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Contractibility as Uniqueness

UCLA Distinguished Lecture Series

An analogy



contractibility $::$ uniqueness

1. Contractibility as Uniqueness
2. Categorifying Uniqueness
3. ∞ -Categorifying Uniqueness



1

Contractibility as Uniqueness

The algebra of paths

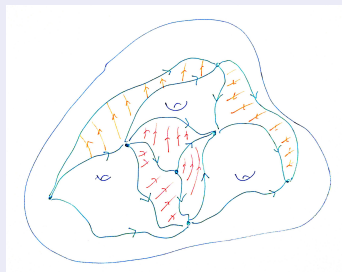
The standard technique used to distinguish your favorite space A from other spaces is to compute an algebraic invariant of the space.

The “algebra of paths” of a space is described in increasing precision by:

- the fundamental group $\pi_1(A, x)$ of loops in A based at x up to homotopy
- the fundamental groupoid $\pi_1 A$ of paths in A up to homotopy
- the fundamental ∞ -groupoid $\pi_\infty A$ of paths in A

$\pi_\infty A$ has:

- points of A as objects
- paths of A as 1-arrows
- paths between paths in A as 2-arrows
- paths between paths between paths in A as 3-arrows, and so on ...

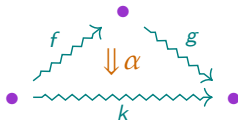


Witnesses to composition

Q: How do we define the **composite** of two paths?

A: We don't!

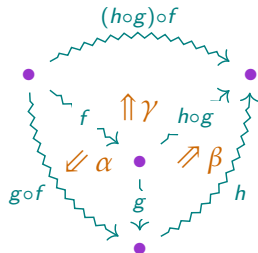
Instead of a composition **operation**, composites of paths are **witnessed** by higher paths.



Q: How unique is path composition?

Partial A: Unique enough for **associativity**.

Given composable paths f , g , h and specified higher paths α , β , γ witnessing composition relations, these higher paths compose. More precisely, a 3-arrow expresses a coherence between compositions witnessed by 2-arrows.



Homotopical uniqueness of path composition



Theorem. The space of composites of two paths f and g in A is contractible.

Proof: The space of composites of paths f and g in A is defined by the pullback:

$$\begin{array}{ccc}
 S^{n-1} & \longrightarrow & \text{Comp}(f, g) \hookrightarrow A^\Delta \\
 \downarrow & \nearrow & \downarrow \text{restrict} \\
 D^n & \xrightarrow{\quad} & * \xrightarrow{f \wedge g} A^\Delta
 \end{array}
 \quad \Leftrightarrow \quad
 \begin{array}{ccc}
 S^{n-1} \times \Delta \cup_{S^{n-1} \times \Delta} D^n \times \Delta & \longrightarrow & A \\
 \downarrow & \nearrow & \\
 D^n \times \Delta & &
 \end{array}$$

A space is **contractible** just when any sphere S^{n-1} can be filled to a disk D^n for $n \geq 1$. This filling problem transposes to an extension problem, and the extension exists since the inclusion admits a continuous deformation retract. \square

Summary



- In a **group(oid)**, composable arrows have a **unique** composite.
- In a **∞ -group(oid)**, composable arrows have a **contractible space** of composites.

The analogy

ordinary mathematics	::	higher mathematics
uniqueness	::	contractibility

can be made even tighter.

Aim: Express the classical notion of uniqueness more categorically.



2

Categorifying Uniqueness

Uniqueness



To say C has a **unique** element means

$$\exists x \in C, \forall y \in C, x = y$$

Here “ $x = y$ ” is a **predicate** — a mathematical statement that is either true or false, depending on two free variables $x, y : C$.

In **proof-relevant mathematics**, we interpret “ $x = y$ ” as the set of all **proofs** that x equals y (which is empty if x and y are not equal).

Then we can form the set $\sum_{x \in C} \prod_{y \in C} x = y$

inspired by a notational analogy with the sentence $\exists x \in C, \forall y \in C, x = y$.

The set $\sum_{x \in C} \prod_{y \in C} x = y$ is also a set of proofs

— but proofs of what?

