

# EE3204/EE3204E Computer Communication Networks I (Part 1)

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## Performance of CSMA/CD using a simplified model

Let  $T_f$  be the frame transmission time.

Let  $T_p$  be the end-to-end one way propagation delay.

Let  $\tau$  be the duration of a (contention) slot which is given by  $2 T_p$ .

Assume that time on medium is organized into cycles each of which consists of a collision/idle period ( $T_c$ ) followed by a transmission period ( $T_f$ ). A collision/idle period consists of a sequence of slots during which there are either collisions or no transmissions. Frame transmission occurs during the transmission period.

Therefore, the utilization (or effective throughput or efficiency)  $U$  is given by  $U = \frac{T_f}{T_c + T_f}$

Let there be  $N$  nodes each of which attempts to transmit during a slot with probability  $p$ .

Let  $A$  be the probability that exactly one node transmits during a slot.

Let  $n_c$  be the mean number of slots in a collision/idle period.

Let  $Pr[i]$  be the probability that the collision/idle period consists of  $i$  slots.

$$A = Np(1-p)^{N-1}$$

The mean number of slots in collision/idle period,  $n_c$ , is calculated as follows.

$$\begin{aligned} n_c &= \sum_{i=0}^{\infty} i Pr[i] \\ &= \sum_{i=1}^{\infty} i(1-A)^i A \end{aligned}$$

$$n_c = \frac{1}{A} - 1$$

The utilization  $U$  is given by

$$\begin{aligned}
 \text{Utilization } U &= \frac{T_f}{T_c + T_f} \\
 &= \frac{T_f}{n_c \tau + T_f} \\
 &= \frac{T_f}{2n_c T_p + T_f} \\
 &= \frac{1}{1 + 2an_c} \\
 &= \frac{1}{1 + 2a\left(\frac{1}{A} - 1\right)}
 \end{aligned}$$

### Calculation of Maximum Utilization

$$A = Np(1 - p)^{N-1}$$

Maximum utilization is achieved when  $A$  is maximized.

$A_{\max}$  can be determined by equating  $dA/dp$  to 0.

$$dA/dp = N[-p(N-1)(1-p)^{N-2} + (1-p)^{N-1}] = 0$$

$$\Rightarrow p(N-1) = 1-p$$

$$\Rightarrow pN = 1$$

$$\Rightarrow p = \frac{1}{N} \quad (\text{it can be shown that for this value of } p, d^2A/dp^2 < 0)$$

Therefore  $A$  is maximized when  $p = \frac{1}{N}$  and  $A_{\max}$  is given by

$$A_{\max} = \left(1 - \frac{1}{N}\right)^{N-1}$$

The maximum utilization  $U_{\max}$  is given by

$$U_{\max} = \frac{1}{1 + 2a\left(\frac{1}{A_{\max}} - 1\right)}$$

When  $N \rightarrow \infty$  (i.e. large number of active nodes)

$\lim_{N \rightarrow \infty} A_{\max} = \lim_{N \rightarrow \infty} \left(1 - \frac{1}{N}\right)^{N-1} = \frac{1}{e}$ , where  $e = 2.72$  and the maximum utilization is given by

$$\begin{aligned} U_{\max} &= \frac{1}{1 + 2a \left( \frac{1}{A_{\max}} - 1 \right)} \\ &= \frac{1}{1 + 2a(e - 1)} \\ &= \frac{1}{1 + 3.44a} \end{aligned}$$