

CS4212 - Compiler Design

Exceptions and OOL

Outline

- Exceptions
 - Checked
 - Unchecked
 - `finally` clauses
- Object Oriented Languages
 - Polymorphism
 - Inheritance

General Facts about Exceptions

- Useful for error handling
- Two parts:
 - `try statement with catch/finally clauses`
 - `throw/raise statement: execution jumps to the catch clause that can handle the exception`
- Without function calls: similar to a labeled `break`
- With function calls: non-local returns

Exception Example

[C]

```
try {  
    ...  
    throw exc ;  
    ...  
} catch (Exc1 e) {  
    ...  
} catch (Exc2 e) {  
    ...  
} finally {  
    ...  
}
```

```
{ struct Exc { exc_type t ; ... } E ;  
  E.t = no_exception ;  
  ...  
  E.t = exc_type ;  
  // Fill E with more relevant data  
  goto catch_clauses ;  
  ...  
catch_clauses:  
  switch ( E.t ) {  
  case exc1_type:  
    // catch Exc1 code  
    goto finally_clause ;  
  
  case exc2_type:  
    // catch Exc2 code  
    goto finally_clause ; // may be optimized  
  
  case no_exception:  
  finally_clause:  
    // finally code  
    break ;  
  }  
}
```

Function Calls

- Java distinguishes between checked and unchecked (runtime)
- Checked Exceptions
 - return individually from each called function
 - return structure contains possibly thrown exception
 - the caller must check return structure
 - no exception: proceed as usual
 - exception raised: re-raise
 - fill up return structure and return immediately

Function Calls

- Unchecked exceptions
 - `try` statement saves processor "state", including the *current stack configuration*
 - Keep a global variable whose type is a structure to record the currently thrown exception
 - `throw` statement
 - fills up global variable with exception details
 - restores "state", performing a long return into the `try` statement
 - global variable helps select correct `catch` clause

Unchecked Exception - Refinement ^[C]

- `try` statements may be nested
- saved state encoded as element to be placed on *exception stack*
- each `try` pushes current state on exception stack
- `throw`: restores state at top of exception stack
- if restored `try` cannot handle current exception, `re-throw`
- before re-throwing, execute `finally`, if one exists

Example

```

int f() {
    ...
    try {
        ...
        g() ;
        ...
    } catch (E1 e) {
        ...
    }
    ...
    return ... ;
}

int g() {
    ...
    h() ;
    ...
    return ... ;
}

```

```

int h() {
    ...
    try {
        ...
        E1 e1 ;
        ...
        throw e1 ;
        ...
        E2 e2 ;
        ...
        throw e2 ;
        ...
    } catch (E2 e2) {
        ...
    } finally {
        ...
    }
    ...
    return ... ;
}

```


Checked Exc: Translation of f in C

```

int f() {
    ...
    try {
        ...
        g() ;
        ...
    } catch (E1 e) {
        ...
    }
    ...
    return R ;
}

```

```

struct E { exc_type T ; exc_val V ;
           int retVal ; } ;

struct E f() {
    ...
    { struct E gRet ;
      ...
      gRet = g() ;
      switch ( gRet.T ) {
        case E1 :
            ... // handle exception E1
            goto finally ;
        case NOEXCEPTION :
            finally :
                ... // finally block, if any, otherwise empty
                break ;
        default : // other exception, rethrow
            return gRet ;
      }
    }
    ...
    struct fRet ;
    fRet.T = NOEXCEPTION ;
    fRet.retVal = R ;
    return fRet ;
}

```

Checked Exc: Translation of g in C

```
int g() {  
    ...  
    h() ;  
    ...  
    return R ;  
}
```

```
struct E g() {  
    ...  
    struct E hRet ;  
    hRet = h() ;  
    if ( hRet.T != NOEXCEPTION )  
        return hRet ;  
    ...  
    struct E gRet ;  
    gRet.T = NOEXCEPTION ;  
    gRet.RetVal = R ;  
    return gRet ;  
}
```

