

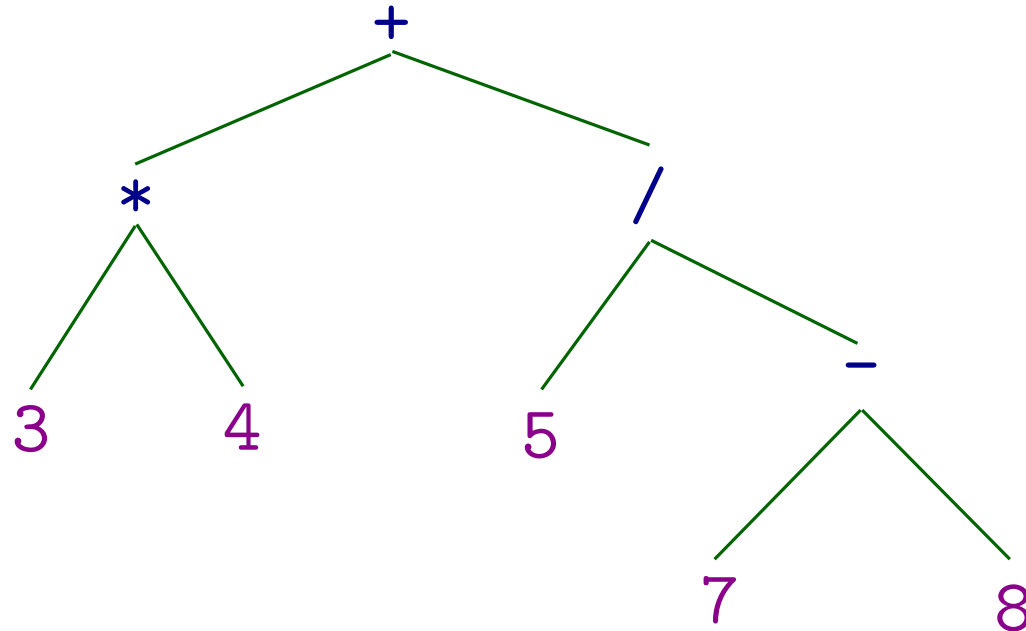
Compilation of Expressions

CS4212 – Lecture 3 (live)

Outline

- Version 1: simple binary expressions
- Version 2: add conditional operators
- Version 3: better engineering
- Version 4: adding assignments

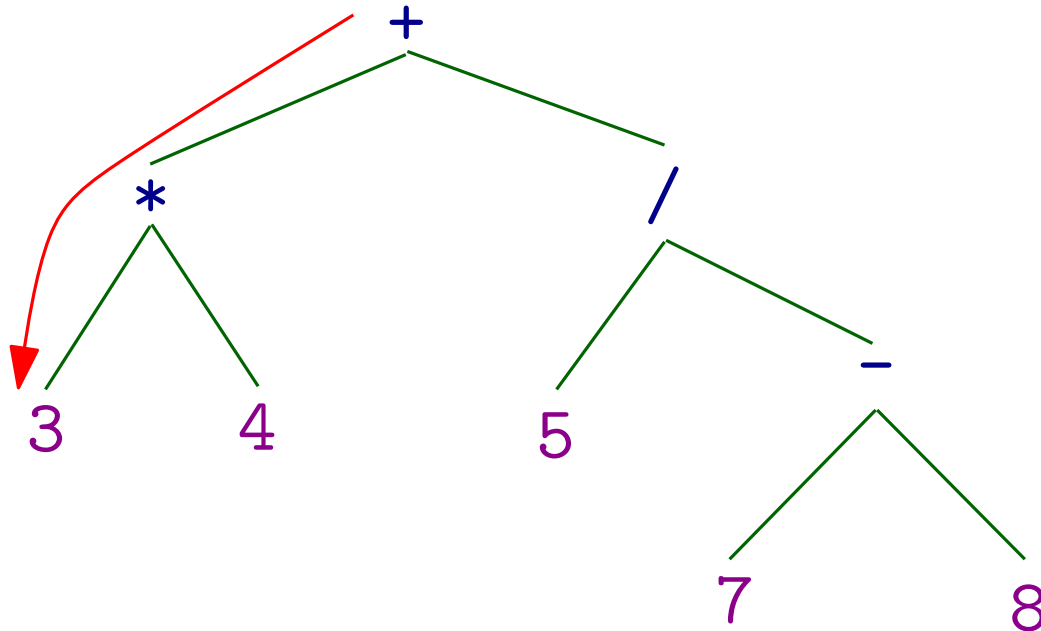
Compilation by Post-Order Traversal



```

pushl $3
pushl $4
popl %ebx
popl %eax
imull %ebx,%eax
pushl %eax
pushl $5
pushl $7
pushl $8
popl %ebx
popl %eax
sub %ebx,%eax
pushl %eax
popl %ebx
popl %eax
cdq
idivl %ebx
pushl %eax
popl %ebx
popl %eax
addl %ebx,%eax
pushl %eax
  
