Chapter 7: Routing Dynamically

Objectives

After completion of this chapter, our goals are for you to be able to answer the following questions:

- What is the purpose of dynamic routing protocols?
- Comparing dynamic routing vs static routing?
- How dynamic routing protocols share route information and achieve convergence?
- Differences between the categories of dynamic routing protocols?
- The different types of distance vector routing protocols?
- The RIP routing protocol configuration?
- How Algorithm used by link-state routing protocols determine the best path?
- What are the advantages and disadvantages of using link-state routing protocols?
- Know how to determine the source route, administrative distance, and metric for a given route?
- How do you describe the differences between the IPv4 route lookup process and the IPv6 route lookup process?
- Can you determine which route will be used to forward a packet upon analyzing a routing table?

1. Introduction

A dynamic routing protocol is the communication used between routers. It allows one router to share information with other routers regarding the networks it knows about as well as its proximity to other routers. The information a router gets from another router, using a routing protocol, is used to build and maintain a routing table.

In a large network with numerous networks and subnets, configuring and maintaining static routes between these networks requires a great deal of administrative and operational overhead. This operational overhead is especially cumbersome when changes to the network occur, such as a down link or implementing a new subnet. Implementing dynamic routing protocols can ease the burden of configuration and maintenance tasks and give the network scalability.

Chapter 7 introduces dynamic routing protocols. It explores the benefits of using dynamic routing protocols, how different routing protocols are classified, and the metrics routing protocols use to determine the best path for network traffic. Other topics covered in this chapter include the characteristics of dynamic routing protocols and how the various routing protocols differ. Network professionals must understand the different routing protocols available in order to make informed decisions about when to use static or dynamic routing. They also need to know which dynamic routing protocol is most appropriate in a particular network environment.

2. Dynamic Routing Protocols

Dynamic routing protocols play an important role in today's networks. The following sections describe several important benefits that dynamic routing protocols provide. In many networks, dynamic routing protocols are typically used with static routes.

The Evolution of Dynamic Routing Protocols

Dynamic routing protocols have been used in networks since the late 1980s. One of the first routing protocols was ***Routing Information Protocol (RIP)***. RIP version 1 (RIPv1) was released in 1988, but some of the basic algorithms within the protocol were used on the ***Advanced Research Projects Agency Network (ARPANET)*** as early as 1969.

As networks evolved and became more complex, new routing protocols emerged. The RIP routing protocol was updated to accommodate growth in the network environment, into RIPv2. However, the newer version of RIP still does not scale to the larger network implementations of today. To address the needs of larger networks, two advanced routing protocols were developed: ***Open Shortest Path First (OSPF)*** and ***Intermediate System-to-Intermediate System (IS-IS)***. Cisco developed the ***Interior Gateway Routing Protocol (IGRP)*** and ***Enhanced IGRP (EIGRP)***, which also scales well in larger network implementations.

Additionally, there was the need to connect different internetworks and provide routing between them. The ***Border Gateway Protocol (BGP)*** is now used between Internet service providers (ISPs). BGP is also used between ISPs and their larger private clients to exchange routing information.

Table 7-1 classifies the protocols.

Table 7-1 Routing Protocol Classification

	Interior Gateway Protocols				**Exterior Gateway Protocols**
	Distance Vector		Link-State		Path Vector
IPv4	RIPv2	EIGRP	OSPFv2	IS-IS	BGP-4
IPv6	RIPng	EIGRP for IPv6	OSPFv3	IS-IS for IPv6	MBGP

With the advent of numerous consumer devices using IP, the IPv4 addressing space is nearly exhausted; thus, IPv6 has emerged. To support the communication based on IPv6, newer versions of the IP routing protocols have been developed, as shown by the IPv6 row in Table 7-1.

RIP is the simplest of dynamic routing protocols and is used in this section to provide a basic level of routing protocol understanding.

Purpose of Dynamic Routing Protocols

Routing protocols are used to facilitate the exchange of routing information between routers. A routing protocol is a set of processes, algorithms, and messages that are used to exchange routing information and populate the routing table with the routing protocol's choice of best paths. The purpose of dynamic routing protocols includes:

- Discovery of remote networks
- Maintaining up-to-date routing information
- Choosing the best path to destination networks
- Ability to find a new best path if the current path is no longer available

The main components of dynamic routing protocols include:

- ***Data structures***: Routing protocols typically use tables or databases for their operations. This information is kept in RAM.
- ***Routing protocol messages***: Routing protocols use various types of messages to discover neighboring routers, exchange routing information, and perform other tasks to learn and maintain accurate information about the network.
- ***Algorithm***: An algorithm is a finite list of steps used to accomplish a task. Routing protocols use algorithms for facilitating routing information and for best path determination.

Figure 7-1 highlights the data structures, routing protocol messages, and routing algorithm used by EIGRP.

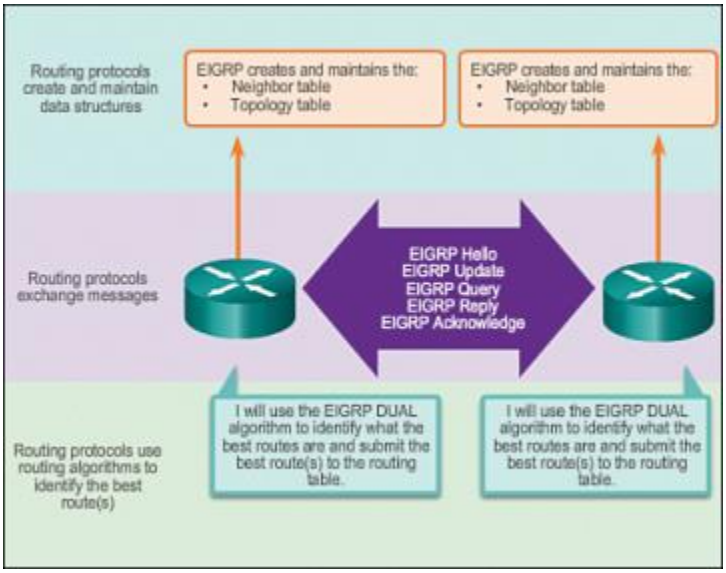


Figure 7-1 Components of Routing Protocols

The Role of Dynamic Routing Protocols

Routing protocols allow routers to dynamically share information about remote networks and automatically add this information to their own routing tables.

Routing protocols determine the best path, or route, to each network. That route is then added to the routing table. A primary benefit of dynamic routing protocols is that routers exchange routing information when there is a topology change. This exchange allows routers to automatically learn about new networks and also to find alternate paths when there is a link failure to a current network.

Compared to static routing, dynamic routing protocols require less administrative overhead. However, the expense of using dynamic routing protocols is dedicating part of a router's resources for protocol operation, including CPU time and network link bandwidth. Despite the benefits of dynamic routing, static routing still has its place. There are times when static routing is more appropriate and other times when dynamic routing is the better choice. Networks with moderate levels of complexity may have both static and dynamic routing configured.

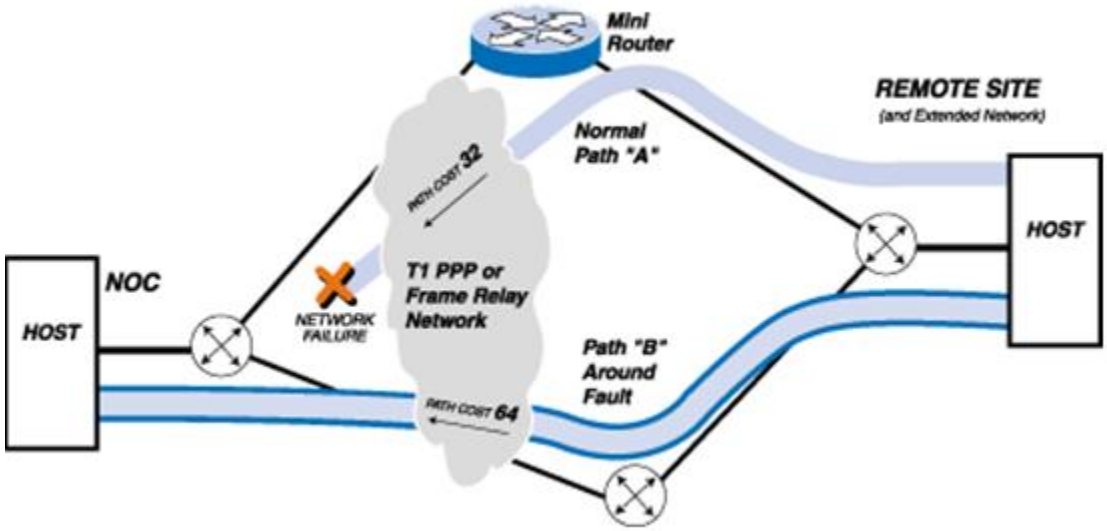
3. Static vs Dynamic Routing

Difference between static and dynamic routing is with regard to the way routing entries enter into the system. Routing in computer networking refers to the process of proper forwarding of packets across computer networks so that finally the packets reach the correct destination. Routing is of two main types as static routing and dynamic routing. In static routing, the network administrator manually sets the routing entries in the routing tables. That is where he manually put sentries that specify which path must be followed in order for a packet to reach a certain destination. On the other hand, in dynamic routing, routing entries are auto-generated using routing protocols automatically without any intervention of the network administrator. The algorithms used are complex but for current networks, which are quite large in size and ones that undergo changes often, dynamic routing is most suitable.

What is Static Routing?

In static routing, network **administrator manually enters the routing entries** to the routing table of each router and computer. A routing entry is an entry that specifies what the gateway that a packet must be forwarded, in order for it to reach a certain destination. On each router or computer, a table called routing table exists that contains a number of routing entries. For a simple small network, entering static routes to each router would be doable but it becomes too tedious with the increase of the size and the complexity of the network. Also, if a change occurs in a network that affects routing (for example, a router is down, or a new router is added), routing

entries must be manually changed. So, in static routing, management of routing tables must also be done by the administrator. The advantage of static routing is that there is not much processing. The only action is to do a lookup on the routing table for a specific destination and hence the routing hardware must not need any sophisticated processors making them cheaper.



A system of dynamic routing for future transportation

What is Dynamic Routing?

In dynamic routing, **routing entries are auto-generated** by routing algorithms. Hence, the administrator does not have to do any manual editing. Routing algorithms are complex mathematical algorithms where the routers advertise about their links and using that information, the most ideal routes are calculated. There are different methods depending on how advertising and calculations happens. **Link state algorithms** and **distance vector algorithms** are two such famous methods. OSPF (Open Shortest Path First) is an algorithm that follows a link state algorithm and RIP (Routing Information Protocol) is an algorithm that uses distance vector algorithm. For modern large networks that involve a lot of changes during operation, dynamic routing is ideal.

In dynamic routing, routing tables are periodically updated and hence, if any change has occurred, the new routing tables would be formed according to them. Another advantage is that in dynamic routing, depending on the congestion, the routing is adapted. That is, if a certain path is too much congested, routing protocols would figure them out and those paths would be avoided in the future routing tables. The drawback of dynamic routing is that the calculation are complex that it would need considerable amount of processing. Therefore, the cost of such routing hardware would be costly.

What is the difference between Static Routing and Dynamic Routing?

• In static routing, network administrator manually enters entries to the routing tables. But in dynamic routing, network administrator does not have to enter any entries as the entries are auto generated.

• In dynamic routing, routing entries are generated using complex routing algorithms. In static routing, no such algorithms are involved.

• For static routing, the action is to just do a lookup on a table and hence does not need any processing making the hardware less costly. But, dynamic routing algorithms involves a lot of calculations. Hence, it requires much processing capabilities. As a result, the hardware would be costly.

• In static routing, routers do not advertise or broadcast any information about the links to other routers. But, in dynamic routing, tables are generated using such information advertised by routers.

• In dynamic routing, routing tables are periodically updated and hence are sensitive to any changes in the network. But, in static routing, the network administrator will have to manually do any changes.

• Static routing can be used for small networks. But, for larger networks, static routing cannot be maintained and hence dynamic routing is used.

• In static routing, if there is a link failure, communication would be affected till the link is up again or the administrator manually sets up an alternate path. But, in dynamic routing, in such an event, the routing table would be updated to have an alternate path.

• Static routing is much secure as no advertisements are sent. But, in dynamic routing, broadcasts and advertisements happens making it less secure.

Summary: Static vs Dynamic Routing

In computer networking, routing is one of the most important things that make a computer network properly work. Static routing is the process where the administrator has to manually setup the routing entries. On the other hand, in dynamic routing, routing tables are automatically generated using algorithms called routing algorithms such as RIP and OSPF. For large complicated networks, using static routing is very tedious and hence one has to go for dynamic routing. The advantage of dynamic routing is that the routing tables will be periodically generated and hence they would comply with any change in the network. But the disadvantage is that the calculations in dynamic routing require more processing power.

4. Routing Protocol Operating Fundamentals

All routing protocols basically perform the same tasks. They all exchange routing updates and converge to build routing tables that are used by the router to make packet forwarding decisions. This section provides an overview of routing protocol fundamentals.

Dynamic Routing Protocol Operation

All routing protocols are designed to learn about remote networks and to quickly adapt whenever there is a change in the topology. The method that a routing protocol uses to accomplish this depends upon the algorithm it uses and the operational characteristics of that protocol.

In general, the operations of a dynamic routing protocol can be described as follows:

1. The router sends and receives routing messages on its interfaces.
2. The router shares routing messages and routing information with other routers that are using the same routing protocol.
3. Routers exchange routing information to learn about remote networks.
4. When a router detects a topology change, the routing protocol can advertise this change to other routers.

Cold Start

All routing protocols follow the same patterns of operation. When a router powers up, it knows nothing about the network topology. It does not even know that there are devices on the other end of its links. The only information that a router has is from its own saved configuration file stored in NVRAM.

After a router boots successfully, it applies the saved configuration. If the IP addressing is configured correctly, then the router initially discovers its own directly connected networks.

Notice how the routers proceed through the boot process and then discover any directly connected networks and subnet masks. This information is added to their routing tables as follows:

- R1 adds the 10.1.0.0 network available through interface FastEthernet 0/0 and adds 10.2.0.0 available through interface Serial 0/0/0.
- R2 adds the 10.2.0.0 network available through interface Serial 0/0/0 and adds 10.3.0.0 available through interface Serial 0/0/1.
- R3 adds the 10.3.0.0 network available through interface Serial 0/0/1 and adds 10.4.0.0 available through interface FastEthernet 0/0.

With this initial information, the routers then proceed to find additional route sources for their routing tables.

Network Discovery

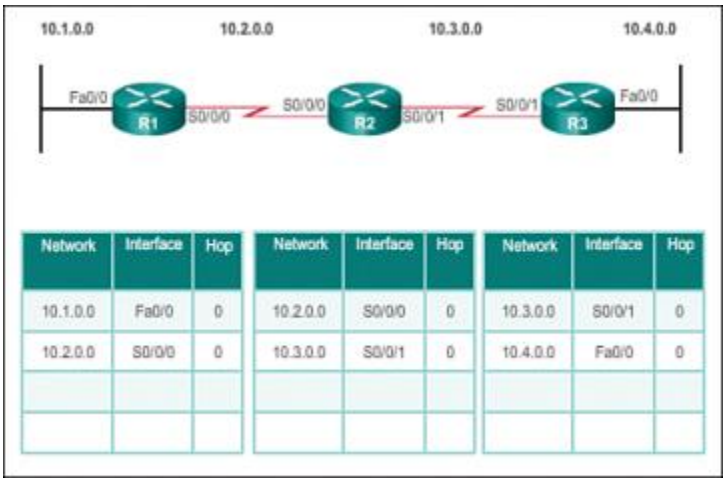
After initial boot up and discovery, the routing table is updated with all directly connected networks and the interfaces those networks reside on.

If a routing protocol is configured, the next step is for the router to begin exchanging routing updates to learn about any remote routes.

The router sends an update packet out all interfaces that are enabled on the router. The update contains the information in the routing table, which currently comprises all directly connected networks.

At the same time, the router also receives and processes similar updates from other connected routers. Upon receiving an update, the router checks it for new network information. Any networks that are not currently listed in the routing table are added.

Figure 7-2 depicts an example topology setup between three routers, R1, R2, and R3. Notice that only the directly connected networks are listed in each router's respective routing table.



Network	Interface	Hop
10.1.0.0	Fa0/0	0
10.2.0.0	S0/0/0	0

Network	Interface	Hop
10.2.0.0	S0/0/0	0
10.3.0.0	S0/0/1	0

Network	Interface	Hop
10.3.0.0	S0/0/1	0
10.4.0.0	Fa0/0	0

Figure 7-2 Initial Routing Table before Exchange

Based on this topology, a listing of the different updates that R1, R2, and R3 send and receive during initial convergence is provided:

R1:

- Sends an update about network 10.1.0.0 out the Serial 0/0/0 interface
- Sends an update about network 10.2.0.0 out the FastEthernet 0/0 interface
- Receives update from R2 about network 10.3.0.0 and increments the hop count by 1
- Stores network 10.3.0.0 in the routing table via Serial 0/0/0 with a metric of 1

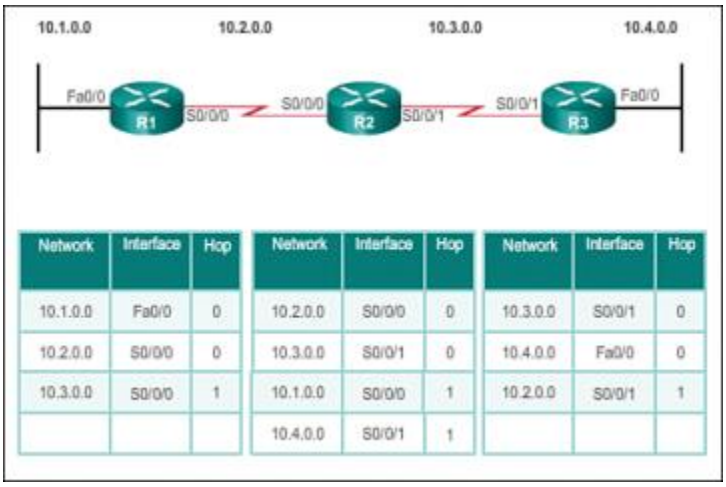
R2:

- Sends an update about network 10.3.0.0 out the Serial 0/0/0 interface
- Sends an update about network 10.2.0.0 out the Serial 0/0/1 interface
- Receives an update from R1 about network 10.1.0.0 and increments the hop count by 1
- Stores network 10.1.0.0 in the routing table via Serial 0/0/0 with a metric of 1
- Receives an update from R3 about network 10.4.0.0 and increments the hop count by 1
- Stores network 10.4.0.0 in the routing table via Serial 0/0/1 with a metric of 1

R3:

- Sends an update about network 10.4.0.0 out the Serial 0/0/1 interface
- Sends an update about network 10.3.0.0 out the FastEthernet 0/0 interface
- Receives an update from R2 about network 10.2.0.0 and increments the hop count by 1
- Stores network 10.2.0.0 in the routing table via Serial 0/0/1 with a metric of 1

Figure 7-3 displays the routing tables after the initial exchange.



Network	Interface	Hop
10.1.0.0	Fa0/0	0
10.2.0.0	S0/0/0	0
10.3.0.0	S0/0/0	1

Network	Interface	Hop
10.2.0.0	S0/0/0	0
10.3.0.0	S0/0/1	0
10.1.0.0	S0/0/0	1
10.4.0.0	S0/0/1	1

Network	Interface	Hop
10.3.0.0	S0/0/1	0
10.4.0.0	Fa0/0	0
10.2.0.0	S0/0/1	1

Figure 7-3 Routing Table After Initial Exchange

After this first round of update exchanges, each router knows about the connected networks of its directly connected neighbors. However, did you notice that R1 does not yet know about 10.4.0.0 and that R3 does not yet know about 10.1.0.0? Full knowledge and a converged network do not take place until there is another exchange of routing information.

Exchanging the Routing Information

At this point the routers have knowledge about their own directly connected networks and about the connected networks of their immediate neighbors. Continuing the journey toward convergence, the routers exchange the next round of periodic updates. Each router again checks the updates for new information.

After initial discovery is complete, each router continues the convergence process by sending and receiving the following updates.

R1:

- Sends an update about network 10.1.0.0 out the Serial 0/0/0 interface
- Sends an update about networks 10.2.0.0 and 10.3.0.0 out the FastEthernet 0/0 interface
- Receives an update from R2 about network 10.4.0.0 and increments the hop count by 1
- Stores network 10.4.0.0 in the routing table via Serial 0/0/0 with a metric of 2
- Same update from R2 contains information about network 10.3.0.0 with a metric of 1. There is no change; therefore, the routing information remains the same.

R2:

- Sends an update about networks 10.3.0.0 and 10.4.0.0 out of Serial 0/0/0 interface
- Sends an update about networks 10.1.0.0 and 10.2.0.0 out of Serial 0/0/1 interface
- Receives an update from R1 about network 10.1.0.0. There is no change; therefore, the routing information remains the same.
- Receives an update from R3 about network 10.4.0.0. There is no change; therefore, the routing information remains the same.

R3:

- Sends an update about network 10.4.0.0 out the Serial 0/0/1 interface
- Sends an update about networks 10.2.0.0 and 10.3.0.0 out the FastEthernet 0/0 interface
- Receives an update from R2 about network 10.1.0.0 and increments the hop count by 1
- Stores network 10.1.0.0 in the routing table via Serial 0/0/1 with a metric of 2
- Same update from R2 contains information about network 10.2.0.0 with a metric of 1. There is no change; therefore, the routing information remains the same.

