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How I became seduced by univalent foundations

2022 Fields Medal Symposium: Akshay Venkatesh

Plan



1. A reintroduction to proofs
2. On the art of giving the same name to different things
3. Contractibility as uniqueness
4. Computer proof assistants that compute



1

A reintroduction to proofs

Conjunction and disjunction



Compare a traditional introduction to the logical operators “and” \wedge and “or” \vee using truth tables with the following:

Conjunction \wedge is the logical operator defined by the rules:

- \wedge intro: If p is true and q is true, then $p \wedge q$ is true.
- \wedge elim₁: If $p \wedge q$ is true, then p is true.
- \wedge elim₂: If $p \wedge q$ is true, then q is true.

Disjunction \vee is the logical operator defined by the rules:

- \vee intro₁: If p is true, then $p \vee q$ is true.
- \vee intro₂: If q is true, then $p \vee q$ is true.
- \vee elim: If $p \vee q$ is true, and if r can be derived from p and from q , then r is true.

— from Clive Newstead’s *An Infinite Descent into Pure Mathematics*.

Implication



The **introduction rules** explain how to prove a proposition involving a particular connective, while the **elimination rules** explain how to use a hypothesis involving a particular connective.

Implication \Rightarrow is the logical operator defined by the rules:

- \Rightarrow **intro**: If q can be derived from the assumption that p is true, then $p \Rightarrow q$ is true.
- \Rightarrow **elim**: If $p \Rightarrow q$ is true and p is true, then q is true.

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Theorem. For any propositions p , q , and r , $((p \Rightarrow q) \wedge (q \Rightarrow r)) \Rightarrow (p \Rightarrow r)$.

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Universal and existential quantification



Let $p : X \rightarrow \{\perp, \top\}$ be an X -indexed family of propositions, a **predicate** on X .

Universal quantification $\forall x \in X, p(x)$ is the logical formula defined by the rules:

- **\forall intro**: If $p(x)$ can be derived from the assumption that x is an arbitrary element of X , then $\forall x \in X, p(x)$ is true.
- **\forall elim**: If $\forall x \in X, p(x)$ is true and $a \in X$, then $p(a)$ is true.

Existential quantification $\exists x \in X, p(x)$ is the logical formula defined by the rules:

- **\exists intro**: If $a \in X$ and $p(a)$ is true, then $\exists x \in X, p(x)$ is true.
- **\exists elim**: If $\exists x \in X, p(x)$ is true and q can be derived from the assumption that $p(a)$ is true for some $a \in X$, then q is true.

Dependent type theory



Dependent type theory is a formal system for mathematical statements and proofs that has the following primitive notions:

- types, e.g., \mathbb{N} , \mathbb{Q} , Group

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- terms, e.g., $17 : \mathbb{N}$, $\sqrt{2} : \mathbb{R}$, $K_4 : \text{Group}$
- dependent types, e.g., $\mathbb{R}^\bullet : \mathbb{N} \rightarrow \text{Type}$, $\text{is-prime} : \mathbb{N} \rightarrow \text{Type}$, $\text{Mat}_{\bullet \times \bullet} : \mathbb{N} \rightarrow \mathbb{N} \rightarrow \text{Type}$

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- dependent terms, e.g., $\vec{0}^\bullet : \prod_{n:\mathbb{N}} \mathbb{R}^n$, $!_\bullet : \prod_{n:\mathbb{N}} \text{Mat}_{n,n}$, $S_\bullet : \prod_{n:\mathbb{N}} \text{Group}$

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all of which can occur in an arbitrary context of variables from previously-defined types.

In a mathematical statement of the form “Let ...be ...then ...” The stuff following the “let” likely declares the names of the variables in the context described after the “be”, while the stuff after the “then” most likely describes a type or term in that context.

Products and coproducts



Given types A and B , the **product type** $A \times B$ is governed by the rules:

- \times **intro**: given terms $a : A$ and $b : B$ there is a term $(a, b) : A \times B$
- \times **elim**: given a term $p : A \times B$ there are terms $\text{pr}_1 p : A$ and $\text{pr}_2 p : B$

plus computation rules that relate pairings and projections.

