

CSC343 Assignment 3

Ben Cheng

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$R_1 = \text{DEFGHIJK}$, $S_1 = \{D \rightarrow FG, E \rightarrow HK, F \rightarrow EIJ, F \rightarrow K\}$

1. $D^+ = \text{DEFGHIJK}$
 $E^+ = \text{EHK}$
 $F^+ = \text{EFHIJK}$

Since E and F are not superkeys of the relation, we know that each of them violates BCNF:

$E \rightarrow H, K$
 $F \rightarrow E, I, J, K$

2. First, we can create the relation $R_1 = \text{EHK}$ with the functional dependency $E \rightarrow HK$. This relation is now in BCNF.

However, the original relation $R_2 = \text{DEFGIJ}$ is still not in BCNF because $F^+ = \text{EFIJ}$. Therefore, we need to further decompose R .

We can create two new relations: $R_3 = \text{EFIJ}$ and $R_4 = \text{DFG}$. R_3 and R_4 is now in BCNF.

The result is $\{R_1 = \text{EHK}, R_3 = \text{EFIJ}, R_4 = \text{DFG}\}$

The relations is $\{\text{DFG}, \text{EFIJ}, \text{EHK}\}$

3. The BCNF decomposition does not guarantee dependencies but dependencies are preserved this time.
Because the removed relation $F \rightarrow K$ can be derived by $F \rightarrow E$ and $E \rightarrow H$.

4. 1.1

| D | E | F | G | H | I | J | K |
|---|---|---|---|---|---|---|---|
| d | e | f | g | h | i | j | k |
| | e | | | h | | | k |
| | w | f | | h | i | j | k |
| | e | f | | | i | j | k |

A whole row is filled, so lossless.

2

$$R_2 = JKLMNOPQ$$

$$S_2 = \{JLM \rightarrow N, K \rightarrow LM, KN \rightarrow JLO, M \rightarrow JKO, N \rightarrow JL\}$$

1. Step 1. Split RHS

- $JLM \rightarrow N$
- $K \rightarrow L$
- $K \rightarrow M$
- $KN \rightarrow J$
- $KN \rightarrow L$
- $KN \rightarrow O$
- $M \rightarrow J$
- $M \rightarrow K$
- $M \rightarrow O$
- $N \rightarrow J$
- $N \rightarrow L$

$JLM \rightarrow N$ can't be saved as $J = J, L = L, M = M, JL = JL, JM = JM$, and $LM = LM$.

$K \rightarrow L, K \rightarrow M, M \rightarrow J, M \rightarrow K, M \rightarrow O$, and $N \rightarrow J, N \rightarrow L$ can't be saved as they are singleton.

For $KN \rightarrow J, KN \rightarrow L$, and $KN \rightarrow O$, $K^* = JKLMNO$ covers all of them, so they can be saved as $K \rightarrow J, K \rightarrow L$, and $K \rightarrow O$.

So we are left with

- $JLM \rightarrow N$
- $K \rightarrow L$
- $K \rightarrow M$
- $K \rightarrow J$
- $K \rightarrow O$

- $M \rightarrow J$
- $M \rightarrow K$
- $M \rightarrow O$
- $N \rightarrow J$
- $N \rightarrow L$

(a)

| | FD | | J | K | L | M | N | O | P | Q | |
|---|---------------------|-------------------|----------|----------|----------|----------|----------|----------|----------|----------|---------|
| a | $JLM \rightarrow N$ | $S_{2-\{a\}}$ | J | K | L | M | | O | | | need |
| b | $K \rightarrow L$ | $S_{2-\{b\}}$ | J | K | | M | | O | | | need |
| c | $K \rightarrow M$ | $S_{2-\{c\}}$ | J | K | L | | | O | | | need |
| d | $K \rightarrow J$ | $S_{2-\{d\}}$ | J | K | L | M | N | O | | | no need |
| e | $K \rightarrow O$ | $S_{2-\{d,e\}}$ | J | K | L | M | N | O | | | no need |
| f | $M \rightarrow J$ | $S_{2-\{d,e,f\}}$ | | K | L | M | | O | | | need |
| g | $M \rightarrow K$ | $S_{2-\{d,e,g\}}$ | J | | | M | | O | | | need |
| h | $M \rightarrow O$ | $S_{2-\{d,e,h\}}$ | J | K | L | M | N | | | | need |
| i | $N \rightarrow J$ | $S_{2-\{d,e,i\}}$ | | | L | | N | | | | need |
| j | $N \rightarrow L$ | $S_{2-\{d,e,j\}}$ | J | | | | N | | | | need |

$$\{JLM \rightarrow N, K \rightarrow LM, M \rightarrow JKO, N \rightarrow JL\}$$

2. Attributes on the left but not right: none
 Attributes on the right but not left: O, it's in no key
 Attributes on both left and right: JKLMN, need to check
 Attributes on neither left and right: P, Q, in every key

(a) :

| J | K | L | M | N | Closure |
|----------|----------|----------|----------|----------|--------------------------------|
| 1 | 0 | 0 | 0 | 0 | $JPQ^* = JPQ(\text{no})$ |
| 0 | 1 | 0 | 0 | 0 | $KPQ^* = JKLMNOPQ(\text{yes})$ |
| 0 | 0 | 1 | 0 | 0 | $LPQ^* = LPQ(\text{no})$ |
| 0 | 0 | 0 | 1 | 0 | $MPQ^* = JKLMNOPQ(\text{yes})$ |
| 0 | 0 | 0 | 0 | 1 | $NPQ^* = JLNPQ(\text{no})$ |
| 1 | 1 | 0 | 0 | 0 | no |
| 1 | 0 | 1 | 0 | 0 | $JLPQ^* = JLPQ(\text{no})$ |
| 1 | 0 | 0 | 1 | 0 | no |
| 1 | 0 | 0 | 0 | 1 | $JNPQ^* = JLPQ(\text{no})$ |
| 0 | 1 | 1 | 0 | 0 | no |
| 0 | 1 | 0 | 1 | 0 | no |
| 0 | 1 | 0 | 0 | 1 | no |
| 0 | 0 | 1 | 1 | 0 | no |
| 0 | 0 | 1 | 0 | 1 | no |
| 0 | 0 | 0 | 1 | 1 | no |

The keys: KPQ, MPQ

3. Relation: JLMN, KLM, MJKO, NJL

But still, no relation is a superkey of P and Q

Take a key KPQ

JLMN, KLM, MJKO, NJL, KPQ Becomes: JLMN, KLM, MJKO, KPQ
as NJL contained in JLMN