

# Week2

# x86-64 General-Purpose Registers

	<div><div>%rax</div><div>%eax</div></div>	<div><div>%r8</div><div>%r8d</div></div>	a5
	<div><div>%rbx</div><div>%ebx</div></div>	<div><div>%r9</div><div>%r9d</div></div>	a6
a4	<div><div>%rcx</div><div>%ecx</div></div>	<div><div>%r10</div><div>%r10d</div></div>	
a3	<div><div>%rdx</div><div>%edx</div></div>	<div><div>%r11</div><div>%r11d</div></div>	
a2	<div><div>%rsi</div><div>%esi</div></div>	<div><div>%r12</div><div>%r12d</div></div>	
a1	<div><div>%rdi</div><div>%edi</div></div>	<div><div>%r13</div><div>%r13d</div></div>	
	<div><div>%rsp</div><div>%esp</div></div>	<div><div>%r14</div><div>%r14d</div></div>	
	<div><div>%rbp</div><div>%ebp</div></div>	<div><div>%r15</div><div>%r15d</div></div>	

– Extend existing registers to 64 bits. Add 8 new ones.

## 32-bit code for swap

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

swap:

```
pushl %ebp
movl  %esp,%ebp
pushl %ebx
```

} Set  
Up

```
movl  8(%ebp), %edx
movl  12(%ebp), %ecx
movl  (%edx), %ebx
movl  (%ecx), %eax
movl  %eax, (%edx)
movl  %ebx, (%ecx)
```

} Body

```
popl  %ebx
popl  %ebp
ret
```

} Finish

## 64-bit code for swap

```
void swap(int *xp, int *yp)
{
    int t0 = *xp;
    int t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

swap:

```
movl    (%rdi), %edx
movl    (%rsi), %eax
movl    %eax, (%rdi)
movl    %edx, (%rsi)
```

ret

} Set  
Up

} Body

} Finish

# Object Code