Proofs of uniqueness



The set $\sum_{x \in C} \prod_{y \in C} x = y$ is also a set of proofs

— but proofs of what?

An element of $\sum_{x \in C} \prod_{y \in C} x = y$ is

- the choice of some element $c \in C$
- together with a proof, for all $z \in C$, that c equals z .

Thus $\sum_{x \in C} \prod_{y \in C} x = y$ is the set of proofs of the sentence $\exists x \in C, \forall y \in C, x = y$

asserting that C has a **unique** element.

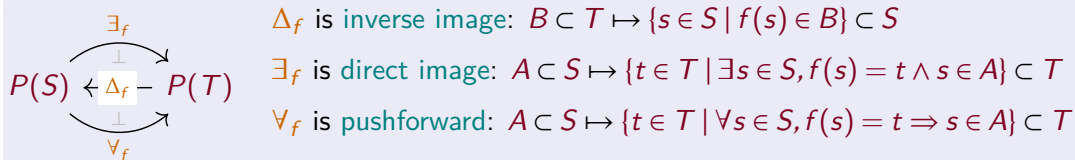
It remains to explain the analogy:

logic	\exists	\forall
sets	\sum	\prod

Digression: quantifiers as adjoints



A set function $f : S \rightarrow T$ induces order-preserving functions between their powersets:



For the unique function $! : S \rightarrow *$ these reduce to

The diagram shows two powersets, $P(S)$ and $P(*)$, connected by three arrows. The top arrow is labeled \exists and points from $P(S)$ to $P(*)$. The bottom arrow is labeled \forall and points from $P(S)$ to $P(*)$. The middle arrow is labeled Δ and points from $P(*)$ to $P(S)$. The arrows are connected by two vertical lines, one on each side, with a horizontal line connecting them in the center, forming a box-like structure.

$P(S) \xleftarrow{\Delta} P(*) = \{*, \emptyset\}$

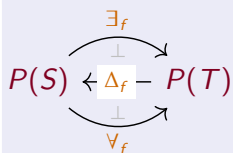
The set $P(S) = \{A \subset S\}$ can be identified with the set of **predicates** $p(s)$ with one free variable $s \in S$ — the corresponding subset is $\{s \in S \mid p(s) \text{ is true}\}$. If we interpret the two elements of $P(*)$ by $* =: \text{"true"}$ and $\emptyset =: \text{"false"}$ then

- \exists is the function that sends the predicate $p(s)$ to the sentence $\exists s \in S, p(s)$
- \forall is the function that sends the predicate $p(s)$ to the sentence $\forall s \in S, p(s)$

Digression: locally cartesian closed categories



For any function $f: S \rightarrow T$ there are functors:

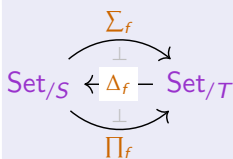


Δ_f is **inverse image**: $B \subset T \mapsto \{s \in S \mid f(s) \in B\} \subset S$

\exists_f is **direct image**: $A \subset S \mapsto \{t \in T \mid \exists s \in S, f(s) = t \wedge s \in A\} \subset T$

\forall_f is **pushforward**: $A \subset S \mapsto \{t \in T \mid \forall s \in S, f(s) = t \Rightarrow s \in A\} \subset T$

In **proof-relevant mathematics**, it is natural to replace the poset $P(S)$ by the category \mathbf{Set}_S of **S -indexed sets**. An object $\{P(s)\}_{s \in S}$ is a family of sets where $P(s)$ can be thought of as the set of proofs of some predicate $p(s)$ on $s \in S$.



