ECOLE POLYTECHNIQUE DE LOUVAIN

LINGI2132 - LANGUAGES AND TRANSLATORS

Assignement 2 - Report

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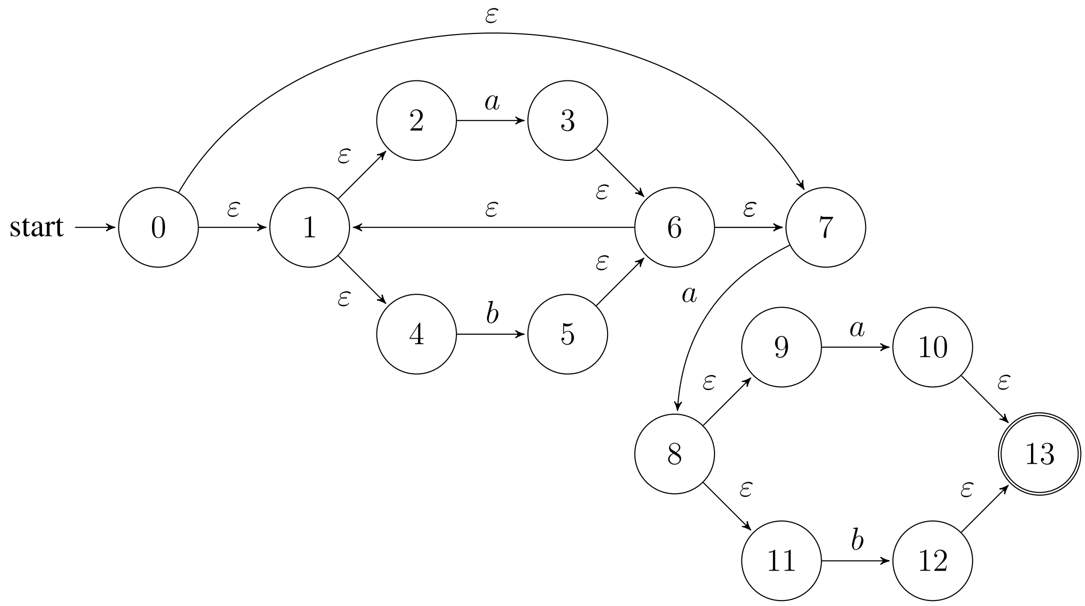
Academic Year 2013-2014

1 Lexical Analysis

1.1 Give 5 different strings belonging to the language described by this regExp.

1. aab
2. bbaa
3. ababab
4. aaaaaaa
5. bbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbab

1.2 Construct the NFA from this regExp (Thompson Construction



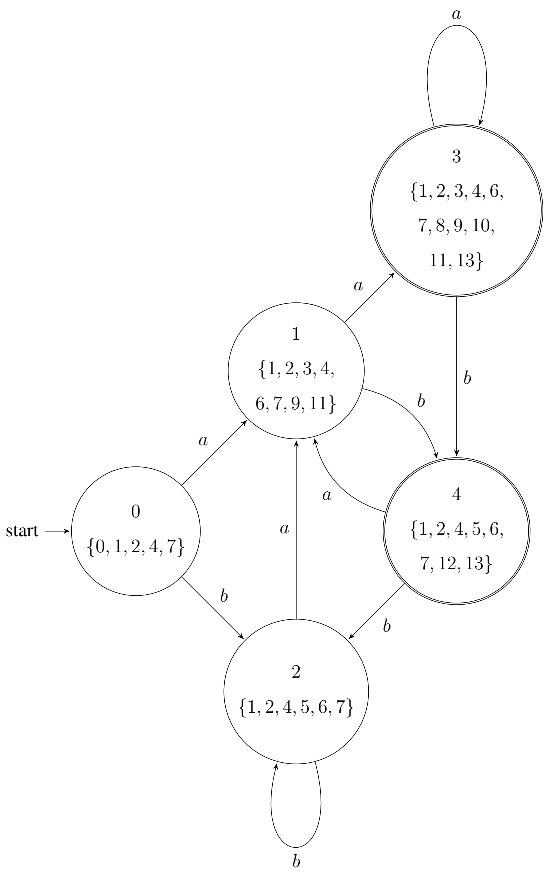
1.3 Transform the NFA into a DFA (justify important steps, ε-closures).

