Internetworking: Packet Switching & Performance Metrics

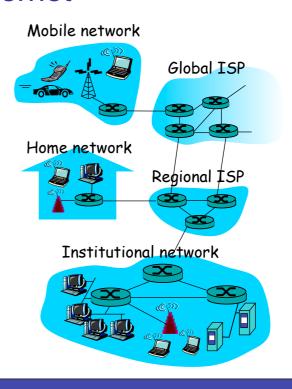
EE3204: Computer Communication Networks I Mehul Motani motani@nus.edu.sg

EE3204 Lecture Notes

Computer Communication Networks I

A look inside the Internet

- millions of connected computing devices running network apps
- communication linksfiber, copper, radio,satellite
- routers: forward packets (chunks of data)

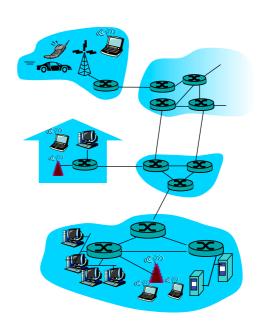


© Mehul Motani

Internetworking: Packet Switching & Performance Metrics 2

A closer look at network structure

- network edge: applications and hosts
- access networks, physical media: wired, wireless communication links
- network core:
 - interconnected routers
 - network of networks



© Mehul Motani

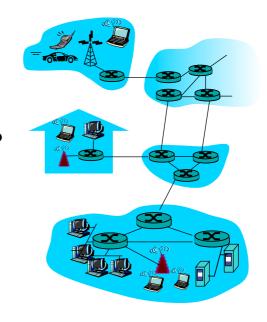
Internetworking: Packet Switching & Performance Metrics 3

EE3204 Lecture Notes

Computer Communication Networks I

The Network Core

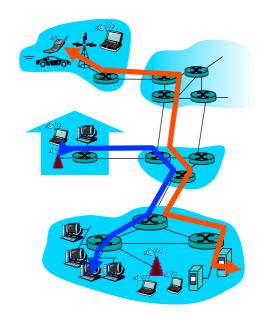
- mesh of interconnected routers
- <u>the</u> fundamental question: how is data transferred through net?
 - circuit switching: dedicated circuit per call: telephone net
 - packet-switching: data sent thru net in discrete "chunks"



Network Core: Circuit Switching

End-end resources reserved for "call"

- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required



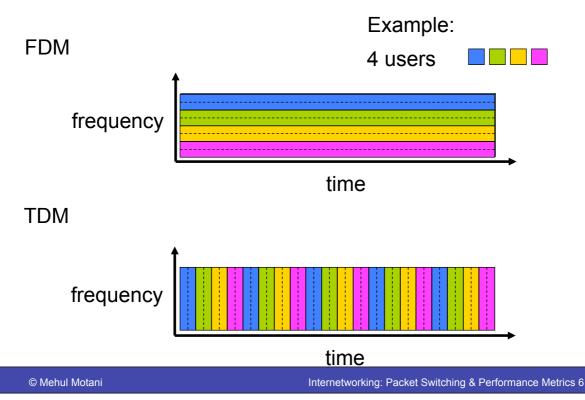
© Mehul Motani

Internetworking: Packet Switching & Performance Metrics 5

EE3204 Lecture Notes

Computer Communication Networks I

Circuit Switching: FDM and TDM



Network Core: Packet Switching

each end-end data stream divided into packets

- user A, B packets share network resources
- each packet uses full link bandwidth
- resources used as needed

resource contention:

- aggregate resource demand can exceed amount available
- congestion: packets queue, wait for link use
- store and forward: packets move one hop at a time
 - Node receives complete packet before forwarding

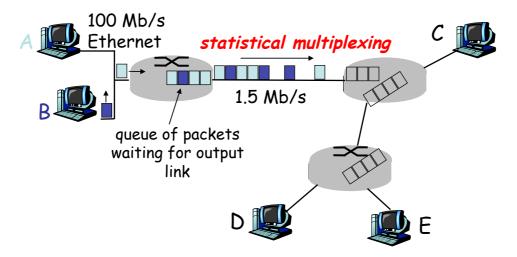
© Mehul Motani

Internetworking: Packet Switching & Performance Metrics 7

EE3204 Lecture Notes

Computer Communication Networks I

Packet Switching: Statistical Multiplexing



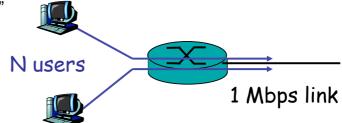
Sequence of A & B packets does not have fixed pattern, bandwidth shared on demand **⇒** *statistical multiplexing*. TDM: each host gets same slot in revolving TDM frame.

