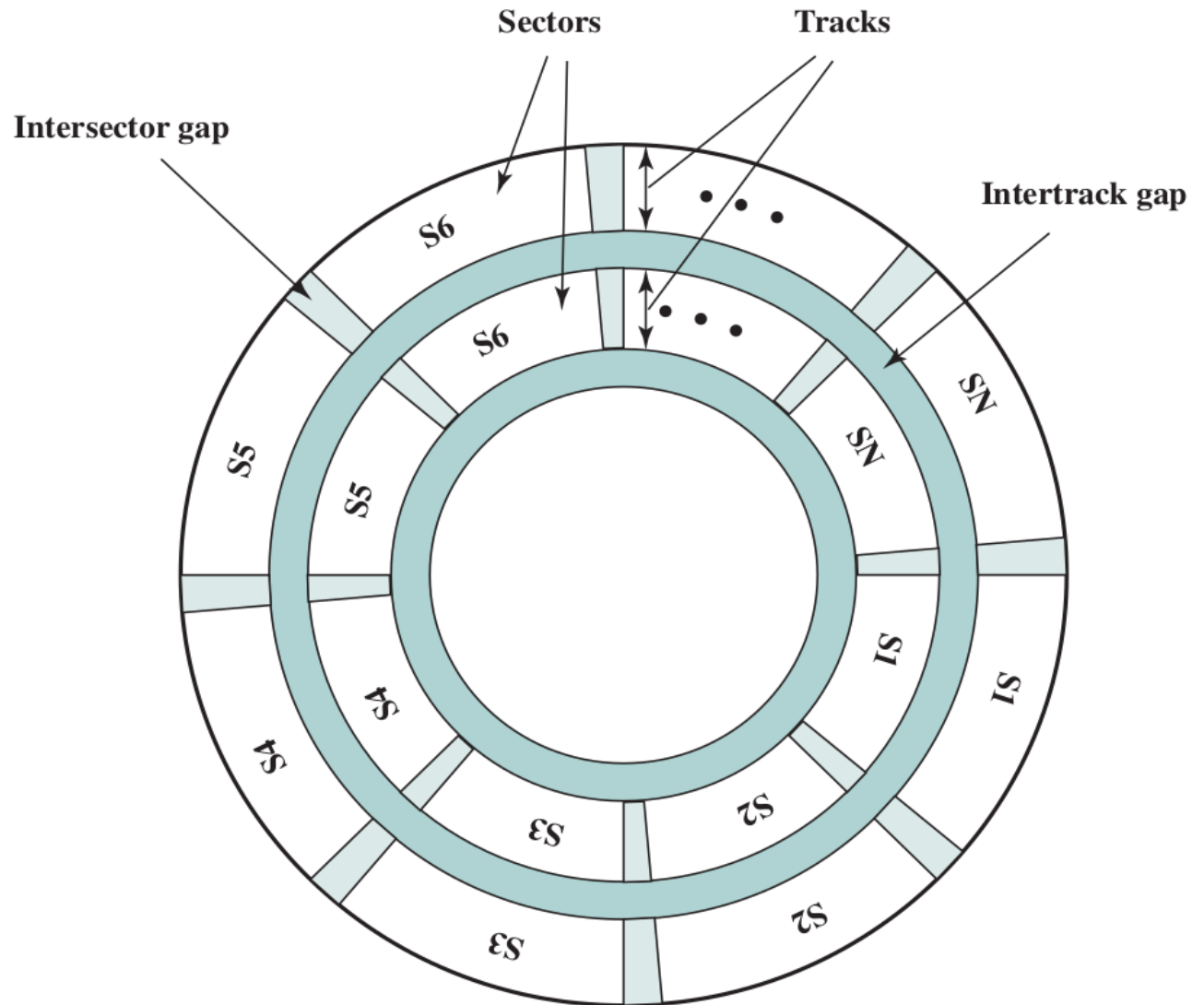
- ☑ DVD

Magnetic Disk

- ⌘ Metal or plastic disk coated, on one or both sides, with ***magnetizable material***
- ⌘ Data read and written through a magnetic head (coil) by means of **induction**