Given types A and B , the **coproduct type** $A + B$ is governed by the rules:

- $+$ **intro**: given a term $a : A$, there is a term $\text{inl}a : A + B$, and
given a term $b : B$, there is a term $\text{inr}b : A + B$
- $+$ **elim**: given a family of types $C : (A + B) \rightarrow \text{Type}$, if there are dependent terms
 $c : \prod_{a:A} C(\text{inl}a)$ and $d : \prod_{b:B} C(\text{inr}b)$, then there is a term $e : \prod_{x:A+B} C(x)$

plus computation rules that relate the inclusions and the elimination.

Functions in set theory vs functions in type theory



In set theory, a function $f: X \rightarrow Y$ is a subset $\Gamma_f \subset X \times Y$ with the property that

$$\forall x \in X, \exists! y \in Y, (x, y) \in \Gamma_f.$$

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Given types A and B , the function type $A \rightarrow B$ is governed by the rules:

- \rightarrow intro: if in the context of a variable $x: A$ there is a term $b_x: B$,
there is a term $\lambda x. b_x: A \rightarrow B$
- \rightarrow elim: given terms $f: A \rightarrow B$ and $a: A$, there is a term $f(a): B$

plus computation rules that relate λ -abstractions and evaluations.

Dependent functions and dependent sums



Let $B : A \rightarrow \text{Type}$ be a family of types over a type A .

The dependent function type $\prod_{x:A} B(x)$ is governed by the rules:

- Π_{intro} : if in the context of a variable $x : A$ there is a term $b_x : B$
there is a term $\lambda x. b_x : \prod_{x:A} B(x)$
- Π_{elim} : given terms $f : \prod_{x:A} B(x)$ and $a : A$ there is a term $f(a) : B(a)$

plus computation rules that relate λ -abstractions and evaluations.

The dependent sum type $\sum_{x:A} B(x)$ is governed by the rules:

- Σ_{intro} : if there are terms $a : A$ and $b : B(a)$, there is a term $(a, b) : \sum_{x:A} B(x)$
- Σ_{elim} : given a type family $C : \sum_{x:A} B(x) \rightarrow \text{Type}$, if there is a term $c : \prod_{a:A} \prod_{b:B(a)} C(a, b)$, there is a term $d : \prod_{z:\sum_{x:A} B(x)} C(z)$

plus computation rules that relate pairings and eliminations.

The natural numbers in set theory



Recall the von Neumann and Zermelo constructions of the natural numbers in set theory:

$$3_{\text{vN}} := \{\{\}, \{\{\}\}, \{\{\}, \{\{\}\}\}\} \quad 3_{\text{Z}} := \{\{\{\{\}\}\}\}$$

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By **Dedekind's categoricity theorem**, the natural numbers \mathbb{N} are characterized by **Peano's postulates**:

- There is a natural number $0 \in \mathbb{N}$.
- Every natural number $n \in \mathbb{N}$ has a successor $\text{suc } n \in \mathbb{N}$.
- 0 is not the successor of any natural number.
- No two natural numbers have the same successor.
- The **principle of mathematical induction**: for all $P : \mathbb{N} \rightarrow \{\perp, \top\}$

$$P(0) \Rightarrow (\forall k \in \mathbb{N}, P(k) \Rightarrow P(\text{suc } k)) \Rightarrow (\forall n \in \mathbb{N}, P(n))$$

The natural numbers in dependent type theory



The natural numbers type \mathbb{N} is governed by the rules:

- $\mathbb{N}_{\text{intro}}$: there is a term $0 : \mathbb{N}$ and for any term $n : \mathbb{N}$ there is a term $\text{succ } n : \mathbb{N}$

The elimination rule strengthens the principle of mathematical induction by replacing the predicate $P : \mathbb{N} \rightarrow \{\perp, \top\}$ by an arbitrary family of types $P : \mathbb{N} \rightarrow \text{Type}$.