```
int sum(int a, int b) {  
    return (a+b) ;  
}
```

<sum>:

55	push	%ebp
89 e5	mov	%esp, %ebp
8b 45 0c	mov	0xc(%ebp), %eax
03 45 08	add	0x8(%ebp), %eax
5d	pop	%ebp
c3	ret	

## Code for sum

0x401040 <sum>:

0x55  
0x89  
0xe5  
0x8b  
0x45  
0x0c  
0x03  
0x45  
0x08  
0x5d  
0xc3

- Total of 11 bytes
- Each instruction: 1, 2, or 3 bytes
- Starts at address 0x401040

# Disassembling Object Code

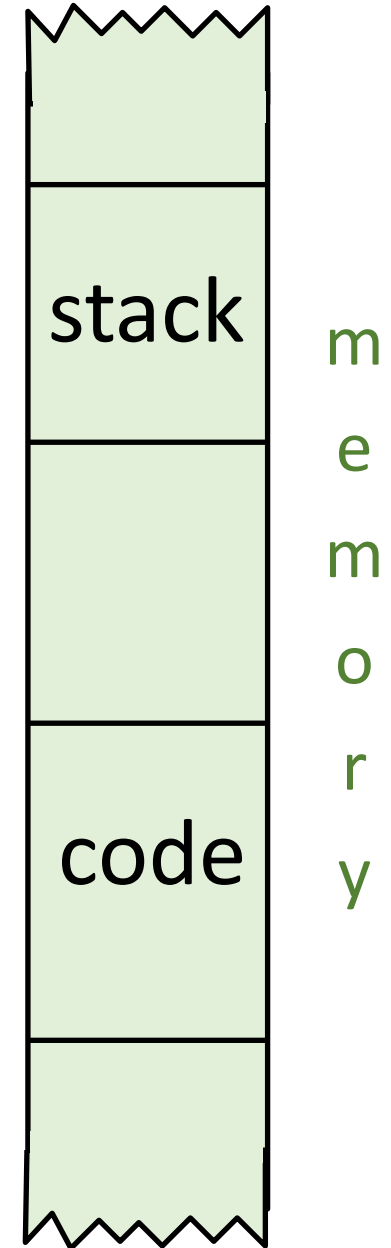
- **Disassembler**
  - `objdump -d <file>`
    - useful tool for examining object code
    - can be run on either executable or object (.o) file

# Today

- Procedures
