Your name Problem Set - 4

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ID: 123456789

Number of hours delay for this Problem Set:

Cumulative number of hours delay so far:

Put hours here
Put hours here

I discussed this homework with:

Put names here

Problem 1 - 15 points

Suppose that we have a production $A \to BCD$. Each of the four nonterminals A, B, C, and D have two attributes: s is a synthesized attribute, and i is an inherited attribute. For each of the sets of rules below, tell whether (i) the rules are consistent with an S-attributed definition (ii) the rules are consistent with an L-attributed definition, and (iii) whether the rules are consistent with any evaluation order at all?

```
a A: s = B: i + C: s.
```

b
$$A: s = B: i + C: s$$
 and $D: i = A: i + B: s$.

c
$$A : s = B : s + D : s$$
.

$$\mbox{d} \ A: s = D: i, \quad B: i = A: s + C: s, \quad C: i = B: s, \mbox{ and } D: i = B: i + C: i.$$

Solution:

a

b

 \mathbf{c}

d

Problem 2 - 15 points

Construct the DAG for the expression ((x+y)-((x+y)*(x-y)))+((x+y)*(x-y))

Solution: Your solutions go here

Problem 3 - 15 points

Translate the arithmetic expression a + (b + c).

- a A syntax tree.
- b Quadruples.
- c Triples.
- d Indirect triples.

Solution:

a
b
\mathbf{c}
d
Problem 4 - 20 points A real array $A[i; j; k]$ has index i ranging from 1 to 4, j ranging from 0 to 4, and k ranging from 5 to 10. Reals take 8 bytes each. If A is stored row-major, starting at byte 0, find the location of:
a $A[3;4;5]$
b $A[1;2;7]$
c A[4;3;9].
Repeat the above if A is stored in column-major order.
Solution: Row-major:
a
b
\mathbf{c}
Column-major:
\mathbf{a}
b
\mathbf{c}

Problem 5 - 20 points

Add rules to the syntax-directed definition of Fig. 1 for the following control-flow constructs:

PRODUCTION	SEMANTIC RULES
$P \rightarrow S$	S.next = newlabel() $P.code = S.code \mid\mid label(S.next)$
$S \rightarrow \mathbf{assign}$	$S.code = \mathbf{assign}.code$
$S \rightarrow \mathbf{if} (B) S_1$	B.true = newlabel() $B.false = S_1.next = S.next$ $S.code = B.code \mid\mid label(B.true) \mid\mid S_1.code$
$S \rightarrow \mathbf{if} (B) S_1 \mathbf{else} S_2$	$B.true = newlabel()$ $B.false = newlabel()$ $S_1.next = S_2.next = S.next$ $S.code = B.code$ $ label(B.true) S_1.code$ $ gen('goto' S.next)$ $ label(B.false) S_2.code$
$S \rightarrow $ while $(B) S_1$	$begin = newlabel()$ $B.true = newlabel()$ $B.false = S.next$ $S_1.next = begin$ $S.code = label(begin) B.code$ $ label(B.true) S_1.code$ $ gen('goto' begin)$
$S \rightarrow S_1 S_2$	$S_1.next = newlabel()$ $S_2.next = S.next$ $S.code = S_1.code \mid\mid label(S_1.next) \mid\mid S_2.code$

Figure 1: Rules to the syntax-directed definition

- \bullet A repeat-statement **repeat** S while B.
- A for-loop for $(S_1; B; S_2)S_3$.

Solution:

Problem 6 - 15 points

Translate the following expressions using the ifFalse mechanism:

```
a if (a == b \ \&\& \ c == d \ || \ e == f) \ x == 1;
b if (a == b \ || \ c == d \ || \ e == f) \ x == 1;
c if (a == b \ \&\& \ c == d \ \&\& \ e == f) \ x == 1;
```

Solution:

a

 \mathbf{c}