```

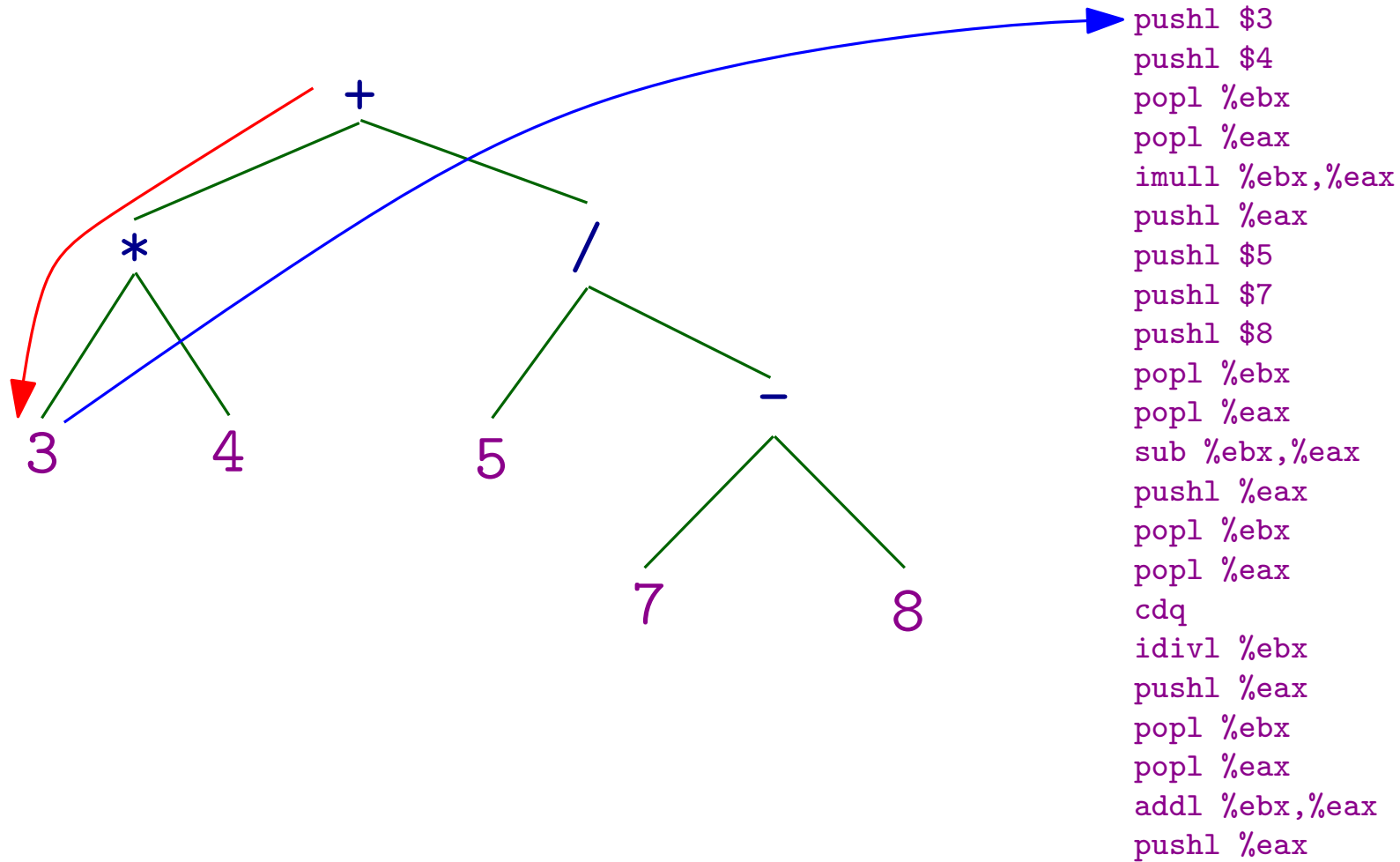
Compilation by Post-Order Traversal



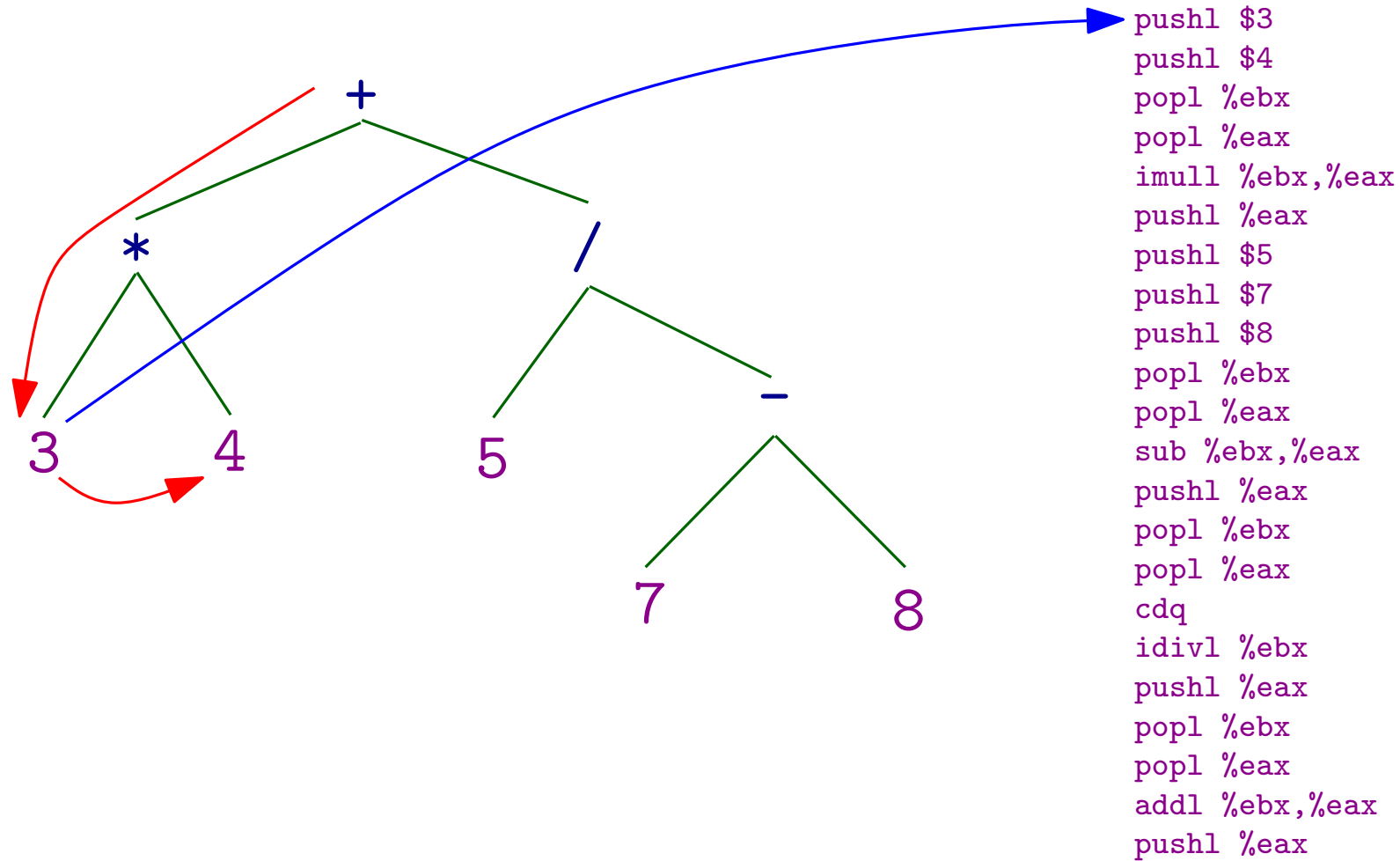
```

pushl $3
pushl $4
popl %ebx
popl %eax
imull %ebx,%eax
pushl %eax
pushl $5
pushl $7
pushl $8
popl %ebx
popl %eax
sub %ebx,%eax
pushl %eax
popl %ebx
popl %eax
cdq
idivl %ebx
pushl %eax
popl %ebx
popl %eax
addl %ebx,%eax
pushl %eax
  