  - Mechanisms
  - **Stack Structure**
  - **Calling Conventions**
    - Passing control
    - Passing data
    - Managing local data
  - **Illustration of Recursion**

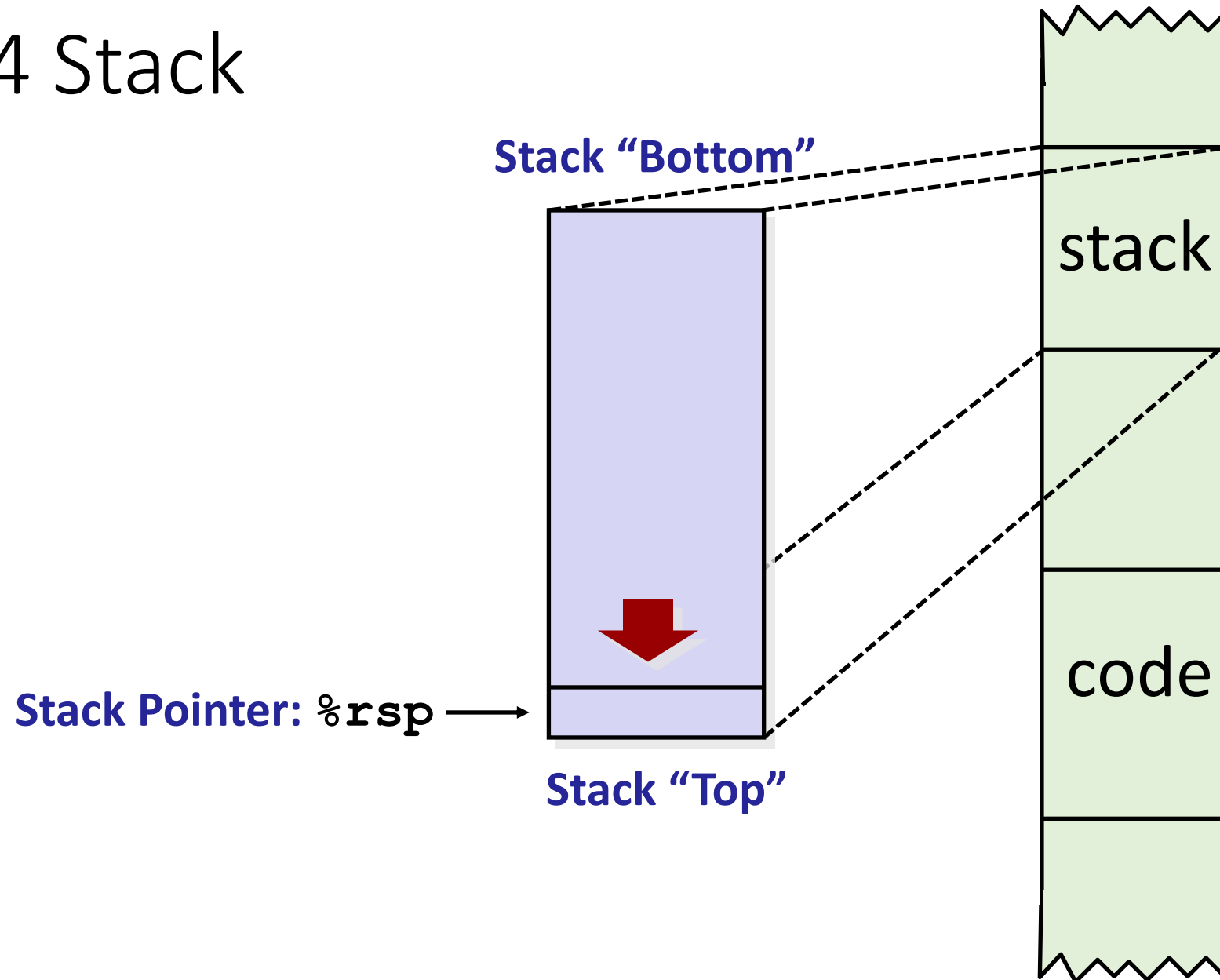
# x86-64 Stack

- Region of memory managed with stack discipline
  - Memory viewed as array of bytes.
  - Different regions have different purposes.



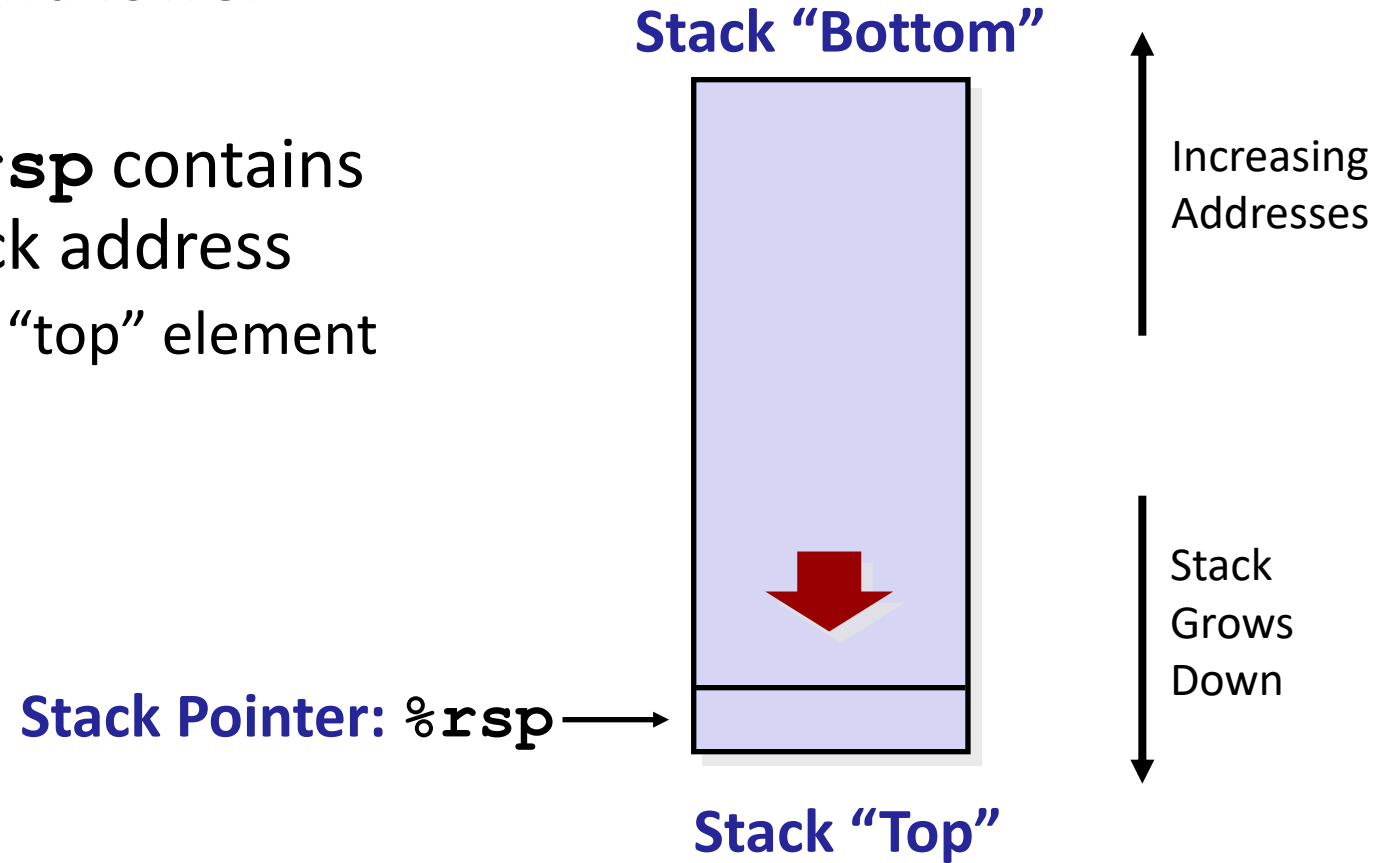


# x86-64 Stack



# x86-64 Stack

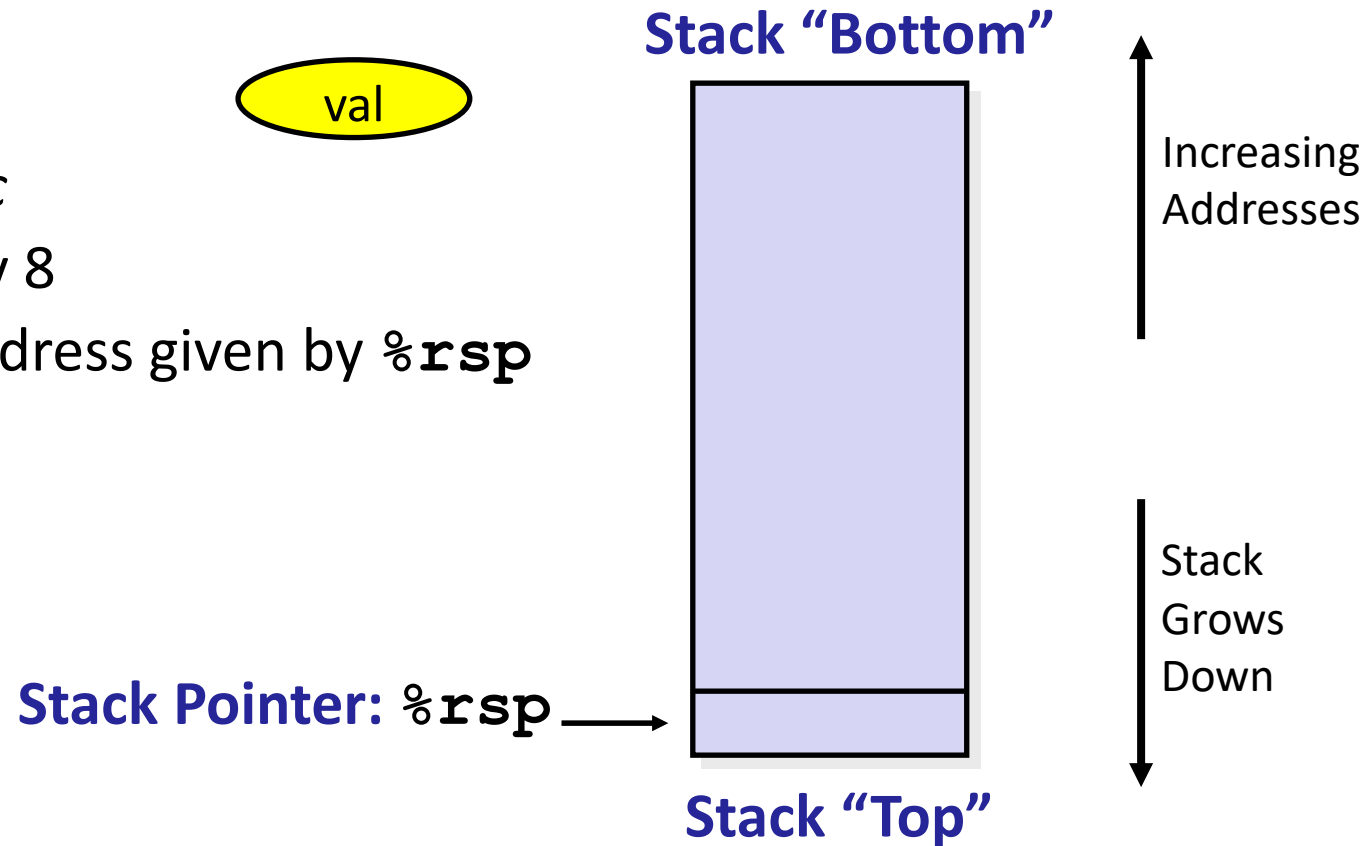
- Grows toward lower addresses
- Register `%rsp` contains lowest stack address
  - address of “top” element



# x86-64 Stack: Push

- **pushq *Src***