- \mathbb{N}_{elim} : for any type family $P : \mathbb{N} \rightarrow \text{Type}$, to prove $p : \prod_{n:\mathbb{N}} P(n)$ it suffices to prove $p_0 : P(0)$ and $p_s : \prod_{k:\mathbb{N}} P(k) \rightarrow P(\text{succ } k)$. That is

$$\mathbb{N}_{\text{ind}} : P(0) \rightarrow \left(\prod_{k \in \mathbb{N}} P(k) \rightarrow P(\text{succ } k) \right) \rightarrow \left(\prod_{n \in \mathbb{N}} P(n) \right)$$

Computation rules establish that p is defined recursively from p_0 and p_s .

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Note the other two Peano postulates are missing because they are provable!



2

On the art of giving the same name to different things

Identity types



The following rules for **identity types** were developed by Martin-Löf:

Given a type A and terms $x, y : A$, the **identity type** $x =_A y$ is governed by the rules:

- **=intro**: given a type A and term $x : A$ there is a term $\text{refl}_x : x =_A x$

The elimination rule for the identity type defines an induction principle analogous to recursion over the natural numbers: it provides sufficient conditions for which to define a dependent function out of the identity type family.

- **=elim**: for any type family $P(x, y, p)$ over $x, y : A$ and $p : x =_A y$, to prove $P(x, y, p)$ for all x, y, p it suffices to assume y is x and p is refl_x . That is

$$\text{=ind} : \left(\prod_{x:A} P(x, x, \text{refl}_x) \right) \rightarrow \left(\prod_{x,y:A} \prod_{p:x=_A y} P(x, y, p) \right)$$

A computation rule establishes that the proof of $P(x, x, \text{refl}_x)$ is the given one.

The homotopical interpretation of dependent type theory



Identity types can be iterated: given $x, y : A$ and $p, q : x =_A y$ there is a type $p =_{x=_A y} q$.

Does this type always have a term? In other words, are identity proofs unique?

The homotopical interpretation of dependent type theory



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Theorem (Lumsdaine, Garner–van den Berg after Hofmann–Streicher). The terms belonging to the iterated identity types of any type A form an ∞ -groupoid.

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The ∞ -groupoid structure of A has

- terms $x : A$ as objects,
- identifications $p : x =_A y$ as 1-morphisms aka **paths**,
- higher identifications $h : p =_{x=_A y} q$ as 2-morphisms aka **homotopies**, ...

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The required structures are proven from **=elim**:

- **constant paths** (reflexivity) $\text{refl}_x : x = x$
- **reversal** (symmetry) $p : x = y$ yields $p^{-1} : y = x$
- **concatenation** (transitivity) $p : x = y$ and $q : y = z$ yield $p * q : x = z$

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and furthermore concatenation is associative and unital, the associators are coherent ...

Leibniz' indiscernibility of identicals as path lifting



The elimination rule for the identity types is also called **path induction**:

- **=elim**: for any type family $P(x, y, p)$ over $x, y : A$ and $p : x =_A y$, to prove $P(x, y, p)$ for all x, y, p it suffices to assume y is x and p is refl_x . That is

$$\text{=ind} : \left(\prod_{x:A} P(x, x, \text{refl}_x) \right) \rightarrow \left(\prod_{x,y:A} \prod_{p:x=_A y} P(x, y, p) \right)$$

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In first-order logic, one axiom for the equality relation is **indiscernibility of identicals**:

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Proposition. Let $P : A \rightarrow \text{Type}$ be any family of types. For any $x, y : A$ and $p : x =_A y$, there is a transport function $\text{tr}_{P,p} : P(x) \rightarrow P(y)$.

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Proof: By **=elim**, to define $\text{tr}_P : \prod_{x,y:A} \prod_{p:x=_A y} P(x) \rightarrow P(y)$ it suffices to define a term of type $\prod_{x:A} P(x) \rightarrow P(x)$, for which we take the identity function $\lambda x. \lambda q. q$. \square

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Corollary. For any $P : A \rightarrow \text{Type}$, $x, y : A$, and $p : x =_A y$, then $P(x) \simeq P(y)$.