```

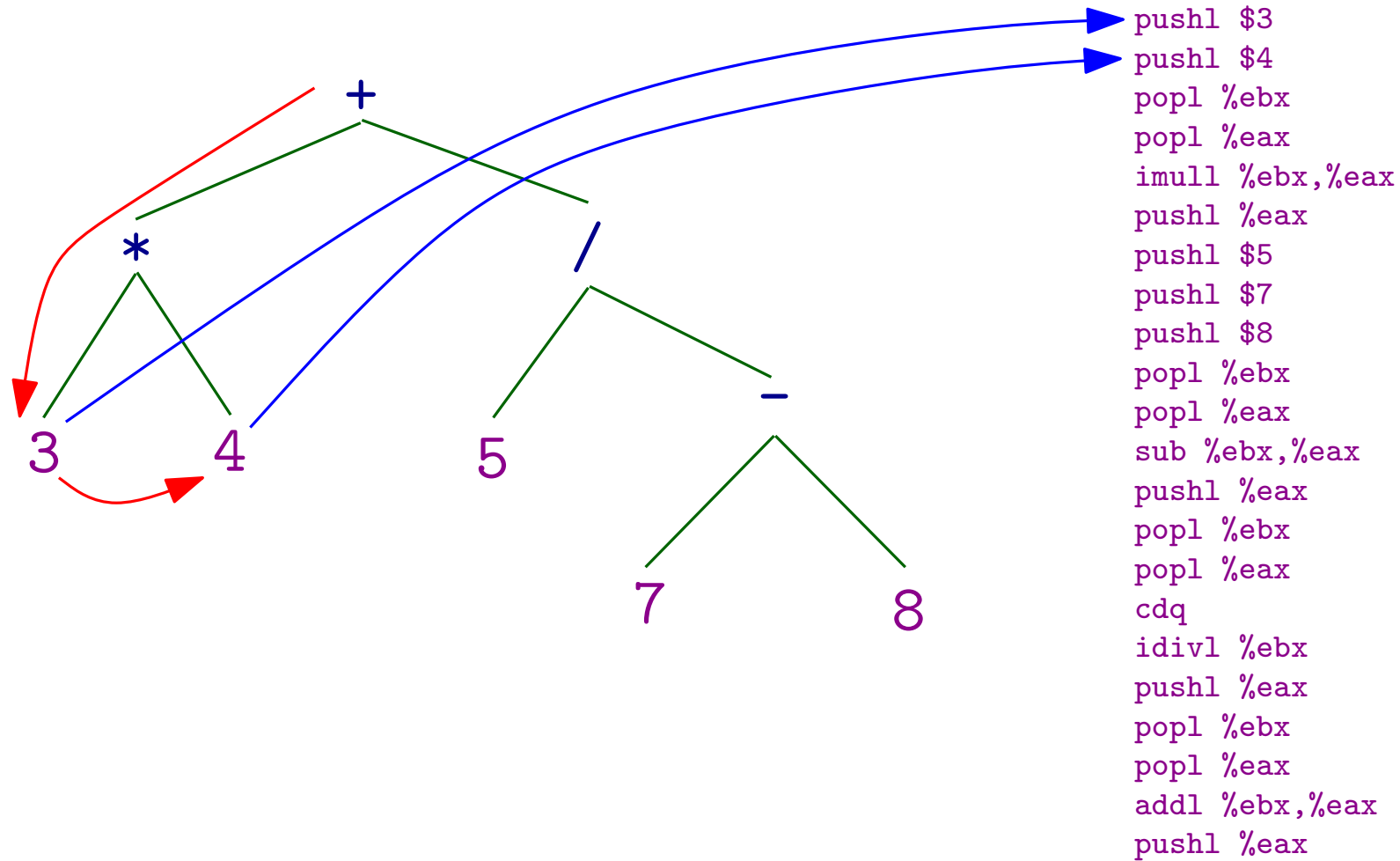
Compilation by Post-Order Traversal



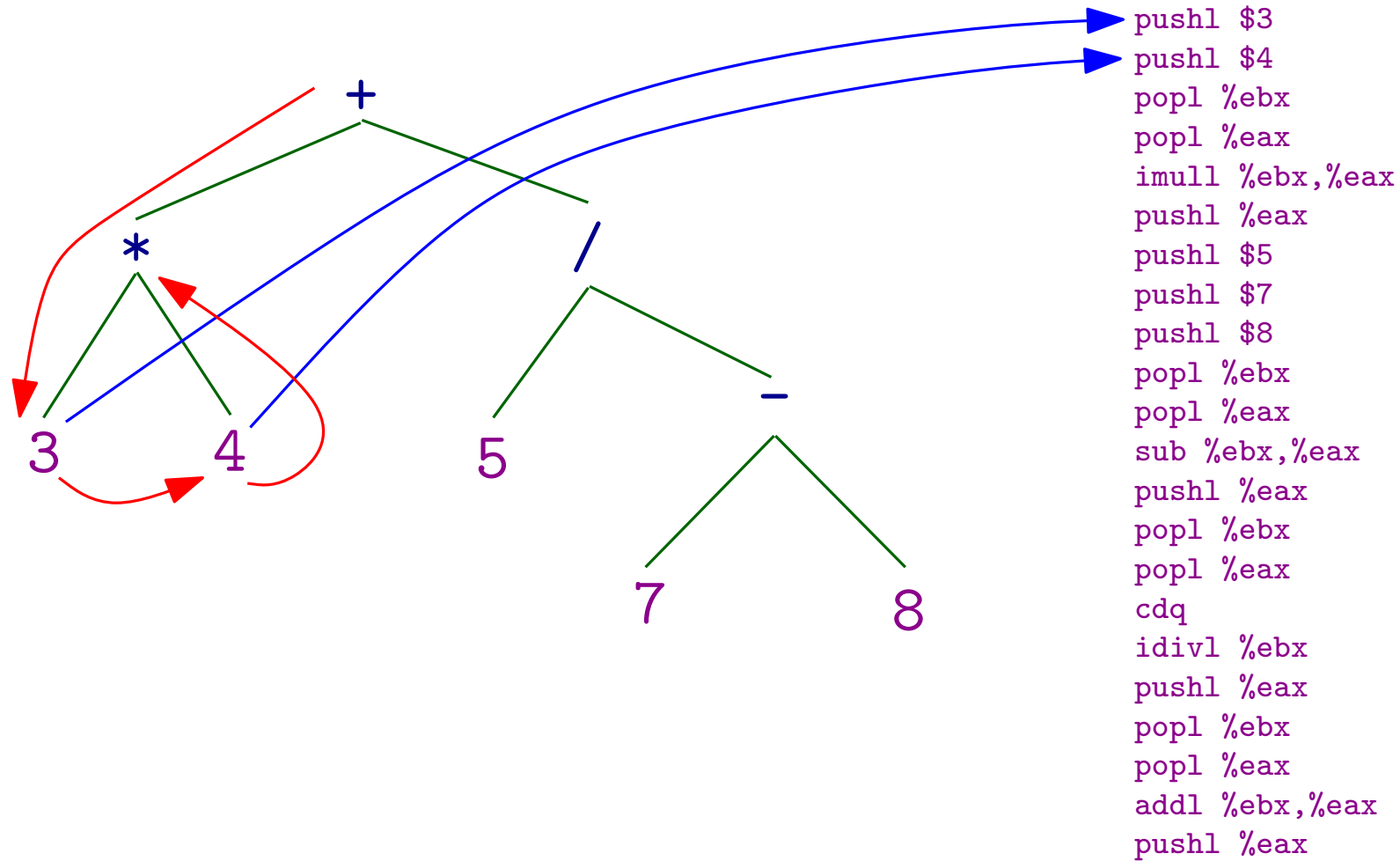
Compilation by Post-Order Traversal



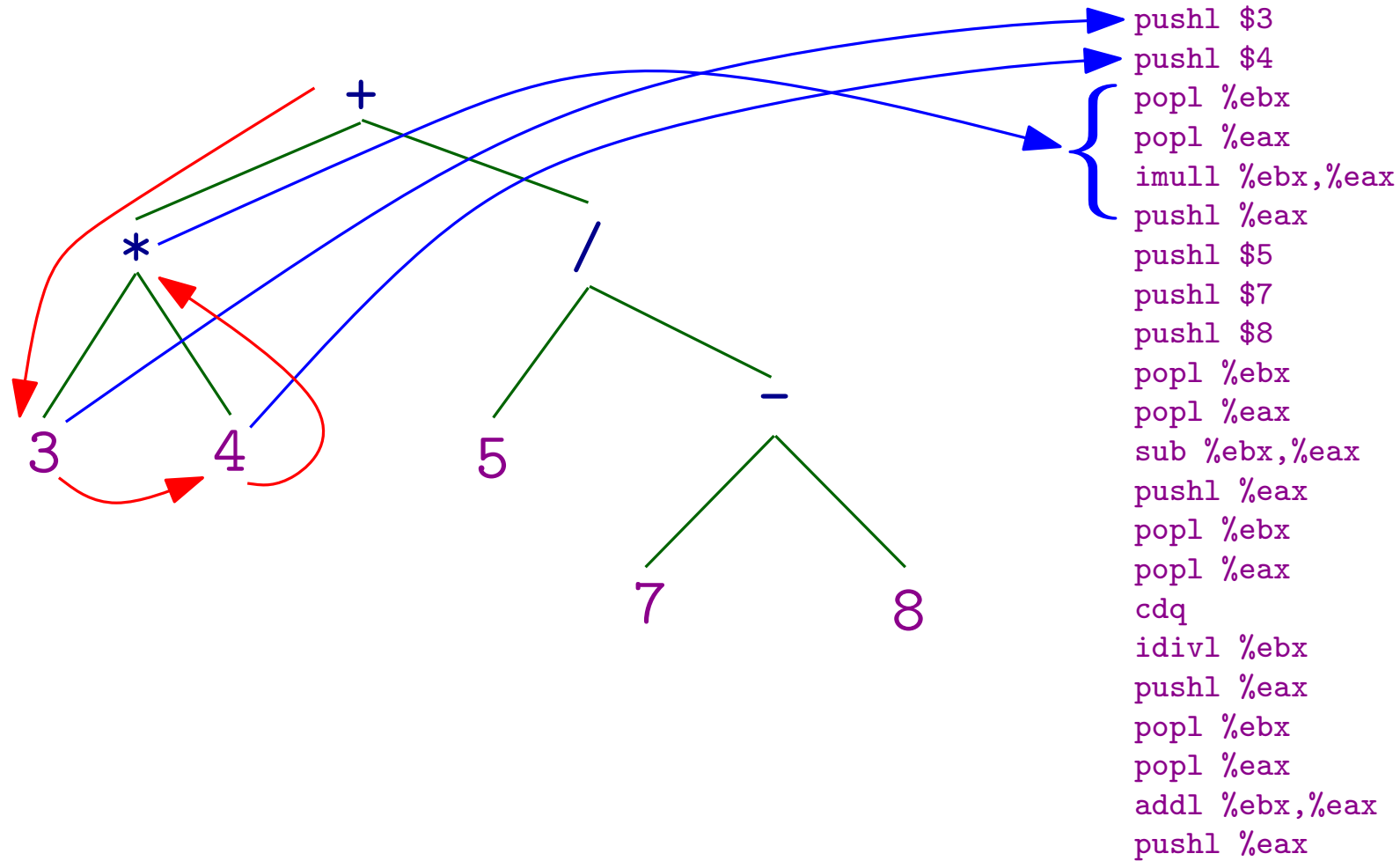
Compilation by Post-Order Traversal



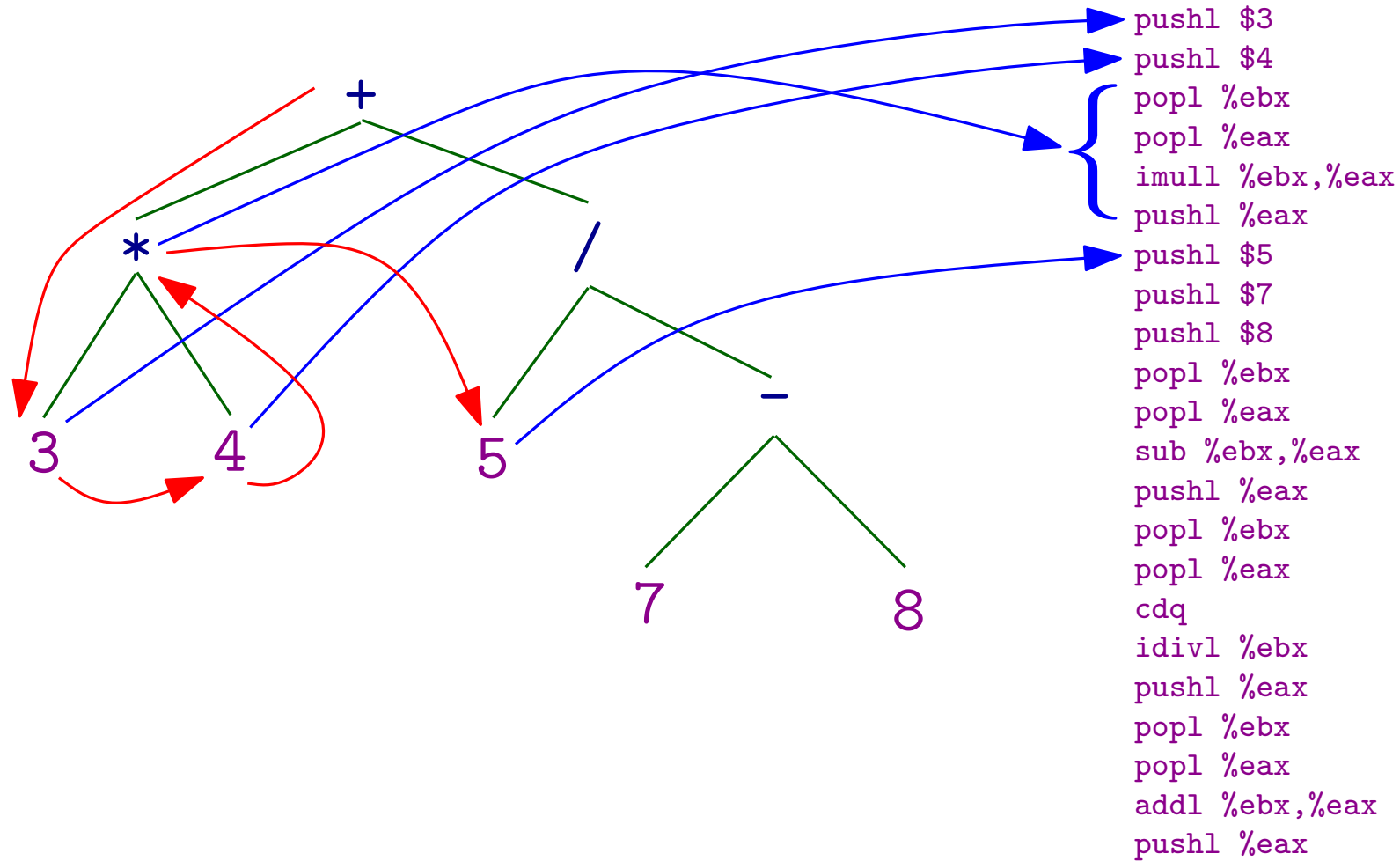
Compilation by Post-Order Traversal



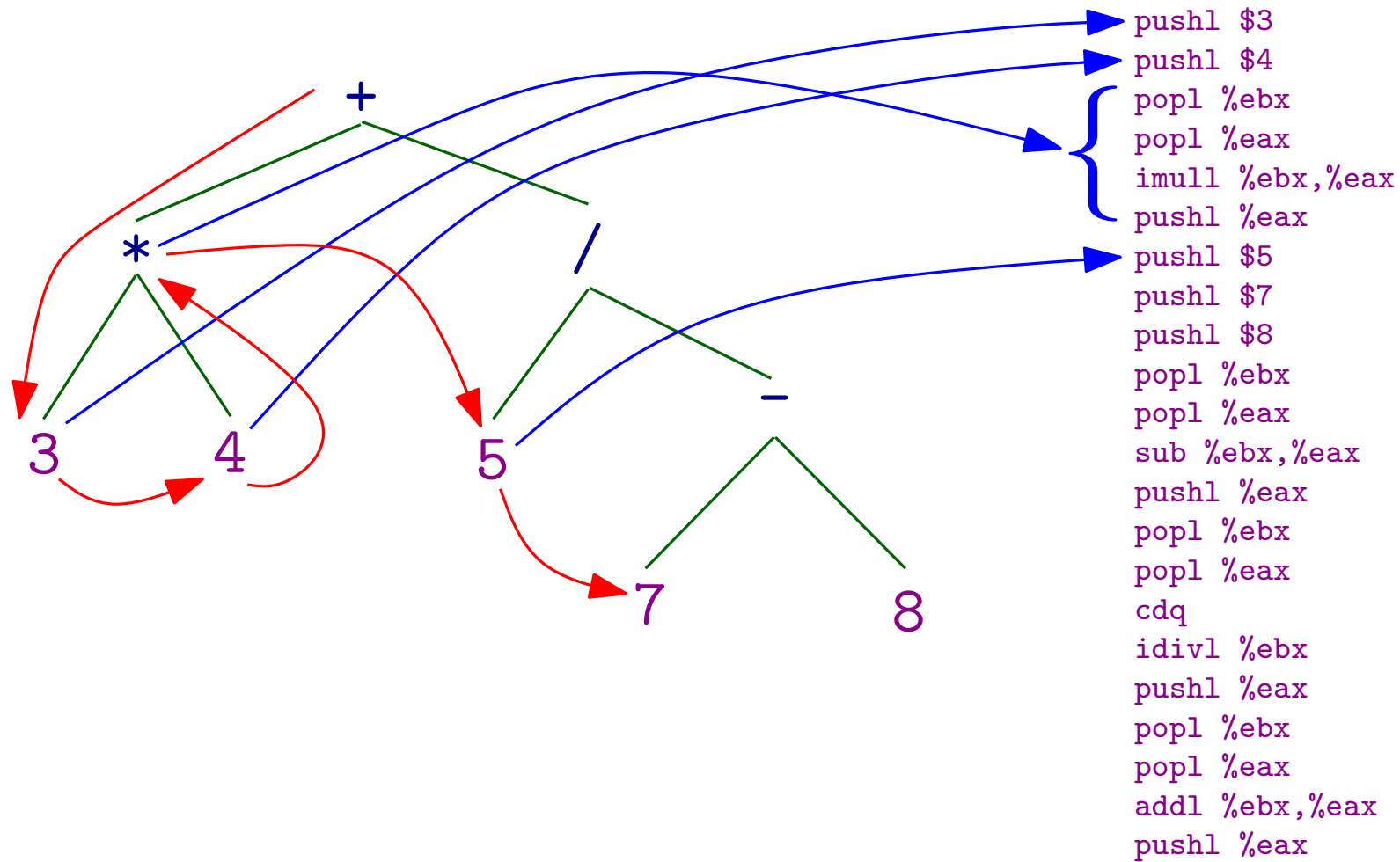
Compilation by Post-Order Traversal



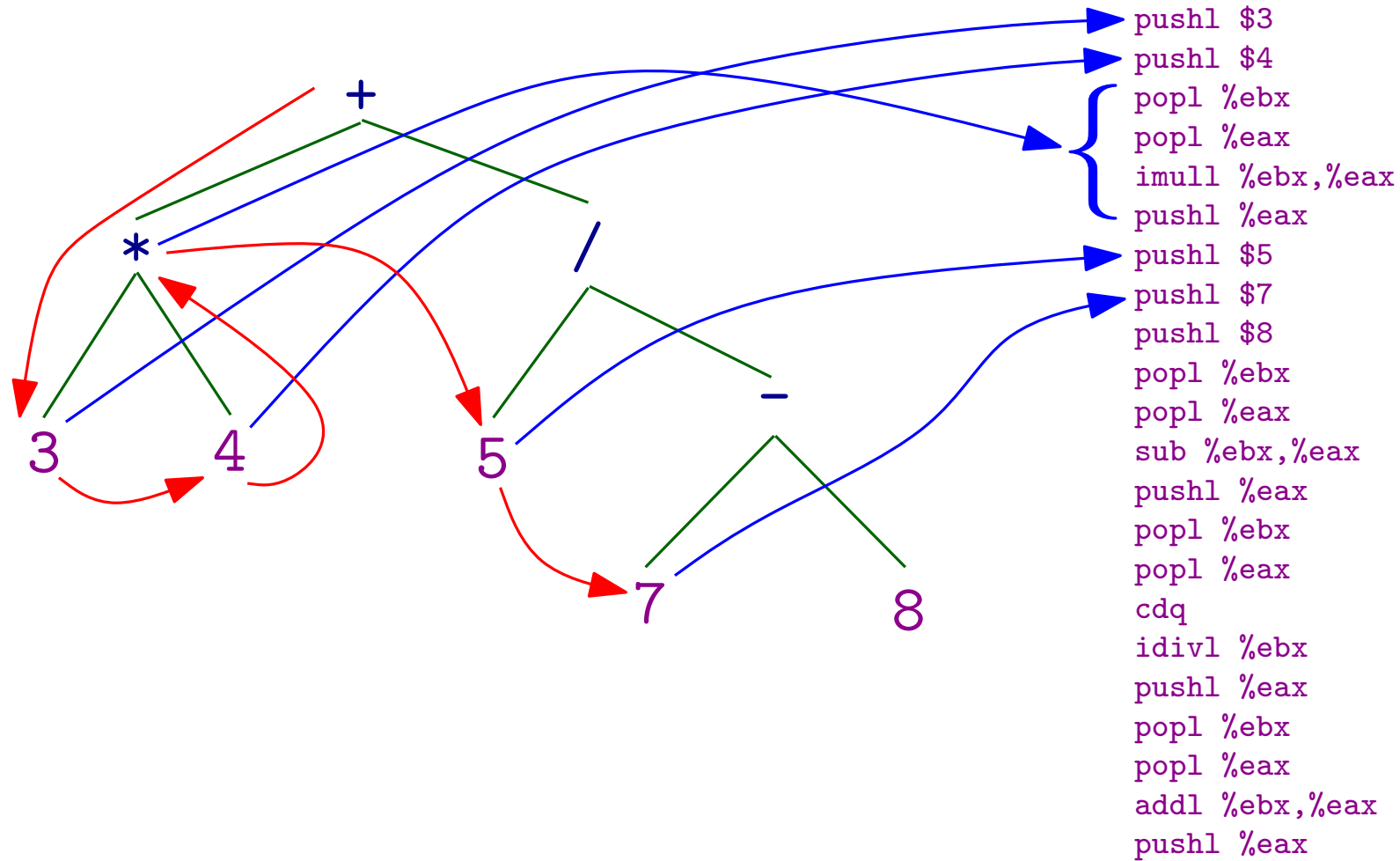
Compilation by Post-Order Traversal



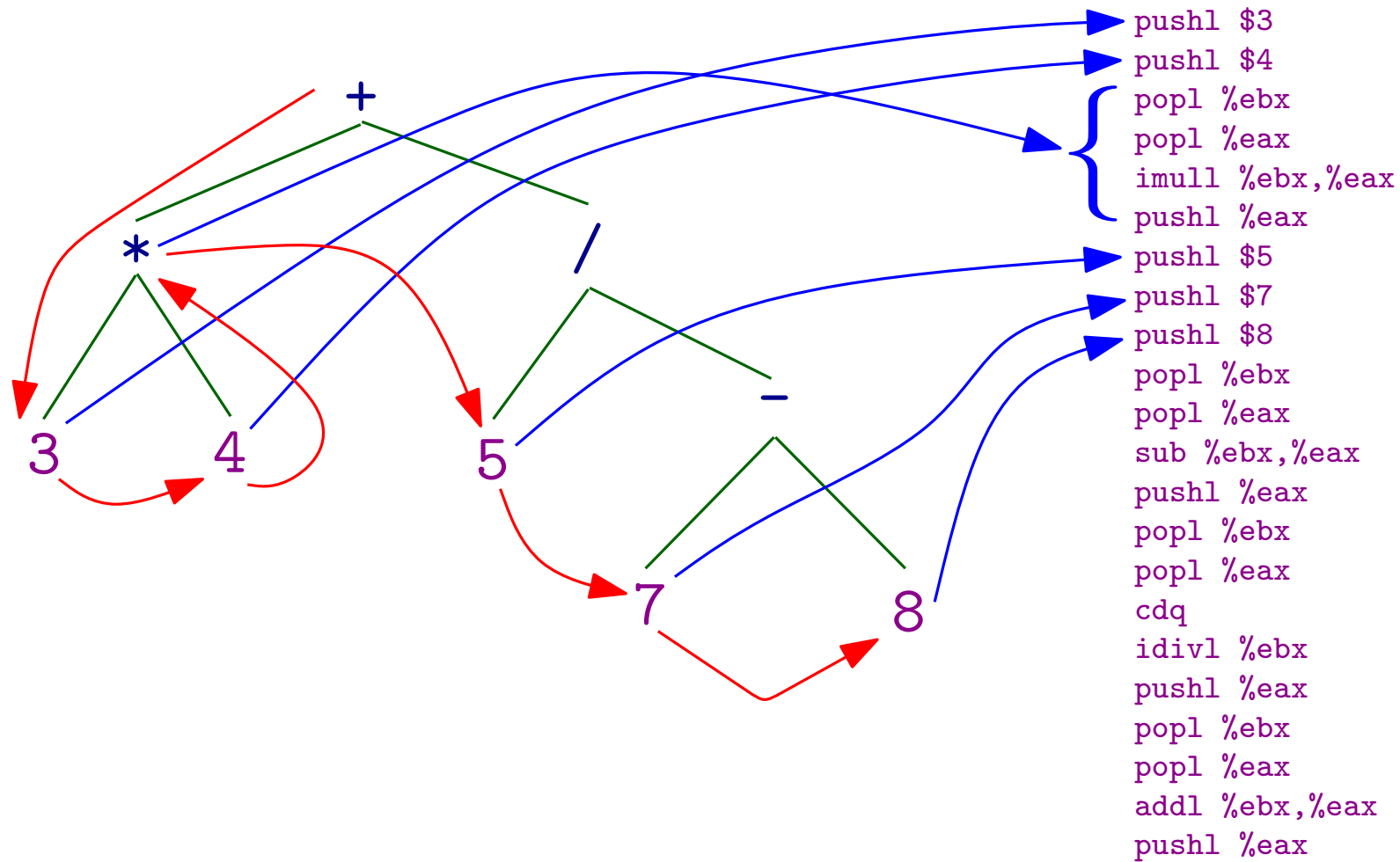
Compilation by Post-Order Traversal



Compilation by Post-Order Traversal

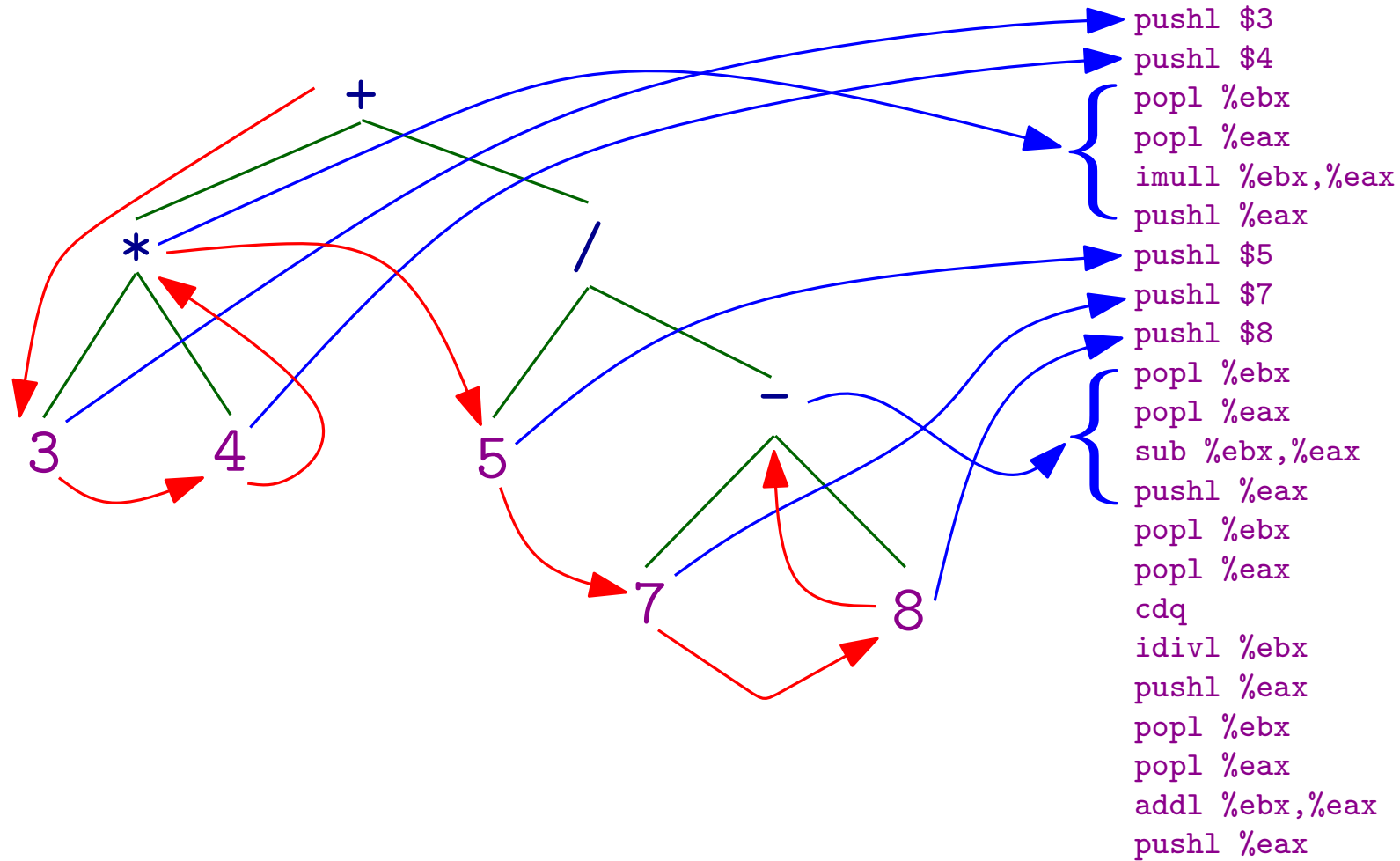


Compilation by Post-Order Traversal

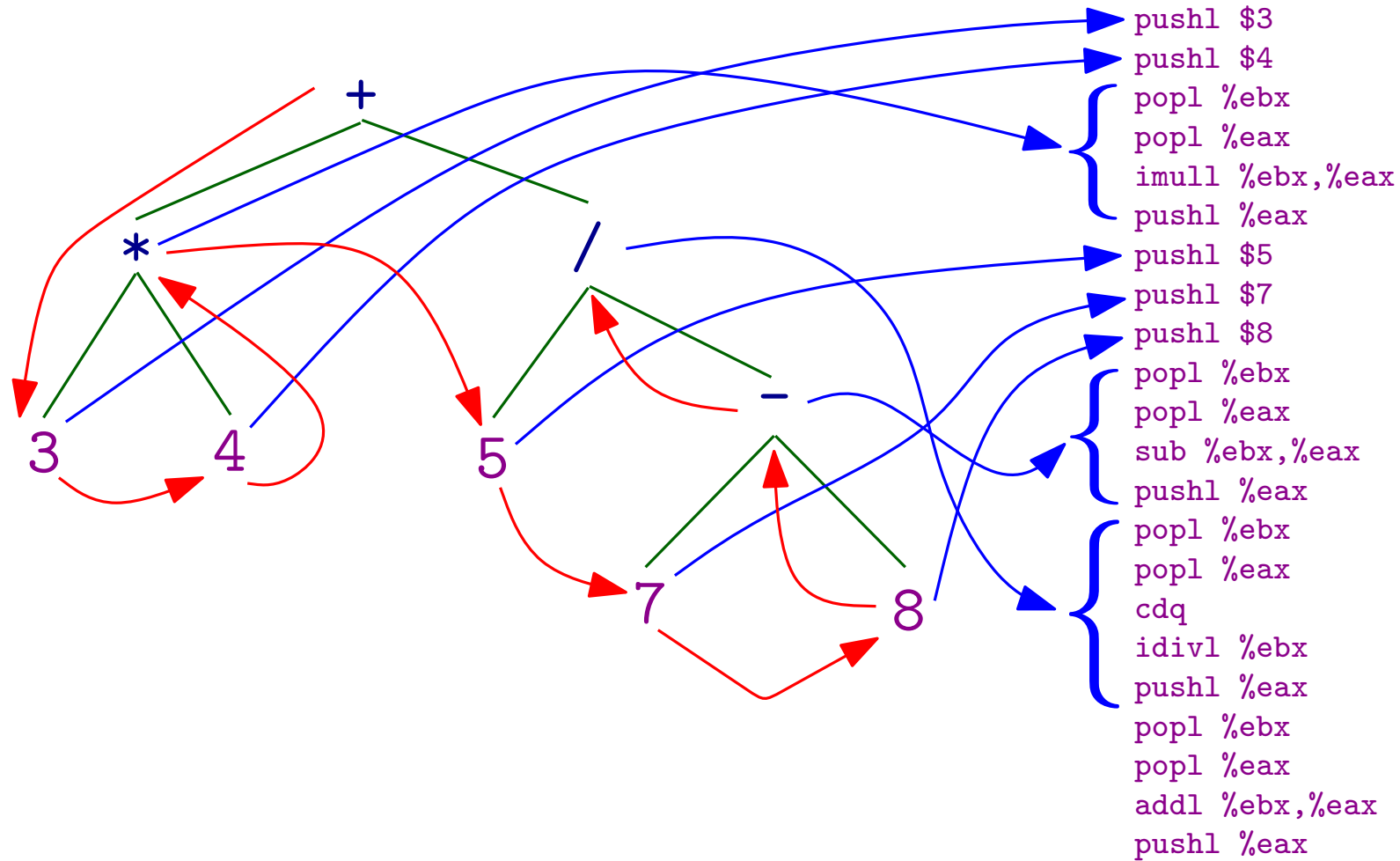