Figure 7-4 displays the routing tables after the routers have converged.

Network	Interface	Hop
10.1.0.0	Fa0/0	0
10.2.0.0	S0/0/0	0
10.3.0.0	S0/0/0	1
10.4.0.0	S0/0/0	2

Network	Interface	Hop
10.2.0.0	S0/0/0	0
10.3.0.0	S0/0/1	0
10.1.0.0	S0/0/0	1
10.4.0.0	S0/0/1	1

Network	Interface	Hop
10.3.0.0	S0/0/1	0
10.4.0.0	Fa0/0	0
10.2.0.0	S0/0/1	1
10.1.0.0	S0/0/1	2

Figure 7-4 Routing Table After Convergence

Distance vector routing protocols typically implement a routing loop prevention technique known as split horizon. Split horizon prevents information from being sent out the same interface from which it was received. For example, R2 does not send an update containing the network 10.1.0.0 out of Serial 0/0/0, because R2 learned about network 10.1.0.0 through Serial 0/0/0.

After routers within a network have converged, the router can then use the information within the route table to determine the best path to reach a destination. Different routing protocols have different ways of calculating the best path.

Achieving Convergence

The network has converged when all routers have complete and accurate information about the entire network, as shown in Figure 7-4. ***Convergence*** is the time it takes routers to share information, calculate best paths, and update their routing tables. A network is not completely operable until the network has converged; therefore, most networks require short convergence times.

Convergence is both collaborative and independent. The routers share information with each other, but must independently calculate the impacts of the topology change on their own routes. Because they develop an agreement with the new topology independently, they are said to converge on this consensus.

Convergence properties include the speed of propagation of routing information and the calculation of optimal paths. The speed of propagation refers to the amount of time it takes for routers within the network to forward routing information.

As shown in Figure 7-5, routing protocols can be rated based on the speed to convergence; the faster the convergence, the better the routing protocol. Generally, older protocols, such as RIP, are slow to converge, whereas modern protocols, such as EIGRP and OSPF, converge more quickly.

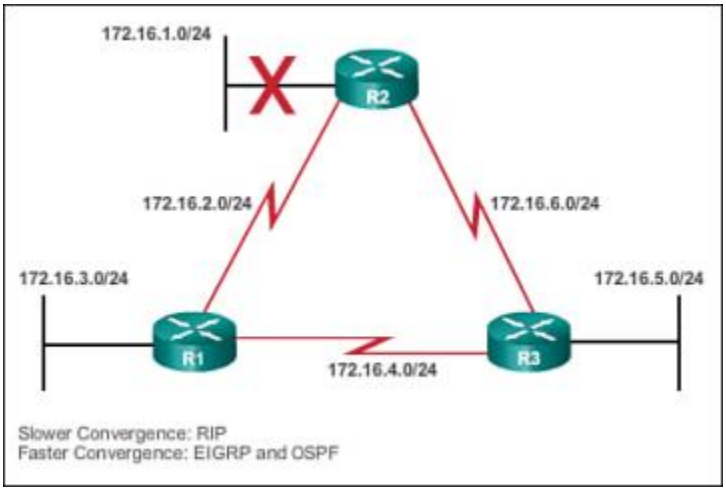


Figure 7-5 Converging

5. Types of Routing Protocols

Table 3-1 showed how routing protocols can be classified according to various characteristics. This section gives an overview of the most common IP routing protocols. Most of these routing protocols will be examined in detail in other chapters. For now, this section gives a very brief overview of each protocol.

Classifying Routing Protocols

Routing protocols can be classified into different groups according to their characteristics. Specifically, routing protocols can be classified by their:

- **Purpose**: Interior Gateway Protocol (IGP) or Exterior Gateway Protocol (EGP)
- **Operation**: Distance vector protocol, link-state protocol, or path-vector protocol
- **Behavior**: Classful (legacy) or classless protocol

For example, IPv4 routing protocols are classified as follows:

- **RIPv1 (legacy)**: IGP, distance vector, classful protocol
- **IGRP (legacy)**: IGP, distance vector, classful protocol developed by Cisco (deprecated from 12.2 IOS and later)
- **RIPv2**: IGP, distance vector, classless protocol
- **EIGRP**: IGP, distance vector, classless protocol developed by Cisco
- **OSPF**: IGP, link-state, classless protocol
- **IS-IS**: IGP, link-state, classless protocol
- **BGP**: EGP, path-vector, classless protocol

The ***classful routing protocols***, RIPv1 and IGRP, are legacy protocols and are only used in older networks. These routing protocols have evolved into the ***classless routing protocols***, RIPv2 and EIGRP, respectively. Link-state routing protocols are classless by nature.

Figure 7-6 displays a hierarchical view of dynamic routing protocol classification.

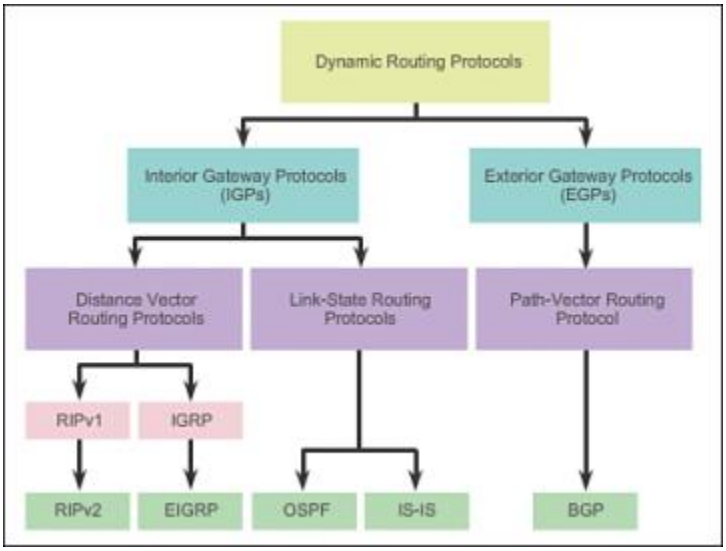


Figure 7-6 Routing Protocol Classification

IGP and EGP Routing Protocols

An ***autonomous system (AS)*** is a collection of routers under a common administration such as a company or an organization. An AS is also known as a routing domain. Typical examples of an AS are a company's internal network and an ISP's network.

The Internet is based on the AS concept; therefore, two types of routing protocols are required:

- ***Interior Gateway Protocols (IGP)***: Used for routing within an AS. It is also referred to as intra-AS routing. Companies, organizations, and even service providers use an IGP on their internal networks. IGPs include RIP, EIGRP, OSPF, and IS-IS.
- ***Exterior Gateway Protocols (EGP)***: Used for routing between autonomous systems. It is also referred to as inter-AS routing. Service providers and large companies may interconnect using an EGP. The Border Gateway Protocol (BGP) is the only currently viable EGP and is the official routing protocol used by the Internet.

Because BGP is the only EGP available, the term EGP is rarely used; instead, most engineers simply refer to BGP.

The example in Figure 7-7 provides simple scenarios highlighting the deployment of IGPs, BGP, and static routing.

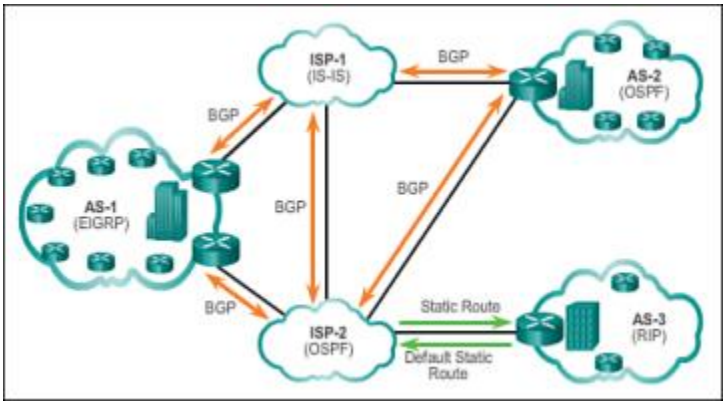


Figure 7-7 IGP versus EGP Routing Protocols

There are five individual autonomous systems in the scenario:

- **ISP-1**: This is an AS and it uses IS-IS as the IGP. It interconnects with other autonomous systems and service providers using BGP to explicitly control how traffic is routed.
- **ISP-2**: This is an AS and it uses OSPF as the IGP. It interconnects with other autonomous systems and service providers using BGP to explicitly control how traffic is routed.
- **AS-1**: This is a large organization and it uses EIGRP as the IGP. Because it is multihomed (i.e., connects to two different service providers), it uses BGP to explicitly control how traffic enters and leaves the AS.
- **AS-2**: This is a medium-sized organization and it uses OSPF as the IGP. It is also multihomed; therefore, it uses BGP to explicitly control how traffic enters and leaves the AS.
- **AS-3**: This is a small organization with older routers within the AS; it uses RIP as the IGP. BGP is not required because it is single-homed (i.e., connects to one service provider). Instead, static routing is implemented between the AS and the service provider.

Distance Vector Routing Protocols

Distance vector means that routes are advertised by providing two characteristics:

- **Distance**: Identifies how far it is to the destination network and is based on a metric such as the hop count, cost, bandwidth, delay, and more
- **Vector**: Specifies the direction of the next-hop router or exit interface to reach the destination

For example, in Figure 7-8, R1 knows that the distance to reach network 172.16.3.0/24 is one hop and that the direction is out of the interface Serial 0/0/0 toward R2.

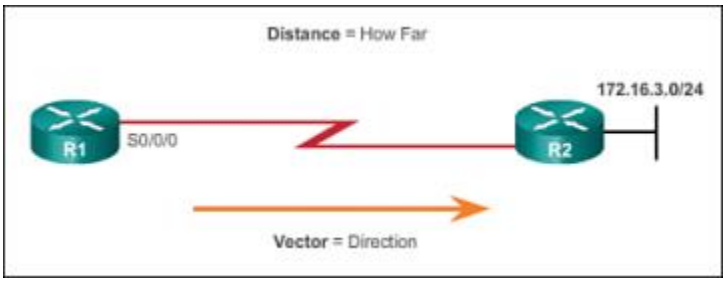


Figure 7-8 The Meaning of Distance Vector

A router using a ***distance vector routing protocol*** does not have the knowledge of the entire path to a destination network. Distance vector protocols use routers as sign posts along the path to the final destination. The only information a router knows about a remote network is the distance or metric to reach that network and which path or interface to use to get there. Distance vector routing protocols do not have an actual map of the network topology.

There are four distance vector IPv4 IGPs:

- **RIPv1**: First generation legacy protocol
- **RIPv2**: Simple distance vector routing protocol
- **IGRP**: First generation Cisco proprietary protocol (obsolete and replaced by EIGRP)
- **EIGRP**: Advanced version of distance vector routing

Link-State Routing Protocols

In contrast to distance vector routing protocol operation, a router configured with a ***link-state routing protocol*** can create a complete view or topology of the network by gathering information from all of the other routers.

To continue our analogy of sign posts, using a link-state routing protocol is like having a complete map of the network topology. The sign posts along the way from source to destination are not necessary, because all link-state routers are using an identical map of the network. A link-state router uses the link-state information to create a topology map and to select the best path to all destination networks in the topology.

RIP-enabled routers send periodic updates of their routing information to their neighbors. Link-state routing protocols do not use periodic updates. After the network has converged, a link-state update is only sent when there is a change in the topology. For example, in Figure 7-9, the link-state update is sent when the 172.16.3.0 network goes down.

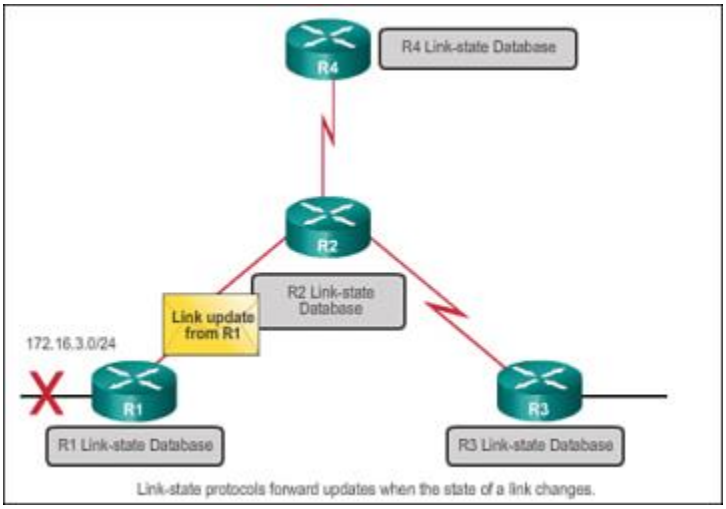


Figure 7-9 Link-State Protocol Operation

Link-state protocols work best in situations where:

- The network design is hierarchical, usually occurring in large networks
- Fast convergence of the network is crucial
- The administrators have good knowledge of the implemented link-state routing protocol

There are two link-state IPv4 IGPs:

- **OSPF**: Popular standards-based routing protocol
- **IS-IS**: Popular in provider networks

Classful Routing Protocols

The biggest distinction between classful and classless routing protocols is that classful routing protocols do not send subnet mask information in their routing updates. Classless routing protocols include subnet mask information in the routing updates.

The two original IPv4 routing protocols developed were RIPv1 and IGRP. They were created when network addresses were allocated based on classes (i.e., class A, B, or C). At that time, a routing protocol did not need to include the subnet mask in the routing update, because the network mask could be determined based on the first octet of the network address.

Only RIPv1 and IGRP are classful. All other IPv4 and IPv6 routing protocols are classless. Classful addressing has never been a part of IPv6.

The fact that RIPv1 and IGRP do not include subnet mask information in their updates means that they cannot provide variable-length subnet masks (VLSMs) and Classless Inter-Domain Routing (CIDR).

Classful routing protocols also create problems in discontiguous networks. A discontiguous network is when subnets from the same classful major network address are separated by a different classful network address.

To illustrate the shortcoming of classful routing, refer to the topology in Figure 7-10.

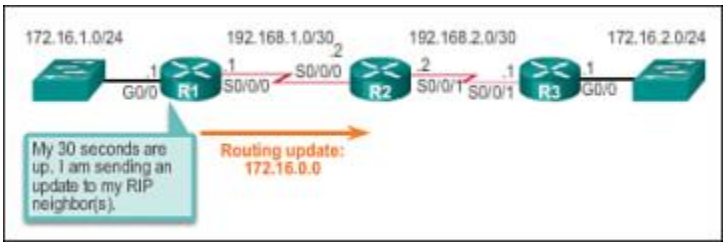


Figure 7-10 R1 Forwards a Classful Update to R2

Notice that the LANs of R1 (172.16.1.0/24) and R3 (172.16.2.0/24) are both subnets of the same class B network (172.16.0.0/16). They are separated by different classful network addresses (192.168.1.0/30 and 192.168.2.0/30).

When R1 forwards an update to R2, RIPv1 does not include the subnet mask information with the update; it only forwards the class B network address 172.16.0.0.

R2 receives and processes the update. It then creates and adds an entry for the class B 172.16.0.0/16 network in the routing table, as shown in Figure 7-11.

Figure 7-11. R2 Adds the Entry for 172.16.0.0 via R1

When R3 forwards an update to R2, it also does not include the subnet mask information and therefore only forwards the classful network address 172.16.0.0.

R2 receives and processes the update and adds another entry for the classful network address 172.16.0.0/16 to its routing table, as shown in Figure 7-12. When there are two entries with identical metrics in the routing table, the router shares the load of the traffic equally among the two links. This is known as load balancing.

Figure 7-12 R2 Adds the Entry for 172.16.0.0 via R3

Discontiguous networks have a negative impact on a network. For example, a ping to 172.16.1.1 would return "U.U.U" because R2 would forward the first ping out its Serial 0/0/1 interface toward R3, and R3 would return a Destination Unreachable (U) error code to R2. The second ping would exit out of R2's Serial 0/0/0 interface toward R1, and R1 would return a successful code (.). This pattern would continue until the **ping** command is done.

Classless Routing Protocols

Modern networks no longer use classful IP addressing and the subnet mask cannot be determined by the value of the first octet. The classless IPv4 routing protocols (RIPv2, EIGRP, OSPF, and IS-IS) all include the subnet mask information with the network address in routing updates. Classless routing protocols support VLSM and CIDR.

IPv6 routing protocols are classless. The distinction whether a routing protocol is classful or classless typically only applies to IPv4 routing protocols. All IPv6 routing protocols are considered classless because they include the prefix-length with the IPv6 address.

Figures 7-13 through 7-15 illustrate how classless routing solves the issues created with classful routing.

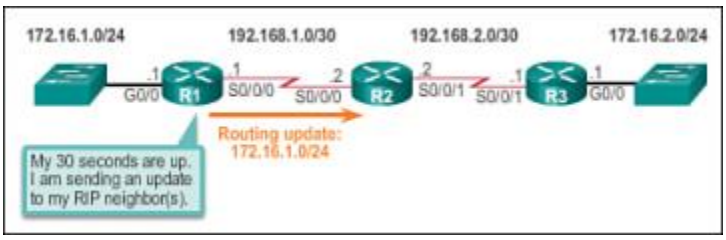


Figure 7-13 R1 Forwards a Classless Update to R2

```
R2# show ip route | begin Gateway
Gateway of last resort is not set

      172.16.0.0/24 is subnetted, 1 subnets
R        172.16.1.0 [120/1] via 192.168.1.1, 00:00:06,
         Serial0/0/0
      192.168.1.0/24 is variably subnetted, 2 subnets,
      2 masks
C        192.168.1.0/30 is directly connected, Serial0/0/0
L        192.168.1.2/32 is directly connected, Serial0/0/0
R2#
```

Figure 7-14 R2 Adds the Entry for the 172.16.1.0/24 Network via R1

```
R2# show ip route | begin Gateway
Gateway of last resort is not set

     172.16.0.0/24 is subnetted, 2 subnets
R       172.16.1.0 [120/1] via 192.168.1.1, 00:00:03,
        Serial0/0/0
R       172.16.2.0 [120/1] via 192.168.2.1, 00:00:03,
        Serial0/0/1
     192.168.1.0/24 is variably subnetted, 2 subnets,
     2 masks
C       192.168.1.0/30 is directly connected, Serial0/0/0
L       192.168.1.2/32 is directly connected, Serial0/0/0
     192.168.2.0/24 is variably subnetted, 2 subnets,
     2 masks
C       192.168.2.0/30 is directly connected, Serial0/0/1
L       192.168.2.2/32 is directly connected, Serial0/0/1
R2#
```

Figure 7-15 Entry for the 172.16.2.0/24 Network via R3

In the ***discontiguous network*** design of Figure 7-13, the classless protocol RIPv2 has been implemented on all three routers. When R1 forwards an update to R2, RIPv2 includes the subnet mask information with the update 172.16.1.0/24.

In Figure 7-14, R2 receives, processes, and adds two entries in the routing table. The first line displays the classful network address 172.16.0.0 with the /24 subnet mask of the update. This is known as the parent route. The second entry displays the VLSM network address 172.16.1.0 with the exit and next-hop address. This is referred to as the child route. Parent routes never include an exit interface or next-hop IP address.

When R3 forwards an update to R2, RIPv2 includes the subnet mask information with the update 172.16.2.0/24.

R2 receives, processes, and adds another child route entry 172.16.2.0/24 under the parent route entry 172.16.0.0, as shown in Figure 7-15.

A ping from R2 to 172.16.1.1 would now be successful.

Routing Protocol Characteristics

Routing protocols can be compared based on the following characteristics:

- **Speed of convergence**: Speed of convergence defines how quickly the routers in the network topology share routing information and reach a state of consistent knowledge. The faster the convergence, the more preferable the protocol. Routing loops can occur when inconsistent routing tables are not updated due to slow convergence in a changing network.
- **Scalability**: Scalability defines how large a network can become, based on the routing protocol that is deployed. The larger the network is, the more scalable the routing protocol needs to be.
- **Classful or classless (use of VLSM)**: Classful routing protocols do not include the subnet mask and cannot support ***variable-length subnet mask (VLSM)***. Classless routing protocols include the subnet mask in the updates. Classless routing protocols support VLSM and better route summarization.

- **Resource usage**: Resource usage includes the requirements of a routing protocol such as memory space (RAM), CPU utilization, and link bandwidth utilization. Higher resource requirements necessitate more powerful hardware to support the routing protocol operation, in addition to the packet forwarding processes.
- **Implementation and maintenance**: Implementation and maintenance describes the level of knowledge that is required for a network administrator to implement and maintain the network based on the routing protocol deployed.

Table 7-2 summarizes the characteristics of each routing protocol.

Table 7-2 Comparing Routing Protocols

	RIPv1	RIPv2	IGRP	EIGRP	OSPF	IS-IS
Speed of Convergence	Slow	Slow	Slow	Fast	Fast	Fast
Scalability – Size of Network	Small	Small	Small	Large	Large	Large
Use of VLSM	No	Yes	No	Yes	Yes	Yes
Resource Usage	Low	Low	Low	Medium	High	High
Implementation and Maintenance	Simple	Simple	Simple	Complex	Complex	Complex

Routing Protocol Metrics

There are cases when a routing protocol learns of more than one route to the same destination. To select the best path, the routing protocol must be able to evaluate and differentiate between the available paths. This is accomplished through the use of routing ***metrics***.

