

# Integer Multiplication and Division

COE 301

Computer Organization

Prof. Muhamed Mudawar

College of Computer Sciences and Engineering  
King Fahd University of Petroleum and Minerals

## Presentation Outline

- ❖ Unsigned Integer Multiplication
- ❖ Signed Integer Multiplication
- ❖ Faster Integer Multiplication
- ❖ Integer Division
- ❖ Integer Multiplication and Division in MIPS

## Unsigned Integer Multiplication

❖ Paper and Pencil Example:

**Multiplicand**  $1100_2 = 12$   
**Multiplier**  $\times 1101_2 = 13$

$$\begin{array}{r} 1100 \\ 0000 \\ 1100 \\ 1100 \\ \hline 10011100_2 = 156 \end{array}$$

Binary multiplication is easy

$0 \times \text{multiplicand} = 0$

$1 \times \text{multiplicand} = \text{multiplicand}$

**Product**  $10011100_2 = 156$

❖  $m\text{-bit multiplicand} \times n\text{-bit multiplier} = (m+n)\text{-bit product}$

❖ Accomplished via **shifting** and **addition**

❖ Consumes more time and more chip area than addition

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## Sequential Unsigned Multiplication

❖ Initialize Product = 0

❖ Check each bit of the Multiplier

❖ If Multiplier bit = 1 then **Product = Product + Multiplicand**

❖ Rather than shifting the multiplicand to the left

Instead, **Shift the Product to the Right**

Has the same net effect and produces the same result

Minimizes the hardware resources

❖ One cycle per iteration (for each bit of the Multiplier)

✧ Addition and shifting can be done simultaneously

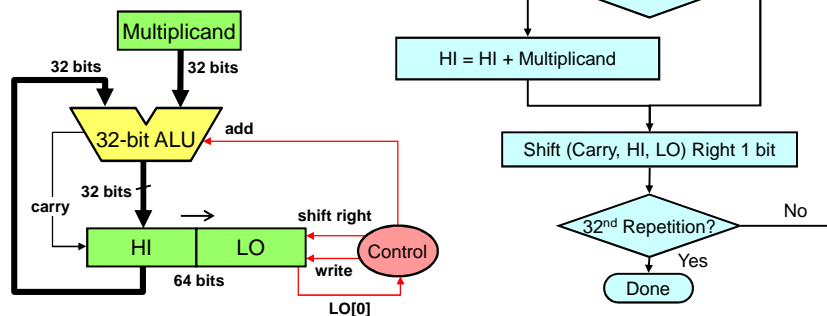
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## Sequential Multiplication Hardware

- ❖ Initialize HI = 0
- ❖ Initialize LO = Multiplier
- ❖ Final Product = HI and LO registers
- ❖ Repeat for each bit of Multiplier



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## Sequential Multiplier Example

- ❖ Consider:  $1100_2 \times 1101_2$ , Product =  $10011100_2$
- ❖ 4-bit multiplicand and multiplier are used in this example
- ❖ 4-bit adder produces a 5-bit sum (with carry)

Iteration		Multiplicand	Carry	Product = HI, LO
0	Initialize (HI = 0, LO = Multiplier)	1 1 0 0		0 0 0 0 1 1 0 1
1	LO[0] = 1 => ADD	1 1 0 0	0	1 1 0 0 1 1 0 1
	Shift Right (Carry, HI, LO) by 1 bit	1 1 0 0		0 1 1 0 0 1 1 0
2	LO[0] = 0 => Do Nothing			
	Shift Right (Carry, HI, LO) by 1 bit	1 1 0 0		0 0 1 1 0 0 1 1
3	LO[0] = 1 => ADD	1 1 0 0	0	1 1 1 1 0 0 1 1
	Shift Right (Carry, HI, LO) by 1 bit	1 1 0 0		0 1 1 1 1 0 0 1
4	LO[0] = 1 => ADD	1 1 0 0	1	0 0 1 1 1 0 0 1
	Shift Right (Carry, HI, LO) by 1 bit	1 1 0 0		1 0 0 1 1 1 0 0

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## Next ...

