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Course: CSDS 337 - Compiler Design

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**Problem Set - 3**

ID: 3559750

Term: Spring 2024

Due Date: 28<sup>th</sup> Feb, 2024

Number of hours delay for this Problem Set:

0

Cumulative number of hours delay so far:

26

I discussed this homework with:

Jackson Schuetzle.

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### Problem 1 - 10 points

Design grammar for the following language:

- The set of all strings of 0s and 1s such that every 0 is immediately followed by at least one 1.

*Solution:*

$S \rightarrow 0x \mid 1S \mid \epsilon$

$x \rightarrow 1S$

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### Problem 2 - 20 points

The following is a grammar for regular expressions over symbols  $a$  and  $b$  only, using  $+$  in place of  $|$  for union, to avoid conflict with the use of vertical bar as a metasymbol in grammars:

$rexpr \rightarrow rexpr + rterm \mid rterm$

$rterm \rightarrow rterm \ rfactor \mid rfactor$

$rfactor \rightarrow rfactor * \mid rprimary$

$rprimary \rightarrow a|b$

- Left factor this grammar.
- Does left factoring make the grammar suitable for top-down parsing?
- In addition to left factoring, eliminate left recursion from the original grammar.
- Is the resulting grammar suitable for top-down parsing?

*Solution:* Your solutions go here

- The grammar is already left factored. No nonterminal has two or more productions beginning with the same symbol.
- No because the grammar is left recursive.

c  $rexpr \rightarrow A_1$

$A_1 \rightarrow +rterm \ A_1 \mid \epsilon$

$rterm \rightarrow rfactor \ A_2$

$A_2 \rightarrow rfactor \ A_2 \mid \epsilon$

$rfactor \rightarrow rprimary \ A_3$

$A_3 \rightarrow *A_3 \mid \epsilon$

$rprimary \rightarrow a \mid b$

d Yes because it is left-factored and right-recursive.

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**Problem 3 - 20 points**

Consider the grammar for  $S \rightarrow S + S | SS | (S) | S * | a$  and the string  $(a + a) * a$ .

- a Devise a predictive parser and show the parsing tables. You may use left-factor and/or eliminate left-recursion from your grammar first.
- b Compute FIRST and FOLLOW for your grammar.

*Solution:*

- a Left-factor the grammar:

$$\begin{aligned} S &\rightarrow SS' \mid (S) \mid a \\ S' &\rightarrow +S \mid S \mid * \end{aligned}$$

Eliminate left recursion:

$$\begin{aligned} S &\rightarrow (S)A \mid aA \\ A &\rightarrow S'A \mid \epsilon \\ S' &\rightarrow +S \mid S \mid * \end{aligned}$$

b

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**Problem 4 - 10 points**

Give the bottom-up parses for the following input string and grammar:  $aaa * a + +$  and  $S \rightarrow SS + | SS * | a$ .

*Solution:*

$\rightarrow \underline{aaa} * a + +$   
 $\rightarrow S \underline{aa} * a + +$   
 $\rightarrow SS \underline{a} * a + +$   
 $\rightarrow SSS \underline{*} a + +$   
 $\rightarrow SS \underline{a} + +$   
 $\rightarrow SSS \underline{+} +$   
 $\rightarrow SS \underline{+}$   
 $\rightarrow S$

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**Problem 5 - 20 points**

Construct the SLR sets of items for the (augmented) grammar  $S \rightarrow SS + | SS * | a$ . Compute the GOTO function for these sets of items. Show the parsing table for this grammar. Is this grammar SLR?

*Solution:*

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**Problem 6 - 20 points**

Construct the canonical parsing table for the following augmented grammar:

$$\begin{aligned} S' &\rightarrow S \\ S &\rightarrow AA \\ A &\rightarrow aA | b \end{aligned}$$

*Solution:*