# **CG2271 Real Time Operating Systems**

### **Lab 6 - Synchronization Mechanisms**

#### **Answer Book**

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Lab day: Friday Lab time: 4pm ~ 6pm

# Question 1 (2 marks)

Output is pasted below:

ADD: New value of ctr is 1 PRINT: Current value of ctr is 1

PRINT: ctr is too small. Zzzz.. g'night!

ADD: New value of ctr is 2 ADD: New value of ctr is 3 ADD: New value of ctr is 4

ADD: New value of ctr is 5

ADD: New value of ctr is 6

ADD: New value of ctr is 7

ADD: New value of ctr is 8 ADD: New value of ctr is 9

ADD: New value of ctr is 10

ADD: Reached limit of 10! Waking up print thread

PRINT: ctr is now 10! Exiting thread.

This program would first create two threads: add and print.

When print is executed, it would first lock ctr\_mutex, then check if ctr reaches the MAX\_COUNT; if not, it would be blocked, unlock the ctr\_mutex, and wait for the ctr cond signal.

In the add function, it would add ctr each time in the while loop; if ctr reaches hte MAX\_COUNT, it would signal "ctr\_cond" so that the blocked print can be executed.

### Question 2 (3 marks)

This program does not deadlock because the "pthread\_cond\_wait" would automatically unlock the "ctr\_mutex".

#### Question 3 (2 marks)

At least one child thread does not block because the initial value of "sema1" is 1; therefore when the sem\_wait is first called, it's positive therefore no need to be blocked.

### Question 4 (2 marks)

Each time the main thread calls sem post, 1 thread is woken.

# Question 5 (2 marks)

POSIX semaphores are different from mutexes because semaphore is implemented as a counter – it would count up when there's a "post" signal, and count-down when there's a "wait" signal – by doing so, it allows mulitple threads to be in "run-state" even if the "wait" is being called in each thread; while for the mutex, it's a global boolean variable – it would have no use when one tries to unlock a unlocked process. (but for semophere by doing a "post" would count up, therefore in the future this count-up can cancel a "wait" signal.)

### Question 6 (3 marks)

The changes we made were to use a semophare to control the order of execution.

1. add in the inlcude file for semophare:

```
#include <semaphore.h>
```

2. add in a semaphore variable:

```
sem_t sema1;
```

3. in the "add" function, wait for semal at the start:

```
sem_wait(&sema1);
```

4. in the "print" function, after the printf statement, add in the "sem\_post(&sema1); " to wake up add function

## Question 7 (2 marks)

```
The final sum is 0 + 1 + 2 + ... + 49 = 1225
```

# Question 8 (3 marks)

The sum is incorrect. This is because the function reader executes before writter has written anything into the queue – meaning, even if there's nothing to read from the queue

## Question 9 (6 marks)

The changes we made were:

- 1. add in "semaphore.h" library to pcomm.h
- 2. in the struct t (struct for the queue), add in two semaphore:

```
1. sem_t sema_full; // wait on full push
```

2. sem\_t sema\_empty; // wait on empty get

- 3. in pq\_create, initialize the two semaphores:
  - 1. sem\_init(&(tmp->sema\_full),1,0);
  - 2. sem init(&(tmp->sema empty),1,0);
- 4. in pq\_destory, destory the two semaphores:
  - sem\_destroy(&(q->sema\_full));
  - sem\_destroy(&(q->sema\_empty));
- 5. In pq\_put, add in wait for semaphore when the queue is full:

```
// Wait if full.
while(q->count == q->size){
    sem_wait(&(q->sema_full));
}
```

6. also in this function, when the number of elements changed from 0 to one, give a post to sema\_empty:

```
if(q->count == 1) sem_post(&q->sema_empty);
```

- 7. Similarly in pq\_get function, add in this
  - 1. // Wait on empty -- a while loop is needed to use up all the positive posts

```
while(q->count == 0) sem wait(\&(q->sema empty));
```

2. // check to wake up the put, if it's no longer full

```
if(q->count == q->size -1) sem_post(&q->sema_full);
```

## Question 10 (5 marks)

The changes we made, in the file "lab6\_3.c" were:

- 1. include the library "semaphore.h"
- 2. create a global variable (semaphore): "sem t all completed"
- 3. in the main function:
  - initialize this semaphore: sem\_init(&all\_completed,1,0);
  - 2. before destorying the queue, wait 6 times for all programs to complete:

 $for(i=0;\ i<6;\ i++) \quad sem\_wait(\&all\_completed);$ 

4. at the end of the writter function and the end of the reader function, post to semaphore all\_complete:

sem\_post(&all\_completed); // finish execution, post to allow main to go

ANSWER BOOK TOTAL: \_\_\_\_/30

**DEMO:** \_\_\_\_/5

TOTAL: \_\_\_\_/35