The homotopy type theoretic Rosetta stone



type theory	logic	set theory	homotopy theory
A	proposition	set	space
$x : A$	proof	element	point
$\emptyset, 1$	\perp, \top	$\{\}, \{\{\}\}$	$\emptyset, *$
$A \times B$	A and B	set of pairs	product space
$A + B$	A or B	disjoint union	coproduct
$A \rightarrow B$	A implies B	set of functions	function space
$P : A \rightarrow \text{Type}$	predicate	family of sets	fibration
$f : \prod_{x:A} P(x)$	conditional proof	family of elements	section
$\prod_{x:A} P(x)$	$\forall x. P(x)$	product	space of sections
$\sum_{x:A} P(x)$	$\exists x. P(x)$	disjoint union	total space
$p : x =_A y$	proof of equality	$x = y$	path from x to y
$\sum_{x,y:A} x =_A y$	equality relation	diagonal	path space for A

The homotopy interpretation, developed by [Awodey–Warren](#) and [Voevodsky](#), is justified by Voevodsky's model of types as Kan complexes and type families as Kan fibrations.

Contractible types



The homotopical perspective on type theory suggests new definitions:

A type A is **contractible** if it comes with a term of type

$$\text{is-contr}(A) := \sum_{a:A} \prod_{x:A} a =_A x$$

By Σ **elim**, a proof of contractibility provides:

- a term $c : A$ called the **center of contraction** and
- a dependent function $h : \prod_{x:A} c =_A x$ called the **contracting homotopy**, which can be thought of as a continuous choice of paths $h_x : c =_A x$ for each $x : A$.

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By Σ **elim**, a proof of contractibility provides:

- a term $c : A$ called the **center of contraction** and
- a dependent function $h : \prod_{x:A} c =_A x$ called the **contracting homotopy**, which can be thought of as a continuous choice of paths $h_x : c =_A x$ for each $x : A$.

If A is contractible and $x, y : A$, then there is a term $p : x =_A y$, so x and y are **indiscernible**. Thus, contractible types behave as if they have a single term.

The hierarchy of types



Contractible types, those types A for which the type

$$\text{is-contr}(A) := \sum_{a:A} \prod_{x:A} a =_A x$$

has a term, form the bottom level of Voevodsky's hierarchy of types.

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Traditional set theory and logic are encoded by layers formed by the 0-types and -1-types.

Equivalences



Similarly, homotopy theory suggests definitions of when two types A and B are **equivalent** or when a function $f : A \rightarrow B$ is an **equivalence**:

An **equivalence** between types A and B is a term of type:

$$A \simeq B := \sum_{f:A \rightarrow B} \left(\sum_{g:B \rightarrow A} \prod_{a:A} g(f(a)) =_A a \right) \times \left(\sum_{h:B \rightarrow A} \prod_{b:B} f(h(b)) =_B b \right)$$

By Σ_{elim} and \times_{elim} , a term of type $A \simeq B$ provides:

- functions $f : A \rightarrow B$ and $g, h : B \rightarrow A$ and
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This type is not a **proposition** and may have non-trivial higher structure.

The univalence axiom



Another notion of sameness between types is provided by the **universe** \mathcal{U} of types, which has (small) types A, B as its terms $\rightsquigarrow A, B : \mathcal{U}$.

Q: How do the types $A =_{\mathcal{U}} B$ and $A \simeq B$ compare?

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By **=elim**, there is a canonical function

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Univalence Axiom: $\text{id-to-equiv} : (A =_{\mathcal{U}} B) \rightarrow (A \simeq B)$ is an equivalence for all $A, B : \mathcal{U}$.

“Identity is equivalent to equivalence.”

$$(A =_{\mathcal{U}} B) \simeq (A \simeq B)$$

Consequences of univalence

There are myriad consequences of the univalence axiom $(A =_{\mathcal{U}} B) \simeq (A \simeq B)$:



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- **Function extensionality**: for any $f, g : A \rightarrow B$, the canonical function defines an equivalence between the identity type and the type of homotopies:

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- By **indiscernibility of identicals**, if $x, y : A$ and $p : x =_A y$ then $P(x) \simeq P(y)$ for any $P : A \rightarrow \mathbf{Type}$. By univalence, whenever $X \simeq Y$ then $X =_{\mathcal{U}} Y$ and thus any type constructed from X is equivalent to the corresponding type constructed from Y .

Voevodsky's univalence axiom — which is justified by the homotopical model of type theory — captures the common mathematical practice of transporting results proven about one object to any other object that is equivalent to it!



