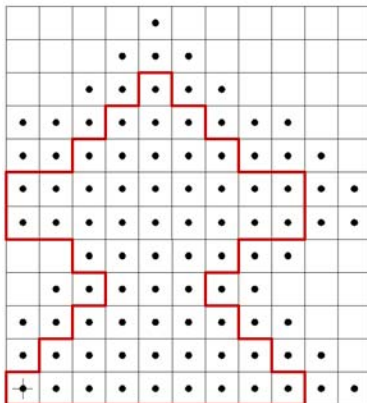
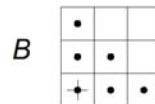
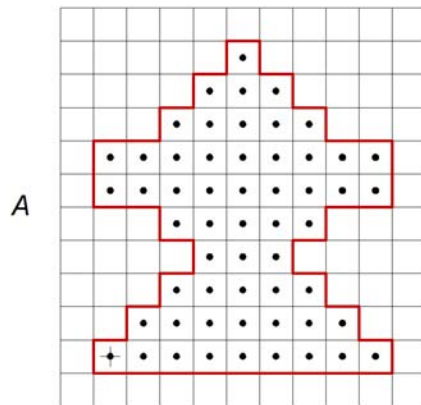
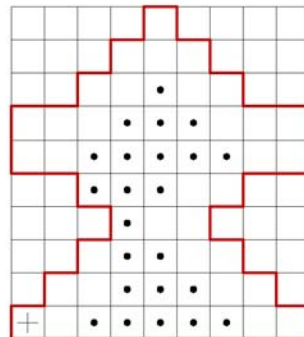
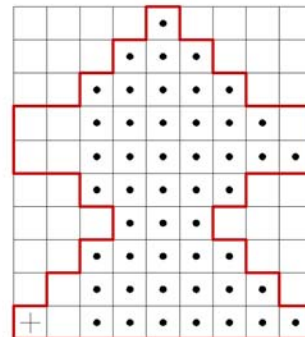
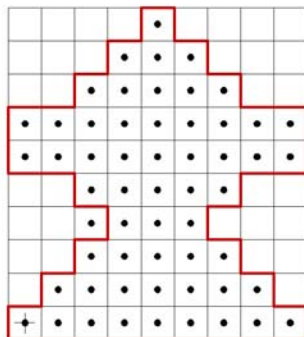
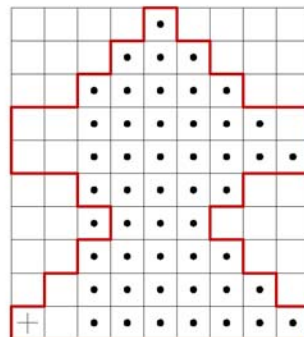
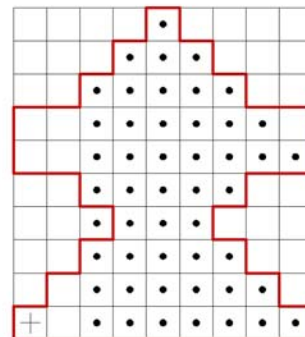