- ❖ Unsigned Integer Multiplication
- ❖ Signed Integer Multiplication
- ❖ Faster Integer Multiplication
- ❖ Integer Division
- ❖ Integer Multiplication and Division in MIPS

## Signed Integer Multiplication

- ❖ So far, we have dealt with unsigned integer multiplication
- ❖ First Attempt:
  - ✧ Convert multiplier and multiplicand into positive numbers
    - If negative then obtain the 2's complement and remember the sign
  - ✧ Perform unsigned multiplication
  - ✧ Compute the sign of the product
  - ✧ If product sign < 0 then obtain the 2's complement of the product
- ❖ Better Version:
  - ✧ Use the unsigned multiplication hardware
  - ✧ When shifting right, **extend the sign** of the product
  - ✧ If multiplier is negative, the **last step** should be a **subtract**

## Signed Multiplication (Pencil & Paper)

### ❖ Case 1: Positive Multiplier

$$\begin{array}{r}
 \text{Multiplicand} \quad 1100_2 = -4 \\
 \text{Multiplier} \quad \times 0101_2 = +5 \\
 \hline
 \text{Sign-extension} \left\{ \begin{array}{l} \boxed{1111}1100 \\ \boxed{11}1100 \end{array} \right. \\
 \hline
 \text{Product} \quad 11101100_2 = -20
 \end{array}$$

### ❖ Case 2: Negative Multiplier

$$\begin{array}{r}
 \text{Multiplicand} \quad 1100_2 = -4 \\
 \text{Multiplier} \quad \times 1101_2 = -3 \\
 \hline
 \text{Sign-extension} \left\{ \begin{array}{l} \boxed{1111}1100 \\ \boxed{11}1100 \end{array} \right. \\
 \quad \quad \quad 00100 \quad (\text{2's complement of } 1100) \\
 \hline
 \text{Product} \quad 00001100_2 = +12
 \end{array}$$

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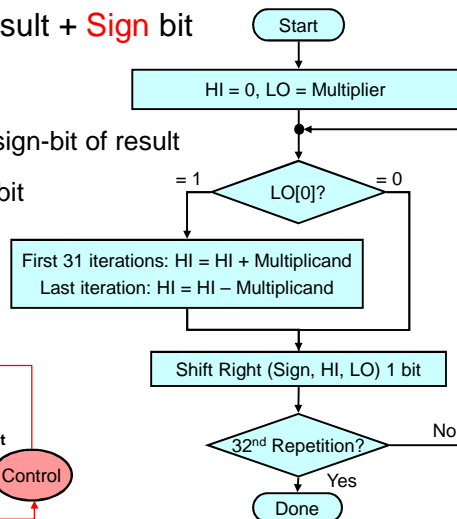
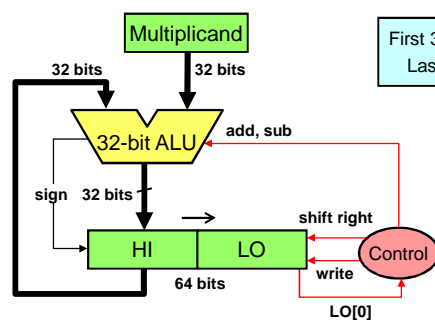
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## Sequential Signed Multiplier

### ❖ ALU produces 32-bit result + Sign bit

### ❖ Check for overflow

- ❖ No overflow → Extend sign-bit of result
- ❖ Overflow → Invert sign bit



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## Signed Multiplication Example

- ❖ Consider:  $1100_2$  (-4)  $\times$   $1101_2$  (-3), Product =  $00001100_2$
- ❖ Check for overflow: No overflow  $\rightarrow$  Extend sign bit
- ❖ Last iteration: add 2's complement of Multiplicand