- Fetch operand at *Src*
- Decrement `%rsp` by 8
- Write operand at address given by `%rsp`

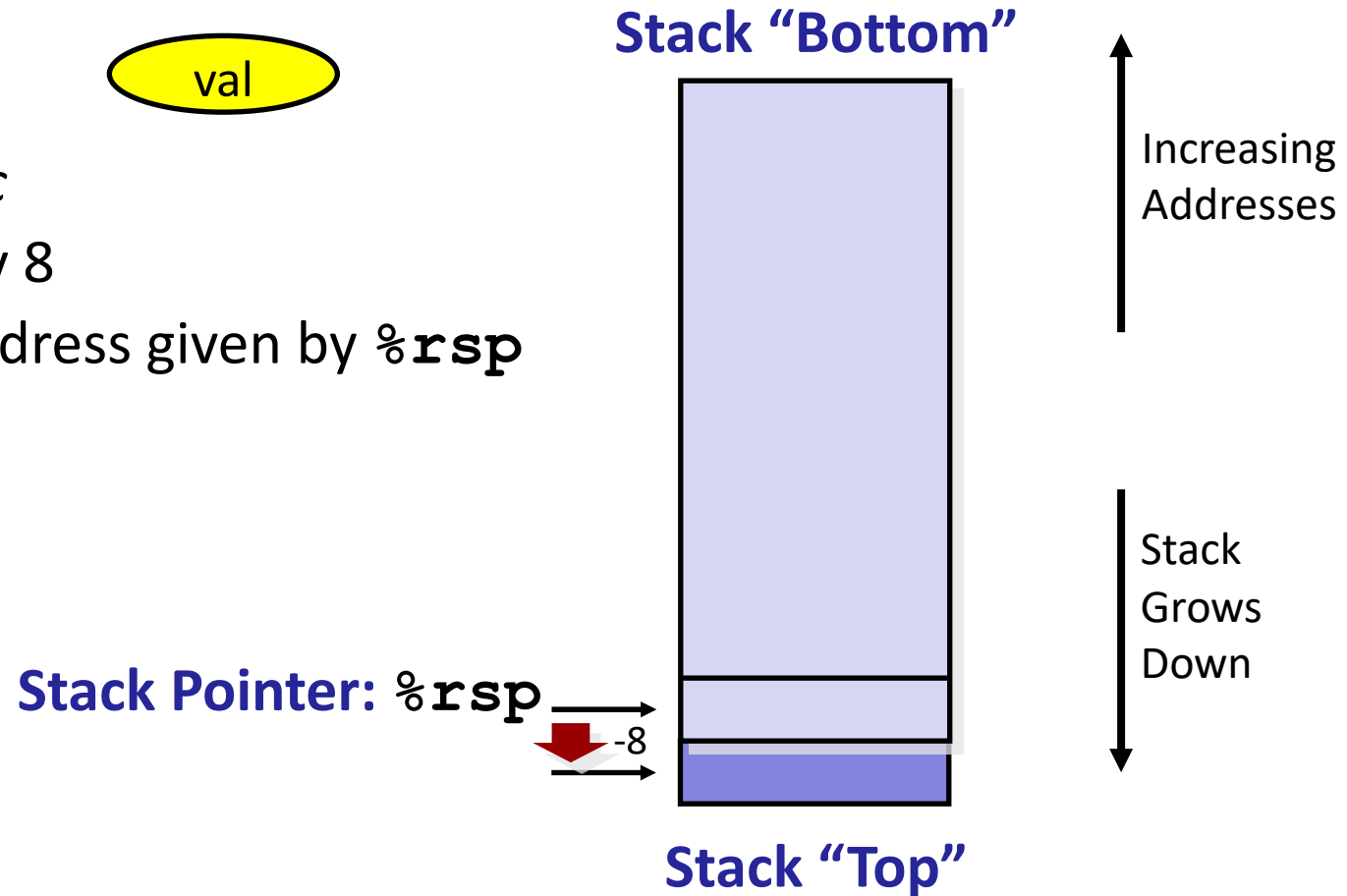


# x86-64 Stack: Push

- **pushq *Src***

- Fetch operand at *Src*
- Decrement `%rsp` by 8
- Write operand at address given by `%rsp`

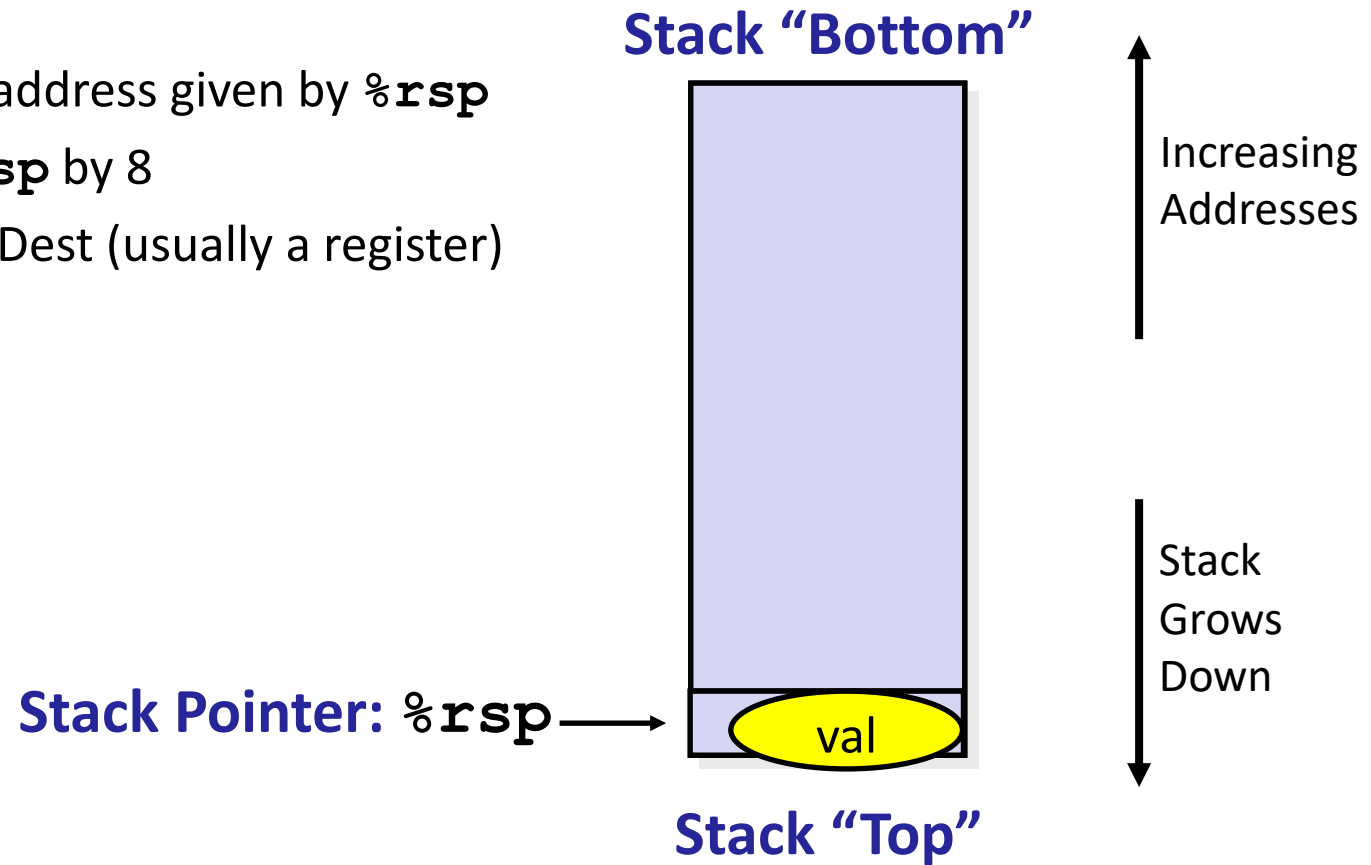
val



# x86-64 Stack: Pop

## ■ `popq Dest`

- Read value at address given by `%rsp`
- Increment `%rsp` by 8
- Store value at `Dest` (usually a register)



# x86-64 Stack: Pop

## ■ `popq Dest`

- Read value at address given by `%rsp`
- Increment `%rsp` by 8
- Store value at `Dest` (usually a register)

val

Stack Pointer: `%rsp`

+8

Stack "Bottom"



Stack "Top"

Increasing  
Addresses

Stack  
Grows  
Down

# x86-64 Stack: Pop

## ■ `popq Dest`

- Read value at address given by `%rsp`
- Increment `%rsp` by 8
- Store value at `Dest` (usually a register)

Stack Pointer: `%rsp` →

Stack “Bottom”



Increasing  
Addresses

Stack  
Grows  
Down

(The memory doesn't change,  
only the value of `%rsp`)

# Today

- Procedures
  - Mechanisms
  - Stack Structure
  - Calling Conventions
    - Passing control
    - Passing data
    - Managing local data
  - Illustration of Recursion



# Code Examples