Checked Exc: Translation of h in C ^[C]

```
int h() {  
    ...  
    try {  
        ...  
        E1 e1 ;  
        ...  
        throw e1 ;  
        ...  
        E2 e2 ;  
        ...  
        throw e2 ;  
        ...  
    } catch (E2 e2) {  
        ...  
    } finally {  
        ...  
    }  
    ...  
    return R ;  
}  
  
struct E h() {  
    ...  
    { struct E hRet ;  
        ...  
        hRet.T = E1 ; hRet.V = ... ; goto catch ;  
        ...  
        hRet.T = E2 ; hRet.V = ... ; goto catch ;  
        ...  
    catch:  
        switch ( hRet.T ) {  
            case E2 :  
                ... // handle E2  
                hRet.T = NOEXCEPTION ;  
                goto finally ;  
            default:  
            finally:  
                ...  
                if ( hRet.T != NOEXCEPTION ) return hRet ;  
                break ;  
        }  
        ...  
        hRet.T = NOEXCEPTION ;  
        hRet.RetVal = R ;  
        return hRet ;  
    }  
}
```

setjmp/longjmp in C

- `int setjmp(jmp_buf env)`
 - Sets up the local `jmp_buf` buffer and initializes it for the jump.
 - Saves the program's calling environment in the environment buffer `env`.
 - Direct invocation: `setjmp` returns 0.
 - Return from call to `longjmp`: returns nonzero.
- `void longjmp(jmp_buf env, int value)`
 - Restores context of environment buffer `env`.
 - The value specified by `value` is passed from `longjmp` to `setjmp`.
 - Program execution continues as if the corresponding invocation of `setjmp` had just returned.
 - `value != 0 -> setjmp` returns 1 ; otherwise returns `value`.

Example

```
#include <stdio.h>
#include <setjmp.h>

static jmp_buf buf;

void second(void) {
    printf("second\n");           // prints
    longjmp(buf,1);              // jumps back to where setjmp was called -
    making setjmp now return 1
}

void first(void) {
    second();                    // does not print
    printf("first\n");
}

int main() {
    if ( ! setjmp(buf) ) {
        first();                // when executed, setjmp returns 0
    } else {                    // when longjmp jumps back, setjmp returns 1
        printf("main\n");       // prints
    }

    return 0;
}
```

Unchecked Exceptions: f

```
int f() {
    ...
    try {
        ...
        g() ;
        ...
    } catch (E1 e) {
        ...
    }
    ...
    return R ;
}
```

g remains the same!

```
extern struct E exception ;
extern jmp_buf push(), pop() ;

int f() {
    ...
    if ( setjmp(push()) ) { // try
        ...
        g() ;
        ...
        pop();
    } else {
        switch ( exception.T ) {
            case E1 :
                ...
                exception.T = NOEXCEPTION ;
                goto finally ;
            default:
            finally:
                ... // may be empty
                if ( exception.T != NOEXCEPTION )
                    longjmp(pop()) ;
        }
    }
    ...
    return R ;
}
```

Unchecked Exceptions: h

[A]

```
int h() {  
    ...  
    try {  
        ...  
        E1 e1 ;  
        ...  
        throw e1 ;  
        ...  
        E2 e2 ;  
        ...  
        throw e2 ;  
        ...  
    } catch (E2 e2) {  
        ...  
    } finally {  
        ...  
    }  
    ...  
    return R ;  
}
```

```
int h() {  
    ...  
    if (setjmp(push)) {  
        ...  
        exception.T = E1 ;  
        exception.V = e1 ;  
        longjmp(pop()) ;  
        ...  
        exception.T = E2 ;  
        exception.V = e2 ;  
        longjmp(pop()) ;  
        ...  
        pop();  
    } else {  
        switch ( exception.T ) {  
            case E2 : // handle E2  
                ...  
                exception.T = NOEXCEPTION ;  
                goto finally ;  
            default:  
                finally:  
                ...  
                if ( exception.T != NOEXCEPTION )  
                    longjmp(pop()) ;  
        }  
    }  
    ...  
    return R ;  
}
```

Discussion

- **Checked exceptions**
 - easier to optimize, more compositional
 - more amenable to debugging
 - less flexible: require intermediate functions (like *g*) to be aware that an exception may pass through
- **Unchecked exceptions**
 - more flexible: intermediate functions need not be changed
 - less amenable to debugging
 - more difficult to optimize