Iteration		Multiplicand	Sign	Product = HI, LO
0	Initialize (HI = 0, LO = Multiplier)	1 1 0 0		0 0 0 0 1 1 0 <b>1</b>
1	LO[0] = 1 $\Rightarrow$ ADD		$\rightarrow +$	<b>1</b> 1 1 0 0 1 1 0 1
	Shift (Sign, HI, LO) right 1 bit	1 1 0 0		1 1 1 0 0 1 1 <b>0</b>
2	LO[0] = 0 $\Rightarrow$ Do Nothing			
	Shift (Sign, HI, LO) right 1 bit	1 1 0 0		1 1 1 1 0 0 1 <b>1</b>
3	LO[0] = 1 $\Rightarrow$ ADD		$\rightarrow +$	<b>1</b> 1 0 1 1 0 0 1 1
	Shift (Sign, HI, LO) right 1 bit	1 1 0 0		1 1 0 1 1 0 0 <b>1</b>
4	LO[0] = 1 $\Rightarrow$ SUB (ADD 2's compl)	0 1 0 0	$\rightarrow +$	<b>0</b> 0 0 0 1 1 0 0 1
	Shift (Sign, HI, LO) right 1 bit			0 0 0 0 1 1 0 0

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## Next ...

- ❖ Unsigned Integer Multiplication
- ❖ Signed Integer Multiplication
- ❖ **Faster Integer Multiplication**
- ❖ Integer Division
- ❖ Integer Multiplication and Division in MIPS

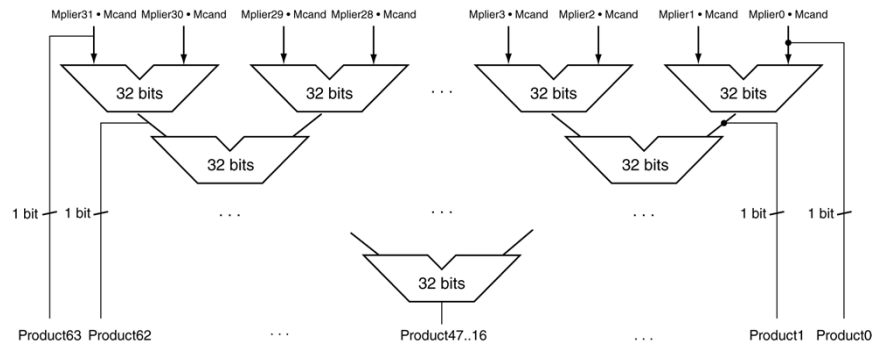
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## Faster Integer Multiplier

❖ Uses Multiple Adders (Cost vs. Performance)



■ Can be pipelined

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## Using Multiple Adders

❖ 32-bit adder for each bit of the multiplier

- ✧ AND multiplicand with each bit of multiplier
- ✧ Product = accumulated shifted sum

❖ Each adder produces a 33-bit output

- ✧ Most significant bit is a carry bit

❖ Array multiplier can be optimized

- ✧ Additions can be done in parallel
- ✧ Multiple-level tree reduction to produce final product
- ✧ Carry save adders reduce delays

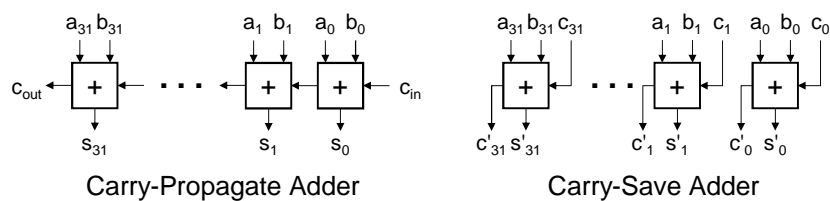
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## Carry Save Adders

- ❖ Used when adding multiple numbers (as in multipliers)
- ❖ All the bits of a carry-save adder work in parallel
  - ✧ The carry does not propagate as in a carry-propagate adder
  - ✧ This is why a carry-save is faster than a carry-propagate adder
- ❖ A carry-save adder has 3 inputs and produces two outputs
  - ✧ It adds 3 numbers and produces partial sum and carry bits



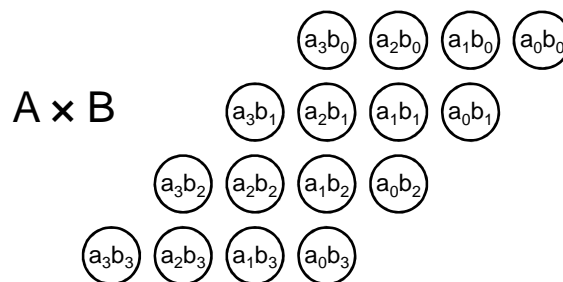
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## Tree Multiplier - 1 of 2