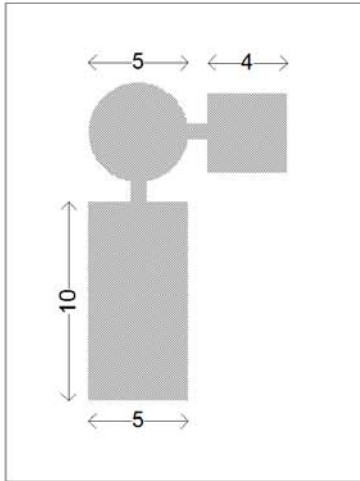
# EE3206/EE3206E INTRODUCTION TO COMPUTER VISION AND IMAGE PROCESSING

## Tutorial Set G – Solutions

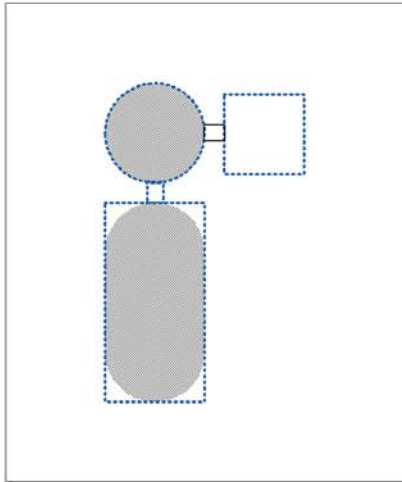
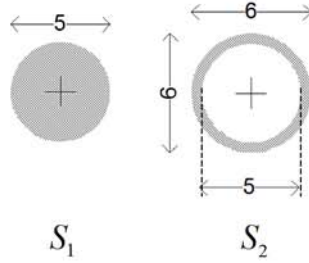
### Question 1


