

## CS3230 : Tutorial - 9

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Please drop your answer sheets in Bakh's office or Rahul's CQT office by 1 pm Tuesday, 23rd October, 2012.

1. We are given a sequence of  $n$  non-negative real numbers  $p_1, p_2, \dots, p_n$ . We are supposed to determine  $\max\{p(j) - p(i) \mid 1 \leq i < j \leq n\}$ . Write a dynamic-programming algorithm (idea and pseudocode) running in linear time for this task.
2. Write an algorithm which takes as input array  $d$  of  $n$  denominations with  $d[1] > d[2] > \dots > d[n] = 1$  and an amount  $A$ . It outputs array  $C$  such that  $C[j]$  (for  $1 \leq j \leq A$ ) is the minimum number of coins used for amount  $j$  (using denominations in  $d$ ). The algorithm must use storage only  $O(n)$  (that is the storage used should not depend on  $A$ ). Write the idea of the algorithm and then the pseudocode. What is the running time of your algorithm?

3. Consider the following algorithm:

```
power2(k) {  
  result = 2  
  for i = 1 to k  
    result = result * result  
  return result  
}
```

Is this a polynomial time algorithm if  $k$  is given in unary ? Is this a polynomial time algorithm if  $k$  is given in binary ?

4. Show that if a decision problem  $A$  is in  $\mathcal{P}$  (the class of decision problems solvable in polynomial time), then for every other decision problem  $B$ , we have  $A \leq_P B$  using a Karp reduction (a.k.a polynomial time many-one reduction), unless  $B$  or  $\bar{B}$  is empty.
5. Take it granted that there exists an (infinite) listing of all polynomial time algorithms. Let  $P_k$  be the  $k$ -th algorithm in this listing. Take it granted that there exists a universal program  $U$  such that  $U(x, k) = P_k(x)$  (that is the output of  $U$  on input  $(x, k)$  is the same as the output of  $P$  on input  $x$ ) for every binary string  $x$  and number natural number  $k$ . The running time of  $U$  on input  $(x, k)$  is polynomial in  $|x| + k$ . Consider the following algorithm:

```
diagp(k) {  
  if (U(k, k) == yes)
```

```
    return no
else
    return yes
}
```

Show that *diagp* is an exponential time algorithm and the language it decides is different from any language in  $P$ . From this conclude that  $\mathcal{P} \neq \mathcal{EXP}$ , where  $\mathcal{EXP}$  is the class of decision problems that can be decided in exponential time. Above whenever number  $k$  is given as input to any program it is given using binary encoding.