Δ_f is **substitution**: $\{Q(t)\}_{t \in T} \mapsto \{Q(f(s))\}_{s \in S}$

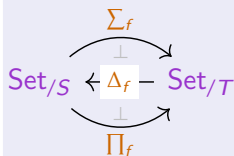
Σ_f is **sum**: $\{P(s)\}_{s \in S} \mapsto \{\sum_{s \in f^{-1}(t)} P(s)\}_{t \in T}$

Π_f is **product**: $\{P(s)\}_{s \in S} \mapsto \{\prod_{s \in f^{-1}(t)} P(s)\}_{t \in T}$

Summary



The triple of adjoint functors



Δ_f is **substitution**: $\{Q(t)\}_{t \in T} \mapsto \{Q(f(s))\}_{s \in S}$

Σ_f is **sum**: $\{P(s)\}_{s \in S} \mapsto \{\sum_{s \in f^{-1}(t)} P(s)\}_{t \in T}$

Π_f is **product**: $\{P(s)\}_{s \in S} \mapsto \{\prod_{s \in f^{-1}(t)} P(s)\}_{t \in T}$

gives a more formal way to understand the set

$$\sum_{x \in C} \prod_{y \in C} x = y$$

~> The set of proofs “ $x = y$ ” defines an indexed set $\{x = y\}_{x, y \in C} \in \text{Set}_{C \times C}$

~> Product along the projection $\pi_1 : C \times C \rightarrow C$ gives $\{\prod_{y \in C} x = y\}_{x \in C} \in \text{Set}_C$

~> Sum along $! : C \rightarrow *$ gives the set $\sum_{x \in C} \prod_{y \in C} x = y \in \text{Set}_* = \text{Set}$



3

∞ -Categorifying Uniqueness

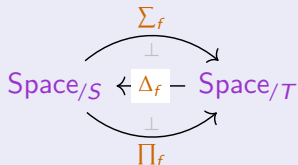
A convenient category of spaces



Now replace **Set** by a convenient category **Space** of spaces and continuous maps.

A **family of spaces** $\{E_b\}_{b \in B} \in \mathbf{Space}/_B$ is a continuous map $\pi: E \rightarrow B$, where the space E_b is the fiber over a point $b \in B$ in the base space while the total space $E \simeq \sum_{b \in B} E_b$.

Any continuous $f: S \rightarrow T$ gives rise to an adjoint triple:



Δ_f is **pullback**

Σ_f is **composition**

Π_f is **pushforward**

Identifications as paths



Q: For a space C , how to interpret the family of spaces $\{x = y\}_{x,y \in C} \in \text{Space}_{C \times C}$?

First guess: $\Delta \downarrow \in \text{Space}_{C \times C}$ — but a better choice is the **path space** $C' \downarrow \in \text{Space}_{C \times C}$.

New idea:

A point $p \in x = y$ is a **path** from x to y in C , providing a **proof** that x equals y .

A space of proofs



What is a point in the space $\sum_{x \in C} \prod_{y \in C} x = y$?

The functor $\sum_B : \mathbf{Space}_B \rightarrow \mathbf{Space}$ takes $\{E_b\}_{b \in B}$ to the total space $\sum_{b \in B} E_b$.

\rightsquigarrow a point in $\sum_{b \in B} E_b$ is a pair (a, e_a) of a point $a \in B$ and a point $e_a \in E_a$

The functor $\prod_B : \mathbf{Space}_B \rightarrow \mathbf{Space}$ takes $\{E_b\}_{b \in B}$ to the space of sections $\prod_{b \in B} E_b$.

\rightsquigarrow a point in $\prod_{b \in B} E_b$ is section $s : B \rightarrow \sum_{b \in B} E_b$ of the projection to B

- So a point in $\sum_{x \in C} \prod_{y \in C} x = y$ is a pair (c, h) where $c \in C$ and $h \in \prod_{y \in C} c = y$.
- The point $h \in \prod_{y \in C} c = y$ is a section $h : C \rightarrow \sum_{y \in C} c = y$ to the projection.

Together $(c, h) \in \sum_{x \in C} \prod_{y \in C} x = y$ defines:

- a center of contraction c and
- a contracting homotopy h ,

proving that the space C is contractible!

Contractibility as uniqueness



In summary, a point in the set

$$\sum_{x \in C} \prod_{y \in C} x = y$$

is a proof that C is **unique**, while a point in the space

$$\sum_{x \in C} \prod_{y \in C} x = y$$

is a proof that C is **contractible**.

Next time: If identifications $p \in x = y$ are paths, which may carry data, what strategies exist to prove a theorem involving a hypothesis of the form $x = y$? We'll introduce one powerful proof technique: the **principle of path induction** in **homotopy type theory**.

Thank you!