 $A \oplus B$ 

 $A \ominus B$ 

 $A \circ B$ 

 $A \bullet B$ 

 $(A \circ B) \bullet B$ 

 $(A \bullet B) \circ B$

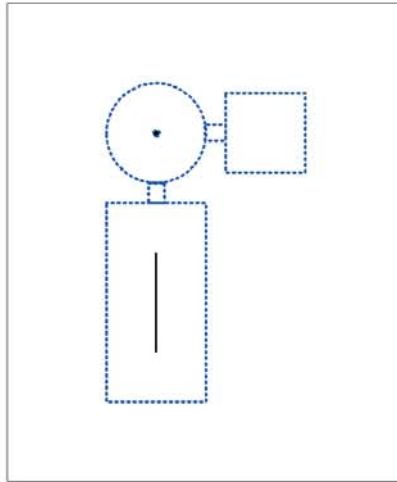
## Question 2



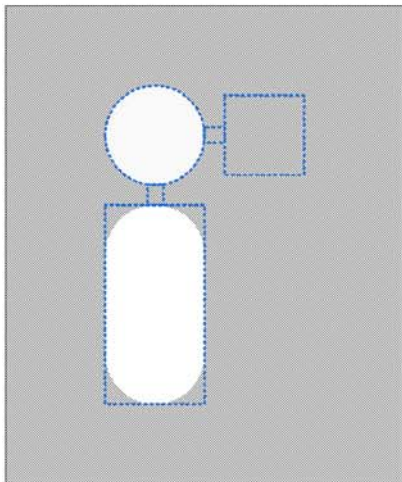
$I$



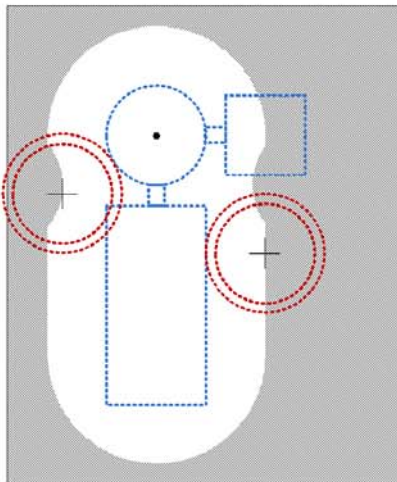
$I_1 = I \circ S_1$



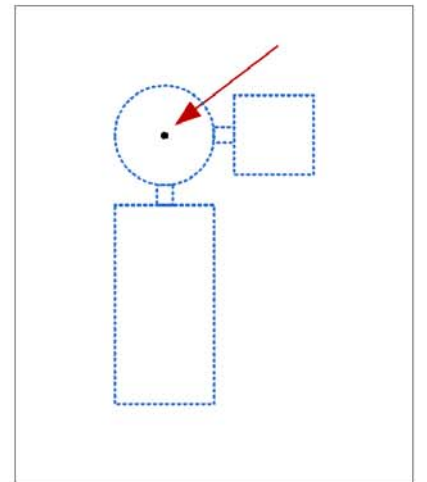
$I_1 \ominus S_1$



$I_1^c$



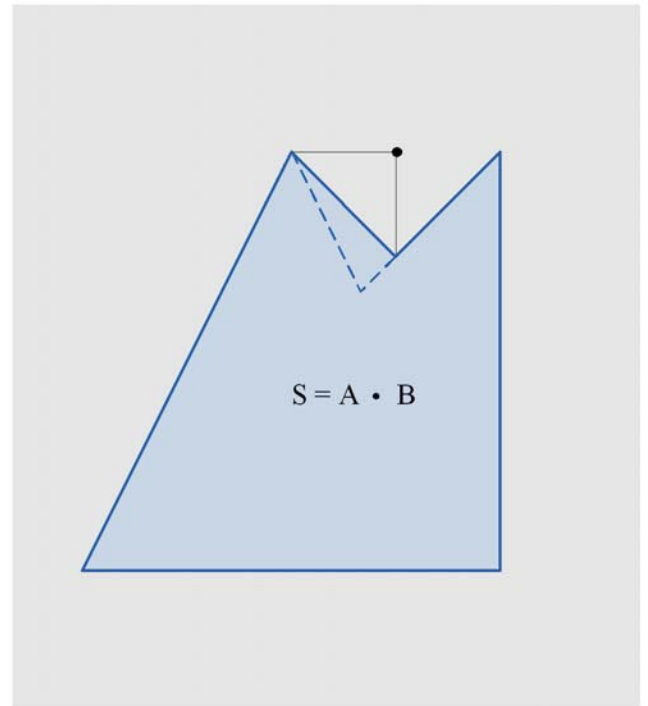
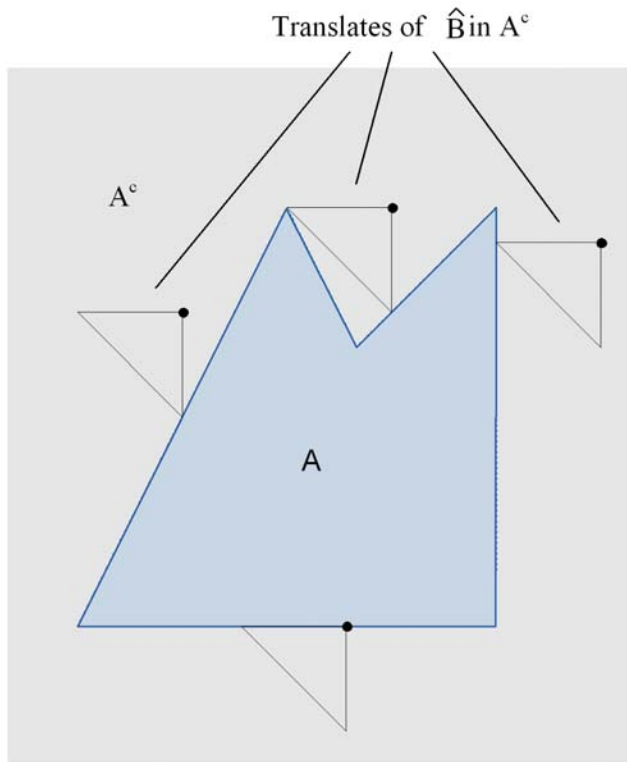
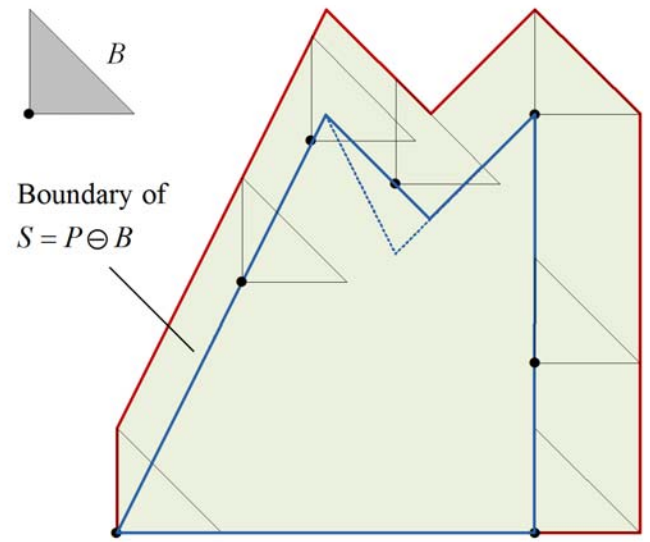
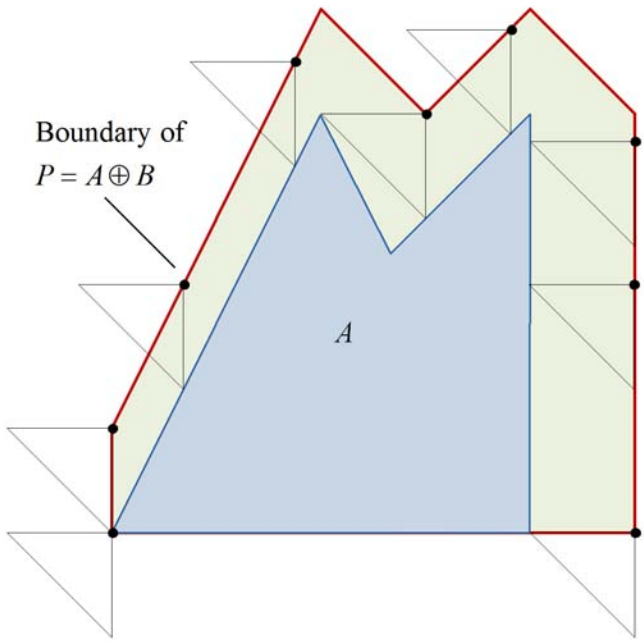
$I_1^c \ominus S_2$