A metric is a measurable value that is assigned by the routing protocol to different routes based on the usefulness of that route. In situations where there are multiple paths to the same remote network, the routing metrics are used to determine the overall "cost" of a path from source to destination. Routing protocols determine the best path based on the route with the lowest cost.

Different routing protocols use different metrics. The metric used by one routing protocol is not comparable to the metric used by another routing protocol. Two different routing protocols might choose different paths to the same destination.

For example, assume that PC1 wants to send a packet to PC2. In Figure 7-16, the RIP routing protocol has been enabled on all routers and the network has converged. RIP makes a routing protocol decision based on the least number of hops. Therefore, when the packet arrives on R1, the best route to reach the PC2 network would be to send it directly to R2 even though the link is much slower that all other links.

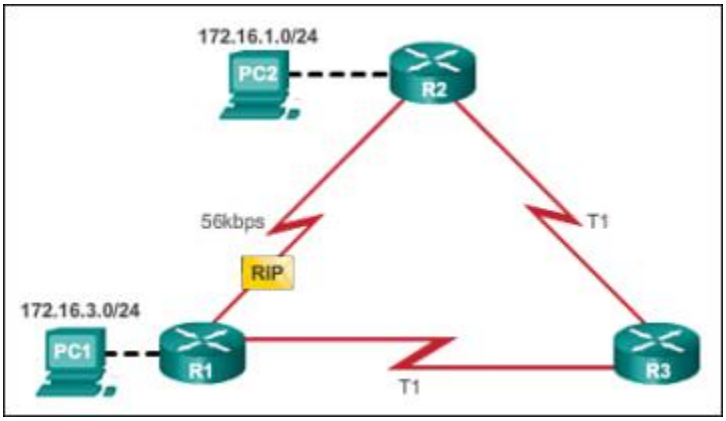


Figure 7-16 RIP Uses Shortest Hop Count Path

In Figure 7-17, the OSPF routing protocol has been enabled on all routers and the network has converged. OSPF makes a routing protocol decision based on the best bandwidth. Therefore, when the packet arrives on R1, the best route to reach the PC2 network would be to send it to R3, which would then forward it to R2.

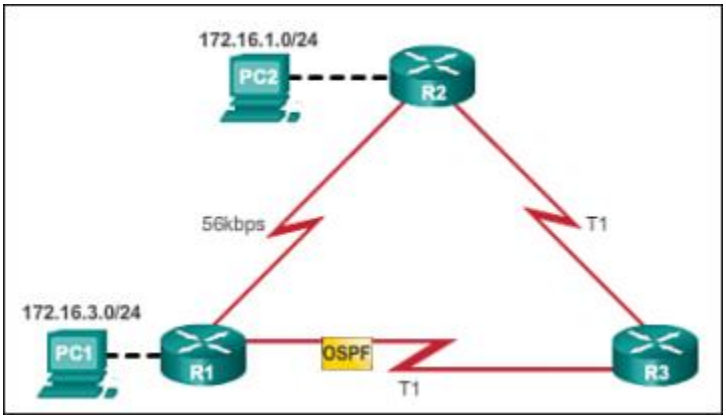


Figure 7-17 OSPF Uses Faster Links

6. Distance Vector Dynamic Routing

This section describes the characteristics, operations, and functionality of distance vector routing protocols. Understanding the operation of distance vector routing is critical to enabling, verifying, and troubleshooting these protocols.

Distance Vector Technologies

Distance vector routing protocols share updates between neighbors. Neighbors are routers that share a link and are configured to use the same routing protocol. The router is only aware of the network addresses of its own interfaces and the remote network addresses it can reach through its neighbors. Routers using distance vector routing are not aware of the network topology.

Some distance vector routing protocols send periodic updates. For example, RIP sends a periodic update to all of its neighbors every 30 seconds. RIP does this even if the topology has not changed; it continues to send updates. RIPv1 reaches all of its neighbors by sending updates to the all-hosts IPv4 address of 255.255.255.255, a broadcast.

The broadcasting of periodic updates is inefficient because the updates consume bandwidth and consume network device CPU resources. Every network device has to process a broadcast message. RIPv2 and EIGRP, instead, use multicast addresses so that only neighbors that need updates will receive them. EIGRP can also send a unicast message to only the affected neighbor. Additionally, EIGRP only sends an update when needed, instead of periodically.

As shown in Figure 7-18, the two modern IPv4 distance vector routing protocols are RIPv2 and EIGRP. RIPv1 and IGRP are listed only for historical accuracy.

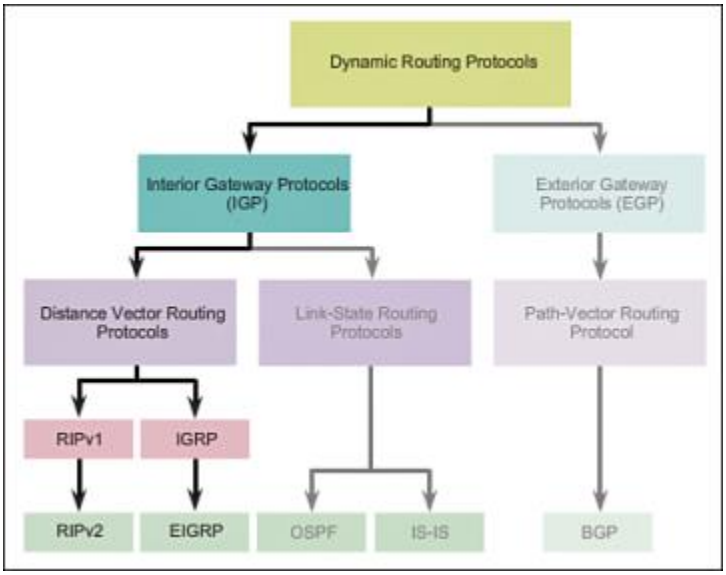


Figure 7-18 Distance Vector Routing Protocols

Distance Vector Algorithm

At the core of the distance vector protocol is the routing algorithm. The algorithm is used to calculate the best paths and then send that information to the neighbors.

The algorithm used for the routing protocols defines the following processes:

- Mechanism for sending and receiving routing information
- Mechanism for calculating the best paths and installing routes in the routing table
- Mechanism for detecting and reacting to topology changes

In the animation in the online course, R1 and R2 are configured with the RIP routing protocol. The algorithm sends and receives updates. Both R1 and R2 then glean new information from the update. In this case, each router learns about a new network. The algorithm on each router makes its calculations independently and updates the routing table with the new information. When the LAN on R2 goes down, the algorithm constructs a triggered update and sends it to R1. R1 then removes the network from the routing table.

Different routing protocols use different algorithms to install routes in the routing table, send updates to neighbors, and make path determination decisions. For example:

- RIP uses the Bellman-Ford algorithm as its routing algorithm. It is based on two algorithms developed in 1958 and 1956 by Richard Bellman and Lester Ford, Jr.
- IGRP and EIGRP use the Diffusing Update Algorithm (DUAL) routing algorithm developed by Dr. J.J. Garcia-Luna-Aceves at SRI International.

7. Types of Distance Vector Routing Protocols

There are two main distance vector routing protocols. This section highlights similarities and differences between RIP and EIGRP.

Routing Information Protocol

The Routing Information Protocol (RIP) was a first generation routing protocol for IPv4 originally specified in RFC 1058. It is easy to configure, making it a good choice for small networks.

RIPv1 has the following key characteristics:

- Routing updates are broadcasted (255.255.255.255) every 30 seconds.
- The hop count is used as the metric for path selection.
- A hop count greater than 15 hops is deemed infinite (too far). That 15th hop router would not propagate the routing update to the next router.

In 1993, RIPv1 evolved to a classless routing protocol known as RIP version 2 (RIPv2). RIPv2 introduced the following improvements:

- **Classless routing protocol**: It supports VLSM and CIDR, because it includes the subnet mask in the routing updates.
- **Increased efficiency**: It forwards updates to multicast address 224.0.0.9, instead of the broadcast address 255.255.255.255.
- **Reduced routing entries**: It supports manual route summarization on any interface.
- **Secure**: It supports an authentication mechanism to secure routing table updates between neighbors.

Table 7-3 summarizes the differences between RIPv1 and RIPv2.

Table 7-3 RIPv1 versus RIPv2

Characteristics and Features	RIPv1	RIPv2
Metric	Both use hop count as a simple metric. The maximum number of hops is 15.	
Updates Forwarded to Address	255.255.255.255	224.0.0.9
Supports VLSM	No	Yes
Supports CIDR	No	Yes
Supports Summarization	No	Yes
Supports Authentication	No	Yes

RIP updates are encapsulated into a UDP segment, with both source and destination port numbers set to UDP port 520.

In 1997, the IPv6-enabled version of RIP was released. RIPng is based on RIPv2. It still has a 15-hop limitation and the administrative distance is 120.

Enhanced Interior Gateway Routing Protocol

The Interior Gateway Routing Protocol (IGRP) was the first proprietary IPv4 routing protocol, developed by Cisco in 1984. It used the following design characteristics:

- Bandwidth, delay, load, and reliability are used to create a composite metric.
- Routing updates are broadcast every 90 seconds, by default.

In 1992, IGRP was replaced by Enhanced IGRP (EIGRP). Like RIPv2, EIGRP also introduced support for VLSM and CIDR. EIGRP increases efficiency, reduces routing updates, and supports secure message exchange.

Table 7-4 summarizes the differences between IGRP and EIGRP.

Table 7-4 IGRP versus EIGRP

Characteristics and Features	IGRP	EIGRP
Metric	Both use a composite metric based on bandwidth and delay. Reliability and load can also be included in the metric calculation if configured.	
Updates Forwarded to Address	255.255.255.255	224.0.0.10
Supports VLSM	No	Yes

Characteristics and Features	IGRP	EIGRP
Supports CIDR	No	Yes
Supports Summarization	No	Yes
Supports Authentication	No	Yes

EIGRP also introduced:

- ***Bounded triggered updates***: It does not send periodic updates. Only routing table changes are propagated, whenever a change occurs. This reduces the amount of load the routing protocol places on the network. Bounded triggered updates means that EIGRP only sends to the neighbors that need it. It uses less bandwidth, especially in large networks with many routes.
- ***Hello keepalive mechanism***: A small Hello message is periodically exchanged to maintain adjacencies with neighboring routers. This means a very low usage of network resources during normal operation, instead of the periodic updates.
- **Maintains a topology table**: Maintains all the routes received from neighbors (not only the best paths) in a topology table. DUAL can insert backup routes into the EIGRP topology table.
- **Rapid convergence**: In most cases, it is the fastest IGP to converge because it maintains alternate routes, enabling almost instantaneous convergence. If a primary route fails, the router can use the alternate route identified. The switchover to the alternate route is immediate and does not involve interaction with other routers.
- **Multiple network layer protocol support**: EIGRP uses Protocol Dependent Modules (PDM), which means that it is the only protocol to include support for protocols other than IPv4 and IPv6, such as legacy IPX and AppleTalk.

8. RIP and RIPng Routing

Although the use of RIP has decreased in the past decade, it is still important to your networking studies because it might be encountered in a network implementation. As well, understanding how RIP operates and knowing its implementation will make learning other routing protocols easier.

Configuring the RIP Protocol

In this section, you will learn how to configure, verify, and troubleshoot RIPv2.

Router RIP Configuration Mode

Although RIP is rarely used in modern networks, it is useful as a foundation for understanding basic network routing. For this reason, this section provides a brief overview of how to configure basic RIP settings and to verify RIPv2.

Refer to the reference topology in Figure 7-19 and the addressing table in Table 3-7.

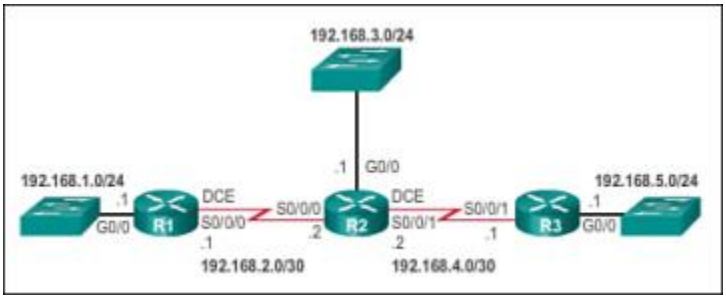


Figure 7-19 RIP Reference Topology

Table 7-5 Addressing Table

Device	Interface	IP Address	Subnet Mask
R1	G0/0	192.168.1.1	255.255.255.0
	S0/0/0	192.168.2.1	255.255.255.0
R2	G0/0	192.168.3.1	255.255.255.0
	S0/0/0	192.168.2.2	255.255.255.0
	S0/0/1	192.168.4.2	255.255.255.0
R3	G0/0	192.168.5.1	255.255.255.0
	S0/0/1	192.168.4.1	255.255.255.0

In this scenario, all routers have been configured with basic management features and all interfaces identified in the reference topology are configured and enabled. There are no static routes configured and no routing protocols enabled; therefore, remote network access is currently impossible. RIPv2 is used as the dynamic routing protocol.

To enable RIP, use the **router rip** command to enter router configuration mode, as shown in the following output. This command does not directly start the RIP process. Instead, it provides access to the router configuration mode where the RIP routing settings are configured.

```
R1# conf t

Enter configuration commands, one per line. End with CNTL/Z.

R1(config)# router rip

R1(config-router)#
```

To disable and eliminate RIP, use the **no router rip** global configuration command. This command stops the RIP process and erases all existing RIP configurations.

Figure 7-20 displays a partial list of the various RIP commands that can be configured. This section covers the two highlighted commands as well as **network**, **passive-interface**, and **version**.

```
R1(config-router)# ?
Router configuration commands:
  address-family          Enter Address Family command mode
  auto-summary            Enable automatic network number
                          summarization
  default                 Set a command to its defaults
  default-information     Control distribution of default
                          information
  default-metric          Set metric of redistributed routes
  distance                Define an administrative distance
  distribute-list         Filter networks in routing updates
  exit                    Exit from routing protocol
                          configuration mode
  flash-update-threshold  Specify flash update threshold in
                          second
  help                    Description of the interactive help
                          system
  input-queue             Specify input queue depth
```

Figure 7-20 RIP Configuration Options

Advertising Networks

By entering the RIP router configuration mode, the router is instructed to run RIP. But the router still needs to know which local interfaces it should use for communication with other routers, as well as which locally connected networks it should advertise to those routers.

To enable RIP routing for a network, use the **network** *network-address* router configuration mode command. Enter the classful network address for each directly connected network. This command:

- Enables RIP on all interfaces that belong to a specific network. Associated interfaces now both send and receive RIP updates.
- Advertises the specified network in RIP routing updates sent to other routers every 30 seconds.

If a subnet address is entered, the IOS automatically converts it to the classful network address. Remember RIPv1 is a classful routing protocol for IPv4. For example, entering the **network 192.168.1.32** command would automatically be converted to **network 192.168.1.0** in the running configuration file. The IOS does not give an error message, but instead corrects the input and enters the classful network address.

In the following command sequence, the **network** command is used to advertise the R1 directly connected networks.

```
R1(config)# router rip

R1(config-router)# network 192.168.1.0

R1(config-router)# network 192.168.2.0

R1(config-router)#
```

Examining Default RIP Settings

The output of the **show ip protocols** command in Figure 7-21 displays the IPv4 routing protocol settings currently configured on the router.

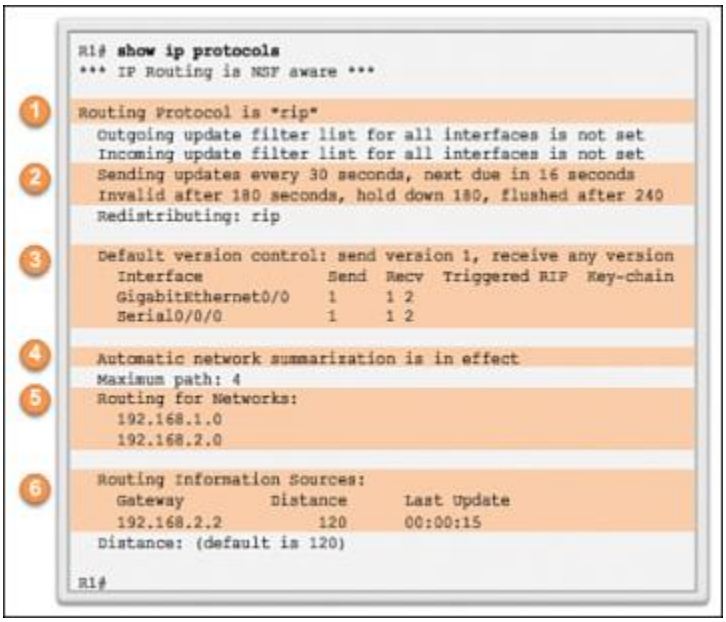


Figure 7-21 Verifying RIP Settings on R1

This output confirms that:

1. RIP routing is configured and running on router R1.
2. The values of various timers; for example, the next routing update is sent by R1 in 16 seconds.
3. The version of RIP configured is currently RIPv1.
4. R1 is currently summarizing at the classful network boundary.
5. The classful networks are advertised by R1. These are the networks that R1 includes in its RIP updates.
6. The RIP neighbors are listed, including their next-hop IP address, the associated AD that R2 uses for updates sent by this neighbor, and when the last update was received from this neighbor.

This command is also very useful when verifying the operations of other routing protocols (i.e., EIGRP and OSPF).

The **show ip route** command displays the RIP routes installed in the routing table. InFigure 7-22, R1 now knows about the highlighted networks.

```
R1# show ip route | begin Gateway
Gateway of last resort is not set

      192.168.1.0/24 is variably subnetted, 2 subnets, 2 masks
C        192.168.1.0/24 is directly connected, GigabitEthernet0/0
L        192.168.1.1/32 is directly connected, GigabitEthernet0/0
      192.168.2.0/24 is variably subnetted, 2 subnets, 2 masks
C        192.168.2.0/24 is directly connected, Serial0/0/0
L        192.168.2.1/32 is directly connected, Serial0/0/0
R     192.168.3.0/24 [120/1] via 192.168.2.2, 00:00:24, Serial0/0/0
R     192.168.4.0/24 [120/1] via 192.168.2.2, 00:00:24, Serial0/0/0
R     192.168.5.0/24 [120/2] via 192.168.2.2, 00:00:24, Serial0/0/0
R1#
```

Figure 7-22 Verifying RIP Routes on R1

Enabling RIPv2

By default, when a RIP process is configured on a Cisco router, it is running RIPv1, as shown in the following output:

```
R1# show ip protocols

*** IP Routing is NSF aware ***


Routing Protocol is "rip"

  Outgoing update filter list for all interfaces is not set

  Incoming update filter list for all interfaces is not set

  Sending updates every 30 seconds, next due in 16 seconds

  Invalid after 180 seconds, hold down 180, flushed after 240

  Redistributing: rip

  Default version control: send version 1, receive any version

    Interface              Send  Recv  Triggered RIP  Key-chain
```

```
    GigabitEthernet0/0    1      1 2
    Serial0/0/0           1      1 2
  Automatic network summarization is in effect
  Maximum path: 4
  Routing for Networks:
    192.168.1.0
    192.168.2.0


Routing Information Sources:
    Gateway         Distance      Last Update
    192.168.2.2          120      00:00:15
  Distance: (default is 120)

R1#
```

However, even though the router only sends RIPv1 messages, it can interpret both RIPv1 and RIPv2 messages. A RIPv1 router ignores the RIPv2 fields in the route entry.

Use the **version 2** router configuration mode command to enable RIPv2, as shown in Figure 7-23.

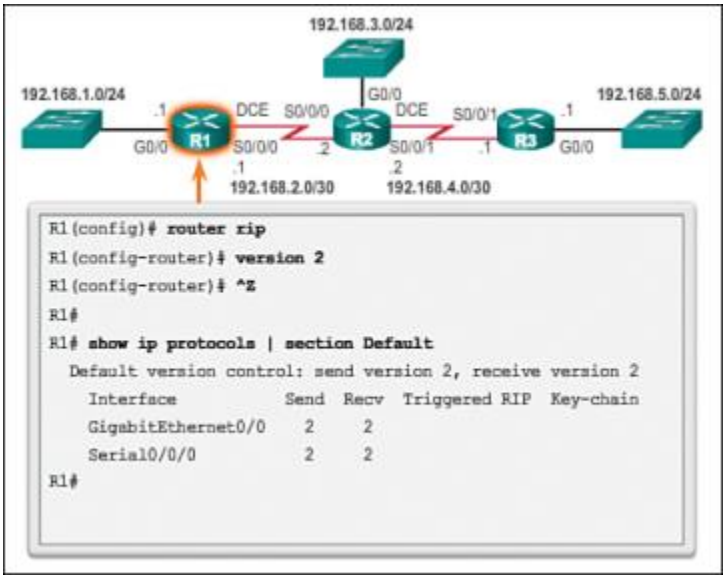


Figure 7-23 Enable and Verify RIPv2 on R1

Notice how the **show ip protocols** command verifies that R2 is now configured to send and receive version 2 messages only. The RIP process now includes the subnet mask in all updates, making RIPv2 a classless routing protocol.