```
void multstore(long x, long y, long *dest)
{
    long t = mult2(x, y);
    *dest = t;
}
```

```
00000000000400540 <multstore>:
400540: push    %rbx                # Save %rbx
400541: mov     %rdx,%rbx           # Save dest
400544: callq   400550 <mult2>      # mult2(x,y)
400549: mov     %rax, (%rbx)         # Save at dest
40054c: pop     %rbx                # Restore %rbx
40054d: retq                      # Return
```

```
long mult2(long a, long b)
{
    long s = a * b;
    return s;
}
```

```
00000000000400550 <mult2>:
400550: mov     %rdi,%rax           # a
400553: imul    %rsi,%rax           # a * b
400557: retq                      # Return
```

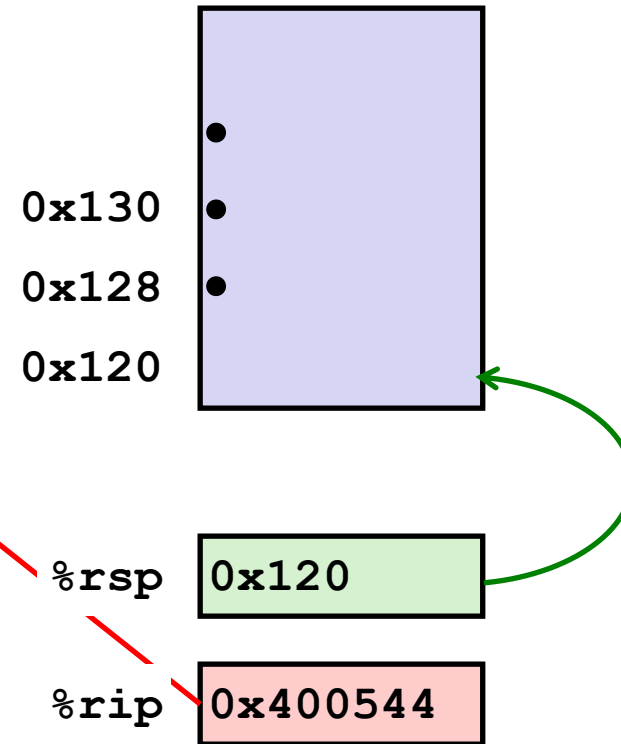
# Procedure Control Flow

- Use stack to support procedure call and return
- **Procedure call: `call label`**
  - Push return address on stack
  - Jump to ***label***
- Return address:
  - Address of the next instruction right after call
- **Procedure return: `ret`**
  - Pop address from stack
  - Jump to address

# Control Flow Example #1