$I_2 = (I_1 \ominus S_1) \cap (I_1^c \ominus S_2)$

### Question 3

$$\begin{aligned}
 (A \bullet B)^c &= [(A \oplus B) \ominus B]^c && \text{(definition of closing)} \\
 &= (A \oplus B)^c \oplus \hat{B} && \text{(duality of erosion and dilation)} \\
 &= (A^c \ominus \hat{B}) \oplus \hat{B} && \text{(duality of erosion and dilation)} \\
 &= A^c \circ \hat{B} && \text{(definition of opening)}
 \end{aligned}$$

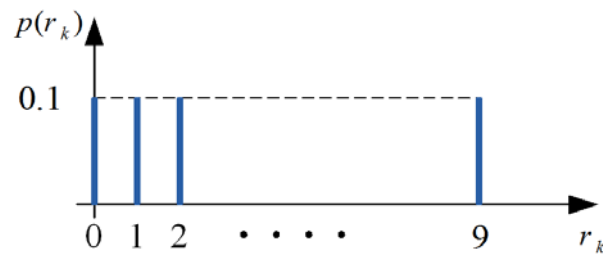


## Question 4

### Part (a)

$$\begin{aligned} I(E) &= -\log_r P(E) \\ &= -\log_2(0.2) = 2.322 \text{ bits} \\ &= -\log_{10}(0.2) = 0.699 \text{ Hartleys} \end{aligned}$$

### Part(b)

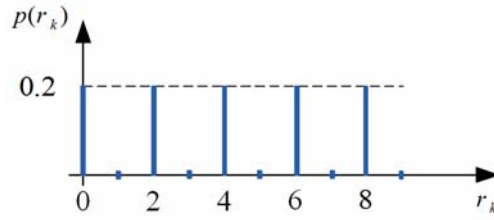


Probability values :  $p(r_k) = 0.1, \quad r_k = 0, 1, 2, \dots, 9$

The entropy is

$$\begin{aligned} H &= -\sum_{k=0}^9 p(r_k) \log p(r_k) \\ &= -\sum_{k=0}^9 0.1 \log(0.1) \\ &= 0.1 \log(10) \times 10 \\ &= 3.322 \text{ bits} \end{aligned}$$

**Part (c)**

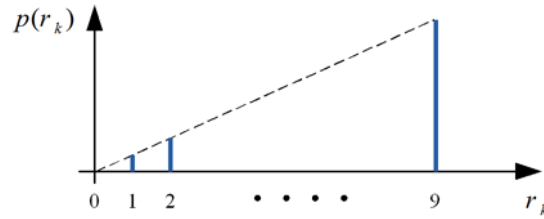


Probability values :  $p(r_k) = 0.2, \quad r_k = 0, 2, 4, 6, 8$

The entropy is

$$\begin{aligned}
 H &= - \sum_{k=0}^9 p(r_k) \log p(r_k) \\
 &= - \sum_{k=0}^9 0.2 \log(0.2) \\
 &= -0.2 \log(0.2) \times 5 \\
 &= \log(5) \\
 &= 2.322 \text{ bits}
 \end{aligned}$$