Configuring **version 1** enables RIPv1 only, while configuring **no version** returns the router to the default setting of sending version 1 updates but listening for version 1 or version 2 updates.

The following output verifies that there are no RIP routes still in the routing table:

```
R1# show ip route | begin Gateway

Gateway of last resort is not set


      192.168.1.0/24 is variably subnetted, 2 subnets, 2 masks

C        192.168.1.0/24 is directly connected, GigabitEthernet0/0

L        192.168.1.1/32 is directly connected, GigabitEthernet0/0

      192.168.2.0/24 is variably subnetted, 2 subnets, 2 masks
```

```
C        192.168.2.0/24 is directly connected, Serial0/0/0

L        192.168.2.1/32 is directly connected, Serial0/0/0

R1#
```

There are no RIP routes because R1 is now only listening for RIPv2 updates. R2 and R3 are still sending RIPv1 updates. Therefore, the **version 2** command must be configured on all routers in the routing domain.

Disabling Auto Summarization

As shown in Figure 7-24, RIPv2 automatically summarizes networks at major network boundaries by default, just like RIPv1.

```
R1# show ip protocols
*** IP Routing is NSF aware ***

Routing Protocol is "rip"
 Outgoing update filter list for all interfaces is not set
 Incoming update filter list for all interfaces is not set
 Sending updates every 30 seconds, next due in 16 seconds
 Invalid after 180 seconds, hold down 180, flushed after
 240
 Redistributing: rip
 Default version control: send version 2, receive version 2
    Interface         Send  Recv  Triggered RIP  Key-chain
    GigabitEthernet0/0    1     1 2
    Serial0/0/0           1     1 2
 Automatic network summarization is in effect
 Maximum path: 4
 Routing for Networks:
    192.168.1.0
    192.168.2.0
 Routing Information Sources:
    Gateway         Distance      Last Update
    192.168.2.2          120      00:00:15
 Distance: (default is 120)
R1#
```

Figure 7-24 Verify RIPv2 Route Summarization

To modify the default RIPv2 behavior of ***automatic summarization***, use the **no auto-summary** router configuration mode command as shown in the following command sequence:

```
R1(config)# router rip

R1(config-router)# no auto-summary

R1(config-router)# end

R1#

*Mar 10 14:11:49.659: %SYS-5-CONFIG_I: Configured from console by console

R1# show ip protocols | section Automatic
```

```
  Automatic network summarization is not in effect

R1#
```

This command has no effect when using RIPv1. When automatic summarization has been disabled, RIPv2 no longer summarizes networks to their classful address at boundary routers. RIPv2 now includes all subnets and their appropriate masks in its routing updates. The **show ip protocols** output now states that automatic network summarization is not in effect.

RIPv2 must be enabled before automatic summarization is disabled.

Configuring Passive Interfaces

By default, RIP updates are forwarded out all RIP-enabled interfaces. However, RIP updates really only need to be sent out interfaces connecting to other RIP-enabled routers.

For instance, refer to the topology in Figure 7-24. RIP sends updates out of its Gigabit Ethernet 0/0 interface even though no RIP device exists on that LAN. R1 has no way of knowing this and, as a result, sends an update every 30 seconds. Sending out unneeded updates on a LAN impacts the network in three ways:

- **Wasted bandwidth**: Bandwidth is used to transport unnecessary updates. Because RIP updates are either broadcasted or multicasted, switches also forward the updates out all ports.
- **Wasted resources**: All devices on the LAN must process the update up to the transport layers, at which point the devices will discard the update.
- **Security risk**: Advertising updates on a broadcast network is a security risk. RIP updates can be intercepted with packet sniffing software. Routing updates can be modified and sent back to the router, corrupting the routing table with false metrics that misdirect traffic.

To address these problems, an interface can be configured to stop sending routing updates. This is referred to as configuring a ***passive interface***. Use the **passive-interface** router configuration command to prevent the transmission of routing updates through a router interface but still allow that network to be advertised to other routers. The command stops routing updates out the specified interface. However, the network that the specified interface belongs to is still advertised in routing updates that are sent out other interfaces.

There is no need for R1, R2, and R3 to forward RIP updates out of their LAN interfaces. The configuration in Figure 7-25 identifies the R1 Gigabit Ethernet 0/0 interface as passive.

```
R1(config)# router rip
R1(config-router)# passive-interface g0/0
R1(config-router)# end
R1#
R1# show ip protocols | begin Default
  Default version control: send version 2, receive version 2
    Interface              Send  Recv  Triggered RIP  Key-chain
    Serial0/0/0            2     2
  Automatic network summarization is not in effect
  Maximum path: 4
  Routing for Networks:
    192.168.1.0
    192.168.2.0
  Passive Interface(s):
    GigabitEthernet0/0
  Routing Information Sources:
    Gateway         Distance      Last Update
    192.168.2.2           120      00:00:06
  Distance: (default is 120)

R1#
```

Figure 7-25 Configuring and Verifying a Passive Interface on R1

The **show ip protocols** command is then used to verify that the Gigabit Ethernet interface was passive. Notice that the Gigabit Ethernet 0/0 interface is no longer listed as sending or receiving version 2 updates, but instead is now listed under the Passive Interface(s) section. Also notice that the network 192.168.1.0 is still listed under Routing for Networks, which means that this network is still included as a route entry in RIP updates that are sent to R2.

All routing protocols support the **passive-interface** command.

As an alternative, all interfaces can be made passive using the **passive-interface default** command. Interfaces that should not be passive can be re-enabled using the **no passive-interface** command.

Propagating a Default Route

In the topology in Figure 7-26, R1 is single-homed to a service provider. Therefore, all that is required for R1 to reach the Internet is a ***default static route*** going out of the Serial 0/0/1 interface.

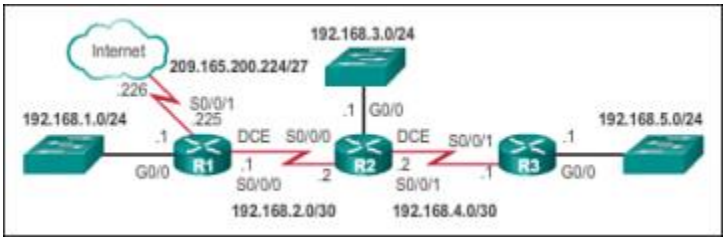


Figure 7-26 Propagating a Default Route on R1

Similar default static routes could be configured on R2 and R3, but it is much more scalable to enter it one time on the edge router R1 and then have R1 propagate it to all other routers using RIP. To provide Internet connectivity to all other networks in the RIP routing domain, the default static route needs to be advertised to all other routers that use the dynamic routing protocol.

To propagate a default route, the edge router must be configured with:

- A default static route using the **ip route 0.0.0.0 0.0.0.0** *exit-intf next-hop-ip* command.
- The **default-information originate** router configuration command. This instructs R1 to originate default information, by propagating the static default route in RIP updates.

The example in Figure 7-27 configures a fully specified default static route to the service provider, and then the route is propagated by RIP. Notice that R1 now has a Gateway of Last Resort and default route installed in its routing table.

```
R1(config)# ip route 0.0.0.0 0.0.0.0 S0/0/1 209.165.200.226
R1(config)# router rip
R1(config-router)# default-information originate
R1(config-router)# ^Z
R1#
*Mar 10 23:33:51.801: %SYS-5-CONFIG_I: Configured from console by
console
R1# show ip route | begin Gateway
Gateway of last resort is 209.165.200.226 to network 0.0.0.0

S*    0.0.0.0/0 [1/0] via 209.165.200.226, Serial0/0/1
      192.168.1.0/24 is variably subnetted, 2 subnets, 2 masks
C        192.168.1.0/24 is directly connected, GigabitEthernet0/0
L        192.168.1.1/32 is directly connected, GigabitEthernet0/0
      192.168.2.0/24 is variably subnetted, 2 subnets, 2 masks
C        192.168.2.0/24 is directly connected, Serial0/0/0
L        192.168.2.1/32 is directly connected, Serial0/0/0
R     192.168.3.0/24 [120/1] via 192.168.2.2, 00:00:08,
Serial0/0/0
R     192.168.4.0/24 [120/1] via 192.168.2.2, 00:00:08,
Serial0/0/0
R     192.168.5.0/24 [120/2] via 192.168.2.2, 00:00:08,
Serial0/0/0
      209.165.200.0/24 is variably subnetted, 2 subnets, 2 masks
C        209.165.200.0/24 is directly connected, Serial0/0/1
```

Figure 7-27 Configuring and Verifying a Default Route on R1

Although RIP is rarely used in modern networks, it is useful as a foundation for understanding basic network routing. In this activity, you will configure a default route, configure RIP version 2 with appropriate network statements and passive interfaces, and verify full connectivity.

Configuring the RIPng Protocol

In this section, you will learn how to configure, verify, and troubleshoot RIPng.

Advertising IPv6 Networks

As with its IPv4 counterpart, RIPng is rarely used in modern networks. It is also useful as a foundation for understanding basic network routing. For this reason, this section provides a brief overview of how to configure basic ***RIPng***.

Refer to the reference topology in Figure 7-28.

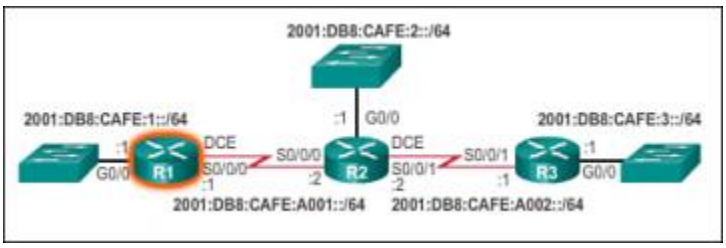


Figure 7-28 Enabling RIPng on the R1 Interfaces

In this scenario, all routers have been configured with basic management features and all interfaces identified in the reference topology are configured and enabled. There are no static routes configured and no routing protocols enabled; therefore, remote network access is currently impossible.

To enable an IPv6 router to forward IPv6 packets, **ipv6 unicast-routing** must be configured.

Unlike RIPv2, RIPng is enabled on an interface and not in router configuration mode. In fact, there is no **network** *network-address* command available in RIPng. Instead, use the **ipv6 rip** *domain-name* **enable** interface configuration command.

In the following output, IPv6 unicast routing is enabled and the Gigabit Ethernet 0/0 and Serial 0/0/0 interfaces are enabled for RIPng using the domain name RIP-AS:

```
R1(config)# ipv6 unicast-routing

R1(config)#

R1(config)# interface gigabitethernet 0/0

R1(config-if)# ipv6 rip RIP-AS enable

R1(config-if)# exit

R1(config)#
```

```
R1(config)# interface serial 0/0/0

R1(config-if)# ipv6 rip RIP-AS enable

R1(config-if)# no shutdown

R1(config-if)#
```

The process to propagate a default route in RIPng is identical to RIPv2 except that an IPv6 default static route must be specified. For example, assume that R1 had an Internet connection from a Serial 0/0/1 interface to IP address 2001:DB8:FEED:1::1/64. To propagate a default route, R1 would have to be configured with:

- A default static route using the **ipv6 route 0::/0 2001:DB8:FEED:1::1** global configuration command.
- The **ipv6 rip** *domain-name* **default-information originate** interface configuration mode command. For example, the Serial 0/0/1 interface of R1 would have to be configured with the **ipv6 rip RIP-AS default-information originate** command. This would instruct R1 to be the source of the default route information and propagate the default static route in RIPng updates sent out of the RIPng-enabled interfaces.

Examining the RIPng Configuration

In Figure 7-29, the **show ipv6 protocols** command does not provide the same amount of information as its IPv4 counterpart.

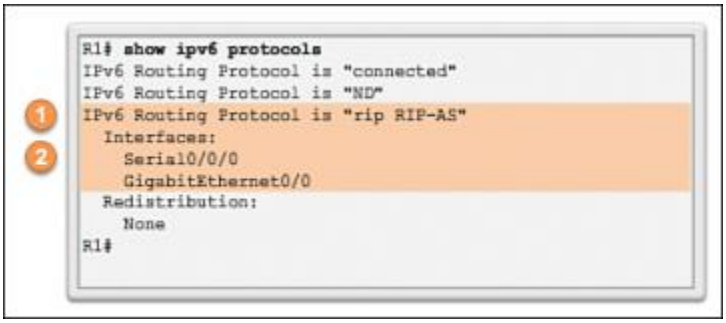


Figure 7-29 Verifying RIPng Settings on R1

However, the command does confirm the following parameters:

1. That RIPng routing is configured and running on router R1.
2. The interfaces configured with RIPng.

The **show ipv6 route** command displays the routes installed in the routing table as shown in Figure 7-30. The output confirms that R1 now knows about the highlighted RIPng networks.

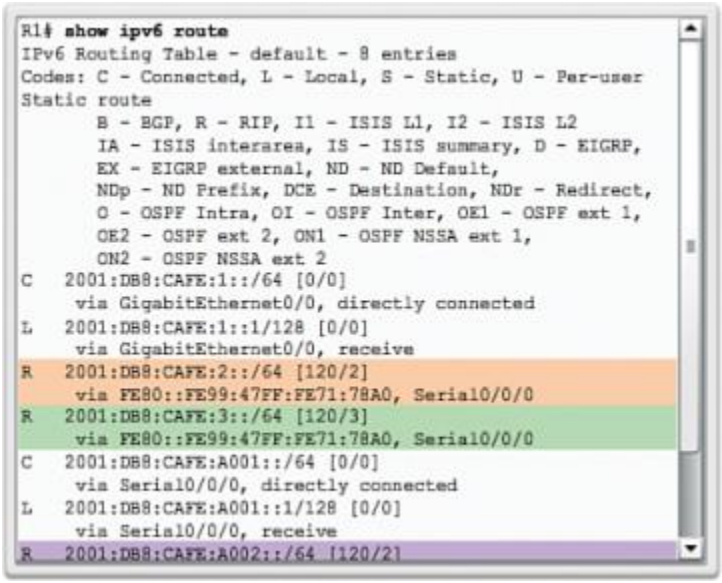


Figure 7-30 Verifying Routes on R1

Notice that the R2 LAN is advertised as two hops away. This is because there is a difference in the way RIPv2 and RIPng calculate the hop counts. With RIPv2 (and RIPv1), the metric to the R2 LAN would be one hop. This is because the metric (hop count) that is displayed in the IPv4 routing table is the number of hops required to reach the remote network (counting the next-hop router as the first hop). In RIPng, the sending router already considers itself to be one hop away; therefore, R2 advertises its LAN with a metric of 1. When R1 receives the update, it adds another hop count of 1 to the metric. Therefore, R1 considers the R2 LAN to be two hops away. Similarly it considers the R3 LAN to be three hops away.

Appending the **rip** keyword to the command as shown in Figure 7-31 only lists RIPng networks.

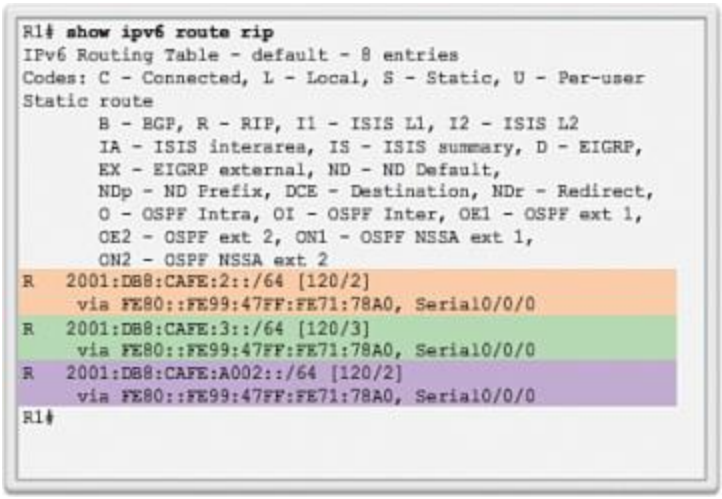


Figure 7-31 Verifying RIPng Routes on R1

RIPng (RIP Next Generation) is a distance vector routing protocol for routing IPv6 addresses. RIPng is based on RIPv2 and has the same administrative distance and 15-hop limitation. This activity will help you become more familiar with RIPng.

9. Link-State Dynamic Routing

Distance vector routing protocols are thought to be simple to understand, whereas link-state routing protocols have the reputation of being very complex, even intimidating. However, link-state routing protocols and concepts are not difficult to understand. In many ways, the link-state process is simpler to understand than distance vector concepts.

Link-State Routing Protocol Operation

This section describes the characteristics, operations, and functionality of link-state routing protocols. Understanding the operation of link-state routing is critical to enabling, verifying, and troubleshooting these protocols.

Shortest Path First Protocols

Link-state routing protocols are also known as shortest path first protocols and are built around Edsger Dijkstra's shortest path first (SPF) algorithm. The SPF algorithm is discussed in more detail in a later section.

The IPv4 link-state routing protocols are shown Figure 7-32:

- Open Shortest Path First (OSPF)
- Intermediate System-to-Intermediate System (IS-IS)

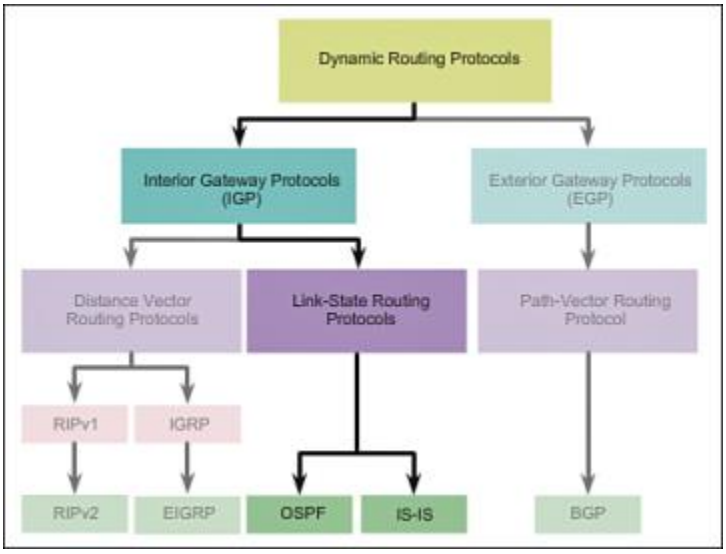


Figure 7-32 Link-State Routing Protocols

Link-state routing protocols have the reputation of being much more complex than their distance vector counterparts. However, the basic functionality and configuration of link-state routing protocols is equally straightforward.

Just like RIP and EIGRP, basic OSPF operations can be configured using the:

- **router ospf** *process-id* global configuration command
- **network** command to advertise networks

Dijkstra's Algorithm

All link-state routing protocols apply ***Dijkstra's algorithm*** to calculate the best path route. The algorithm is commonly referred to as the ***shortest path first (SPF)*** algorithm. This algorithm uses accumulated costs along each path, from source to destination, to determine the total cost of a route.

In Figure 7-33, each path is labeled with an arbitrary value for cost.

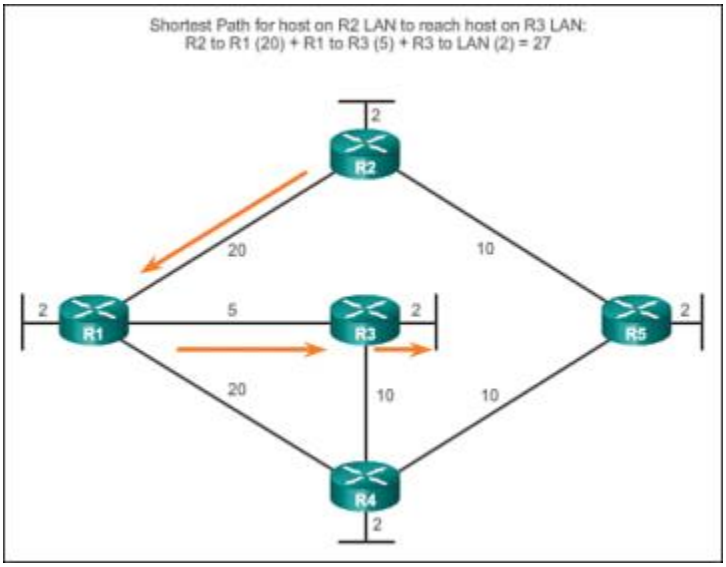


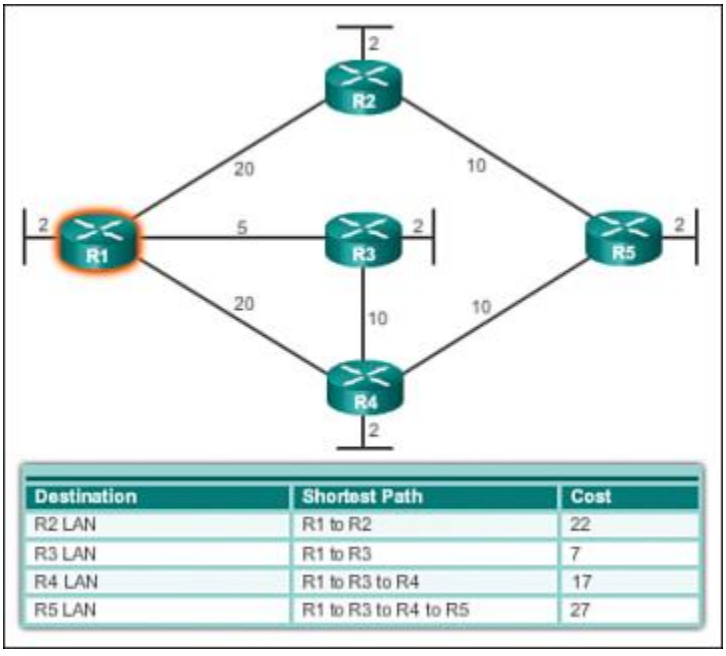
Figure 7-33 Dijkstra's Shortest Path First Algorithm

The cost of the shortest path for R2 to send packets to the LAN attached to R3 is 27. Specifically, the cost is R2 to R1 (20) plus R1 to R3 (5) plus R3 to LAN (2). Each router determines its own cost to each destination in the topology. In other words, each router calculates the SPF algorithm and determines the cost from its own perspective.

