

2W9

- Cristian's Algorithm

$$T_s = \text{Server time} + \frac{T_i - T_o - I}{2}$$

- Berkeley's algorithm

$$T_{\text{master}} = T_m + \text{average time differences}$$

↳ All slaves change time to master

$$\Delta T_n = T_m - T_n$$

- Averaging Algorithm (all process change)

$$\Delta M_0 = \frac{M_0 + M_1 + M_2 + M_3}{4} - M_0$$

- Lamports: Happens before, transitive

- WFL, resource allocation situations

2W10

- MUTEX: Critical section: coordinator, distributed, ring token request, grant, release, ok signal, check sent vs recv time only 1 process token

- Deadlock:
 - 1) Mutual Exclusion: Resource held by at most 1 process
 - 2) Hold and Wait: Process holds resource, waits for another
 - 3) Non-preemption: Resource cannot be pre-empted. Must be released
 - 4) Circular-wait: $P_0 \rightarrow P_1 \rightarrow \dots \rightarrow P_0$

Knots / cycles → one process/resources
↳ multiple resources

- Deadlock detection:

- Ostrich
false deadlock from delays → - Coordinator merge resource graphs
- CM1: probe (original, sender, recv)

> Election

- Bully:
 - Process sends election msg
 - ↳ if higher process, response, new leader
 - ↳ cont. until no response / timeout = new leader
- Ring:
 - Send election msg until reaches sender
 - ↳ Higher process ignores msg
 - ↓ coordinate msg sent

> Concurrency control

- lock based: Shared / Exclusive locks
 - read
 - write

can be read by other shared/unlocked
cannot be written by anything

RW conflicts

- 2-phase locking:
 - 1) Get all locks
 - 2) Release all locks



- Timestamp:
 - Read: $T \geq x_{\text{write}}$ time
 - Write: $T \geq x_{\text{write}}$ AND $T \geq x_{\text{read}}$

- Optimistic:
 - Check conflicts @ end
 - 1) read
 - 2) validate
 - 3) write
 - Validate:
 - 1) T_i finish 3 phase $\leftarrow T_j$
 - OR 2) T_i finish $\leftarrow T_j$ write
 - OR 3) T_i finish read $\leftarrow T_j$ read

7W 11

- ## > Solving Matrix input time computations
- ↳ Row based
 - ↳ Tile based

Question 6 (a) Calculations

1) a) $24 \times 8 \times 64 + 16 \times 16 \times 12.6 + 6462.7 \times 2$
 $= 28439 \text{ GFLOPS}$

