

Counting inversions problem

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Counting inversions

Let a_1, a_2, \dots, a_n be
a list of integers.

Two indices i and j
form an inversion if

$$(1) \quad i < j$$

$$(2) \quad a_i > a_j$$

For instance, 2, 4, 1, 3, 5
has three inversions.

The problem

Input: A list A of integers

$$A = a_1, a_2, \dots, a_n.$$

Output:

The number of inversions of the list A .

There is a "brute force" algorithm that solves the problem:

Initialize $\text{Count} = 0$.

For each (i, j) if i and j form an inversion increment Count .

There is a better way to solve the problem.

Idea:

Embed the Mergesort algorithm into our solution.

How can this be done?

On input $A = a_1, a_2, \dots, a_n$

- (1) Divide A into two equal sized lists X, Y .
- (2) Count inversions in X And Y .
- (3) Sort X AND Y
- (4) Count inversions a_i and a_j such that $a_i \in X$ And $a_j \in Y$.
- (5) Output the sum of the inversions

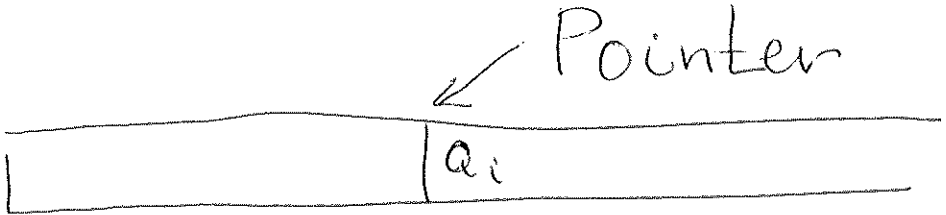
Merge-and-Count(A,B) algorithm.

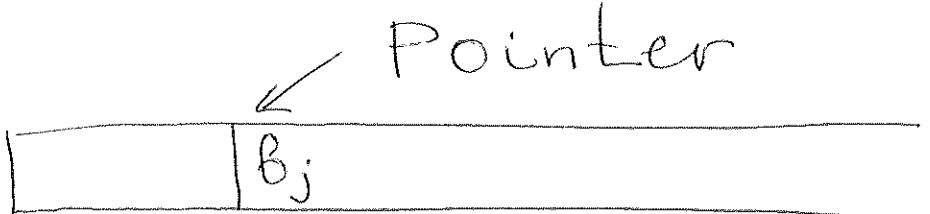
Ingredients:

A, B sorted lists,

Counter

Current pointer

A: A horizontal array with a vertical line near the right end. The element to the right of the line is labeled a_i . An arrow labeled "Pointer" points to the vertical line.

B: A horizontal array with a vertical line near the right end. The element to the right of the line is labeled b_j . An arrow labeled "Pointer" points to the vertical line.

While both A, B nonempty

Append the smaller of
 a_i, b_j to new list C .

If $b_j < a_i$

$Count \leftarrow Count + \text{the}$
remaining items in A .

Advance the pointer in
the list from which
the smaller element was
selected.

Sort-and-Count (L):

If $|L| \leq 1$, then there are no inversions.

Divide L into two equal halves:
 A and B .

$(r_A, A) \leftarrow \text{Sort-and-count}(A)$

$(r_B, B) \leftarrow \text{Sort-and-count}(B)$

$(r, L) \leftarrow \text{Merge-and-Count}(A, B)$

Output $r_A + r_B + r$.

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Correctness is proved by induction on $|L|$.