The focus of this section is on cost, which is determined by the SPF tree. For this reason, the graphics throughout this section show the connections of the SPF tree, not the topology. All links are represented with a solid black line.

SPF Example

The table in Figure 7-34 displays the shortest path and the accumulated cost to reach the identified destination networks from the perspective of R1.



Destination	Shortest Path	Cost
R2 LAN	R1 to R2	22
R3 LAN	R1 to R3	7
R4 LAN	R1 to R3 to R4	17
R5 LAN	R1 to R3 to R4 to R5	27

Figure 7-34 R1 SPF Tree

The shortest path is not necessarily the path with the least number of hops. For example, look at the path to the R5 LAN. It might be assumed that R1 would send directly to R4 instead of to R3. However, the cost to reach R4 directly (22) is higher than the cost to reach R4 through R3 (17).

Observe the shortest path for each router to reach each of the LANs, as shown in Tables 7-6 through 7-8.

Table 7-6 R2 SPF Tree

Destination	Shortest Path	Cost
R1 LAN	R2 to R1	22
R3 LAN	R2 to R1 to R3	27
R4 LAN	R2 to R5 to R4	22
R5 LAN	R2 to R5	12

Table 7-7 R3 SPF Tree

Destination	Shortest Path	Cost
R1 LAN	R3 to R1	7
R2 LAN	R3 to R1 to R2	27
R4 LAN	R3 to R4	12
R5 LAN	R3 to R4 to R5	22

Table 7-7 R4 SPF Tree

Destination	Shortest Path	Cost
R1 LAN	R4 to R3 to R1	17
R2 LAN	R4 to R5 to R2	22
R3 LAN	RR4 to R3	12
R5 LAN	R4 to R5	12

Table 7-8 R5 SPF Tree

Destination	Shortest Path	Cost
R1 LAN	R5 to R4 to R3 to R1	27
R2 LAN	R5 to R2	12
R3 LAN	R5 to R4 to R3	22
R4 LAN	R5 to R4	12

Link-State Updates

Link-state updates (LSUs) are the packets used for OSPF routing updates. This section discusses how OSPF exchanges LSUs to discover the best routes.

Link-State Routing Process

So exactly how does a link-state routing protocol work? With link-state routing protocols, a link is an interface on a router. Information about the state of those links is known as link-states.

All routers in an OSPF area will complete the following generic link-state routing process to reach a state of convergence:

1. Each router learns about its own links and its own directly connected networks. This is done by detecting that an interface is in the up state.
2. Each router is responsible for meeting its neighbors on directly connected networks. Link-state routers do this by exchanging Hello packets with other link-state routers on directly connected networks.
3. Each router builds a ***link-state packet (LSP)***containing the state of each directly connected link. This is done by recording all the pertinent information about each neighbor, including neighbor ID, link type, and bandwidth.
4. Each router floods the LSP to all neighbors. Those neighbors store all LSPs received in a database. They then flood the LSPs to their neighbors until all routers in the area have received the LSPs. Each router stores a copy of each LSP received from its neighbors in a local database.
5. Each router uses the database to construct a complete map of the topology and computes the best path to each destination network. Like having a road map, the router now has a complete map of all destinations in the topology and the routes to reach them. The SPF algorithm is used to construct the map of the topology and to determine the best path to each network.

This process is the same for both OSPF for IPv4 and OSPF for IPv6. The examples in this section refer to OSPF for IPv4.

Link and Link-State

The first step in the link-state routing process is that each router learns about its own links, its own directly connected networks. When a router interface is configured with an IP address and subnet mask, the interface becomes part of that network.

Refer to the topology in Figure 7-35. For purposes of this discussion, assume that R1 was previously configured and had full connectivity to all neighbors. However, R1 lost power briefly and had to restart.

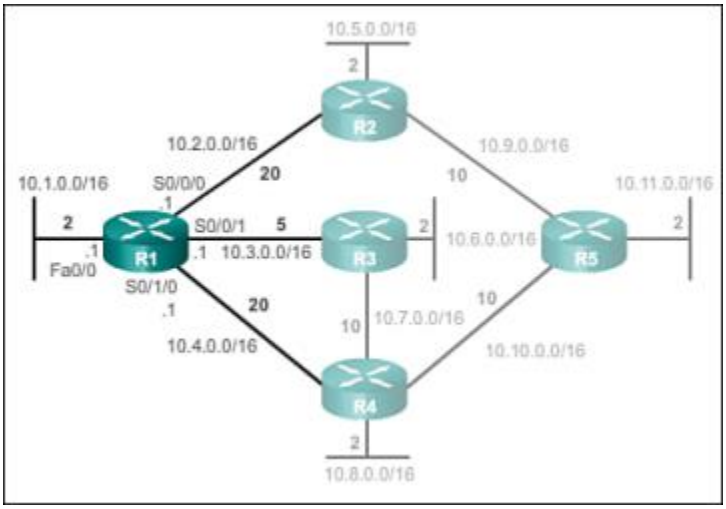


Figure 7-35 R1 Links

During boot up R1 loads the saved startup configuration file. As the previously configured interfaces become active, R1 learns about its own directly connected networks. Regardless of the routing protocols used, these directly connected networks are now entries in the routing table.

As with distance vector protocols and static routes, the interface must be properly configured with an IPv4 address and subnet mask, and the link must be in the up state before the link-state routing protocol can learn about a link. Also, like distance vector protocols, the interface must be included in one of the **network** router configuration statements before it can participate in the link-state routing process.

Figure 7-35 shows R1 linked to four directly connected networks:

- FastEthernet 0/0: 10.1.0.0/16
- Serial 0/0/0: 10.2.0.0/16
- Serial 0/0/1: 10.3.0.0/16
- Serial 0/1/0: 10.4.0.0/16

As shown in Figures 7-36 through 7-39, the link-state information includes:

- The interface's IPv4 address and subnet mask
- The type of network, such as Ethernet (broadcast) or Serial point-to-point link
- The cost of that link
- Any neighbor routers on that link

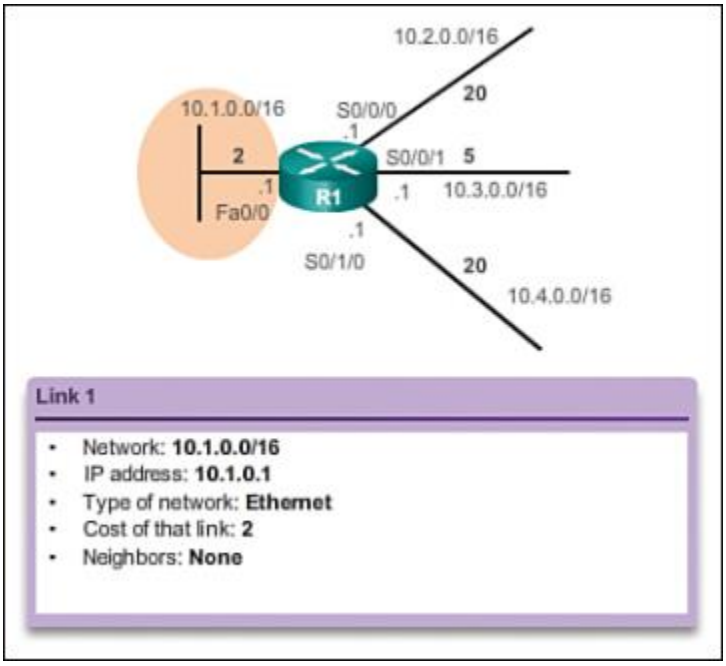


Figure 7-36 Link-State of Interface Fa0/0

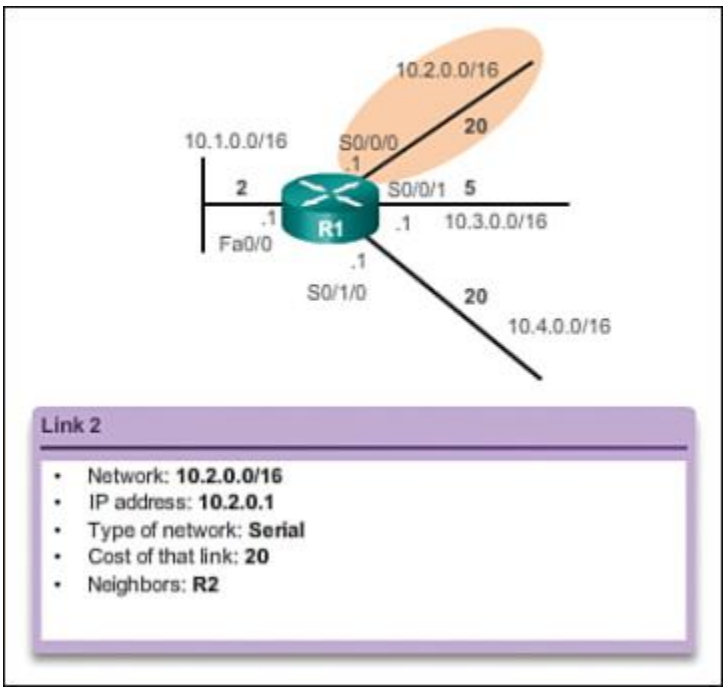


Figure 7-37 Link-State of Interface S0/0/0

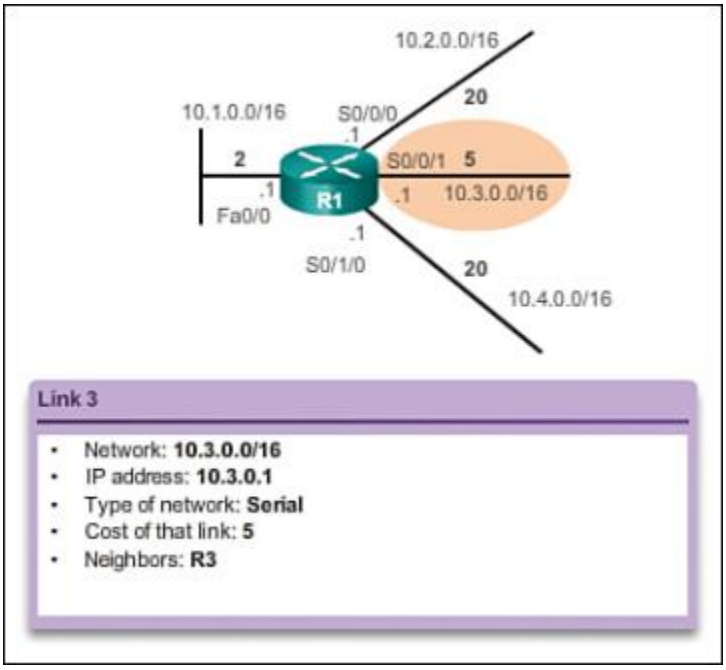


Figure 7-38 Link-State of Interface S0/0/1

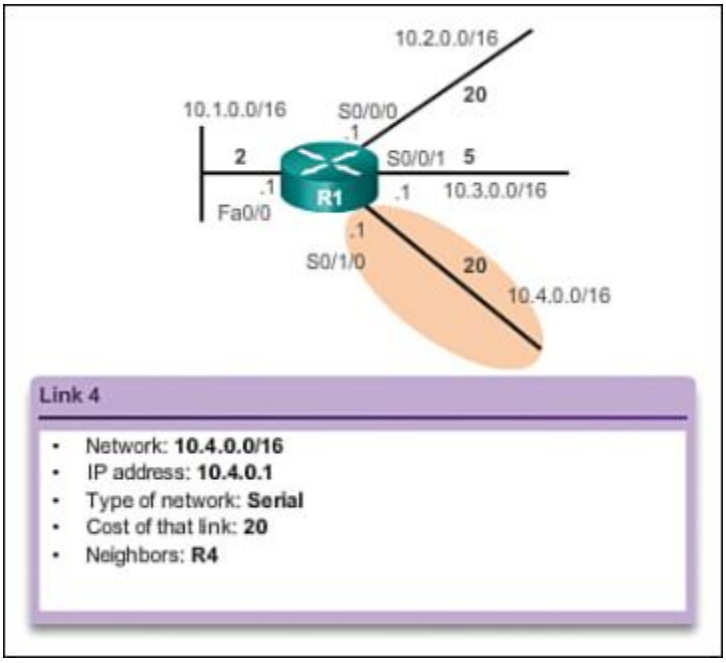


Figure 7-39 Link-State of Interface S0/1/0

Cisco's implementation of OSPF specifies the OSPF routing metric as the cost of the link based on the bandwidth of the outgoing interface. For the purposes of this chapter, we are using arbitrary cost values to simplify the demonstration.

Say Hello

The second step in the link-state routing process is that each router is responsible for meeting its neighbors on directly connected networks.

Routers with link-state routing protocols use a Hello protocol to discover any neighbors on their links. A neighbor is any other router that is enabled with the same link-state routing protocol.

In Figure 7-40, R1 sends Hello packets out its links (interfaces) to discover if there are any neighbors.

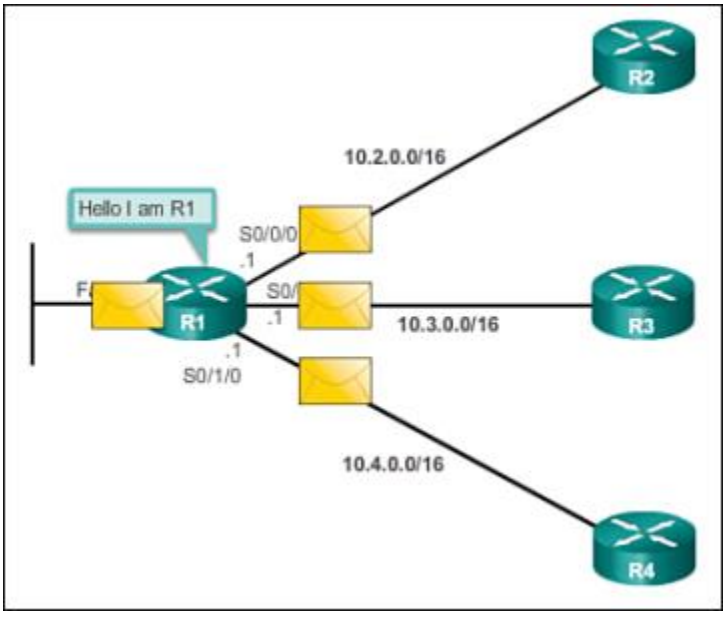


Figure 7-40 R1 Sends Hello Packets

In Figure 7-41, R2, R3, and R4 reply to the Hello packet with their own Hello packets because these routers are configured with the same link-state routing protocol. There are no neighbors out the FastEthernet 0/0 interface. Because R1 does not receive a Hello on this interface, it does not continue with the link-state routing process steps for the FastEthernet 0/0 link.

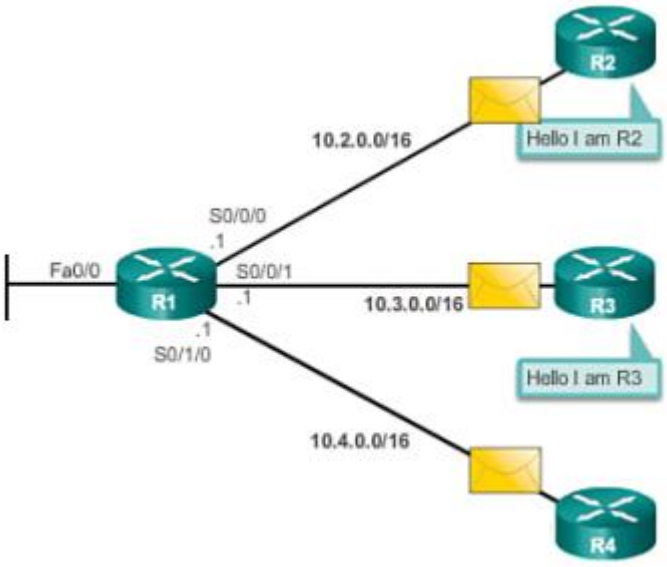


Figure 7-41 R2, R3, and R4 Reply with Hello Packets

When two link-state routers learn that they are neighbors, they form an adjacency. These small Hello packets continue to be exchanged between two adjacent neighbors and serve as a keepalive function to monitor the state of the neighbor. If a router stops receiving Hello packets from a neighbor, that neighbor is considered unreachable and the adjacency is broken. In OSPF, for broadcast multiacess and point-to-point links, the default is 10 seconds. On NBMA, the default is 30 seconds.

Building the Link-State Packet

The third step in the link-state routing process is that each router builds an LSP containing the state of each directly connected link.

After a router has established its adjacencies, it can build its LSPs that contain the link-state information about its links. A simplified version of the LSP from R1 displayed in Figure 7-42 would contain the following:

1. R1; Ethernet network 10.1.0.0/16; Cost 2
2. R1 -> R2; Serial point-to-point network; 10.2.0.0/16; Cost 20
3. R1 -> R3; Serial point-to-point network; 10.3.0.0/16; Cost 5
4. R1 -> R4; Serial point-to-point network; 10.4.0.0/16; Cost 20

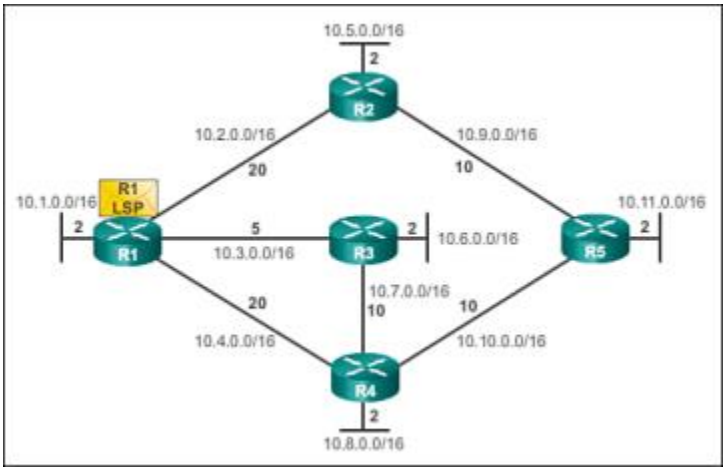


Figure 7-42 Building the LSP

Flooding the LSP

The fourth step in the link-state routing process is that each router floods the LSP to all neighbors, who then store all LSPs received in a database.

Each router floods its link-state information to all other link-state routers in the routing area as shown in Figure 7-43.

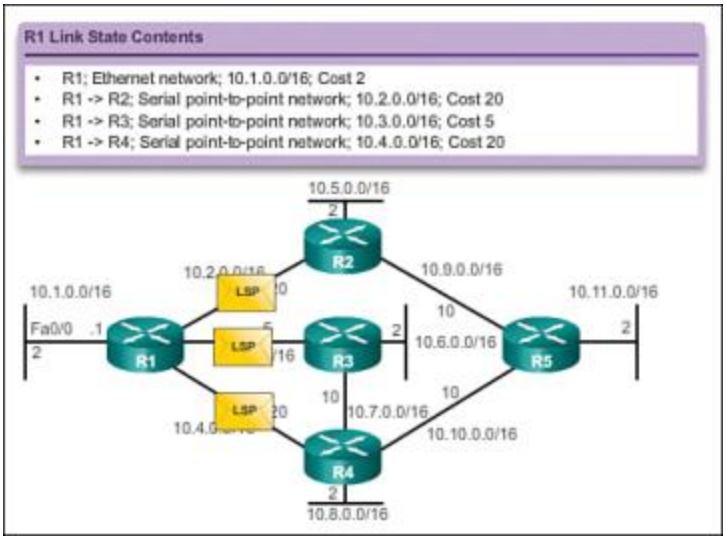


Figure 7-43 R1 Floods Its LSP

Whenever a router receives an LSP from a neighboring router, it immediately sends that LSP out all other interfaces except the interface that received the LSP. This process creates a flooding effect of LSPs from all routers throughout the routing area.

Building the Link-State Database

The final step in the link-state routing process is that each router uses the database to construct a complete map of the topology and computes the best path to each destination network.

Eventually, all routers receive an LSP from every other link-state router in the routing area. These LSPs are stored in the link-state database.

Table 7-9 displays the link-state database content of R1.