b) $8 \times 64 + 6462.7 = 6974.7$ ✓

1) a) $24 \times 8 \times 10 + 12^2 \times 5.6 + 9322.46 \times 2$
 $= 21371.32$

b) $8 \times 10 + 9322.46 = 9402.46$ ✓

2) $10^{10} \times 32 \text{ bits}$

$$t = \frac{10^{10} \times 32}{8 \times 1024^3 \times 256} = 0.146s$$

2) $24000^2 \times 16 \text{ bits}$

$$t = \frac{24000^2 \times 16}{8 \times 1024^3 \times 32} = 0.0335s$$

3) Krishan's

S

$$\frac{75.95}{2} = 37.975$$

$$15:55 + 37.975$$

$$= 16:32.975$$

$$16:21 \rightarrow 16:32.975$$

$$= +11.975s$$

$$T_s = S + \frac{T_1 - T_0 - I}{2}$$

$$= 15:55 + \frac{16:21 - 15:05 - 0.05}{2}$$

$$3) = 27:11 + \frac{27:21 - 26:15 - 0.04}{2}$$

↙

32.98

$$= 27:43.980$$

$$43.98 - 21 = +22.985$$



4) Averaging Algorithm

$$\Delta M_1 = \frac{M_1 + M_2 + M_3 + M_4}{4} - M_1$$

$$1 = \frac{-20 - 10 + M_4}{4}$$

$$4 + 30 = M_4$$

$$M_4 = +34$$

$$M_4 = 12:00:34$$



$$4) 5 = \frac{-5:14 - 2:10 + M_4}{4}$$

$$5 = \frac{-444 + M_4}{4}$$

$$M_4 = 10:07:44$$

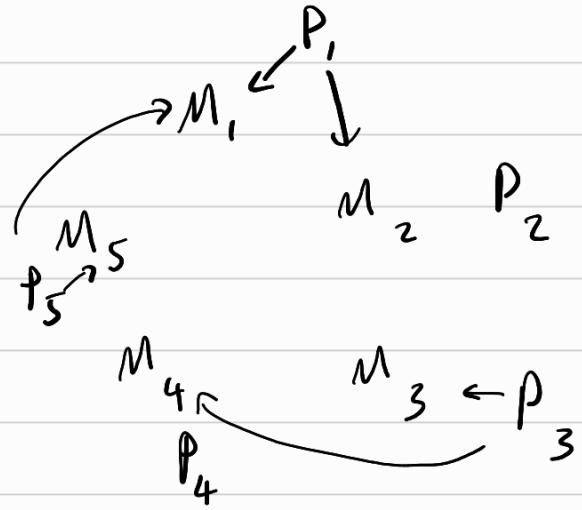
$$M_4 = +464s$$

$$+7:44$$



5) Ring Topology

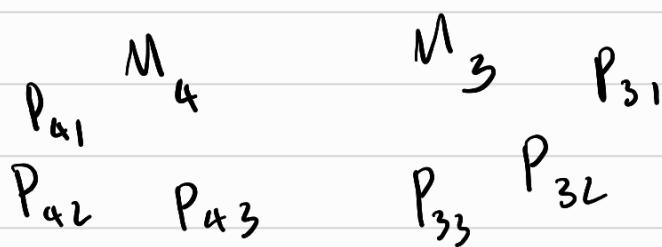
④



5)