3

Contractibility as uniqueness

Rebuilding the pragmatic foundations for higher structures



I am pretty strongly convinced that there is an ongoing reversal in the collective consciousness of mathematicians: the homotopical picture of the world becomes the basic intuition, and if you want to get a discrete set, then you pass to the set of connected components of a space defined only up to homotopy ... Cantor's problems of the infinite recede to the background: from the very start, our images are so infinite that if you want to make something finite out of them, you must divide them by another infinity.

— Yuri Manin “We do not choose mathematics as our profession, it chooses us: Interview with Yuri Manin” by Mikhail Gelfand

∞ -categories in set theory



Essentially, ∞ -categories are 1-categories in which all the **sets** have been replaced by **∞ -groupoids** aka **homotopy types**:

sets $:: \infty$ -groupoids
categories $:: \infty$ -categories

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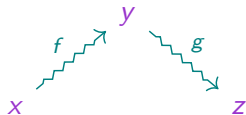
This is why ∞ -categories are so difficult to model within set theory.

Composing paths



In the **total singular complex** aka the **fundamental ∞ -groupoid** aka the **anima** or “soul”

of a space X , composites of paths are witnessed by higher paths:

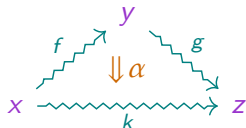


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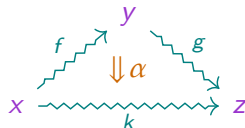
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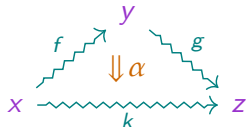


Theorem. The space of composites of two paths f and g in X is contractible.

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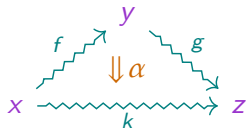
Proof: The **space of composites** of paths f and g in X is defined by the pullback:

$$\begin{array}{ccc} \text{Comp}(f, g) & \hookrightarrow & \text{Map}(\Delta, X) \\ \downarrow & \lrcorner & \downarrow \text{restrict} \\ * & \xrightarrow{f \wedge g} & \text{Map}(\Lambda, X) \end{array}$$

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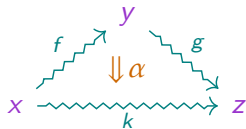
$$\begin{array}{ccccc}
 S^{n-1} & \longrightarrow & \text{Comp}(f, g) & \hookrightarrow & \text{Map}(\Delta, X) \\
 \downarrow & \nearrow & \downarrow & \lrcorner & \downarrow \text{restrict} \\
 D^n & \longrightarrow & * & \xrightarrow{f \wedge g} & \text{Map}(\Lambda, X)
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A space is **contractible** just when any sphere S^{n-1} can be filled to a disk D^n .

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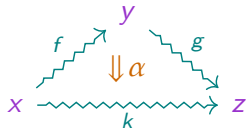
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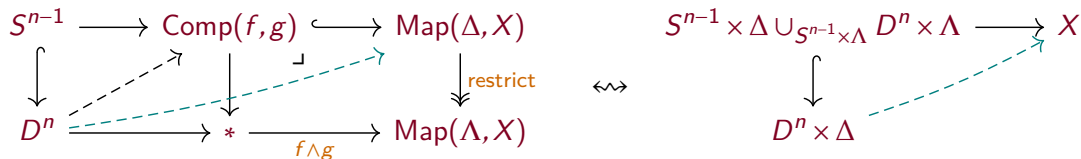
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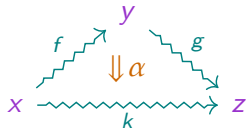
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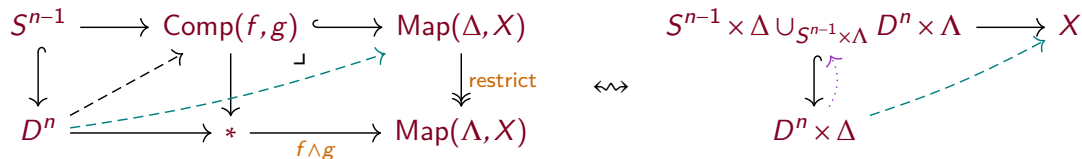
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A space is **contractible** just when any sphere S^{n-1} can be filled to a disk D^n . The extension exists since the inclusion admits a continuous deformation retract. □

∞ -categories in homotopy type theory



The identity type family gives each type the structure of an ∞ -groupoid: each type A has a family of identity types $x =_A y$ over $x, y : A$ whose terms $p : x =_A y$ are called **paths**.