Disk Data Layout



Data Organization and Formatting

⌘ Concentric rings or tracks

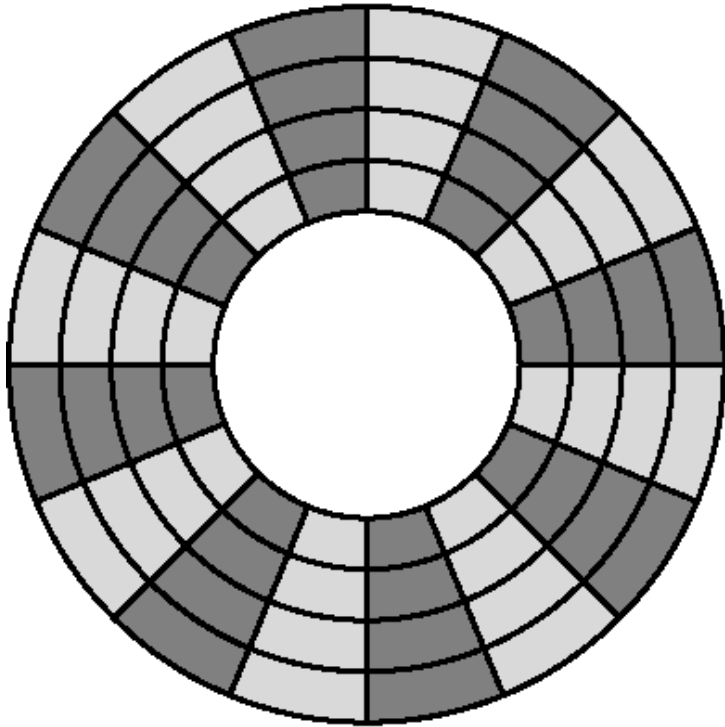
- ☒ Gaps between tracks
- ☒ Reduce gap to increase capacity
- ☒ Same number of bits per track
- ☒ Constant angular velocity