© Mehul Motani

Packet switching versus circuit switching

Packet switching allows more users to use network!

- > 1 Mb/s link
- each user:
 - > 100 kb/s when "active"
 - active 10% of time
- > circuit-switching:
 - > 10 users
- packet switching:
 - with 35 users, probability10 active at same timeis less than .0004



Q: how did we get value 0.0004?

© Mehul Motani

Internetworking: Packet Switching & Performance Metrics 9

EE3204 Lecture Notes

Computer Communication Networks I

Packet switching versus circuit switching

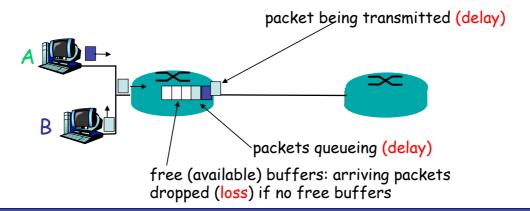
Is packet switching a "slam dunk winner?"

- great for bursty data
 - > resource sharing
 - > simpler, no call setup
- excessive congestion: packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
 - > bandwidth guarantees needed for audio/video apps
 - > still an unsolved problem

How do loss and delay occur?

Packets queue in router buffers in network core

- packet arrival rate to link exceeds output link capacity
- packets queue, wait for turn



© Mehul Motani

Internetworking: Packet Switching & Performance Metrics 11

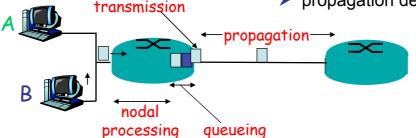
EE3204 Lecture Notes

Computer Communication Networks I

Delay in packet-switched networks

- 1. nodal processing:
- check bit errors
- determine output link
- 2. queueing
- > time waiting at output link for 4. Propagation delay: transmission
- depends on congestion level > s = propagation speed in of router

- 3. Transmission delay:
- R=link bandwidth (bps)
- L=packet length (bits)
- time to send bits into link = L/R
- d = length of physical link
- medium (~2x108 m/sec)
- propagation delay = d/s



Note: s and R are very different quantities!

© Mehul Motani

Internetworking: Packet Switching & Performance Metrics 12

Nodal delay

$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

- → d_{proc} = processing delay
 - > typically a few microseconds or less
- d_{queue} = queuing delay
 - > depends on congestion
- - > = L/R, significant for low-speed links
- d_{prop} = propagation delay
 - a few microsecs to hundreds of msecs

© Mehul Motani

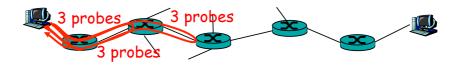
Internetworking: Packet Switching & Performance Metrics 13

EE3204 Lecture Notes

Computer Communication Networks I

"Real" Internet delays and routes

- What do "real" Internet delay & loss look like?
- Traceroute program: provides delay measurement from source to router along endend Internet path towards destination. For all i:
 - > sends three packets that will reach router *i* on path towards destination
 - router *i* will return packets to sender
 - > sender times interval between transmission and reply.



"Real" Internet delays and routes

traceroute: gaia.cs.umass.edu to www.eurecom.fr

```
Three delay measurements from gaia.cs.umass.edu to cs-gw.cs.umass.edu

1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms

2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms

3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms

4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms

5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms

6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms

7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms

7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms

8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms

9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms

10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms

11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms

12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms

13 nice.cssi.renater.fr (195.220.98.10) 126 ms 126 ms 124 ms

14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms

15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms

16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms

17 ***

** means no response (probe lost, router not replying)

19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms
```

© Mehul Motani

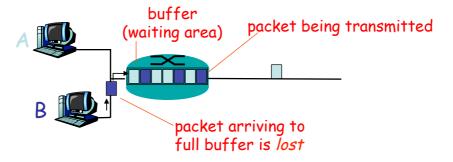
Internetworking: Packet Switching & Performance Metrics 15