Table 7-9 Link-State Database

R1 Link-states:
Connected to network 10.1.0.0/16, cost = 2
Connected to R2 on network 10.2.0.0/16, cost = 20
Connected to R3 on network 10.2.0.0/16, cost = 5
Connected to R4 on network 10.3.0.0/16, cost = 20
R2 Link-states:
Connected to network 10.5.0.0/16, cost = 2
Connected to R1 on network 10.2.0.0/16, cost = 20
Connected to R5 on network 10.9.0.0/16, cost = 10
R3 Link-states:
Connected to network 10.6.0.0/16, cost = 2
Connected to R1 on network 10.3.0.0/16, cost = 5
Connected to R4 on network 10.7.0.0/16, cost = 10
R4 Link-states:
Connected to network 10.8.0.0/16, cost = 2
Connected to R1 on network 10.4.0.0/16, cost = 20
Connected to R3 on network 10.7.0.0/16, cost = 10
Connected to R5 on network 10.10.0.0/16, cost = 10

R1 Link-states:

Connected to network 10.1.0.0/16, cost = 2

Connected to R2 on network 10.2.0.0/16, cost = 20

Connected to R3 on network 10.2.0.0/16, cost = 5

Connected to R4 on network 10.3.0.0/16, cost = 20

R5 Link-states:

Connected to network 10.11.0.0/16, cost = 2

Connected to R2 on network 10.9.0.0/16, cost = 10

Connected to R4 on network 10.10.0.0/16, cost = 10

As a result of the flooding process, R1 has learned the link-state information for each router in its routing area. Notice that R1 also includes its own link-state information in the link-state database.

With a complete link-state database, R1 can now use the database and the shortest path first (SPF) algorithm to calculate the preferred path or shortest path to each network, resulting in the SPF tree.

Building the SPF Tree

Each router in the routing area uses the link-state database and SPF algorithm to construct the ***SPF tree***.

For example, using the link-state information from all other routers, R1 can now begin to construct an SPF tree of the network. To begin, the SPF algorithm interprets each router's LSP to identify networks and associated costs.

The SPF algorithm then calculates the shortest paths to reach each individual network, resulting in the SPF tree as shown in Figure 7-44. R1 now has a complete topology view of the link-state area.

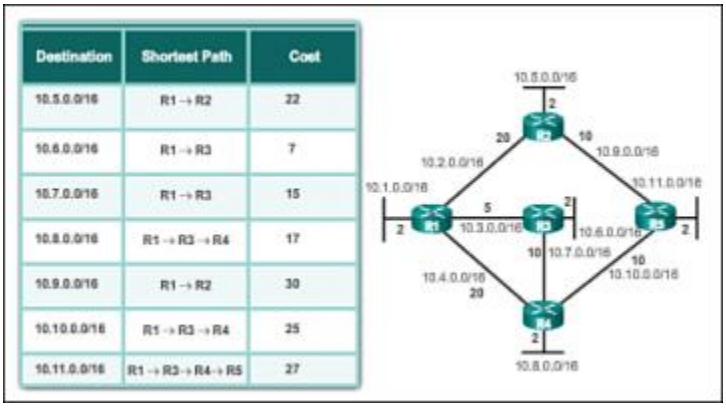


Figure 7-44 Resulting SPF Tree of R1

The entire process can be viewed in the online course on page 3.4.2.7 in Figures 1 through 6.

Each router constructs its own SPF tree independently from all other routers. To ensure proper routing, the link-state databases used to construct those trees must be identical on all routers.

Adding OSPF Routes to the Routing Table

Using the shortest path information determined by the SPF algorithm, these paths can now be added to the routing table. Figure 7-45 shows the routes that have now been added to R1's IPv4 routing table.

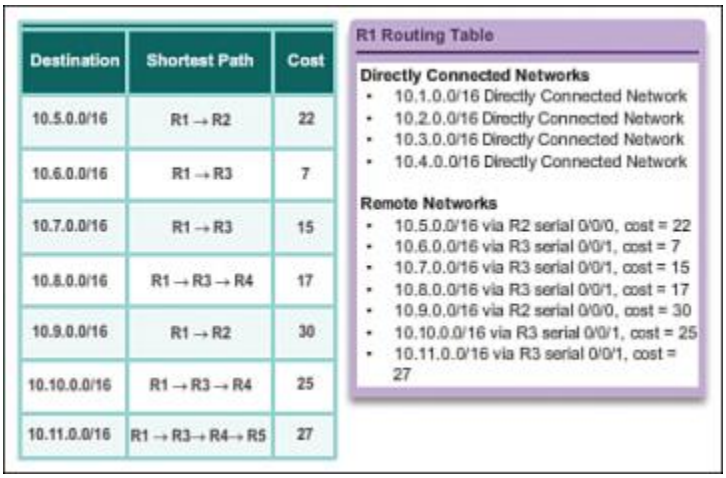


Figure 7-45 Populate the Routing Table

The routing table also includes all directly connected networks and routes from any other sources, such as static routes. Packets are now forwarded according to these entries in the routing table.

Why Use Link-State Routing Protocols?

This section discusses the advantages of using link-state routing protocols and compares the two types of link-state routing protocols.

Why Use Link-State Protocols?

There are several advantages of link-state routing protocols compared to distance vector routing protocols.

- **Builds a topological map**: Link-state routing protocols create a topological map, or SPF tree of the network topology. Because link-state routing protocols exchange link-states, the SPF algorithm can build an SPF tree of the network. Using the SPF tree, each router can independently determine the shortest path to every network.
- **Fast convergence**: When receiving an LSP, link-state routing protocols immediately flood the LSP out all interfaces except for the interface from which the LSP was received. In contrast, RIP needs to process each routing update and update its routing table before flooding the routing update out other interfaces.
- ***Event-driven updates***: After the initial flooding of LSPs, link-state routing protocols only send out an LSP when there is a change in the topology. The LSP contains only the information regarding the affected link. Unlike some distance vector routing protocols, link-state routing protocols do not send periodic updates.
- **Hierarchical design**: Link-state routing protocols use the concept of areas. Multiple areas create a hierarchical design to networks, allowing for better route aggregation (summarization) and the isolation of routing issues within an area.

Link-state protocols also have a few disadvantages compared to distance vector routing protocols:

- **Memory requirements**: Link-state protocols require additional memory to create and maintain the link-state database and SPF tree.
- **Processing requirements**: Link-state protocols can also require more CPU processing than distance vector routing protocols. The SPF algorithm requires more CPU time than distance vector algorithms such as Bellman-Ford, because link-state protocols build a complete map of the topology.
- **Bandwidth requirements**: The flooding of link-state packets can adversely affect the available bandwidth on a network. This should only occur during initial startup of routers, but can also be an issue on unstable networks.

Link-State Protocols Support Multiple Areas

Modern link-state routing protocols are designed to minimize the effects on memory, CPU, and bandwidth. The use and configuration of multiple areas can reduce the size of the link-state databases. Multiple areas can also limit the amount of link-state information flooding in a routing domain and send LSPs only to those routers that need them. When there is a change in the topology, only those routers in the affected area receive the LSP and run the SPF algorithm. This can help isolate an unstable link to a specific area in the routing domain.

For example, in Figure 7-46, there are three separate routing domains: area 1, area 0, and area 51.

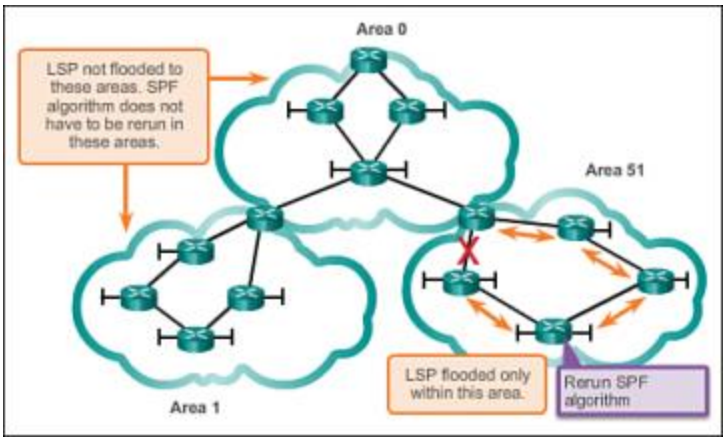


Figure 7-46 Create Areas to Minimize Router Resource Usage

If a network in area 51 goes down, the LSP with the information about this downed link is only flooded to other routers in that area. Only those routers in area 51 need to update their link-state databases, rerun the SPF algorithm, create a new SPF tree, and update their routing tables. Routers in other areas learn that this route is down, but this is done with a type of LSP that does not cause them to rerun their SPF algorithm. Routers in other areas can update their routing tables directly.

Protocols that Use Link-State

There are only two link-state routing protocols, OSPF and IS-IS.

Open Shortest Path First (OSPF) is the most popular implementation. It was designed by the Internet Engineering Task Force (IETF) OSPF Working Group. The development of OSPF began in 1987 and there are two current versions in use:

- **OSPFv2**: OSPF for IPv4 networks (RFC 1247 and RFC 2328)
- ***OSPFv3***: OSPF for IPv6 networks (RFC 2740)

With the OSPFv3 Address Families feature, OSPFv3 includes support for both IPv4 and IPv6.

IS-IS was designed by International Organization for Standardization (ISO) and is described in ISO 10589. The first incarnation of this routing protocol was developed at Digital Equipment Corporation (DEC) and is known as DECnet Phase V. Radia Perlman was the chief designer of the IS-IS routing protocol.

IS-IS was originally designed for the OSI protocol suite and not the TCP/IP protocol suite. Later, Integrated IS-IS, or Dual IS-IS, included support for IP networks. Although IS-IS has been known as the routing protocol used mainly by ISPs and carriers, more enterprise networks are beginning to use IS-IS.

OSPF and IS-IS share many similarities and also have many differences. There are many pro-OSPF and pro-IS-IS factions who discuss and debate the advantages of one routing protocol over the other. Both routing protocols provide the necessary routing functionality.

Things to Remember about Link State Routing

- Link state protocols advertise a large amount of topological information about the network (tells what every metric is for every link in the network)
- Routers must calculate the metric (using Shortest path First Algorithm)
- Routers perform CPU intensive computations on the data.
- Discover neighbors before exchanging information.

Process of Learning Routes:
1. Each router discovers its neighbors on each interface, list kept in **neighbors table**.

2. Each router uses a reliable protocol to exchange topology information in its **topology database**.

3. Each router places the learned topology information in its topology database.

4. Each router then runs the SPF algorithm against its own topology database to calculate the best routes to each subnet in the database.

5. Each router finally places the best route to each subnet in the IP routing table.

OSPF Topology Database: Consists of lists of subnet numbers (links), list of routers (and links they are connected to).

-> Uniquely identifier each router in this database using OSPF Router ID (RID)
To select the RID

- The router **first checks for any loopback** interfaces that are up, and chooses the highest numeric IP address of those.
- **If no loopback exists**, router chooses highest IP address from interfaces that are up and up.

*Note: loopback interface is a virtual interface, configured with

interface loopback [interface #]

Each Router chooses RID when OSPY is initialized (during initial loading of IOS). If other interfaces come up after this, not used unless **clear ip ospf process** is issued.

Meeting OSPF Neighbors: Once router has assigned itself a RID, and some of its interfaces are up, the router is ready to meet its neighbors (connected routers).

- Can become neighbors if connected to same subnet
- Router multicasts OSPF Hello packets out each interface
- Hello message follows IP packet header (**port = 89**)
- Hello packets sent to **224.0.0.5** (all OSPF speaking routers)

- **Routers learn several things from Hello Packets:**
- RID, Area ID, Hello Interval, Dead Interval, router priority, designated router, backup designated router, and a list of neighbors sending router already knew about.

- **To confirm that a Hello Packet was received**, next Hello Message will include the sender's RID within the list of neighbors.

-> **Once router sees** its RID included, two-way state achieved, and more detailed information can be exchanged.

The following must match before routers become neighbors:

1. Subnet mask
2. Hello Interval
3. OSPF Area ID
4. Dead Interval
5. Subnet number (derived using the mask applied to the IP)

Reducing Overhead Using Designated Routers

Sometimes **Designated Routers** (DR) are required before sending **Database Description** (DD) packets.

- DR's always required on a LAN
- Sometimes required with Frame Relay/ATM (depending on topology/config)

After DR is elected, all updates flow through the Designated Router (DR). This means that the DR collects and distributes the routing updates to alleviate OSPF update congestion.

Router decides if it needs to elect a DR depending on the *network type*.

* **Point-to-point** DOES NOT need a DR

* **Broadcast** (for LANs), always needs a DR

- **Non-broadcast Multiaccess** (NBMA), for frame relay, sometimes needs DR, sometimes doesn't. Has 5 different variations, configured with **ip ospf network** [type] command

**** Since DR's are so important**, loss of one could cause delay in convergence, so Backup DR (BDR) is also needed. ****

Electing The Designated Router

To elect, neighboring routers hold an election, and look at two fields in the Hello Packet:

* Router that sends the highest OSPF priority becomes DR

* If there is a tie, the highest RID wins.

* **To elect BDR**, typically the second highest priority is used. *

Other Notes:

* Priority setting of 0 means router will never be DR

* Range of valid priority values is 1-255 (to become a DR)

* If DR is elected, then another router comes online with a higher priority, this router will not become DR until both the DR and BDR fail.

Once DR/BDR is elected:

1. Non-DR send updates to 224.0.0.6 (All OSPF DRs)

2. DR relays these messages to 224.0.0.5 (BDR does not forward, only receives)

3. Once router has exchanged its entire link state database, transition to *Full State*

Steady-State Operation: If Hello Interval is not received for [dead interval] amount of time, the router believes the neighbor has failed.

- Default dead timer is 4 times the hello interval

(10 second hello, 40 second dead timer)

- Router marks as "down" in its neighbor table
- Runs the dijkstra algorithm to calculate new routes, floods to inform other routers of failed link

Loop Avoidance: Link state does not use SPF algorithm, but rather it relies on router broadcasting downed link immediately. This is the main reason for fast convergence time (distance vector uses hold time, split

horizon, etc, while link state does not).

Scaling OSPF: If network has many routers (~50 or more, a few hundred subnets), would result in:
* Slow convergence time
* Memory shortages/processor overloading

Scalability Solutions Include:
* **OSPF Areas:** Break up the network so that routers in one area know less topology information about the subnets in the other area, and don't know about other routers at all.
* **Border Router:** OSPF Area Border Router (ABR), border between 2 different areas (sits in both areas).
* **Makes other routers in same area** view network as if it had fewer routers.
* **Area 0 defined as backbone**, OSPF designs hierarchical

**** Note:** doesn't change of subnets know, just decrease of bytes/require memory to process updates **

Summary of Distance Vector and Link State

Feature	Link State	Distance Vector
Convergence Time	Fast	Slow (loop avoidance features)
Loop avoidance	Built into protocol	Extra features such as route poisoning, split horizon
Memory/CPU	Can be large; good design can minimize	Low
Requires design effort for large networks	Yes	No
Public/Proprietary?	OSPF = public	RIP = public IGRP = Cisco proprietary

Balanced Hybrid Routing Protocol/EIGRP Concepts

EIGRP has some features that act like distance vector protocols, and some that act like link-state protocols.

Feature Comparison with IGRP:

Similarities	Differences
Both Cisco proprietary	EIGRP converges faster
Same logic for equal-cost paths	EIGRP sends routing info once to neighbor, then again only when update occurs.

	IGRP sends every 90 seconds.
Metric's identical (EIGRP just scales by multiplying by 256)	EIGRP can exchange for Novel IPX and AppleTalk, as well as IP

EIGRP Processes and Tables: Follows three general steps to be able to add routes to routing table:

1. **EIGRP neighbor table**: Routers discover other EIGRP routers that are attached to same subnet, form a neighbor relationship and keep a list in this table. a. **show ip eigrp neighbor**
2. **EIGRP topology table**: Exchange of network topology information with known neighbors. a. **show ip eigrp topology**
3. **IP routing table**: EIGRP analyzes topology information, puts lowest metric routes in this table. a. **show ip route** -or- **show ip route eigrp**

** EIGRP could have up to 9 tables, since it supports IP, IPX, and AppleTalk **

Hello Messages: Used to perform neighbor discovery, continually sent to notice when connectivity has failed.

Interval determines how frequently it is sent

- LANs/Point-to-point connections = **5 seconds**
- Multipoint WANS like Frame Relay = **60 seconds**

Update Messages: Conveys topology information to neighbors.

- Sent out multicast address 224.0.0.10 if updating multiple routers
- Sent out Unicast address if single router updated
- Reliable messages sent out **Reliable Transport Protocol** (RTP)

Updating the Routing Table while Avoiding Loops

EIGRP keeps basic topological information (but not full information)

- Routes with ***feasible successor*** can be used immediately after route fails
- Routes without on require EIGRP to perform ***Query and Response*** process to confirm that no loop exists.

Successors are in topology table, and are the best route (the route with lowest metric, which is also in routing table).

Feasible Successors are in topology table, and are placed when the neighbor has a lower metric for its route.

Diffusing Update Algorithm (DUAL) is used in query and reply process, when both successor and feasible successor fail. Sends query to confirm route exists, reply verifies route.

EIGRP Compared

Feature	EIGRP	IGRP	OSPF
Discovers neighbors before exchanging routing information	Y	N	Y
Builds topology table in addition to routing table	Y	N	Y
Converges Quickly	Y	N	Y
Bandwidth/delay metric	Y	Y	N
Sends full routing table during update	N	Y	N
Requires distance vector loop avoidance features	N	Y	N
Public Standard	N	N	Y
Uses DUAL Algorithm	Y	N	N

IP Configuration Commands

Command	Configuration Mode
router ospf *process-id*	Global
network [ip address][wildcard mask] **area** [area id]	Router subcommand
ip ospf cost *interface cost*	Sets cost associated with interface
bandwidth [bandwidth]	Sets interface bandwidth
auto-cost reference bandwidth [number]	Router subcommand that sets the numerator in formula to calculate cost.

ip ospf hello [number]	Interface subcommand that sets Hello interval, and sets dead interval to 4 times this number.
ip ospf network [type]	Interface subcommand that defines the OSPF network type.

IP OSPF Exec Commands

Command	Description
show ip route [ip address]	Shows entire routing table, or subset if parameters entered.
show ip protocols	Shows routing protocol parameters and current timer values.
show ip ospf interface	List the area in which the router resides, and adjacent neighbors.
show ip ospf neighbor	Lists neighbors and current status with neighbors, per interface.
show ip route ospf	Lists routes in routing table learned by ospf.
debug ip ospf events	Issues log messages for each OSPF packet.
debug ip ospf packet	Issues log messages describing the contents of all OSPF packets.
debug ip ospf hello	Issues log messages describing Hellos and Hello failures.

OSPF Single-Area Configuration

```
interface Ethernet 0/0
ip address 10.1.1.1 255.255.255.0
interface serial 0/0
ip address 10.1.4.1 255.255.255.0

router ospf 1
network 10.0.0.0 0.255.255.255 area 0
```

Network : What interfaces you want to include in OSPF configuration . Here 10.0.0.0
Wildcard Mask: If bit set to 1, "don't care" bit (and 0 = include) . Here 0.255.255.255
Area : What area this router is in. Here area 0

OSPF Configuration with Multiple Areas

If router has interfaces in multiple areas:

router ospf 1
network 10.1.1.1 0.0.0.0 area 0
network 10.1.4.1 0.0.0.0 area 1
network 10.1.6.1 0.0.0.0 area 0

Useful Commands
show ip ospf interface ->Details IP address, area , Router ID, Hello/Dead Interval, etc. for all interfaces

show ip route -> Shows all routes known by the router (C – Connected, O – OSPF)

show ip ospf neighbor -> Shows the routers ospf neighbors

Remember that the RID is that router's highest IP address on a physical interface when OSPF starts running. Alternatively, if a loopback interface has been configured, OSPF uses the highest IP address on a loopback interface for the RID, even if that IP address is lower than some physical interface's IP address.

OSPF Troubleshooting

Mismatched Hello Intervals:

* View neighbors:
show ip ospf neighbor ->Output doesn't show neighbors

* Run debugging:
debug ip ospf hello -> Output shows mismatched Hello interval

* To identify the interface:
show ip ospf interface [interface] -> Will give you the hello interval

* To change hello interval for that interface:
configure terminal
interface [interface]
ip ospf hello [count]
exit

EIGRP Configuration

Configured exactly like IGRP, just switch "igrp" with "eigrp" in commands.

IP EIGRP Exec Commands

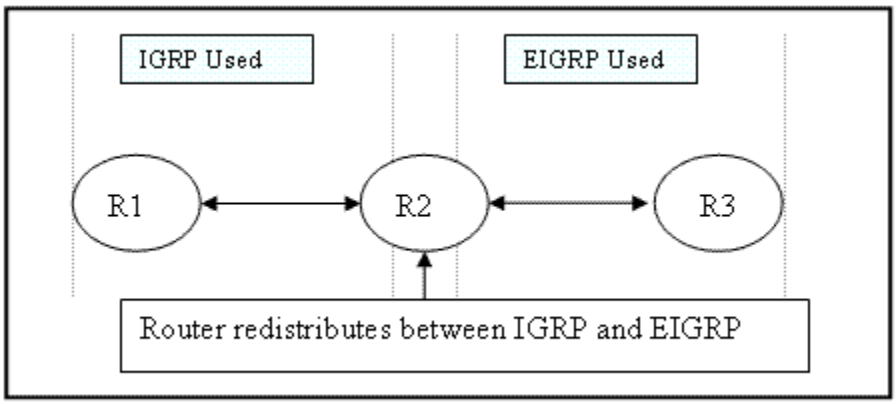
Command	Description
show ip route [ip address]	Shows entire routing table.
show ip eigrp neighbors	Lists EIGRP neighbors and status.
show ip eigrp topology	Lists RIGRP topology table, including feasible successors/successors.
show ip route eigrp	Lists only EIGRP-learned routes
show ip eigrp traffic	Lists traffic statistics about EIGRP

Other Key Points

- Letter "D" signifies EIGRP-learned routes
- All routers must be in same AS number (**network x.x.x.x** [AS number])

IGRP to EIGRP Migration

Feature of EIGRP called Automatic Redistribution



- Border router must be configured for both IGRP and EIGRP
- Both must use same AS number

10. The Routing Table

It is important to know the routing table in depth when troubleshooting network issues. Understanding the structure and lookup process of the routing table will help you diagnose any routing table issue, regardless of your level of familiarity with a particular routing protocol. For example, you might encounter a situation in which the routing table has all of the routes you would expect to see, but packet forwarding is not performing as expected. Knowing how to step through the lookup process of a destination IP address for a packet will enable you to determine whether the packet is being forwarded as expected, if and why the packet is being sent elsewhere, or whether the packet has been discarded.