Compilation by Post-Order Traversal

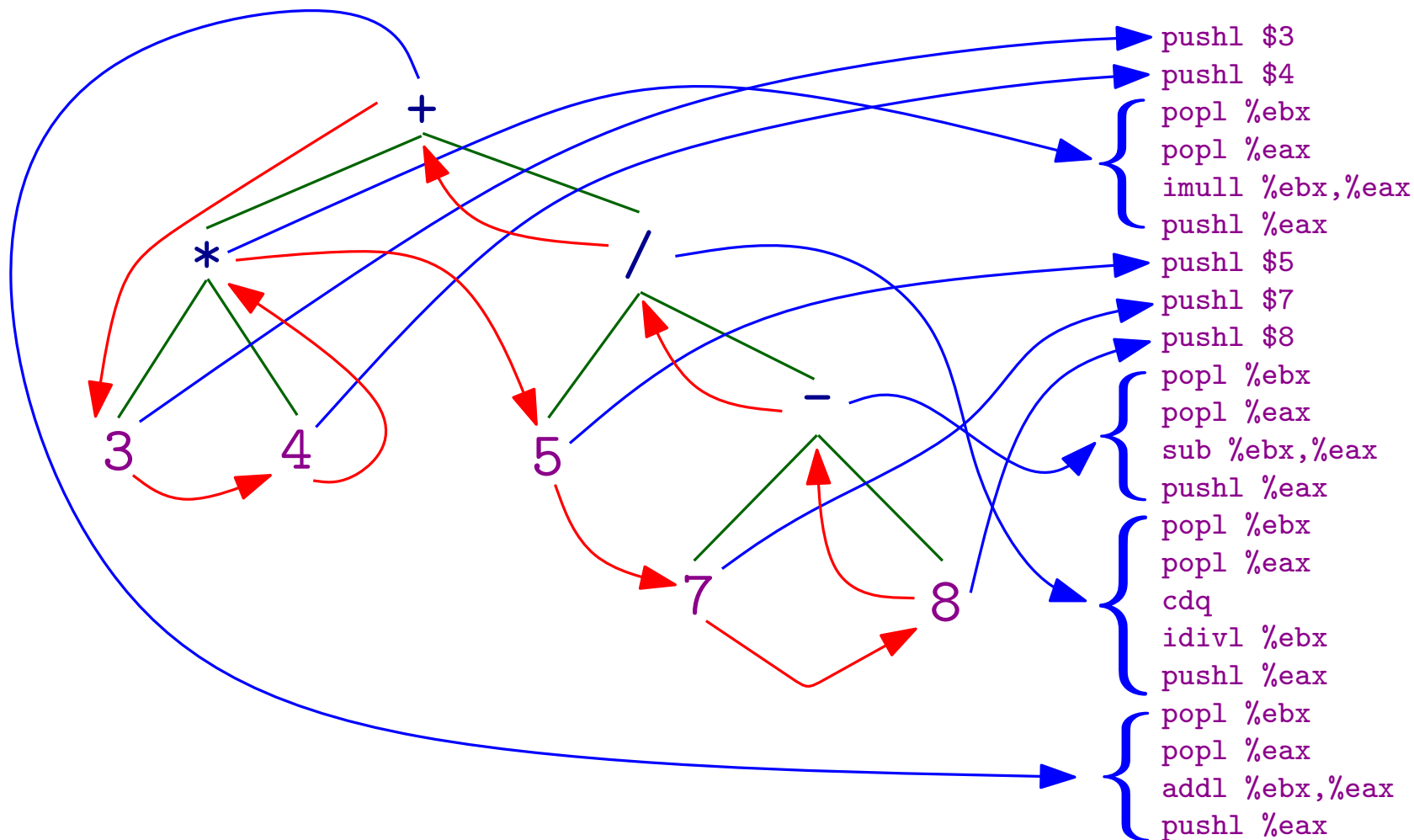
[C]



Compilation by Post-Order Traversal



Compilation by Post-Order Traversal



- For each constant, we push its value on the stack.
- Each operator pulls 2 operands from the stack, computes the result, and pushes it on the stack.
- Some instructions (IDIVL) have peculiar operands


```
ce(C,[Instr]) :-  
    integer(C),  
    !,  
    atomic_list_concat(['pushl $',C],Instr).  
  
ce(E,Code) :-  
    E =.. [Op,E1,E2],  
    member(Op,[+,-,*,/,rem]),  
    !,  
    ce(E1,C1),  
    ce(E2,C2),  
    cop(Op,Cop),  
    append([C1,C2,Cop],Code).
```

```
ce(C,[Instr]) :-  
    integer(C),  
    !,  