⌘ Tracks divided into sectors

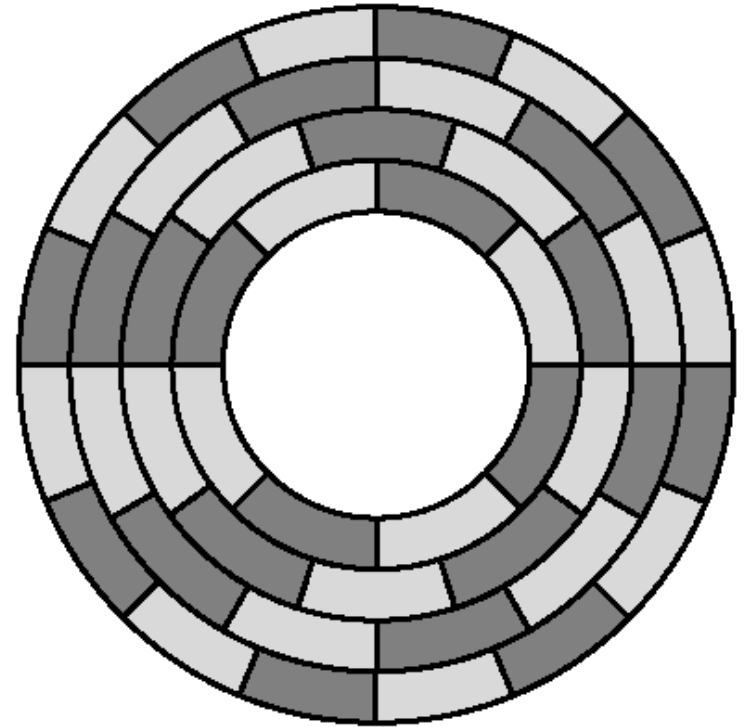
⌘ Data read/written in blocks

- ☒ Minimum block size is one sector
- ☒ May have more than one sector per block

Comparison of variable/fixed density



(a) Constant angular velocity

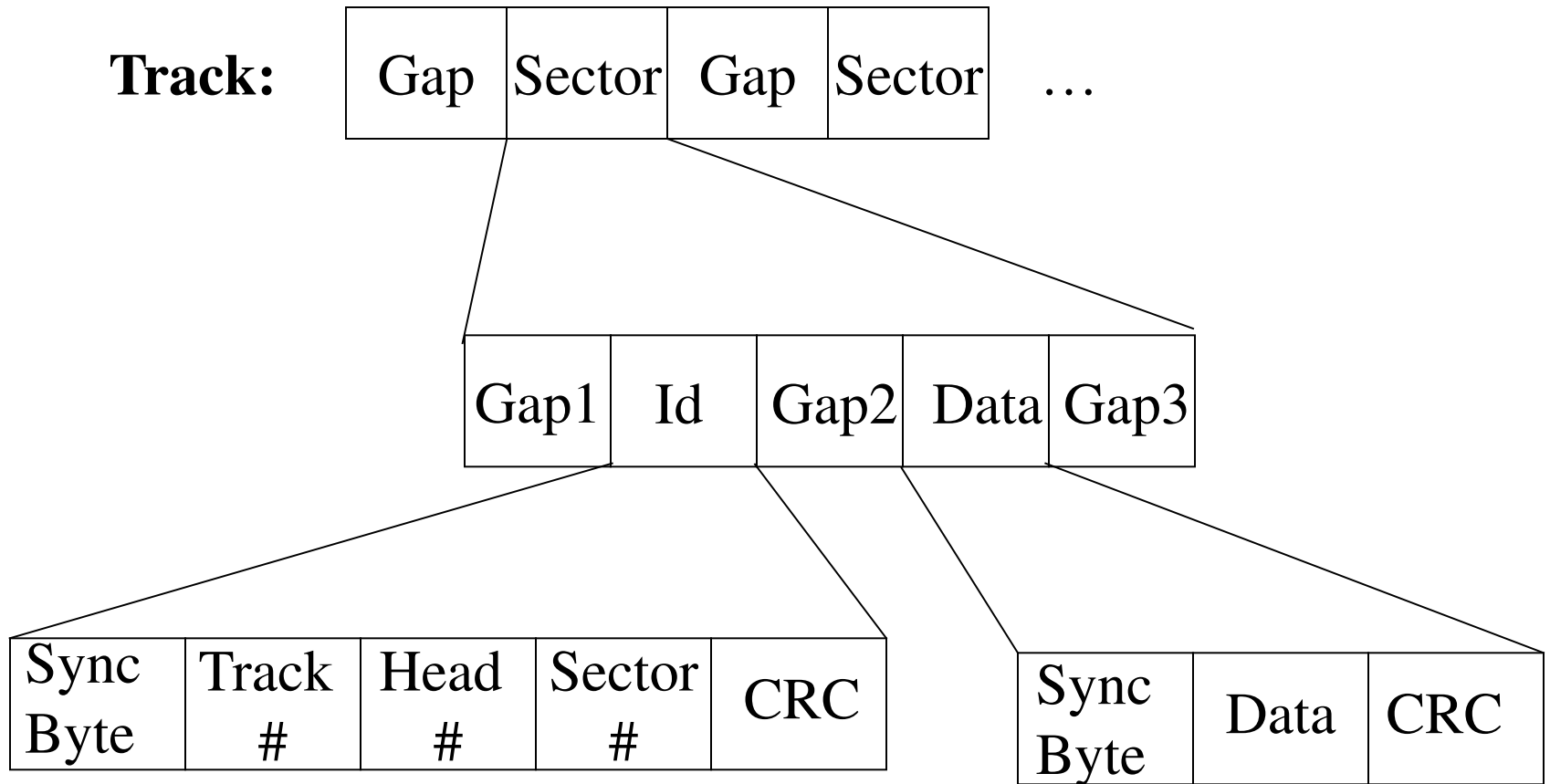


(b) Constant linear velocity

Finding Sectors

- ⌘ Must be able to identify start of track and sector
- ⌘ Format disk
 - ☐ Additional information not available to user
 - ☐ Marks tracks and sectors

An example format



Characteristics of magnetic disks

- ⌘ Single or double (usually) sided
- ⌘ Removable or fixed
- ⌘ Fixed or movable head
- ⌘ Single or multiple platter
- ⌘ Head mechanism
- ⌘ Speed

Removable or Not

⌘ Removable disk

- ☑ Can be removed from drive and replaced with another disk
- ☑ Provides unlimited storage capacity (by changing disk)
- ☑ Easy data transfer between systems