**Part (d)**



Probability values :  $p(r_k) = Kr_k, \quad r_k = 0, 1, 2, \dots, 9$

$$\sum_k p(r_k) = 1 \Rightarrow K = \frac{1}{45}$$

Hence,

$$p(r_k) = \frac{1}{45} r_k$$

The entropy is

$$\begin{aligned}
 H &= - \sum_{k=0}^9 p(r_k) \log p(r_k) \\
 &= - \sum_{k=0}^9 \frac{r_k}{45} \log \frac{r_k}{45} \\
 &= 2.96 \text{ bits}
 \end{aligned}$$

## Question 5

Symbol :	$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$a_5$	$a_6$	$a_7$
Gray level :	0	1	2	3	4	5	6	7
Probability :	0.4	0.08	0.08	0.2	0.12	0.08	0.03	0.01

### Part (a)

The entropy is

$$\begin{aligned}
 H &= -\sum P(a_i) \log_2 P(a_i) \\
 &= -0.4 \log 0.4 - 3 \times 0.08 \log 0.08 - 0.12 \log 0.12 \\
 &\quad - 0.2 \log 0.2 - 0.03 \log 0.03 - 0.01 \log 0.01 \\
 &= 2.453 \text{ bits}
 \end{aligned}$$

Coding efficiency using the natural binary code is

$$2.453/3 = 81.8\%$$

### Part (b)

Original source		Source reduction									
Symbol	Prob.	1		2		3		4		5	6
$a_0$	0.4	<b>1</b>	0.4	<b>1</b>	0.4	<b>1</b>	0.4	<b>1</b>	0.4	<b>1</b>	0.6 <b>0</b>
$a_3$	0.2	<b>000</b>	0.2	<b>000</b>	0.2	<b>000</b>	0.2	<b>000</b>	0.24	<b>01</b>	0.36 <b>00</b> 0.4 <b>1</b>
$a_4$	0.12	<b>010</b>	0.12	<b>010</b>	0.12	<b>010</b>	0.16	<b>001</b>	0.2	<b>000</b>	0.24 <b>01</b>
$a_1$	0.08	<b>0010</b>	0.08	<b>0010</b>	0.12	<b>011</b>	0.12	<b>010</b>	0.16	<b>001</b>	
$a_2$	0.08	<b>0011</b>	0.08	<b>0011</b>	0.08	<b>0010</b>	0.12	<b>011</b>			
$a_5$	0.08	<b>0110</b>	0.08	<b>0110</b>	0.08	<b>0011</b>					
$a_6$	0.03	<b>01110</b>	0.04	<b>0111</b>							
$a_7$	0.01	<b>01111</b>									

Gray Level	Prob.	Straight binary code	Huffman code	$L$
0	0.4	000	1	1
1	0.08	001	0010	4
2	0.08	010	0011	4
3	0.2	011	000	3
4	0.12	100	010	3
5	0.08	101	0110	4
6	0.03	110	01110	5
7	0.01	111	01111	6

Average code length for the Huffman code is

$$\begin{aligned}
\bar{L} &= (1 \times 0.4) + (4 \times 0.08) + (4 \times 0.08) + (3 \times 0.2) \\
&\quad + (3 \times 0.12) + (4 \times 0.08) + (5 \times 0.03) + (5 \times 0.01) \\
&= 2.520 \text{ bits}
\end{aligned}$$