Steps from NFA to DFA :

- $s_0 = \epsilon - closure(\{0\}) = \{0, 1, 2, 4, 7\}$
- $m(s_0, a) = s_1$, where $s_1 = \epsilon - closure(\{3, 8\}) = \{1, 2, 3, 4, 6, 7, 9, 11\}$

- $m(s_0, b) = s_2$, where $s_2 = \epsilon - closure(\{5\}) = \{1, 2, 4, 5, 6, 7\}$
- $m(s_1, a) = s_3$, where $s_3 = \epsilon - closure(\{3, 8, 10\}) = \{1, 2, 3, 4, 6, 7, 8, 9, 11, 13\}$
- $m(s_1, b) = s_4$, where $s_4 = \epsilon - closure(\{5, 12\}) = \{1, 2, 4, 5, 6, 7, 12, 13\}$
- $m(s_2, a) = s_1$
- $m(s_2, b) = s_2$
- $m(s_3, a) = s_3$
- $m(s_3, b) = s_4$
- $m(s_4, a) = s_1$
- $m(s_4, b) = s_2$

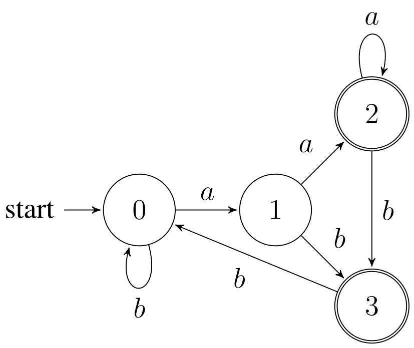
Since s_3 and s_4 contain state 13 which was final in the NFA, they are both final state too.



1.4 Minimize the DFA (Hopcroft Algorithm).

The DFA could be divided in four partitions :

- Two partitions for final states
- A partition which contains previous states $\{0, 2\}$ because :
  - $m(0, a) = 1$
  - $m(2, a) = 1$
  - $m(0, b) = 2$
  - $m(2, b) = 2$
- A partition which contains previous state $\{1\}$:
  - $m(1, a) = 3$
  - $m(1, b) = 4$



2 Parsing

2.1

This grammar is not LL(1) because it has a rule with left recursion: $B := Bv$ and $B := w$.

2.2

We replace the left recursion with a right recursion to avoid the problem indicated in the first exercice.

$$B := wB'$$
$$B' := vB'$$
$$B' := \epsilon$$

Thus we remove from the grammar the rules $B := Bv$ and $B := w$ and we introduce the rules above.

2.3

$first(u) = \{u\}$

$first(v) = \{v\}$

$first(w) = \{w\}$

$first(x) = \{x\}$

$first(y) = \{y\}$

$first(z) = \{z\}$

$\epsilon \in first(E), first(F), first(B')$

$Rule1 : S :== uBDz$

$first(uBDz) = \{u\}$

$\rightarrow u \in first(S)$

$Rule2 : B ::= wB'$

$first(wB') = \{w\}$

$\rightarrow w \in first(B)$

$Rule3 : B' ::= wB'$

$first(vB') = \{v\}$

$\rightarrow v \in first(B')$

$Rule5 : D ::= EF$

$first(EF) = first(E) \backslash \{\epsilon\} \cup first(F)$ (car $E \in first(E)$)= $\{x, y, E\}$

$\rightarrow x, y, E \subset first(D)$

$Rule6 : E ::= y$

$y \in first(E)$

$Rule8 : F ::= x$

$x \in first(F)$

$first(S) = \{u\}$

$first(B) = \{w\}$

$first(B') = \{v, \epsilon\}$

$first(D) = \{x, y, \epsilon\}$

$first(E) = \{y, \epsilon\}$

$first(F) = \{x, \epsilon\}$

Follow

$\#E\, follow(S)$

$Rule1 : S ::= uBDz$

$\rightarrow \{x, y, z\} \subset Bfollow(B)$

$and\, z \in follow(D)$

$Rule2 :\ B ::= wB'$

$\rightarrow follow(B) = \{x, y, z\} \subset follow(B')$

$Rule5 : D ::= EF$

$\rightarrow x \in follow(E)$

$follow(D) = \{z\} \subset follow(E)$

$and\, follow(D) = \{z\} \subset follow(F)$

$follow(S) = \{\#\}$

$follow(B) = \{x, y, z\}$

$follow(B') = \{x, y, z\}$

$follow(D) = \{z\}$

$follow(E) = \{x, z\}$

$follow(F) = \{z\}$

2.4

$Rule1 : S ::= uBDz$

$\rightarrow table[S, u] = 1$

$Rule2 : B ::= wB'$

$\rightarrow table[B, w] = 2$

$Rule3 : B' ::= vB'$

$\rightarrow table[B', v] = 3$

$Rule4 : B' ::= \epsilon$

$\rightarrow table[B', x] = 4$

$\rightarrow table[B', y] = 4$

$\rightarrow table[B', z] = 4$

$Rule5 : D ::= EF$

$\rightarrow table[D, x] = 5$

$\rightarrow table[D, y] = 5$

$\rightarrow table[D, z] = 5$

$Rule6 : E ::= y$

$\rightarrow table[E, y] = 6$

$Rule7 : E ::= \epsilon$

$\rightarrow table[E, x] = 7$

$\rightarrow talbe[E, z] = 7$

$Rule8 : F ::= x$

$\rightarrow table[F, x] = 8$

$Rule9 : F ::= \epsilon$

$\rightarrow table[F, z] = 9$

	u	v	w	x	y	z
S	1					
B			2			
B’		3		4	4	4
D				5	5	5
E				7	6	7
F				8		9