    atomic_list_concat(['pushl $',C],Instr).  
  
ce(E,Code) :-  
    E =.. [Op,E1,E2],  
    member(Op,[+,-,*,/,rem]),  
    !,  
    ce(E1,C1),  
    ce(E2,C2),  
    cop(Op,Cop),  
    append([C1,C2,Cop],Code).
```

Post-order traversal

Code Issued for Operators

```
cop(+,['popl %ebx', 'popl %eax', 'addl %ebx,%eax', 'pushl %eax']).  
cop(-,['popl %ebx', 'popl %eax', 'subl %ebx,%eax', 'pushl %eax']).  
cop(*,['popl %ebx', 'popl %eax', 'imull %ebx,%eax', 'pushl %eax']).  
cop(/,['popl %ebx', 'popl %eax', 'cdq', 'idiv %ebx', 'pushl %eax']).  
cop(rem,['popl %ebx', 'popl %eax', 'cdq', 'idiv %ebx', 'pushl %edx']).
```

Main Predicate

```
out(E) :-
    ce(E,Code),
    Pre = [ '.section .text',
             '.globl _start',
             '_start:',
             'pushl %ebp',
             'movl %esp,%ebp'],
    Post = ['popl %eax',
            'movl %ebp,%esp',
            'popl %ebp',
            'ret'],
    append([Pre,Code,Post],All),
    atomic_list_concat([''|All],'\n\t',AllWritable),
    write(AllWritable).
```

```
out(E) :-  
    ce(E,Code),  
    Pre = [ '.section .text',  
            '.globl _start',  
            '_start:',  
            'pushl %ebp',  
            'movl %esp,%ebp'],  
    Post = ['popl %eax',  
            'movl %ebp,%esp',  
            'popl %ebp',  
            'ret'],  
    append([Pre,Code,Post],All),  
    atomic_list_concat([''|All],'\n\t',AllWritable),  
    write(AllWritable).
```