Parts of an IPv4 Route Entry

A routing table consists of directly connected networks and routes learned statically or dynamically. This section examines these two types of routing table entries.

Routing Table Entries

The topology displayed in Figure 7-47 is used as the reference topology for this section.

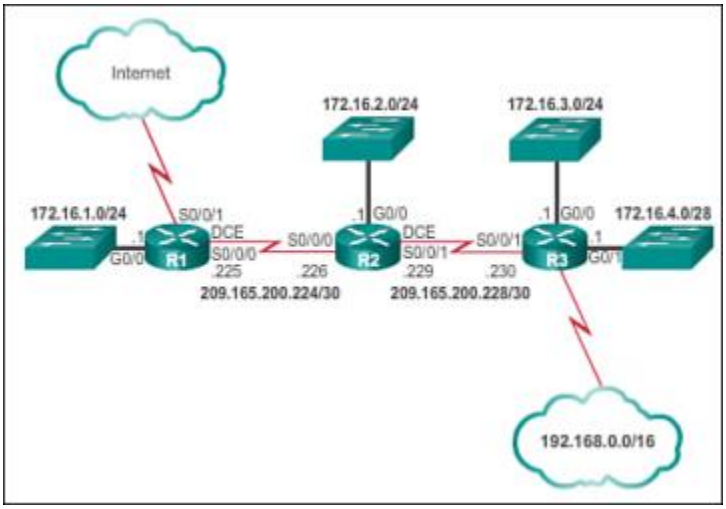


Figure 7-47 Reference Topology

Notice that in the topology:

- R1 is the edge router that connects to the Internet. Therefore, it is propagating a default static route to R2 and R3.
- R1, R2, and R3 contain discontiguous networks separated by another classful network.
- R3 is also introducing a 192.168.0.0/16 supernet route.

Figure 7-48 displays the IPv4 routing table of R1 with directly connected, static, and dynamic routes.

```
R1# show ip route | begin Gateway
Gateway of last resort is 209.165.200.234 to network 0.0.0.0

S* 0.0.0.0/0 [1/0] via 209.165.200.234, Serial0/0/1
                   is directly connected, Serial0/0/1
   172.16.0.0/16 is variably subnetted, 5 subnets, 3 masks
C     172.16.1.0/24 is directly connected, GigabitEthernet0/0
L     172.16.1.1/32 is directly connected, GigabitEthernet0/0
R     172.16.2.0/24 [120/1] via 209.165.200.226, 00:00:12, Serial0/0/0
R     172.16.3.0/24 [120/2] via 209.165.200.226, 00:00:12, Serial0/0/0
R     172.16.4.0/28 [120/2] via 209.165.200.226, 00:00:12, Serial0/0/0
R  192.168.0.0/16 [120/2] via 209.165.200.226, 00:00:03, Serial0/0/0
   209.165.200.0/24 is variably subnetted, 5 subnets, 2 masks
C     209.165.200.224/30 is directly connected, Serial0/0/0
L     209.165.200.225/32 is directly connected, Serial0/0/0
R     209.165.200.228/30 [120/1] via 209.165.200.226, 00:00:12,
                    Serial0/0/0
C     209.165.200.232/30 is directly connected, Serial0/0/1
L     209.165.200.233/30 is directly connected, Serial0/0/1
R1#
```

Figure 7-48 Routing Table of R1

The routing table hierarchy in Cisco IOS was originally implemented with the classful routing scheme. Although the routing table incorporates both classful and classless addressing, the overall structure is still built around this classful scheme.

Directly Connected Entries

As highlighted in Figure 7-49, the routing table of R1 contains three directly connected networks. Notice that two routing table entries are automatically created when an active router interface is configured with an IP address and subnet mask.

```
R1# show ip route | begin Gateway
Gateway of last resort is 209.165.200.234 to network 0.0.0.0

S* 0.0.0.0/0 [1/0] via 209.165.200.234, Serial0/0/1
                   is directly connected, Serial0/0/1
   172.16.0.0/16 is variably subnetted, 5 subnets, 3 masks
C     172.16.1.0/24 is directly connected, GigabitEthernet0/0
L     172.16.1.1/32 is directly connected, GigabitEthernet0/0
R     172.16.2.0/24 [120/1] via 209.165.200.226,00:00:12, Serial0/0/0
R     172.16.3.0/24 [120/2] via 209.165.200.226, 00:00:12, Serial0/0/0
R     172.16.4.0/28 [120/2] via 209.165.200.226, 00:00:12, Serial0/0/0
R     192.168.0.0/16 [120/2] via 209.165.200.226, 00:00:03, Serial0/0/0
   209.165.200.0/24 is variably subnetted, 5 subnets, 2 masks
C     209.165.200.224/30 is directly connected, Serial0/0/0
L     209.165.200.225/32 is directly connected, Serial0/0/0
R     209.165.200.228/30 [120/1] via 209.165.200.226, 00:00:12, Serial0/0/0
C     209.165.200.232/30 is directly connected, Serial0/0/1
L     209.165.200.233/32 is directly connected, Serial0/0/1
R1#
```

Figure 7-49 Directly Connected Interfaces of R1

Figure 7-50 displays one of the routing table entries on R1 for the directly connected network 172.16.1.0. These entries were automatically added to the routing table when the GigabitEthernet 0/0 interface was configured and activated.

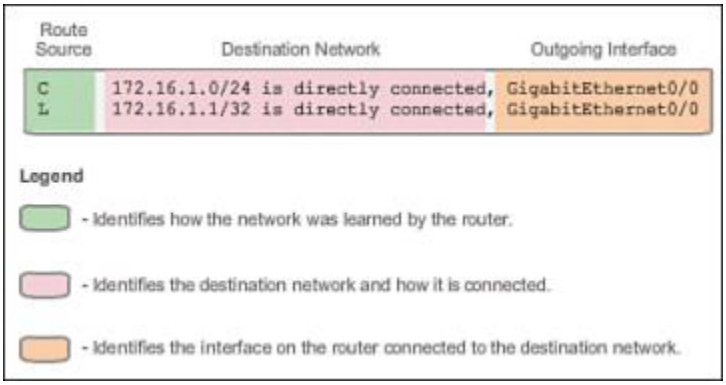


Figure 7-50 Directly Connected Routes of R1

The entries contain the following information:

- **Route source**: Identifies how the route was learned. Directly connected interfaces have two route source codes. **C** identifies a directly connected network. Directly connected networks are automatically created whenever an interface is configured with an IP address and activated. **L** identifies that this is a local route. Local routes are automatically created whenever an interface is configured with an IP address and activated.
- **Destination network**: The address of the remote network and how that network is connected.
- **Outgoing interface**: Identifies the exit interface to use when forwarding packets to the destination network.

Local routing table entries did not appear in routing tables prior to IOS release 15.

A router typically has multiple interfaces configured. The routing table stores information about both directly connected and remote routes. As with directly connected networks, the route source identifies how the route was learned. For instance, common codes for remote networks include:

- **S**: Identifies that the route was manually created by an administrator to reach a specific network. This is known as a static route.
- **D**: Identifies that the route was learned dynamically from another router using the EIGRP routing protocol.
- **O**: Identifies that the route was learned dynamically from another router using the OSPF routing protocol.

- **R**: Identifies that the route was learned dynamically from another router using the RIP routing protocol.

Remote Network Entries

Figure 7-51 displays an IPv4 routing table entry on R1 for the route to remote network 172.16.4.0 on R3.

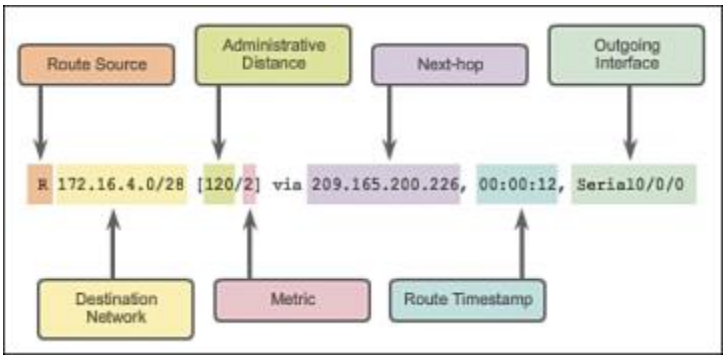


Figure 7-51 Remote Network Route Entry on R1

The entry identifies the following information:

- **Route source**: Identifies how the route was learned.
- **Destination network**: Identifies the address of the remote network.
- **Administrative distance**: Identifies the trustworthiness of the route source.
- **Metric**: Identifies the value assigned to reach the remote network. Lower values indicate preferred routes.
- **Next hop**: Identifies the IPv4 address of the next router to forward the packet to.
- **Route timestamp**: Identifies from when the route was last heard.
- **Outgoing interface**: Identifies the exit interface to use to forward a packet toward the final destination.

Dynamically Learned IPv4 Routes

The structure or format of the routing table might seem obvious until you take a closer look. Understanding the structure of the routing table will help you verify and troubleshoot routing issues because you will understand the routing table lookup process.

Routing Table Terms

A dynamically built routing table provides a great deal of information, as shown in Figure 7-52. Therefore, it is crucial to understand the output generated by the routing table. Special terms are applied when discussing the contents of a routing table.

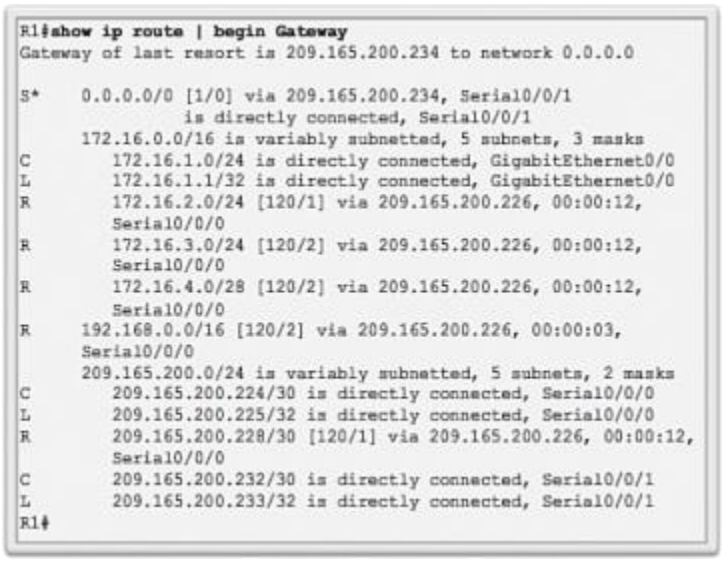


Figure 7-52 Routing Table of R1

The Cisco IP routing table is not a flat database. The routing table is actually a hierarchical structure that is used to speed up the lookup process when locating routes and forwarding packets. Within this structure, the hierarchy includes several levels.

Routes are discussed in terms of:

- Ultimate route
- Level 1 route
- Level 1 parent route
- Level 2 child routes

Ultimate Route

An ***ultimate route*** is a routing table entry that contains either a next-hop IPv4 address or an exit interface. Directly connected, dynamically learned, and local routes are ultimate routes.

In Figure 7-53, the highlighted areas are examples of ultimate routes. Notice that all of these routes specify either a next-hop IPv4 address or an exit interface.

```
R1# show ip route | begin Gateway
Gateway of last resort is 209.165.200.234 to network 0.0.0.0

S*    0.0.0.0/0 [1/0] via 209.165.200.234, Serial0/0/1
                is directly connected, Serial0/0/1
      172.16.0.0/16 is variably subnetted, 5 subnets, 3 masks
C        172.16.1.0/24 is directly connected, GigabitEthernet0/0
L        172.16.1.1/32 is directly connected, GigabitEthernet0/0
R        172.16.2.0/24 [120/1] via 209.165.200.226, 00:00:12,
         Serial0/0/0
R        172.16.3.0/24 [120/2] via 209.165.200.226, 00:00:12,
          Serial0/0/0
R        172.16.4.0/28 [120/2] via 209.165.200.226, 00:00:12,
         Serial0/0/0
R     192.168.0.0/16 [120/2] via 209.165.200.226, 00:00:03,
      Serial0/0/0
      209.165.200.0/24 is variably subnetted, 5 subnets, 2 masks
C        209.165.200.224/30 is directly connected, Serial0/0/0
L        209.165.200.225/32 is directly connected, Serial0/0/0
R        209.165.200.228/30 [120/1] via 209.165.200.226, 00:00:12,
         Serial0/0/0
C        209.165.200.232/30 is directly connected, Serial0/0/1
L        209.165.200.233/32 is directly connected, Serial0/0/1
R1#
```

Figure 7-53 Ultimate Routes of R1

Level 1 Route

A ***level 1 route*** is a route with a subnet mask equal to or less than the classful mask of the network address. Therefore, a level 1 route can be a:

- **Network route**: A network route has a subnet mask equal to that of the classful mask.
- ***Supernet route***: A supernet route is a network address with a subnet mask less than the classful mask, for example, a summary address.
- **Default route**: A default route is a static route with the address 0.0.0.0/0.

The source of the level 1 route can be a directly connected network, static route, or a dynamic routing protocol.

Figure 7-54 highlights how level 1 routes are also ultimate routes.

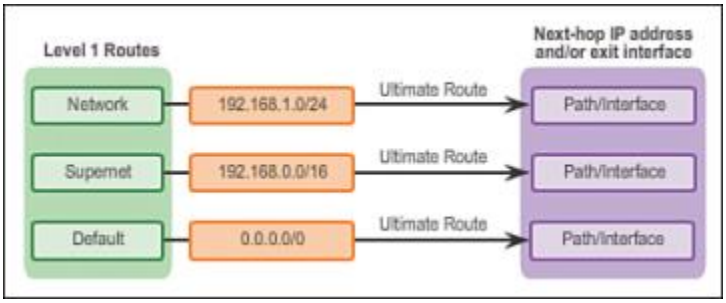


Figure 7-54Sources of Level 1 Routes

Figure 7-55 highlights level 1 routes.

```
R1# show ip route | begin Gateway
Gateway of last resort is 209.165.200.234 to network
0.0.0.0

S*    0.0.0.0/0 [1/0] via 209.165.200.234, Serial0/0/1
                is directly connected, Serial0/0/1
      172.16.0.0/16 is variably subnetted, 5 subnets, 3
masks
C        172.16.1.0/24 is directly connected,
GigabitEthernet0/0
L        172.16.1.1/32 is directly connected,
GigabitEthernet0/0
R        172.16.2.0/24 [120/1] via 209.165.200.226,
00:00:12, Serial0/0/0
R        172.16.3.0/24 [120/2] via 209.165.200.226,
00:00:12, Serial0/0/0
R        172.16.4.0/28 [120/2] via 209.165.200.226,
00:00:12, Serial0/0/0
R     192.168.0.0/16 [120/2] via 209.165.200.226, 00:00:03,
Serial0/0/0
      209.165.200.0/24 is variably subnetted, 5 subnets, 2
masks
C        209.165.200.224/30 is directly connected,
Serial0/0/0
```

Figure 7-55 Example of Level 1 Routes

Level 1 Parent Route

As illustrated in Figure 7-56, a ***level 1 parent route*** is a level 1 network route that is subnetted. A parent route can never be an ultimate route.

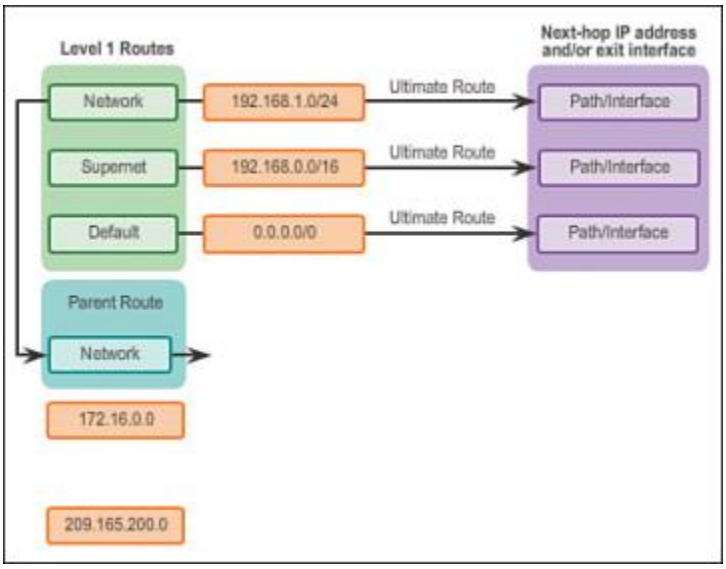


Figure 7-56 Level 1 Parent Route

Figure 7-57 highlights the level 1 parent routes in the routing table of R1. The routing table basically provides a heading for the specific subnets it contains. Each entry displays the classful network address, the number of subnets, and the number of different subnet masks that the classful address has been subdivided into.

```
R1# show ip route | begin Gateway
Gateway of last resort is 209.165.200.234 to network
0.0.0.0

S*    0.0.0.0/0 [1/0] via 209.165.200.234, Serial0/0/1
                is directly connected, Serial0/0/1
      172.16.0.0/16 is variably subnetted, 5 subnets, 3
masks
C        172.16.1.0/24 is directly connected,
GigabitEthernet0/0
L        172.16.1.1/32 is directly connected,
GigabitEthernet0/0
R        172.16.2.0/24 [120/1] via 209.165.200.226,
00:00:12, Serial0/0/0
R        172.16.3.0/24 [120/2] via 209.165.200.226,
00:00:12, Serial0/0/0
R        172.16.4.0/28 [120/2] via 209.165.200.226,
00:00:12, Serial0/0/0
R     192.168.0.0/16 [120/2] via 209.165.200.226, 00:00:03,
Serial0/0/0
      209.165.200.0/24 is variably subnetted, 5 subnets, 2
masks
C        209.165.200.224/30 is directly connected,
Serial0/0/0
```

Figure 7-57 Level 1 Parent Routes of R1

Level 2 Child Route (3.5.2.5)

A ***level 2 child route*** is a route that is a subnet of a classful network address. As illustrated in Figure 7-58, a level 1 parent route is a level 1 network route that is subnetted.

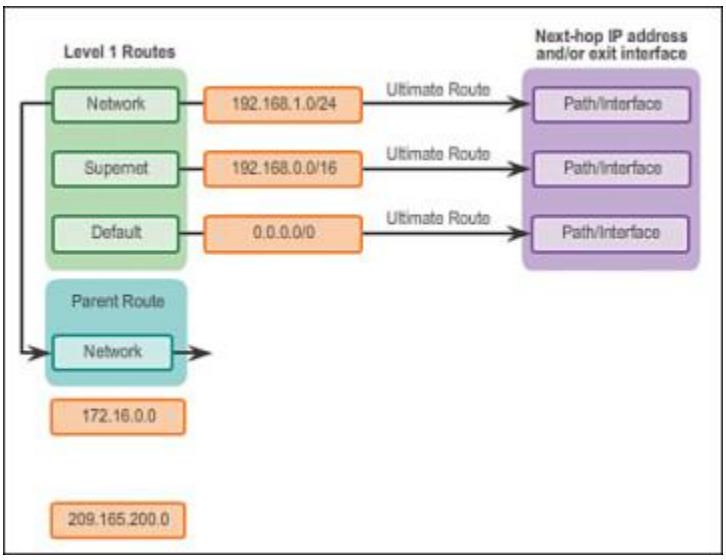


Figure 7-58 Level 2 Child Routes

A level 1 parent route contains level 2 child routes, as shown in Figure 7-59.

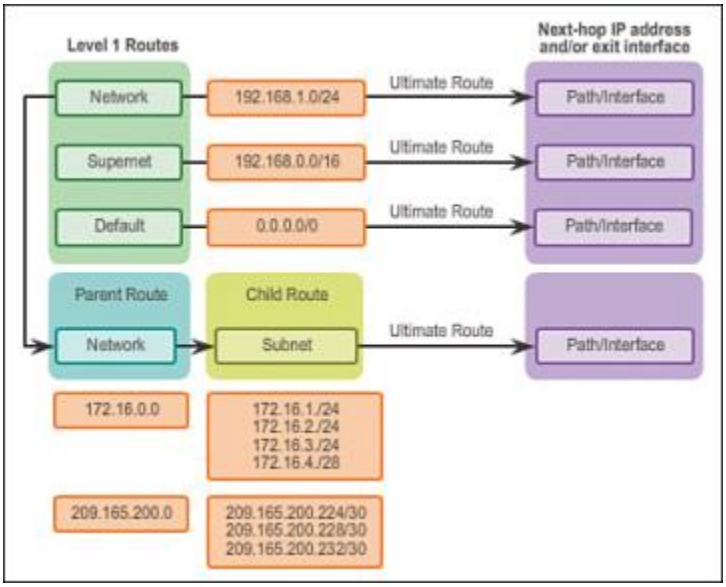


Figure 3-59 Child Routes Are Ultimate Routes

Like a level 1 route, the source of a level 2 route can be a directly connected network, a static route, or a dynamically learned route. Level 2 child routes are also ultimate routes.