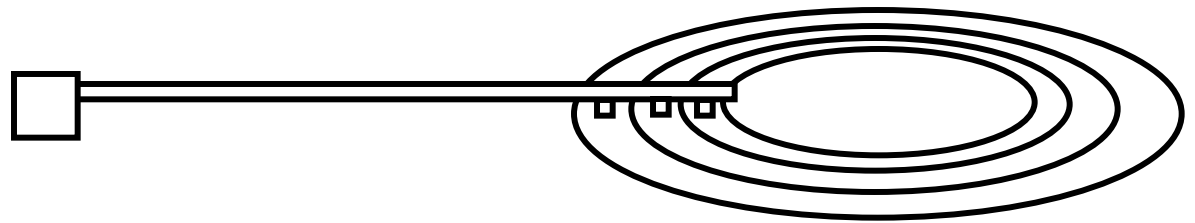
⌘ Nonremovable disk

- ☑ Permanently mounted in the drive

Fixed/Movable Head Disk

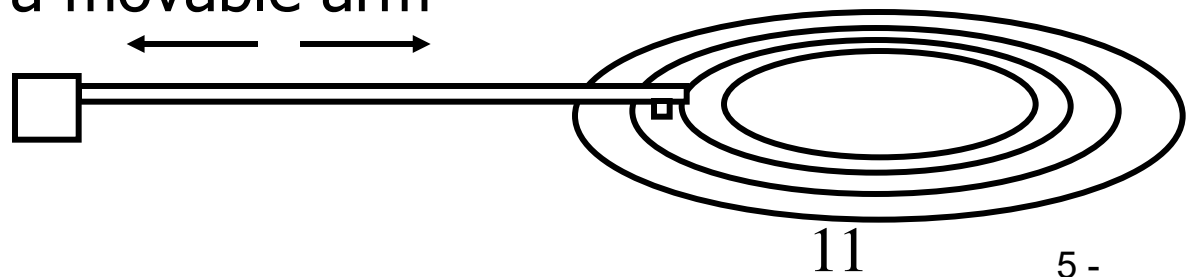
⌘ Fixed head

- ☑ One read/write head per track
- ☑ Heads mounted on a fixed arm



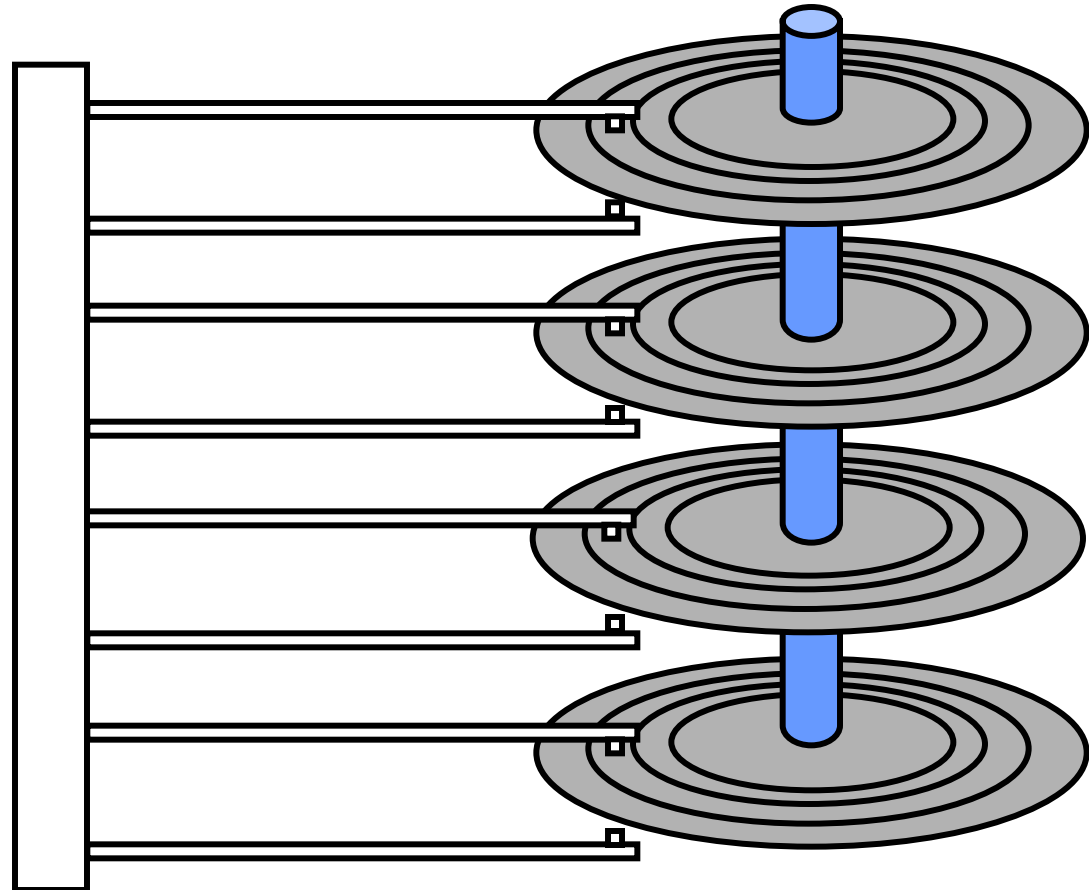
⌘ Movable head

- ☑ One read/write head per side
- ☑ Mounted on a movable arm



Multiple Platters

- ⌘ One head per side
- ⌘ Heads are joined and aligned
- ⌘ Aligned tracks on each platter form cylinders
- ⌘ Data is striped by cylinder
 - ⏏ reduces head movement
 - ⏏ increases speed (transfer rate)