```
0000000000400540 <multstore>:  
.  
.  
400544: callq 400550 <mult2>  
400549: mov  %rax, (%rbx)  
.  
.
```

```
0000000000400550 <mult2>:  
400550: mov  %rdi, %rax  
.  
.  
400557: retq
```



# Control Flow Example #2

```
00000000000400540 <multstore>:
```

•

•

```
400544: callq 400550 <mult2>
```

```
400549: mov    %rax, (%rbx) ←
```

•

•

```
00000000000400550 <mult2>:
```

```
400550: mov    %rdi, %rax ←
```

•

•

```
400557: retq
```

0x130

0x128

0x120

0x118

0x400549

%rsp

0x118

%rip

0x400550

# Control Flow Example #3

```
00000000000400540 <multstore>:
```

•

•

```
400544: callq 400550 <mult2>
```

```
400549: mov    %rax, (%rbx) ←
```

•

•

0x130

0x128

0x120

0x118

0x400549

%rsp

0x118

%rip

0x400557

```
00000000000400550 <mult2>:
```

```
400550: mov    %rdi,%rax
```

•

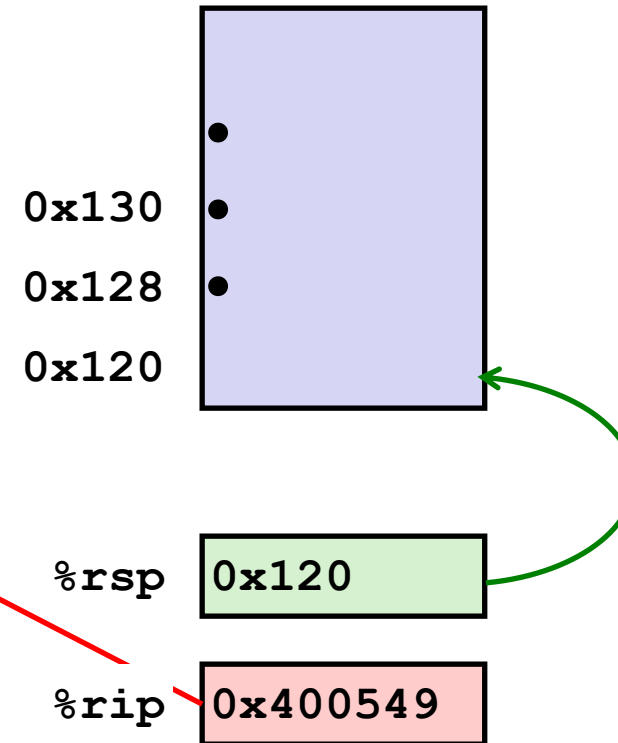
•

```
400557: retq ←
```

# Control Flow Example #4

```
00000000000400540 <multstore>:  
.  
.  
400544: callq 400550 <mult2>  
400549: mov  %rax, (%rbx)  
.  
.
```