- ❖ Suppose we want to multiply two numbers A and B
  - ✧ Example on 4-bit numbers:  $A = a_3 a_2 a_1 a_0$  and  $B = b_3 b_2 b_1 b_0$
- ❖ Step 1: AND (multiply) each bit of A with each bit of B
  - ✧ Requires  $n^2$  AND gates and produces  $n^2$  product bits
  - ✧ Position of  $a_i b_j = (i+j)$ . For example, Position of  $a_2 b_3 = 2+3 = 5$



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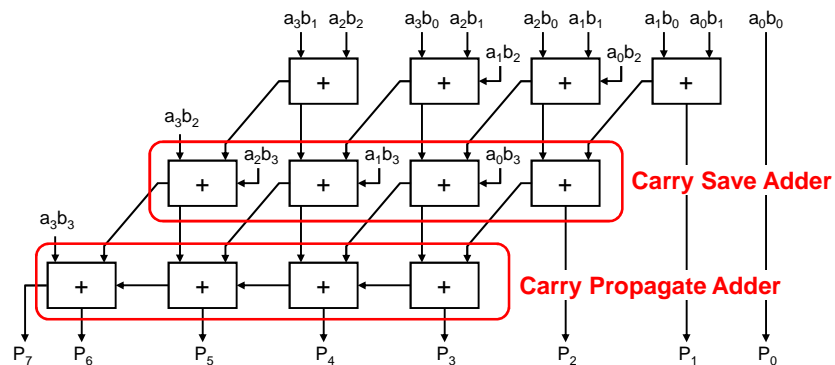


## Tree Multiplier - 2 of 2

Step 2: Use **carry save adders** to add the partial products

✧ Reduce the partial products to just two numbers

Step 3: Add last two numbers using a **carry-propagate adder**



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## Next ...

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- ❖ Faster Integer Multiplication
- ❖ **Integer Division**
- ❖ Integer Multiplication and Division in MIPS

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## Unsigned Division (Paper & Pencil)

**Divisor**  $1011_2$  **Quotient**  $10011_2 = 19$   
**Dividend**  $11011001_2 = 217$

$$\begin{array}{r}
 10011_2 \\
 1011_2 \overline{) 11011001_2} \\
 \underline{-1011} \phantom{000000} \\
 10 \phantom{000000} \\
 101 \phantom{00000} \\
 \underline{1010} \phantom{0000} \\
 10100 \phantom{000} \\
 \underline{-1011} \phantom{000} \\
 1001 \phantom{000} \\
 10011 \phantom{00} \\
 \underline{-1011} \phantom{00} \\
 1000_2 = 8
 \end{array}$$

Dividend =  
 Quotient  $\times$  Divisor  
 + Remainder  
 $217 = 19 \times 11 + 8$

Try to see how big a number can be subtracted, creating a digit of the quotient on each attempt

Binary division is accomplished via shifting and subtraction

**Remainder**  $1000_2 = 8$

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## Sequential Division

- ❖ Uses two registers: HI and LO
- ❖ Initialize: HI = Remainder = 0 and LO = Dividend
- ❖ Shift (HI, LO) LEFT by 1 bit (also Shift Quotient LEFT)
  - ✧ Shift the remainder and dividend registers together LEFT
  - ✧ Has the same net effect of shifting the divisor RIGHT
- ❖ Compute: Difference = Remainder – Divisor
- ❖ If (Difference  $\geq 0$ ) then
  - ✧ Remainder = Difference
  - ✧ Set Least significant Bit of Quotient
- ❖ Observation to Reduce Hardware:
  - ✧ LO register can be also used to store the computed Quotient

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## Sequential Division Hardware

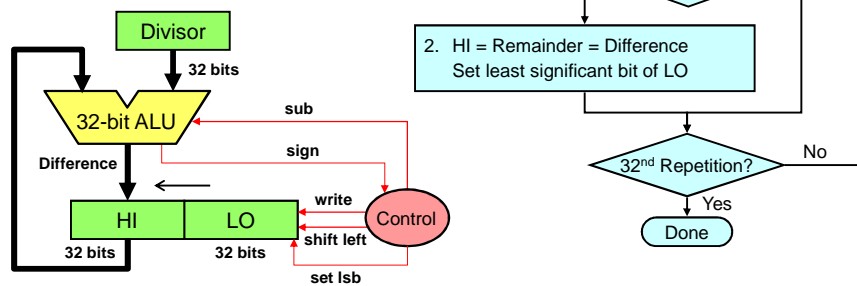
### ❖ Initialize:

❖ HI = 0, LO = Dividend

### ❖ Results:

❖ HI = Remainder

❖ LO = Quotient



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## Unsigned Integer Division Example

❖ Example:  $1110_2 / 0011_2$  (4-bit dividend & divisor)

❖ Result Quotient =  $0100_2$  and Remainder =  $0010_2$

❖ 4-bit registers for Remainder and Divisor (4-bit ALU)

Iteration		HI	LO	Divisor	Difference
0	Initialize	0 0 0 0	1 1 1 0	0 0 1 1	
1	1: Shift Left, Diff = HI - Divisor	0 0 0 1	← 1 1 0 0	0 0 1 1	1 1 1 0
	2: Diff < 0 => Do Nothing				
2	1: Shift Left, Diff = HI - Divisor	0 0 1 1	← 1 0 0 0	0 0 1 1	0 0 0 0
	2: Rem = Diff, set <b>lsb</b> of LO	0 0 0 0	1 0 0 1		
3	1: Shift Left, Diff = HI - Divisor	0 0 0 1	← 0 0 1 0	0 0 1 1	1 1 1 0
	2: Diff < 0 => Do Nothing				
4	1: Shift Left, Diff = HI - Divisor	0 0 1 0	← 0 1 0 0	0 0 1 1	1 1 1 1
	2: Diff < 0 => Do Nothing				

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## Signed Integer Division

- ❖ Simplest way is to remember the signs
- ❖ Convert the dividend and divisor to positive
  - ✧ Obtain the 2's complement if they are negative
- ❖ Do the unsigned division
- ❖ Compute the signs of the quotient and remainder
  - ✧ Quotient sign = Dividend sign XOR Divisor sign
  - ✧ Remainder sign = Dividend sign
- ❖ Negate the quotient and remainder if their sign is negative
  - ✧ Obtain the 2's complement to convert them to negative

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## Signed Integer Division Examples

1. **Positive** Dividend and **Positive** Divisor
  - ✧ Example:  $+17 / +3$       Quotient = +5    Remainder = +2
2. **Positive** Dividend and **Negative** Divisor
  - ✧ Example:  $+17 / -3$       Quotient = -5    Remainder = +2
3. **Negative** Dividend and **Positive** Divisor
  - ✧ Example:  $-17 / +3$       Quotient = -5    Remainder = -2
4. **Negative** Dividend and **Negative** Divisor
  - ✧ Example:  $-17 / -3$       Quotient = +5    Remainder = -2

The following equation must always hold:

$$\text{Dividend} = \text{Quotient} \times \text{Divisor} + \text{Remainder}$$

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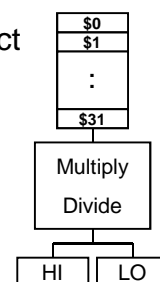
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## Next ...

- ❖ Unsigned Integer Multiplication
- ❖ Signed Integer Multiplication
- ❖ Faster Multiplication
- ❖ Integer Division
- ❖ Integer Multiplication and Division in MIPS

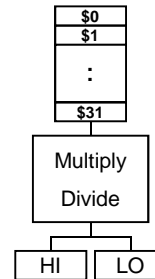
## Integer Multiplication in MIPS

- ❖ Multiply instructions
  - ❖ `mult $s1,$s2`      **Signed multiplication**
  - ❖ `multu $s1,$s2`      **Unsigned multiplication**
- ❖ 32-bit multiplication produces a 64-bit Product
- ❖ Separate pair of 32-bit registers
  - ❖ **HI = high-order 32-bit of product**
  - ❖ **LO = low-order 32-bit of product**
- ❖ MIPS also has a special `mul` instruction
  - ❖ `mul $s0,$s1,$s2`       **$\$s0 = \$s1 \times \$s2$**
  - ❖ **Put low-order 32 bits into destination register**
  - ❖ **HI & LO are undefined**