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Emily Riehl and Michael Shulman, *A type theory for synthetic ∞ -categories*,
Higher Structures 1(1):116–193, 2017

each type A also has a family of hom types $\text{Hom}_A(x, y)$ over $x, y : A$ whose terms $f : \text{Hom}_A(x, y)$ are called **arrows**.

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With more of the work being done by the foundation system, perhaps someday ∞ -category theory will be easy enough to teach to undergraduates?



4

Computer proof assistants that compute

Formalizing univalent foundations



Today we face a problem that involves two difficult to satisfy conditions. On the one hand we have to find a way for computer assisted verification of mathematical proofs. This is necessary, first of all, because we have to stop the dissolution of the concept of proof in mathematics. On the other hand we have to preserve the intimate connection between mathematics and the world of human intuition. This connection is what moves mathematics forward and what we often experience as the beauty of mathematics.

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Thus, the homotopy type theory libraries have been built in **Coq**, **Agda** (`--without-K`), **Lean v2**, and new experimental systems ... though much of the work thusfar has been metatheoretic, developing univalent formal systems and their semantics.

Computing the Brunerie number



Guillaume Brunerie's 2016 PhD thesis contained a proof in homotopy type theory that

there exists $n : \mathbb{Z}$ such that $\pi_4(S^3) \simeq \mathbb{Z}/n\mathbb{Z}$

using the Hopf fibration, the long exact sequence of homotopy groups of a fibration, the Freudenthal suspension theorem, the James construction, the Blakers-Massey theorem, and Whitehead products — theorems of **synthetic homotopy theory** valid in any ∞ -topos.

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This result is quite remarkable in that even though it is a constructive proof, it is not at all obvious how to actually compute this n . At the time of writing, we still haven't managed to extract its value from its definition. A complete and concise definition of this number n is presented in appendix B, for the benefit of someone wanting to implement it in a prospective proof assistant.

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Guillaume Brunerie's 2016 PhD thesis contained a proof in homotopy type theory that
there exists $n : \mathbb{Z}$ such that $\pi_4(S^3) \simeq \mathbb{Z}/n\mathbb{Z}$

using the Hopf fibration, the long exact sequence of homotopy groups of a fibration, the Freudenthal suspension theorem, the James construction, the Blakers-Massey theorem, and Whitehead products — theorems of **synthetic homotopy theory** valid in any ∞ -topos.

This result is quite remarkable in that even though it is a constructive proof, it is not at all obvious how to actually compute this n . At the time of writing, we still haven't managed to extract its value from its definition. A complete and concise definition of this number n is presented in appendix B, for the benefit of someone wanting to implement it in a prospective proof assistant.

When dependent type theory is a foundation for **constructive mathematics**, the univalence axiom posits an inverse equivalence to **id-to-equiv** which is “stuck.” Voevodsky's simplicial set based model proves that it is consistent to assume univalence but does not provide a construction of the inverse equivalence.

A constructive foundation for univalent mathematics



In the last decade, constructive proofs of univalence have been discovered in **cubical** variants of homotopy type theory with semantics in various categories of cubical sets:

Bezem-Coquand-Huber, Cohen-Coquand-Huber-Mörtberg, Awodey,
Angiuli-Brunerie-Coquand-Favonia-Harper-Licata,
Awodey-Cavallo-Coquand-Riehl-Sattler,...

In parallel, a variety of experimental **cubical proof assistants** have been developed —
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On the art of giving the same name to different things:

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Computer proof assistants that compute:

- Axel Ljungström, [The Brunerie Number Is -2](#)
- Anders Mörtberg (j.w.w. Axel Ljungström), [Formalizing \$\pi_4\(S^3\)\$ and computing a Brunerie number in Cubical Agda](#)

Thank you!