EE3204 Lecture Notes

Computer Communication Networks I

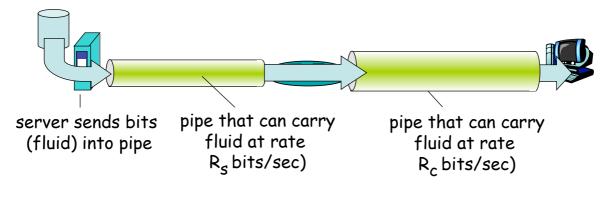
Packet loss

- queue (aka buffer) preceding link in buffer has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not at all



Throughput

- throughput: rate (bits/time) at which bits transferred between sender/receiver
 - *▶instantaneous:* rate at given point in time
 - >average: rate over long(er) period of time



© Mehul Motani

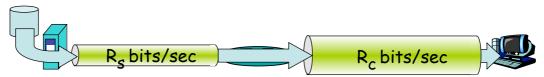
Internetworking: Packet Switching & Performance Metrics 17

EE3204 Lecture Notes

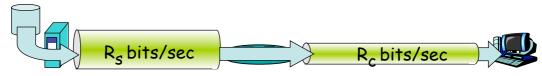
Computer Communication Networks I

Throughput (more)

 $> R_s < R_c$ What is average end-end throughput?



 $> R_s > R_c$ What is average end-end throughput?



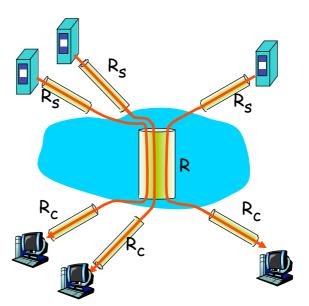
bottleneck link

link on end-end path that constrains end-end throughput

© Mehul Motani

Throughput: Internet scenario

- per-connection end-end throughput: min(R_c,R_s,R/10)
- in practice: R_c or R_s is often bottleneck



10 connections (fairly) share backbone bottleneck link R bits/sec

© Mehul Motani

Internetworking: Packet Switching & Performance Metrics 19

EE3204 Lecture Notes

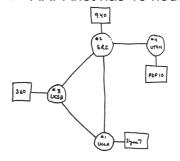
Computer Communication Networks I

Internet History

1961-1972: Early packet-switching principles

- 1961: Kleinrock queueing theory shows effectiveness of packet-switching
- 1964: Baran packetswitching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational

- **1972**:
 - ARPAnet public demonstration
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program
 - ARPAnet has 15 nodes



THE ARPA NETWORK

Internet History

1972-1980: Internetworking, proprietary networks

- > 1970: ALOHAnet satellite network in Hawaii
- 1974: Cerf and Kahn architecture for interconnecting networks
- > 1976: Ethernet at Xerox PARC
- ate70's: proprietary architectures: DECnet, SNA, XNA
- late 70's: switching fixed length packets (ATM precursor)
- 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet architecture

© Mehul Motani

Internetworking: Packet Switching & Performance Metrics 21

EE3204 Lecture Notes

Computer Communication Networks I

Internet History

1980-1990: new protocols and new networks

- ➤ 1983: deployment of TCP/IP
- 1982: smtp e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- > 1985: ftp protocol defined
- ➤ 1988: TCP congestion control

- new national networks: Csnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks

Internet History

1990, 2000's: Internet explosion, web apps

- Early 1990's: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - > HTML, HTTP: Berners-Lee
 - > 1994: Mosaic, later Netscape
 - late 1990's: commercialization of the Web

Late 1990's - 2000's:

- more killer apps: instant messaging, P2P file sharing
- network security to forefront
- est. 50 million host, 100 million+ users
- backbone links running at Gbps

© Mehul Motani

Internetworking: Packet Switching & Performance Metrics 23

EE3204 Lecture Notes

Computer Communication Networks I

Internet History

2007 & beyond: rebirth of the web

- Over 500 million hosts
- Voice, Video over IP
- P2P applications: BitTorrent (file sharing) Skype (VoIP), PPLive (video)
- Rich content & applications: Gaming, YouTube, Facebook, Twitter, Google
- Mobile wireless broadband
- > Web 2.0, Internet 2.0
- Clean slate Internet a complete redesign?