Preamble

Postamble

All code as single atom

```
?- out(1+2).
```

```
.section .text
.globl _start
_start:
pushl %ebp
movl %esp,%ebp
pushl $1
pushl $2
popl %ebx
popl %eax
addl %ebx,%eax
pushl %eax
popl %eax
movl %ebp,%esp
popl %ebp
ret
```

```
true.
```

```
#include <stdio.h>
```

```
int start() asm("_start") ;
```

```
int main() {
    printf("Result = %d\n",start()) ;
}
```

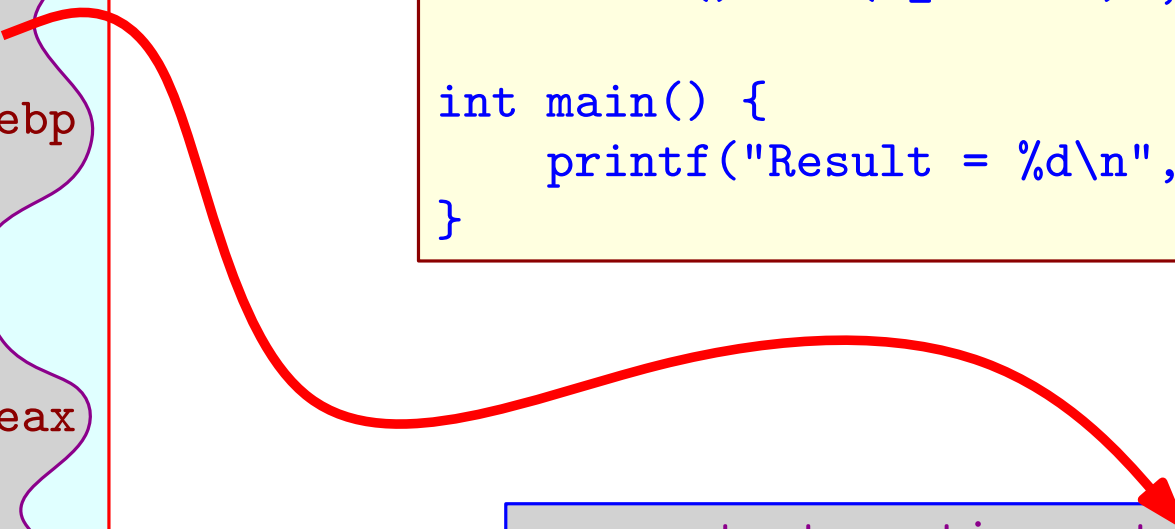
```
gcc -o test runtime.c test.s
```

```
./test
```

```
Result = 5
```

?- out(1+2).

```
.section .text
.globl _start
_start:
pushl %ebp
movl %esp,%ebp
pushl $1
pushl $2
popl %ebx
popl %eax
addl %ebx,%eax
pushl %eax
popl %eax
movl %ebp,%esp
popl %ebp
ret
```



true.

```
#include <stdio.h>
```

```
int start() asm("_start") ;
```

```
int main() {
```

```
    printf("Result = %d\n",start()) ;
```

```
}
```

```
gcc -o test runtime.c test.s
```

```
./test
```

```
Result = 5
```

```
?- out(2*3-4/5+6).
```

```
.section .text
.globl _start
_start:
pushl %ebp
movl %esp,%ebp
pushl $2
pushl $3
popl %ebx
popl %eax
imull %ebx,%eax
pushl %eax
pushl $4
pushl $5
popl %ebx
popl %eax
cdq
idiv %ebx
pushl %eax
```

```
popl %ebx
popl %eax
subl %ebx,%eax
pushl %eax
pushl $6
popl %ebx
popl %eax
addl %ebx,%eax
pushl %eax
popl %eax
movl %ebp,%esp
popl %ebp
ret
```

```
true.
```

Code in `comp_expr_naive_1.pro`


```
?- out(1<2).
```

```
                                .section .text
                                .globl _start
_start:

                                pushl %ebp
                                movl %esp,%ebp
                                pushl $1
                                pushl $2
                                popl %eax
                                popl %ebx
                                cmpl %eax,%ebx
                                jge L0
                                pushl $1
                                jmp L1

L0:

                                pushl $0

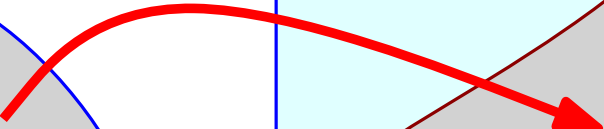
L1:

                                popl %eax
                                movl %ebp,%esp
                                popl %ebp
                                ret

true.
```

- Code generated for comparison operator
- In fact, *code template*
- Each instance of the code needs fresh labels

```
?- out(1<2).  
  
                .section .text  
                .globl _start  
_start:  
  
                pushl %ebp  
                movl %esp,%ebp  
                pushl $1  
                pushl $2  
                popl %eax  
                popl %ebx  
                cmpl %eax,%ebx  
                jge L0  
                pushl $1  
                jmp L1  
  
L0:  
                pushl $0  
  
L1:  
  
                popl %eax  
                movl %ebp,%esp  
                popl %ebp  
                ret  
  
true.
```



Operator Code

```
cop(+, [],  
    ['\n\t\t popl %ebx',  
     '\n\t\t popl %eax',  
     '\n\t\t addl %ebx,%eax',  
     '\n\t\t pushl %eax']).
```

```
cop(<, [L1,L2],  
    ['\n\t\t popl %eax',  
     '\n\t\t popl %ebx',  
     '\n\t\t cmpl %eax,%ebx',  
     '\n\t\t jge ', L1,  
     '\n\t\t pushl $1',  
     '\n\t\t jmp ', L2, '\n',  
L1, ':',  
     '\n\t\t pushl $0', '\n',  
L2, ':',  
    ]).
```

Operator Code

```
cop(+, [],
    ['\n\t\t popl %ebx',
     '\n\t\t popl %eax',
     '\n\t\t addl %ebx,%eax',
     '\n\t\t pushl %eax']).
```

Older operators generate the same code

```
cop(<, [L1,L2],
    ['\n\t\t popl %eax',
     '\n\t\t popl %ebx',
     '\n\t\t cmpl %eax,%ebx',
     '\n\t\t jge ', L1,
     '\n\t\t pushl $1',
     '\n\t\t jmp ', L2, '\n',
    L1, ':',
     '\n\t\t pushl $0', '\n',
    L2, ':',
    ]).
```

Comparison generates a template, with variables L1 and L2 acting as placeholders for fresh variables

Post-Order Traversal

```
ce(C,[Instr],LabelSuffix,LabelSuffix) :-  
    integer(C),!,  
    atomic_list_concat(['\n\t\t pushl $',C],Instr).  
  
ce(E,Code,LabelSuffixIn,LabelSuffixOut) :-  
    E =.. [Op,E1,E2],  
    member(Op,[+,-,*,/,rem,<]),!,  
    cop(Op,LPlaceholders,Cop),  
    generateLabels(LPlaceholders,LabelSuffixIn,LabelSuffixAux1),  
    ce(E1,C1,LabelSuffixAux1,LabelSuffixAux2),  
    ce(E2,C2,LabelSuffixAux2,LabelSuffixOut),  
    append([C1,C2,Cop],Code).
```

Post-Order Traversal

New arguments help generate new labels whenever they are needed. Since assignment is not available, we always need input and output versions.

```
ce(C,[Instr],LabelSuffix,LabelSuffix) :-
    integer(C),!,
    atomic_list_concat(['\n\t\t pushl $',C],Instr).

ce(E,Code,LabelSuffixIn,LabelSuffixOut) :-
    E =.. [Op,E1,E2],
    member(Op,[+,-,*,/,rem,<]),!,
    cop(Op,LPlaceholders,Cop),
    generateLabels(LPlaceholders,LabelSuffixIn,LabelSuffixAux1),
    ce(E1,C1,LabelSuffixAux1,LabelSuffixAux2),
    ce(E2,C2,LabelSuffixAux2,LabelSuffixOut),
    append([C1,C2,Cop],Code).
```

Predicate that bounds placeholders to fresh variables

Label Generator

```
generateLabels([],LabelSuffix,LabelSuffix).  
  
generateLabels([H|T],LabelSuffixIn,LabelSuffixOut) :-  
    atomic_list_concat(['L',LabelSuffixIn],H),  
    LabelSuffixAux #= LabelSuffixIn + 1,  
    generateLabels(T,LabelSuffixAux,LabelSuffixOut).
```

Main Predicate for Version 2

```
out(E) :-
    ce(E,Code,0,_),
    Pre = [ '\n\t\t.section .text',
            '\n\t\t.globl _start',
            '\n_start:',
            '\n\t\tpushl %ebp',
            '\n\t\tmovl %esp,%ebp'],
    Post = ['\n\t\tpopl %eax',
            '\n\t\tmovl %ebp,%esp',
            '\n\t\tpopl %ebp',
            '\n\t\tret'],
    append([Pre,Code,Post],All),
    atomic_list_concat(All,AllWritable),
    write(AllWritable).
```