Code efficiency is

$$\eta = \frac{2.453}{2.520} = 97.3\%$$

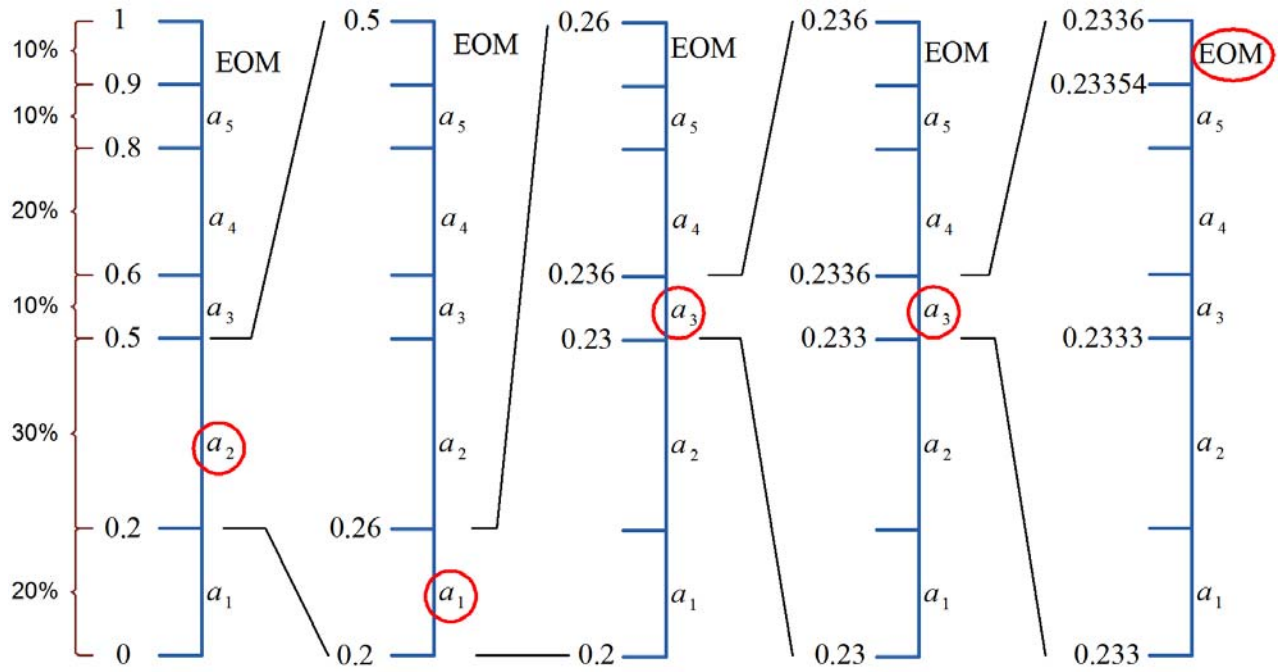
### Part (c)

With 3 bits/pixel, the image occupies  $3 \times 10^4$  bits. Therefore,

$$\text{savings} = (3 - 2.52) \times 10^4 = 4,800 \text{ bits} \quad (16\%)$$

## Question 6

### Part (a)



Hence,  $0.23355 \rightarrow a_2 a_1 a_3 a_3$  (EOM)



**Part (b)**

0	1	2	3	4	5
0	1	2	3	4	5
0	1	2	3	4	5
0	1	2	3	4	5
0	1	2	3	4	5
0	1	2	3	4	5

$I_1$

1	1	3	3	5	5
1	1	3	3	5	5
1	1	3	3	5	5
1	1	3	3	5	5
1	1	3	3	5	5
1	1	3	3	5	5

$I_2$

The symbol probabilities are:

Image $I_1$			Image $I_2$		
Symbol	Gray-level	Prob.	Symbol	Gray-level	Prob.
$a_0$	0	1/6	$a_0$	0	0
$a_1$	1	1/6	$a_1$	1	1/3
$a_2$	2	1/6	$a_2$	2	0
$a_3$	3	1/6	$a_3$	3	1/3
$a_4$	4	1/6	$a_4$	4	0
$a_5$	5	1/6	$a_5$	5	1/3

Image  $I_1$ :

After the first symbol,  $R_1 = \left(\frac{1}{6}\right)$ .

After the second symbol,  $R_2 = \left(\frac{1}{6}\right)^2$ .

...