The routing table hierarchy in Cisco IOS has a classful routing scheme. A level 1 parent route is the classful network address of the subnet route. This is the case even if a classless routing protocol is the source of the subnet route.

Figure 7-60 highlights the level 2 child routes in the routing table of R1.

```
R1#show ip route | begin Gateway
Gateway of last resort is 209.165.200.234 to network
0.0.0.0

S*    0.0.0.0/0 [1/0] via 209.165.200.234, Serial0/0/1
                is directly connected, Serial0/0/1
      172.16.0.0/16 is variably subnetted, 5 subnets, 3
masks
C        172.16.1.0/24 is directly connected,
GigabitEthernet0/0
L        172.16.1.1/32 is directly connected,
GigabitEthernet0/0
R        172.16.2.0/24 [120/1] via 209.165.200.226,
00:00:12, Serial0/0/0
R        172.16.3.0/24 [120/2] via 209.165.200.226,
00:00:12, Serial0/0/0
R        172.16.4.0/28 [120/2] via 209.165.200.226,
00:00:12, Serial0/0/0
R     192.168.0.0/16 [120/2] via 209.165.200.226, 00:00:03,
Serial0/0/0
      209.165.200.0/24 is variably subnetted, 5 subnets, 2
masks
C        209.165.200.224/30 is directly connected,
Serial0/0/0
```

Figure 7-60 Example of Level 2 Child Routes

The IPv4 Route Lookup Process

When a packet arrives on a router interface, the router examines the IPv4 header, identifies the destination IPv4 address, and proceeds through the router lookup process.

In Figure 7-61, the router examines level 1 network routes for the best match with the destination address of the IPv4 packet.

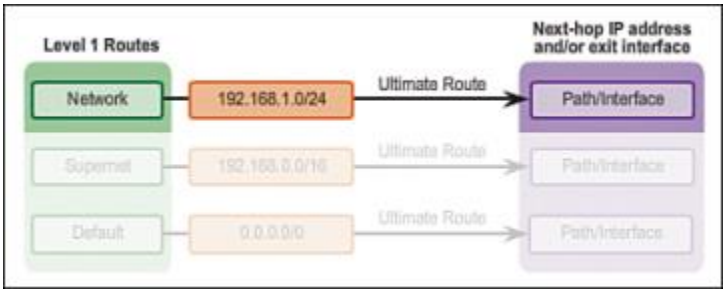


Figure 7-61 Match Level 1 Routes

Specifically, the router proceeds as follows:

1. If the best match is a level 1 ultimate route, then this route is used to forward the packet.
2. If the best match is a level 1 parent route, proceed to the next step.

In Figure 7-62, the router examines child routes (the subnet routes) of the parent route for a best match.

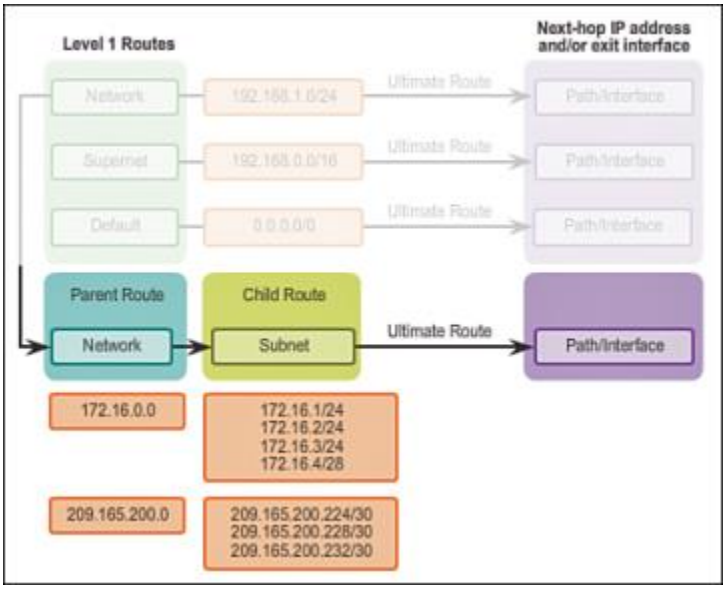


Figure 7-62 Match Level 2 Child Routes

3. If there is a match with a level 2 child route, that subnet is used to forward the packet.
4. If there is not a match with any of the level 2 child routes, proceed to the next step.

In Figure 7-63, the router continues searching level 1 supernet routes in the routing table for a match, including the default route, if there is one.

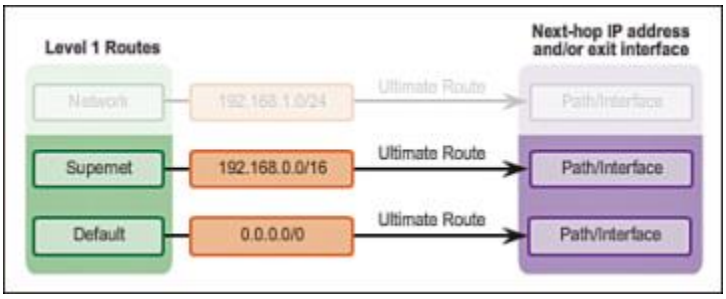


Figure 7-63 Match Supernet and Then Default Route

5. If there is now a lesser match with a level 1 supernet or default routes, the router uses that route to forward the packet.
6. If there is not a match with any route in the routing table, the router drops the packet.

A route referencing only a next-hop IP address and not an exit interface must be resolved to a route with an exit interface. A recursive lookup is performed on the next-hop IP address until the route is resolved to an exit interface.

Best Route = Longest Match

What is meant by the router must find the best match in the routing table? Best match is equal to the longest match.

For there to be a match between the destination IPv4 address of a packet and a route in the routing table, a minimum number of far left bits must match between the IPv4 address of the packet and the route in the routing table. The subnet mask of the route in the routing table is used to determine the minimum number of far left bits that must match. Remember that an IPv4 packet only contains the IPv4 address and not the subnet mask.

The best match is the route in the routing table that has the most number of far left matching bits with the destination IPv4 address of the packet. The route with the greatest number of equivalent far left bits, or the longest match, is always the preferred route.

In Figure 7-64, a packet is destined for 172.16.0.10.

The router has three possible routes that match this packet: 172.16.0.0/12, 172.16.0.0/18, and 172.16.0.0/26. Of the three routes, 172.16.0.0/26 has the longest match and is therefore chosen to forward the packet. Remember, for any of these routes to be considered a match there must be at least the number of matching bits indicated by the subnet mask of the route.

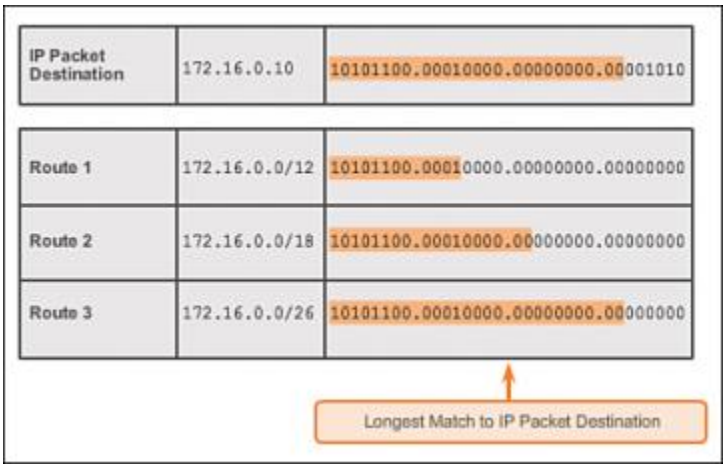


Figure 7-64 Matches for Packets Destined to 172.16.0.10

Analyze an IPv6 Routing Table

The IPv6 routing table shares many similarities with the IPv4 routing table. It also consists of directly connected networks and routes learned statically or dynamically. However, the entries are displayed somewhat differently than IPv4 entries. This section examines the IPv6 routing table.

IPv6 Routing Table Entries

Components of the IPv6 routing table are very similar to the IPv4 routing table. For instance, it is populated using directly connected interfaces, static routes, and dynamically learned routes.

Because IPv6 is classless by design, all routes are effectively level 1 ultimate routes. There is no level 1 parent of level 2 child routes.

The topology displayed in Figure 7-65 is used as the reference topology for this section.

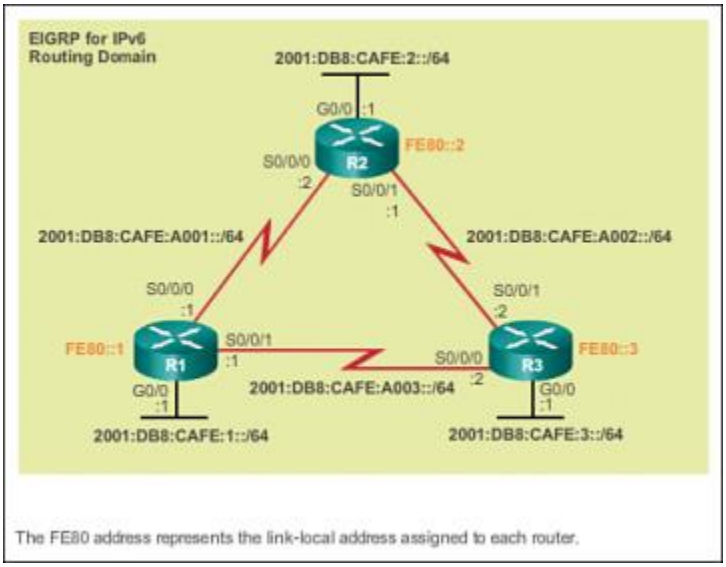


Figure 7-65 Reference IPv6 Topology

Notice that in the topology:

- R1, R2, and R3 are configured in a full mesh topology. All routers have redundant paths to various networks.
- R2 is the edge router and connects to the ISP; however, a default static route is not being advertised.
- EIGRP for IPv6 has been configured on all three routers.

Directly Connected Entries

The routing table of R1 is displayed in Figure 7-66 using the **show ipv6 route** command. Although the command output is displayed slightly differently than in the IPv4 version, it still contains the relevant route information.

```
R1# show ipv6 route
<output omitted>

C   2001:DB8:CAFE:1::/64 [0/0]
     via GigabitEthernet0/0, directly connected
L   2001:DB8:CAFE:1::1/128 [0/0]
     via GigabitEthernet0/0, receive
D   2001:DB8:CAFE:2::/64 [90/3524096]
     via FE80::3, Serial0/0/1
D   2001:DB8:CAFE:3::/64 [90/2170112]
     via FE80::3, Serial0/0/1
C   2001:DB8:CAFE:A001::/64 [0/0]
     via Serial0/0/0, directly connected
L   2001:DB8:CAFE:A001::1/128 [0/0]
     via Serial0/0/0, receive
D   2001:DB8:CAFE:A002::/64 [90/3523840]
     via FE80::3, Serial0/0/1
C   2001:DB8:CAFE:A003::/64 [0/0]
     via Serial0/0/1, directly connected
L   2001:DB8:CAFE:A003::1/128 [0/0]
     via Serial0/0/1, receive
L   FF00::/8 [0/0]
     via Null0, receive
R1#
```

Figure 7-66 IPv6 Routing Table of R1

Figure 7-67 highlights the connected network and local routing table entries of the directly connected interfaces. The three entries were added when the interfaces were configured and activated.

```
R1# show ipv6 route
<output omitted>

C   2001:DB8:CAFE:1::/64 [0/0]
     via GigabitEthernet0/0, directly connected
L   2001:DB8:CAFE:1::1/128 [0/0]
     via GigabitEthernet0/0, receive
D   2001:DB8:CAFE:2::/64 [90/3524096]
     via FE80::3, Serial0/0/1
D   2001:DB8:CAFE:3::/64 [90/2170112]
     via FE80::3, Serial0/0/1
C   2001:DB8:CAFE:A001::/64 [0/0]
     via Serial0/0/0, directly connected
L   2001:DB8:CAFE:A001::1/128 [0/0]
     via Serial0/0/0, receive
D   2001:DB8:CAFE:A002::/64 [90/3523840]
     via FE80::3, Serial0/0/1
C   2001:DB8:CAFE:A003::/64 [0/0]
     via Serial0/0/1, directly connected
L   2001:DB8:CAFE:A003::1/128 [0/0]
     via Serial0/0/1, receive
L   FF00::/8 [0/0]
     via Null0, receive
R1#
```

Figure 7-67 Directly Connected Routes on R1

As shown in Figure 7-68, directly connected route entries display the following information:

- **Route source**: Identifies how the route was learned. Directly connected interfaces have two route source codes (C identifies a directly connected network while L identifies that this is a local route).
- **Directly connected network**: The IPv6 address of the directly connected network.
- **Administrative distance**: Identifies the trustworthiness of the route source. IPv6 uses the same distances as IPv4. A value of 0 indicates the best, most trustworthy source.
- **Metric**: Identifies the value assigned to reach the remote network. Lower values indicate preferred routes.
- **Outgoing interface**: Identifies the exit interface to use when forwarding packets to the destination network.

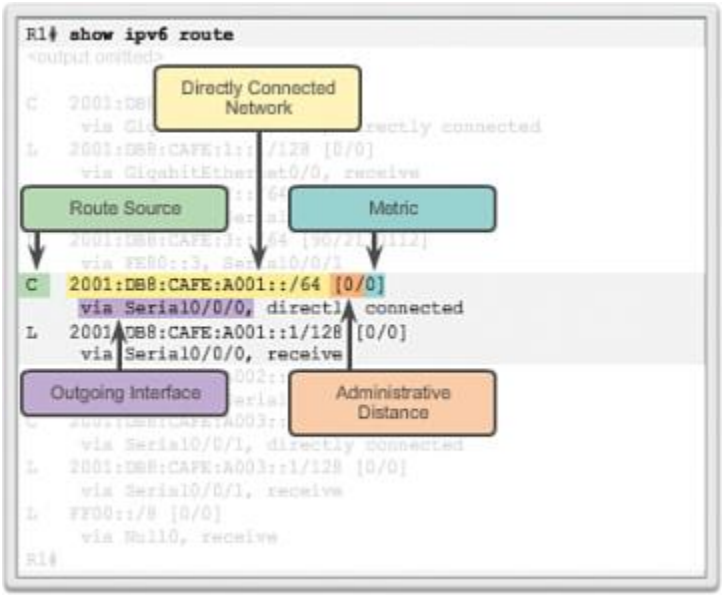


Figure 7-68 Directly Connected Routes on R1

The serial links have reference bandwidths configured to observe how EIGRP metrics select the best route. The reference bandwidth is not a realistic representation of modern networks. It is used only to provide a visual sense of link speed.

Remote IPv6 Network Entries

Figure 7-69 highlights the routing table entries for the three remote networks (i.e., R2 LAN, R3 LAN, and the link between R2 and R3). The three entries were added by the EIGRP.

```
R1# show ipv6 route
<output omitted>

C   2001:DB8:CAFE:1::/64 [0/0]
     via GigabitEthernet0/0, directly connected
L   2001:DB8:CAFE:1::1/128 [0/0]
     via GigabitEthernet0/0, receive
D   2001:DB8:CAFE:2::/64 [90/3524096]
     via FE80::3, Serial0/0/1
D   2001:DB8:CAFE:3::/64 [90/2170112]
     via FE80::3, Serial0/0/1
C   2001:DB8:CAFE:A001::/64 [0/0]
     via Serial0/0/0, directly connected
L   2001:DB8:CAFE:A001::1/128 [0/0]
     via Serial0/0/0, receive
D   2001:DB8:CAFE:A002::/64 [90/3523840]
     via FE80::3, Serial0/0/1
C   2001:DB8:CAFE:A003::/64 [0/0]
     via Serial0/0/1, directly connected
L   2001:DB8:CAFE:A003::1/128 [0/0]
     via Serial0/0/1, receive
L   FF00::/8 [0/0]
     via Null0, receive
R1#
```

Figure 7-69 Remote Networks Entries on R1

Figure 7-70 displays a routing table entry on R1 for the route to remote network 2001:DB8:CAFE:3::/64 on R3.

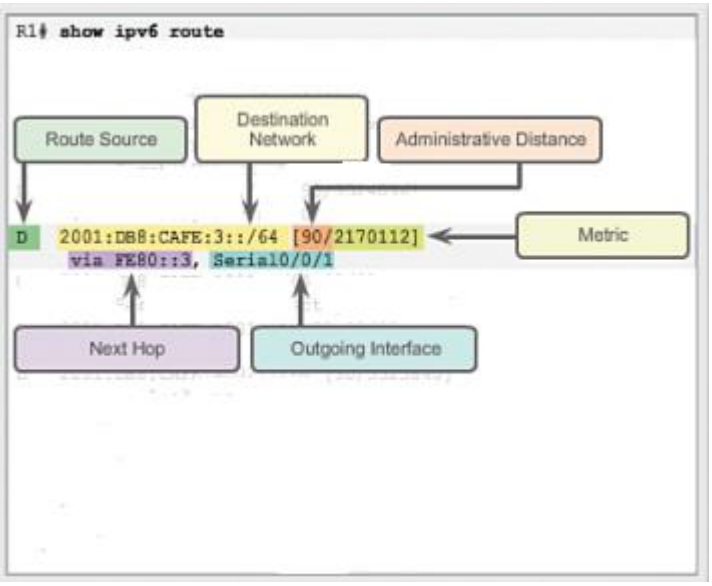


Figure 7-70 Remote Networks Entries on R1

The entry identifies the following information:

- **Route source**: Identifies how the route was learned. Common codes include O (OSPF), D (EIGRP), R (RIP), and S (Static route).
- **Destination network**: Identifies the address of the remote IPv6 network.
- **Administrative distance**: Identifies the trustworthiness of the route source. IPv6 uses the same distances as IPv4.
- **Metric**: Identifies the value assigned to reach the remote network. Lower values indicate preferred routes.
- **Next hop**: Identifies the IPv6 address of the next router to forward the packet to.
- **Outgoing interface**: Identifies the exit interface to use to forward a packet toward the final destination.

When an IPv6 packet arrives on a router interface, the router examines the IPv6 header and identifies the destination IPv6 address. The router then proceeds through the following router lookup process.

The router examines level 1 network routes for the best match with the destination address of the IPv6 packet. Just like IPv4, the longest match is the best match. For example, if there are multiple matches in the routing table, the router chooses the route with the longest match. A match is made by matching the far left bits of the packet's destination IPv6 address with the IPv6 prefix and prefix-length in the IPv6 routing table.

11. Summary

Dynamic routing protocols are used by routers to facilitate the exchange of routing information between routers. The purpose of dynamic routing protocols includes: discovery of remote networks, maintaining up-to-date routing information, choosing the best path to destination networks, and ability to find a new best path if the current path is no longer available. While dynamic routing protocols require less administrative overhead than static routing, they do require dedicating part of a router's resources for protocol operation, including CPU time and network link bandwidth.

Networks typically use a combination of both static and dynamic routing. Dynamic routing is the best choice for large networks and static routing is better for stub networks.

Routing protocols are responsible for discovering remote networks, as well as maintaining accurate network information. When there is a change in the topology, routing protocols propagate that information throughout the routing domain. The process of bringing all routing tables to a state of consistency, where all of the routers in the same routing domain or area have complete and accurate information about the network, is called convergence. Some routing protocols converge faster than others.

Routing protocols can be classified as either classful or classless, as distance vector or link-state, and as an Interior Gateway Protocol or an Exterior Gateway Protocol.

Distance vector protocols use routers as "sign posts" along the path to the final destination. The only information a router knows about a remote network is the distance or metric to reach that network and which path or interface to use to get there. Distance vector routing protocols do not have an actual map of the network topology.

A router configured with a link-state routing protocol can create a complete view or topology of the network by gathering information from all of the other routers.

Metrics are used by routing protocols to determine the best path or shortest path to reach a destination network. Different routing protocols may use different metrics. Typically, a lower metric means a better path. Metrics can be determined by hops, bandwidth, delay, reliability, and load.

Routers sometimes learn about multiple routes to the same network from both static routes and dynamic routing protocols. When a router learns about a destination network from more than one routing source, Cisco routers use the administrative distance value to determine which source to use. Each dynamic routing protocol has a unique administrative value, along with static routes and directly connected networks. The lower the administrative value, the more preferred the route source. A directly connected network is always the preferred source, followed by static routes and then various dynamic routing protocols.

The **show ip protocols** command displays the IPv4 routing protocol settings currently configured on the router. For IPv6, use **show ipv6 protocols**.

With link-state routing protocols such as OSPF, a link is an interface on a router. Information about the state of those links is known as link-states. All link-state routing protocols apply Dijkstra's algorithm to calculate the best path route. The algorithm is commonly referred to as the shortest path first (SPF) algorithm. This algorithm uses accumulated costs along each path, from source to destination, to determine the total cost of a route.

Reference:

- http://www.ciscopress.com/articles/article.asp?p=2180210