## Integer Division in MIPS

- ❖ Divide instructions
  - ✧ `div $s1,$s2`      **Signed division**
  - ✧ `divu $s1,$s2`      **Unsigned division**
- ❖ Division produces quotient and remainder
- ❖ Separate pair of 32-bit registers
  - ✧ **HI = 32-bit remainder**
  - ✧ **LO = 32-bit quotient**
  - ✧ If divisor is 0 then result is **unpredictable**
- ❖ Moving data from HI/LO to MIPS registers
  - ✧ `mfhi Rd` (move from HI to Rd)
  - ✧ `mflo Rd` (move from LO to Rd)



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## Integer Multiply/Divide Instructions

Instruction	Meaning	Format						
<code>mult Rs, Rt</code>	Hi, Lo = $Rs \times Rt$	$op^6 = 0$	$Rs^5$	$Rt^5$	0	0	0	0x18
<code>multu Rs, Rt</code>	Hi, Lo = $Rs \times Rt$	$op^6 = 0$	$Rs^5$	$Rt^5$	0	0	0	0x19
<code>mul Rd, Rs, Rt</code>	$Rd = Rs \times Rt$	0x1c	$Rs^5$	$Rt^5$	$Rd^5$	0	0	0x02
<code>div Rs, Rt</code>	Hi, Lo = $Rs / Rt$	$op^6 = 0$	$Rs^5$	$Rt^5$	0	0	0	0x1a
<code>divu Rs, Rt</code>	Hi, Lo = $Rs / Rt$	$op^6 = 0$	$Rs^5$	$Rt^5$	0	0	0	0x1b
<code>mfhi Rd</code>	$Rd = Hi$	$op^6 = 0$	0	0	$Rd^5$	0	0	0x10
<code>mflo Rd</code>	$Rd = Lo$	$op^6 = 0$	0	0	$Rd^5$	0	0	0x12

- ❖ Signed arithmetic: `mult`, `div` ( $Rs$  and  $Rt$  are signed)
  - ✧ LO = 32-bit low-order and HI = 32-bit high-order of multiplication
  - ✧ LO = 32-bit quotient and HI = 32-bit remainder of division
- ❖ Unsigned arithmetic: `multu`, `divu` ( $Rs$  and  $Rt$  are unsigned)
- ❖ **NO arithmetic exception** can occur

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## Integer to String Conversion

- ❖ Objective: convert an unsigned 32-bit integer to a string
- ❖ How to obtain the decimal digits of the number?
  - ✧ Divide the number by 10, Remainder = decimal digit (0 to 9)
  - ✧ Convert decimal digit into its ASCII representation ('0' to '9')
  - ✧ Repeat the division until the quotient becomes zero
  - ✧ Digits are computed **backwards** from least to most significant
- ❖ Example: convert 2037 to a string
  - ✧ Divide 2037/10    quotient = 203    remainder = 7    char = '7'
  - ✧ Divide 203/10    quotient = 20    remainder = 3    char = '3'
  - ✧ Divide 20/10    quotient = 2    remainder = 0    char = '0'
  - ✧ Divide 2/10    quotient = 0    remainder = 2    char = '2'

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## Integer to String Procedure

```
#-----
# int2str:  Converts an unsigned integer into a string
# Input:   $a0 = unsigned integer
# In/Out:  $a1 = address of string buffer (12 bytes)
#-----

int2str:
    move    $t0, $a0          # $t0 = dividend = unsigned integer
    li      $t1, 10           # $t1 = divisor = 10
    addiu   $a1, $a1, 11      # start at end of string buffer
    sb      $zero, 0($a1)     # store a NULL byte

convert:
    divu    $t0, $t1          # LO = quotient, HI = remainder
    mflo    $t0               # $t0 = quotient
    mfhi    $t2               # $t2 = remainder
    addiu   $t2, $t2, 0x30     # convert digit to a character
    addiu   $a1, $a1, -1       # point to previous byte
    sb      $t2, 0($a1)       # store digit character
    bnez    $t0, convert       # loop if quotient is not 0
    jr      $ra               # return to caller
```

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