After the sixth symbol,  $R_6 = \left(\frac{1}{6}\right)^6 = 2.143 \times 10^{-5}$

Image  $I_2$ :

After the first symbol,  $R_1 = \left(\frac{1}{3}\right)$ .

After the second symbol,  $R_2 = \left(\frac{1}{3}\right)^2$ .

...

After the sixth symbol,  $R_6 = \left(\frac{1}{3}\right)^6 = 1.372 \times 10^{-3}$ .

After the 36th pixel,  $R_{36}$  for  $I_1$  would be much smaller than  $R_{36}$  for  $I_2$ . Since more decimal digits are needed for a smaller range, Image  $I_1$  would require more digits for transmission.

Image  $I_1$

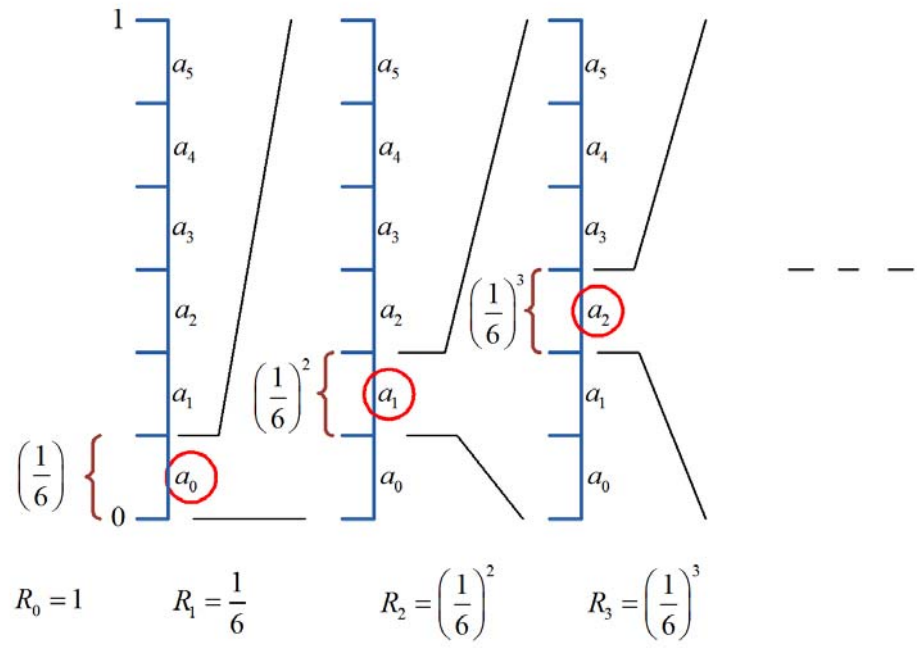
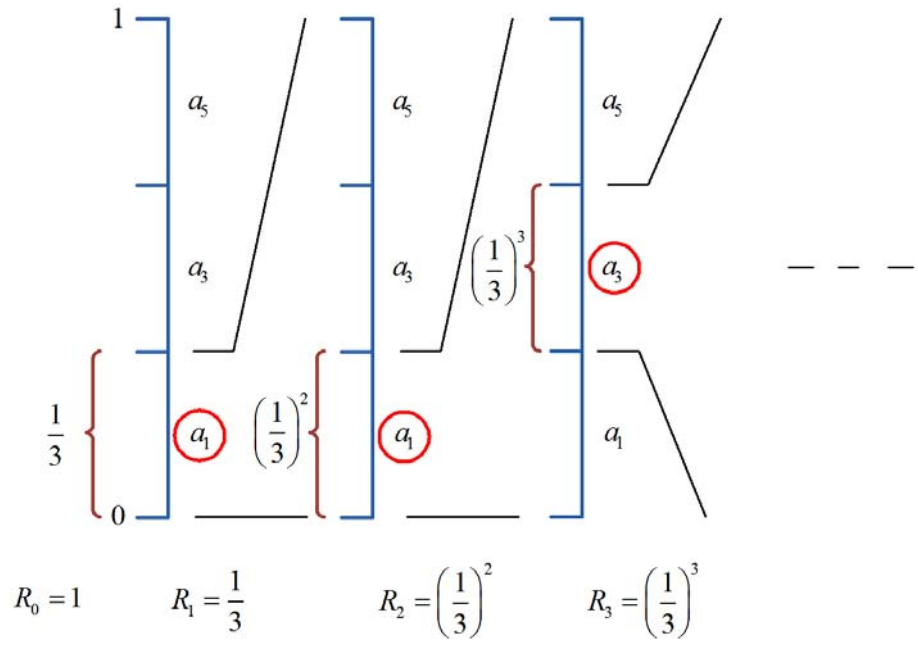