Object Oriented Languages

- **Encapsulation and Data Hiding**
 - static checking of access attributes
 - name mangling
- **Polymorphism**
 - type checking
 - name mangling
- **Inheritance**
 - tables of pointers to functions

Encapsulation and Data Hiding

- Each class translated into a structure
 - data fields are copied into the structure
 - for each method reserve a pointer to that method
- Access modifiers become part of type system
- Data hiding enforced through static checking

Scoping

- Global scope: heap
- Class scope: heap
- Object scope: structure
- Method scope: stack
- Local scope: stack
- Care must be taken in stacking up frames in correct order

Data Access and Polymorphism

- class scopes, return and argument types are used to mangle the name of a static data member, or of a method.
- **Ex:** `int class1::f(char)` translated into `int class1_f_ic(char)`
- Global namespace contains all static references
 - Name mangling ensures that each reference is unique
 - No name clashes should be possible

An OO Example

```
class Speaker {  
    void say(String msg) { System.out.println(msg) ; }  
}  
  
class Lecturer extends Speaker {  
    void lecture(String msg) {  
        say(msg) ;  
        say("You should be taking notes") ;  
    }  
}  
  
class ArrogantLecturer extends Lecturer {  
    void say(String msg) { super.say("It is obvious that" + msg) ; }  
}  
  
...  
  
(new ArrogantLecturer).lecture("The sky is blue") ;
```

Speaker: C Equivalent

```
struct speaker {  
    void (*say)(struct speaker * self, char* msg) ;  
} ;  
  
void speaker_say(struct speaker * self, char* msg) {  
    printf("%s\n",msg) ;  
}  
  
void init_speaker(struct speaker *p) {  
    p->say = speaker_say ;  
}  
  
struct speaker * make_speaker () {  
    struct speaker * retVal = malloc(sizeof(struct speaker));  
    init_speaker(retVal) ;  
    return retVal ;  
}
```

Lecturer: C Equivalent

```
struct lecturer {  
    void (*say)(struct speaker * self, char* msg) ;  
    void (*lecture)(struct lecturer * self, char* msg) ;  
} ;  
  
void lecturer_lecture(struct lecturer * self, char * msg){  
    self->say(self,msg) ;  
    self->say(self,"You should be taking notes") ;  
}  
  
void init_lecturer(struct lecturer *p) {  
    init_speaker(p) ;  
    p->lecture = lecturer_lecture ;  
}  
  
struct lecturer * make_lecturer() {  
    struct lecturer * retVal = malloc(sizeof(struct lecturer));  
    init_lecturer(retVal) ;  
    return retVal ;  
}
```

Arrogant Lecturer: C Equivalent

```
struct alecturer {
    void (*say)(struct speaker * self, char* msg) ;
    void (*lecture)(struct lecturer * self, char* msg) ;
    void (*super_say)(struct speaker * self, char* msg) ;
};

void alecturer_say(struct alecturer * self, char * msg){
    char * p = malloc(200) ;
    *p = '\0' ;
    strcat(p,"It is obvious that " ) ;
    strcat(p,msg) ;
    self->super_say(self,p) ;
}

void init_alecturer(struct alecturer *p) {
    init_lecturer(p) ;
    p->super_say = p->say ;
    p->say = alecturer_say ;
}

struct alecturer * make_alecturer() {
    struct alecturer * retVal = malloc(sizeof(struct alecturer));
    init_alecturer(retVal) ;
    return retVal ;
}
```


Main function

```
int main() {  
    struct alecturer * p = make_alecturer() ;  
    p->lecture(p, "The sky is blue") ;  
}
```

Discussion

- Each class
 - compiled into a structure
 - data members remain the same
 - methods
 - pointer to function member
 - translated into function
- Each method
 - compiled into a function with a mangled name
 - first argument is pointer to self
 - always called via the corresponding pointer

Constructors

- Split into two parts
 - allocation
 - initialization
 - call initializer of extended class
- Return a pointer to the structure representing the object
 - method calls have the pattern
$$p \rightarrow m(p, \dots) ;$$

Class Extension

- Corresponding structure has same prefix as extended class
 - allows polymorphism
- overloaded methods of superclass are saved into separate fields of the structure
 - they can still be called if needed

Optimizations

- Pointers to functions
 - class wide table instead of object membership
- Detect methods that are never overloaded and compile direct calls for them