```
00000000000400550 <mult2>:  
400550: mov  %rdi, %rax  
.  
.  
400557: retq
```



# Today

- Procedures
  - Mechanisms
  - tack Structure
  - Calling Conventions
    - Passing control
    - Passing data
    - Managing local data
  - Illustrations of Recursion & Pointers

# Procedure Data Flow

## Registers

- First 6 arguments

<code>%rdi</code>
<code>%rsi</code>
<code>%rdx</code>
<code>%rcx</code>
<code>%r8</code>
<code>%r9</code>

- Return value

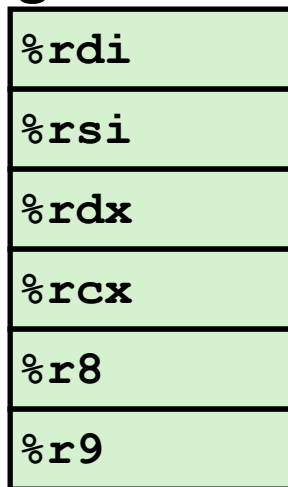
<code>%rax</code>
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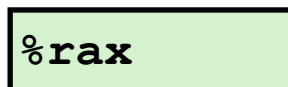
# Procedure Data Flow

## Registers

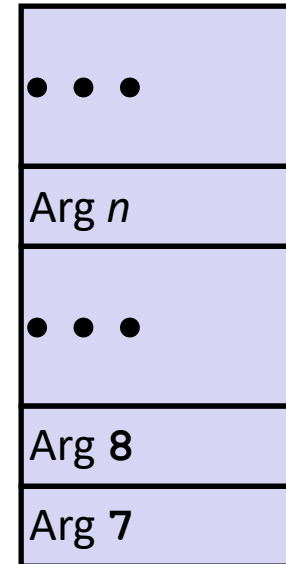
- First 6 arguments



- Return value



## Stack



- Only allocate stack space when needed

# Register Saving Conventions

- Conventions

- ***“Caller Saved”***

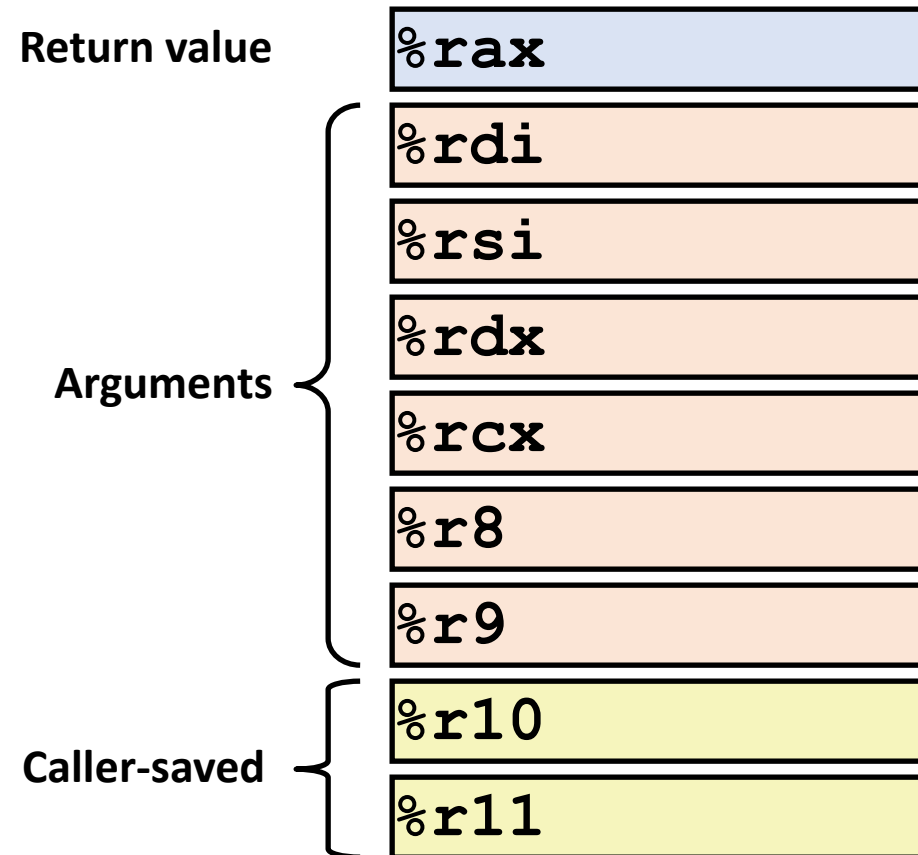
- Caller saves temporary values in its frame before the call
      - push %rax
    - Caller restores them in its frame after the call
      - pop %rax

- ***“Callee Saved”***

- Callee saves temporary values in its frame before using
    - Callee restores them before returning to caller

# x86-64 Linux Register Usage #1

- **%rax**
  - Return value
  - Also caller-saved
  - Can be modified by procedure
- **%rdi, ..., %r9**
  - Arguments
  - Also caller-saved
  - Can be modified by procedure
- **%r10, %r11**
  - Caller-saved
  - Can be modified by procedure



# x86-64 Linux Register Usage #2

- **%rbx, %r12, %r13, %r14**

- Callee-saved
- Callee must save & restore

- **%rbp**

- Callee-saved
- Callee must save & restore
- May be used as frame pointer

- **%rsp**

- Special form of callee save
- Restored to original value upon exit from procedure

