

# Project 9

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Section 1

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## 1 Argue that a DFA and a regular expression are equivalent.

Theorem 1.39:

Every nondeterministic finite automaton has an equivalent deterministic finite automaton.

Lemma 1.55:

If a language is described by a regular language, then it is regular.

Theorem 4.5:

$EQ_{\text{DFA}}$  is a decidable language.

Arguing that a DFA and a regular expression are equivalent.

### 1.1 Express this problem as a language.

Define the language as:

$$M = \{(D, R) \mid D \text{ is a DFA and } R \text{ is a regular expression with } L(D) = L(R)\}$$

Recall that the proof of theorem 4.5 defines a Turing machine  $F$  that decides the language  $EQ_{\text{DFA}} = \{(A, B) \mid A \text{ and } B \text{ are DFA's and } L(A) = L(B)\}$ . Then the following Turing machine  $T$  decides  $M$ .

### 1.2 Show that it is decidable.

$T =$  on input  $(D, R)$ , where  $D$  is a DFA and  $R$  is a regular expression:

- Convert  $R$  into a DFA  $C_R$  using the algorithm in Lemma 1.55.
- Run Turing machine decider  $F$  from Theorem 4.5 on input  $(D, C_R)$ .
- If  $F$  accepts, accept. If  $F$  rejects, reject.

- 2 Suppose  $S$  is the set of all infinite sequences of 0s and 1s.  $001010101\dots$  is a member of  $S$ , for example. Show that  $S$  is uncountable.**

Assume  $S$  is countable  $\Rightarrow \exists f : \mathbb{N} \rightarrow S$  (A bijection)

$n \rightarrow S_n$   
 $f(1) = 10011\dots$  (Pick 1 at index 1, flip to be 0)  
 $f(2) = 01110\dots$  (Pick 1 at index 2, flip to be 0)  
 $f(3) = 11001\dots$  (Pick 0 at index 3, flip to be 1)  
 $\dots$   
 $f(n)$   
 $\dots$

Special sequence  $A$ , it takes the first digit in the first line, second digit in the second line, third digit in third line and so on, but flips each digit.

$A = 001\dots$

$A$  being a sequence of 0s and 1s is in  $S$ , but it is never hit by the bijection.

In other words:

$A \in S$  but  $\forall n \in \mathbb{N}, f \neq A$

- 3 Example 3.9 describes a Turing machine,  $M$ , that decides the language  $w\#w$ , where  $w$  is any string over  $\{0,1\}^*$ . Give the sequence of configurations that  $M$  produces on input  $1\#1$ .**

$q_1 1\#1 \rightarrow q_6 1X1$   
 $Xq_3\#1 \rightarrow q_6 X1X1$   
 $X1q_51 \rightarrow q_6 \sqcup X1X1$   
 $Xq_6X1 \rightarrow q_{\text{reject}} \sqcup X1X1$

- 4 Write a Python program that simulates the Turing Machine  $M_2$ .**

See GitHub Classroom submission Q4.py submitted by Ben10164 (Ben Puryear) in the project-9-Ben10164 repository.

- 5 We know that if  $A \leq_M B$  and  $A$  is undecidable, then  $B$  is undecidable. We use this principle to conclude that an important problem in computational theory is undecidable. What's the problem?**

The halting problem.

- 6 State the problem formally for the previous question.**

$\text{HALT}_{\text{TM}} = \{(M, W) \mid M \text{ is a TM and } M \text{ halts on input } W\}$ .

- 7 What theorem did we call the "gloomiest theorem of them all?"**

Rice's theorem.