Head mechanism

⌘ Contact

- ☒ Floppy

⌘ Fixed gap

⌘ Aerodynamic gap or flying head

- ☒ Winchester

Winchester Hard Disk (1)

- ⌘ Developed by IBM in Winchester (USA)
- ⌘ Sealed unit
- ⌘ One or more platters (disks)
- ⌘ Heads fly on boundary layer of air as disk spins
- ⌘ Very small head-to-disk gap
- ⌘ Getting more robust

Winchester Hard Disk (2)

- ⌘ Universal
- ⌘ Cheap
- ⌘ Fastest external storage
- ⌘ Getting larger all the time
 - ☑ Multiple Gigabyte now usual

Speed

⌘ Seek time

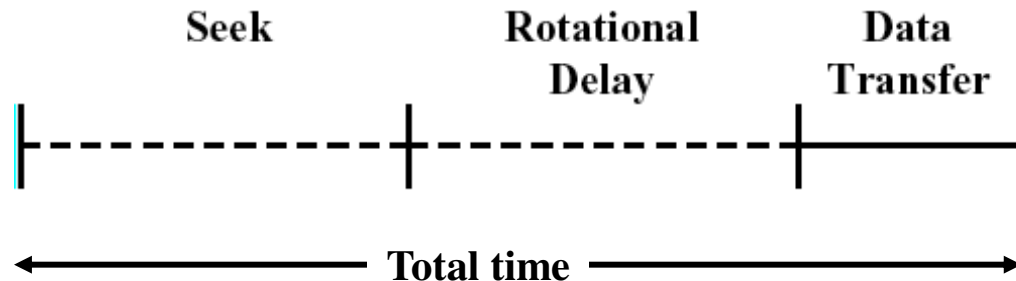
☐ Moving head to the right track

⌘ (Rotational) latency

☐ Waiting for data to rotate under head

⌘ Access time = Seek + Latency

⌘ Transfer rate: speed of copying bytes from disk



Cont'd

TRANSFER TIME The transfer time to or from the disk depends on the rotation speed of the disk in the following fashion:

$$T = \frac{b}{rN}$$

where

T = transfer time

b = number of bytes to be transferred

N = number of bytes on a track

r = rotation speed, in revolutions per second

Thus the total average access time can be expressed as

$$T_a = T_s + \frac{1}{2r} + \frac{b}{rN}$$

RAID

- ⌘ Redundant Array of **Independent** Disks
- ⌘ At least 7 different versions in common use (Not a hierarchy)
- ⌘ Set of physical disks viewed as single logical drive by the operating system
- ⌘ Data distributed (**striped**) across physical drives
- ⌘ Can use redundant capacity to store parity information and provide fault tolerance

Magnetic Tape

- ⌘ Only sequential access
- ⌘ Slower than magnetic and optical disks
- ⌘ Very very cheap
- ⌘ Backup and archive

Optical Storage: CD-ROM

- ⌘ Originally for audio
- ⌘ 650 Mbytes giving over 70 minutes audio
- ⌘ Polycarbonate coated with highly reflective coat, usually aluminum
- ⌘ Data stored as pits
- ⌘ Read by reflecting laser
- ⌘ Audio is single speed
 - ☒ Constant linear velocity
 - ☒ 1.2 m/s
 - ☒ Track (spiral) is 5.27km long
 - ☒ Gives 4391 seconds = 73.2 minutes

Random Access on CD-ROM

- ⌘ Difficult, due to constant density and single track
- ⌘ Move head to rough position
- ⌘ Set correct speed
- ⌘ Read address
- ⌘ Adjust to required location

CD-ROM Pros and Cons

- ⌘ Large capacity
- ⌘ Easy to mass produce
- ⌘ Removable
- ⌘ Robust

- ⌘ Expensive for small runs
- ⌘ Slower than magnetic disk
- ⌘ Read only

Other Optical Storage

⌘ CD-R (for Recordable)

- ☑ Writable, but ... Write Once Read Many (WORM)
- ☑ Now affordable
- ☑ Compatible with CD-ROM drives

⌘ CD-RW (for ReWritable)

- ☑ Erasable, hence writable many times (~1000)
- ☑ Different technology (phase change vs pit)
- ☑ Getting cheaper
- ☑ Mostly, but not always, CD-ROM drive compatible

DVD - Digital Video/Versatile Disk

- ⌘ Optical (CD-sized) disk with a very high capacity:
 - ☒ 4.7 GB per layer (smaller pits and closer tracks)
 - ☒ Up to 2 layers on each of the 2 sides (total 17 GB)
- ⌘ Drives are CD-ROM compatible
- ⌘ Also writable (DVD-R, DVD-RW), but not yet fully standardized