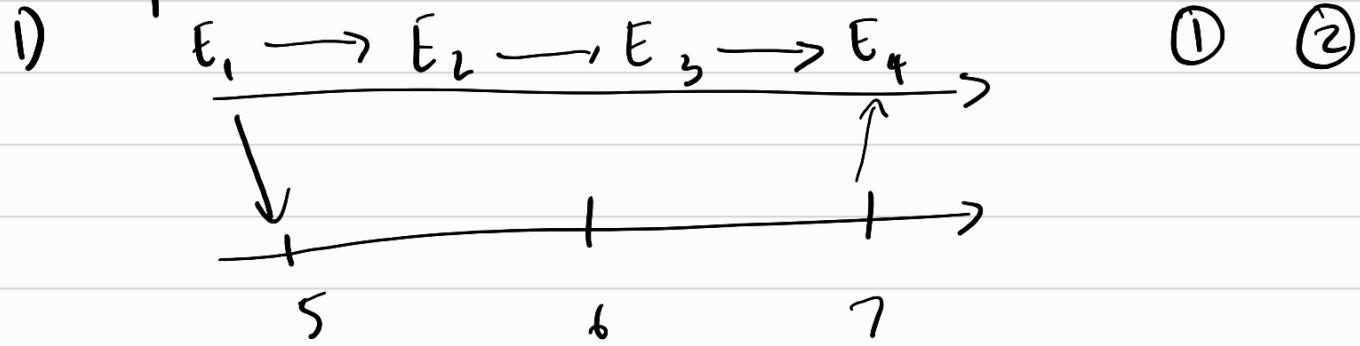


④

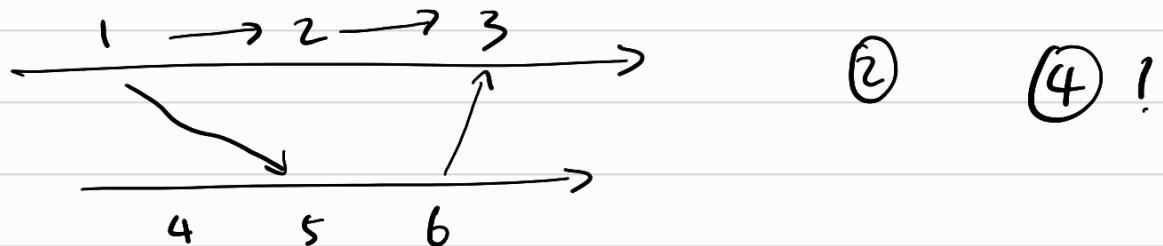


→ W10

Lamport's



1)



②

④ !

Elections

2) ④

2) ①

- Optimistic CC for read
- 3) $T_1 \xrightarrow{\text{r}} x_1 \xrightarrow{\omega} x_2$ $T \geq x_{\text{write}}$
 $T_2 \xrightarrow{\text{r}} x_2 \xrightarrow{\omega} x_3$ for write
 $T_3 \xrightarrow{\text{r}} x_3 \xrightarrow{\omega} x_1$ $T \geq x_{\text{write}}$
 $T \geq x_{\text{read}}$

③

3) ③ ④ ⑤ ⑥

4) Ahmadal's w/ comm

$$S = \frac{1}{0.5 + \frac{0.5}{2}} = 1.33$$

$$\text{orig time} = 80s$$

40s	20s	20s
parallel	serial	comm

$$1) \frac{40}{2} + 20 + 10 \\ = 50s$$

$$2) \frac{40}{2} + 20 + 5 \\ = 45s$$

$$3) 40 + 20 + 10 \\ = 70s$$

$$9) S = \frac{1}{0.5 + \frac{0.5}{2}} \\ = 1.33$$

$$\text{orig}_t = 100s$$

50s	25s	25s
P	S	comm

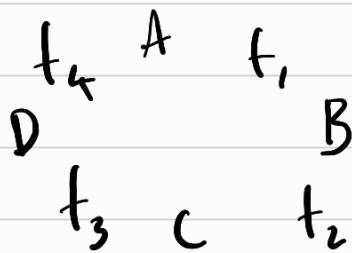
$$a) \frac{50}{2} + 25 + \frac{25}{2^2} \\ = 56.25s$$

$$b) \frac{50}{2} + 25 + \frac{25}{8} \\ = 53.125s$$

$$c) 50 + 25 + \frac{25}{4} \\ = 81.25s$$

> Mutex etc.

10)



③ ④ ⑥ ⑦
⑧

> WII

1) $\frac{600}{2^x} \leq 50 \quad x = 4$

MPI

2) a) $\frac{100}{10} = 5$ rows
 100 columns \rightarrow matrix still needs whole column to perform multiplication

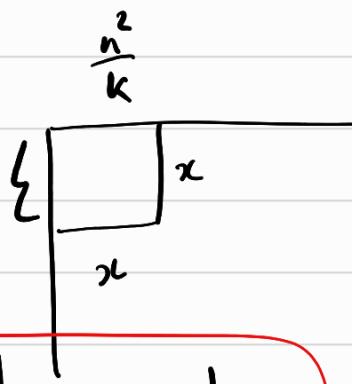
$$105 \text{ sends} \times 100 = 10500 \text{ ms}$$

b) $105 \times 100 \times 19 = 199500 \text{ ms}$
 \downarrow
machines

3) a) tile size = $\frac{100^2}{25} = 400$

$$x = \sqrt{\frac{n^2}{k}}$$

for matrix, we need



x rows and x columns
20 rows, 20 columns
 $\times 100 \times 100$ elements?

$$2000 + 2000 = 4000 \text{ ms}$$

$$4000 \times 24 = 96000 \text{ ms}$$

> W12 +

(1) input vector = 40000×15000