Image  $I_2$



## Question 7

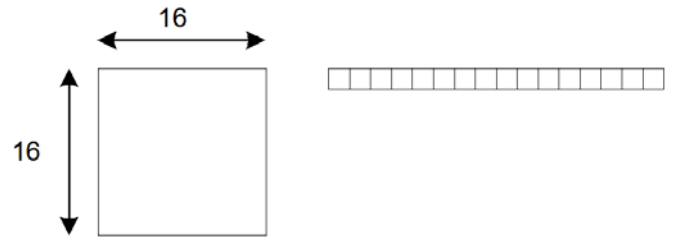
Each pixel is stored as 1 byte. Without run-length coding, the number of bytes required is

$$N_0 = 16^2 = 256$$

Our run-length coding scheme assumes that each row begins with a white pixel, and each run requires 1 byte (to denote the length of the run). For a row starting with a black pixel, an extra byte is needed for the first run (of zero length).

One run per row, 16 rows  
Number of bytes required is

$$\begin{aligned} N_1 &= 16 < N_0 \\ C_R &= 16 \end{aligned}$$

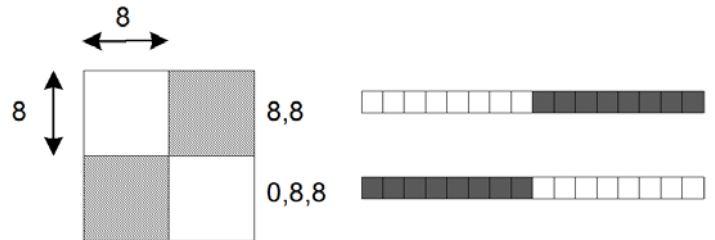


For a row starting with 1, there are two runs

For a row starting with 0, there are three runs

Number of bytes required is

$$\begin{aligned} N_2 &= 8 \times 2 + 8 \times 3 = 40 < N_0 \\ C_R &= 6.4 \end{aligned}$$

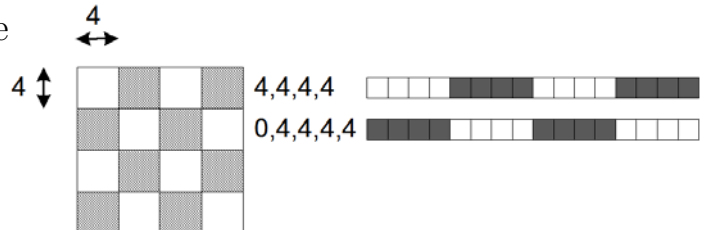


For a row starting with 1, there are four runs

For a row starting with 0, there are five runs

Number of bytes required is

$$\begin{aligned} N_4 &= 8 \times 4 + 8 \times 5 = 72 < N_0 \\ C_R &= 3.6 \end{aligned}$$



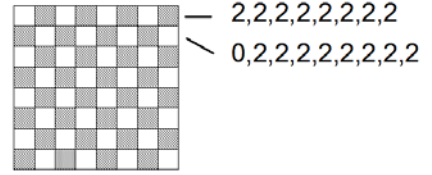
For a row starting with 1, there are eight runs

For a row starting with 0, there are nine runs

Number of bytes required is

$$N_8 = 8 \times 8 + 8 \times 9 = 136 < N_0$$

$$C_R = 1.9$$



For a row starting with 1, there are sixteen runs

For a row starting with 0, there are seventeen runs

Number of bytes required is

$$N_{16} = 8 \times 16 + 8 \times 17 = 264 > N_0$$

$$C_R = 0.97$$