Demo

```
?- out((1<2)*((3<4)+2*4/5)).

        .section .text
        .globl _start
_start:

        pushl %ebp
        movl %esp,%ebp
        pushl $1
        pushl $2
        popl %eax
        popl %ebx
        cmpl %eax,%ebx
        jge L0
        pushl $1
        jmp L1

L0:
        pushl $0

L1:
        pushl $3
        pushl $4
        popl %eax
        popl %ebx
        cmpl %eax,%ebx
        jge L2
        pushl $1
        jmp L3

L2:
        pushl $0

L3:
```

```
        pushl $2
        pushl $4
        popl %ebx
        popl %eax
        imull %ebx,%eax
        pushl %eax
        pushl $5
        popl %ebx
        popl %eax
        cdq
        idiv %ebx
        pushl %eax
        popl %ebx
        popl %eax
        addl %ebx,%eax
        pushl %eax
        popl %ebx
        popl %eax
        imull %ebx,%eax
        pushl %eax
        popl %eax
        movl %ebp,%esp
        popl %ebp
        ret

true.
```

Code in comp_expr_naive_2.pro

Version 3

- We had to add 2 new arguments to the traversal predicate to implement label generation.
 - They are called *attributes*
- In general, as we add more features to the language, more attributes need to be added.
- From software engineering perspective, not a good idea to implement attributes as arguments to the predicate
- Better solution: add a dictionary (input+output versions) to list of arguments; put all attributes in the dictionary as (key,value) pairs

The New Post-Order Traversal Predicate

```
ce(C,[Instr],A,A) :-  
    integer(C),!,  
    atomic_list_concat(['\n\t\t pushl $',C],Instr).  
  
ce(E,Code,AIn,AOut) :-  
    E =.. [Op,E1,E2],  
    member(Op,[+,-,*,/,rem,<]),!,  
    cop(Op,LPlaceholders,Cop),  
    get_assoc(labelsuffix,AIn,LabelSuffixIn,Aaux1,LabelSuffixAux1),  
    generateLabels(LPlaceholders,LabelSuffixIn,LabelSuffixAux1),  
    ce(E1,C1,Aaux1,Aaux2),  
    ce(E2,C2,Aaux2,AOut),  
    append([C1,C2,Cop],Code).
```

The New Main Predicate

```
out(E) :-
    empty_assoc(Empty),
    put_assoc(labelsuffix,Empty,0,A),
    ce(E,Code,A,_),
    Pre = [ '\n\t\t.section .text',
            '\n\t\t.globl _start',
            '\n_start:',
            '\n\t\tpushl %ebp',
            '\n\t\tmovl %esp,%ebp'],
    Post = ['\n\t\tpopl %eax',
            '\n\t\tmovl %ebp,%esp',
            '\n\t\tpopl %ebp',
            '\n\t\tret'],
    append([Pre,Code,Post],All),
    atomic_list_concat(All,AllWritable),
    write(AllWritable).
```

Version 4: Demo

[C]

```
?- out(x=1;y=2;x+y,'test10.s').
true.
```

Code in `comp_expr_naive_3.pro`

```
                .section .text
                .globl _start
_start:
                pushl %ebp
                movl %esp,%ebp
                pushl $1
                popl %eax
                movl %eax,x
                pushl %eax
                popl %eax
                pushl $2
                popl %eax
                movl %eax,y
                pushl %eax
```

```
                popl %eax
                pushl x
                pushl y
                popl %ebx
                popl %eax
                addl %ebx,%eax
                pushl %eax
                popl %eax
                movl %ebp,%esp
                popl %ebp
                ret

                .comm x,4,4
                .comm y,4,4
```

Version 4: Demo

```
?- out(x=1;y=2;x+y,'test10.s').
true.
```

Code in `comp_expr_naive_3.pro`

```

.section .text
.globl _start

_start:
    pushl %ebp
    movl %esp,%ebp
    pushl $1
    = → popl %eax
        movl %eax,x
        pushl %eax
    ; → popl %eax
        pushl $2
    = → popl %eax
        movl %eax,y
        pushl %eax

```

List of variables
must be
collected so as
to reserve space
for them ⇒
another
attribute

```

    popl %eax
    pushl x
    pushl y
    popl %ebx
    popl %eax
    addl %ebx,%eax
    pushl %eax
    popl %eax
    movl %ebp,%esp
    popl %ebp
    ret

    .comm x,4,4
    .comm y,4,4

```

Post-Order Traversal

[C]

```
ce(C,[Instr],A,A) :-
    (    integer(C), P = '$' ; atom(C),P='' ),!,
    atomic_list_concat(['\n\t\t pushl ',P,C],Instr).
ce(E,Code,AIn,AOut) :-
    E =.. [Op,E1,E2],
    member(Op,[+,-,*,/,rem,<,=]),!,
    cop(Op,LPlaceholders,Cop),
    (    Op = (=)
    ->  atom(E1),
        get_assoc(vars,AIn,OldVars,Aaux,NewVars),
        union(OldVars,[E1],NewVars),
        ce(E2,C2,Aaux,AOut),
        LPlaceholders = [E1],
        append([C2,Cop],Code)
    ;   get_assoc(labelsuffix,AIn,LabelSuffixIn,Aaux1,LabelSuffixAux1),
        generateLabels(LPlaceholders,LabelSuffixIn,LabelSuffixAux1),
        ce(E1,C1,Aaux1,Aaux2),
        ce(E2,C2,Aaux2,AOut),
        append([C1,C2,Cop],Code) ).
ce((S1;S2),Code,Ain,Aout) :-
    ce(S1,C1,Ain,Aaux),
    ce(S2,C2,Aaux,Aout),
    append([C1,['\n\t\t popl %eax'],C2], Code).
```

The Operator Code

[C]

```
cop(=, [V],  
    ['\n\t\t popl %eax',  
     '\n\t\t movl %eax,', V,  
     '\n\t\t pushl %eax' ]).
```


The Main Predicate

```
out(E,File) :-
    tell(File),
    empty_assoc(Empty),
    put_assoc(labelsuffix,Empty,0,A1),
    put_assoc(vars,A1,[],A2),
    ce(E,Code,A2,A3),
    Pre = [ '\n\t\t.section .text',
            '\n\t\t.globl _start',
            '\n_start:',
            '\n\t\tpushl %ebp',
            '\n\t\tmovl %esp,%ebp'],
    Post = ['\n\t\tpopl %eax',
            '\n\t\tmovl %ebp,%esp',
            '\n\t\tpopl %ebp',
            '\n\t\tret'],
    append([Pre,Code,Post],All),
    atomic_list_concat(All,AllWritable),
    writeln(AllWritable),
    get_assoc(vars,A3,VarList),
    allocvars(VarList,VarCode),
    atomic_list_concat(VarCode,WritableVars),
    write(WritableVars),
    told.
```

Variable Allocator

[C]

```
allocvars([], []).  
allocvars([V|VT], [D|DT]) :-  
    atomic_list_concat(['\n\t\t .comm ', V, ', ', 4, 4'], D),  
    allocvars(VT, DT).
```

```
?- out(x=1;y=4/2+(0<1);x+2*y,'test0.s').  
true.
```

```
        .section .text  
        .globl _start  
_start:  
        pushl %ebp  
        movl %esp,%ebp  
        pushl $1  
        popl %eax  
        movl %eax,x  
        pushl %eax  
        popl %eax  
        pushl $4  
        pushl $2  
        popl %ebx  
        popl %eax  
        cdq  
        idiv %ebx  
        pushl %eax  
        pushl $0  
        pushl $1  
        popl %eax  
        popl %ebx  
        cmpl %eax,%ebx  
        jge L0  
        pushl $1  
        jmp L1  
  
L0:  
        pushl $0  
  
L1:
```

```
        popl %ebx  
        popl %eax  
        addl %ebx,%eax  
        pushl %eax  
        popl %eax  
        movl %eax,y  
        pushl %eax  
        popl %eax  
        pushl x  
        pushl $2  
        pushl y  
        popl %ebx  
        popl %eax  
        imull %ebx,%eax  
        pushl %eax  
        popl %ebx  
        popl %eax  
        addl %ebx,%eax  
        pushl %eax  
        popl %eax  
        movl %ebp,%esp  
        popl %ebp  
        ret
```

```
        .comm x,4,4  
        .comm y,4,4
```

- Syntax-based processing is achieved by post-order traversal of the AST
 - May require multiple traversals with more complicated language constructs
 - The state of the translation process is recorded in *attributes* (computed or inherited)
 - New features in the language usually require new attributes in the traversal predicate
- The generated code is very inefficient
 - The stack discipline is very simple, but under-utilizes the registers
 - We devise each code template to work independently of the siblings of the current node.
 - Instructions at node boundary may become redundant
 - Hard to optimize in the traversal process; optimization performed later, on the whole generated code.
 - We shall see better utilization of the registers in next recitation.