

Wigner's Semicircular Theorems

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1 Definitions

Assume we have two families of independent and identical real-valued random variables with zero-mean $\{Z_{i,j}\}_{1 \leq i < j}$ and $\{Y_i\}_{1 \leq i}$ such that $\mathbb{E}[Z_{1,2}^2] = 1 = \mathbb{E}[Y_1^2]$ for all integers $k \geq 1$. We also have,

$$r_k := \max \{ \mathbb{E}[|Z_{1,2}^k|], \mathbb{E}[|Y_1^k|] \} < \infty$$

Then we define Wigner matrices as the symmetric $N \times N$ matrices X_N as

$$X_N(i, j) = X_N(j, i) = \begin{cases} \frac{Z_{i,j}}{\sqrt{N}} & \text{if } i < j, \\ \frac{Y_i}{\sqrt{N}} & \text{if } i = j. \end{cases}$$

Let λ_i^N be the real ordered eigenvalues of X_N and we define the empirical distribution of the eigenvalues as the probability measure on \mathbb{R} defined as

$$L_N = \frac{1}{N} \sum_{i=1}^N \delta_{\lambda_i^N}$$

Denote the semicircle distribution as the probability distribution $\sigma(x)dx$ on \mathbb{R} with density

$$\sigma(x) = \frac{1}{2\pi} \sqrt{4 - x^2} \mathbf{1}_{\{|x| \leq 2\}}$$

2 Statement of Theorem

Wigner's Semicircle Theorem for Real Matrices:

For all Wigner matrices, the empirical measure L_N converges in probability to the semicircle distribution

i.e. for all $f \in C_b(\mathbb{R})$ and for all $\epsilon > 0$

$$\lim_{N \rightarrow \infty} \mathbb{P}(|\langle L_N, f \rangle - \langle \sigma, f \rangle| > \epsilon)$$

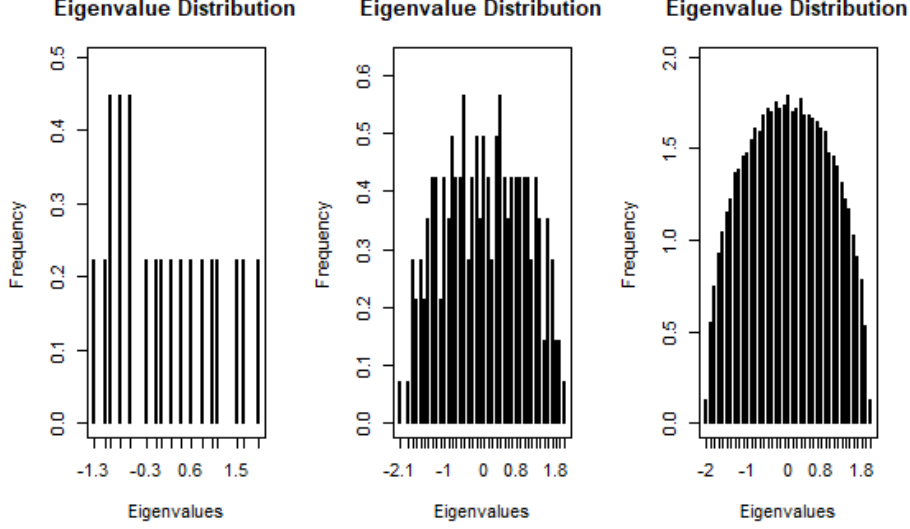


Figure 1: Three realisations of the eigenvalue distribution for $N = 20, 200, 3000$

3 Helpful Lemmas

Lemma 1:

For all $k \in \mathbb{N}$,

$$\lim_{N \rightarrow \infty} \langle \bar{L}_N, x^k \rangle = \langle \sigma, x^k \rangle$$

We denote $m_k := \langle \sigma, x^k \rangle$ and $m_k^N = \langle \bar{L}_N, x^k \rangle$

Proof of Lemma 1:

We start with basic matrix algebra

$$\begin{aligned} \langle \bar{L}_N, x^k \rangle &= \frac{1}{N} \mathbb{E}[\text{tr}(X_N^k)] \\ &= \frac{1}{N} \sum_{i_1, \dots, i_k=1}^N \mathbb{E}[X_N(i_1, i_2) X_N(i_2, i_3) \cdots X_N(i_{k-1}, i_k) X_N(i_k, i_1)] \\ &:= \frac{1}{N} \sum_{i_1, \dots, i_k=1}^N \mathbb{E}[T_{\mathbf{i}}^N] := \frac{1}{N} \sum_{i_1, \dots, i_k=1}^N \bar{T}_{\mathbf{i}}^N \quad \text{where } \mathbf{i} = (i_1, \dots, i_k) \end{aligned} \quad (3.1)$$

We now transfer to the language of combinatorics to complete this proof...

Given a set \mathcal{J} , a \mathcal{J} -letter is an element of \mathcal{J} , and a \mathcal{J} -word is a finite sequence of \mathcal{J} -letters, at least one letter long. A \mathcal{J} -word is closed if the first and last letters are the same, and two \mathcal{J} -words are equivalent, $w_1 \sim w_2$, if there exists a bijection

on \mathcal{J} mapping one to the other. If $\mathcal{J} = \{1, \dots, N\}$, we call it an N -word. The length of a word w , $l(w)$, is the number of letters in w , its weight, $\text{wt}(w)$, is the number of distinct letters in the word, and the support of w , $\text{supp}(w)$, is the set of letters in w .

We can then associate to each word $w = s_1, \dots, s_k$ a graph $G_w = (V_w, E_w)$ where vertices $V_w = \text{supp}(w)$, and edges $E_w = \{\{s_i, s_{i+1}\}; i = 1, \dots, k-1\}$. We call the set of self edges $E_w^s = \{e \in E_w : e = \{u, u\}, u \in V_w\}$, and the set of connecting edges $E_w^c = E_w \setminus E_w^s$. Note that G_w is connected as w defines a path connecting all the vertices, and defines a cycle if the word is closed. Denote N_e^w for $e \in E_w$ the number of times the path w traverses edge e .

Now note that any k -tuple of integers \mathbf{i} defines a closed word $w_{\mathbf{i}} = i_1 i_2 \dots i_k i_1$ with length $k+1$. Then,

$$\bar{T}_{\mathbf{i}}^N = \frac{1}{N^{\frac{k}{2}}} \prod_{e \in E_{w_{\mathbf{i}}}^c} \mathbb{E}[Z_{1,2}^{N_e^{w_{\mathbf{i}}}}] \prod_{e \in E_{w_{\mathbf{i}}}^s} \mathbb{E}[Y_1^{N_e^{w_{\mathbf{i}}}}] \quad (3.2)$$

Hence we see that $\bar{T}_{\mathbf{i}}^N = 0$ unless $N_e^{w_{\mathbf{i}}} \geq 2$ for all $e \in E_{w_{\mathbf{i}}}$. Furthermore, $\text{wt}(w_{\mathbf{i}}) \leq \frac{k}{2} + 1$

Note that if $N \geq t$, then there are $C_{N,t} = N(N-1)\dots(N-t+1)$ N -words equivalent to a given N -word with weight t .

We define $\mathcal{W}_{k,t}$ to be the set of representatives for equivalency classes of closed N -words w with length $k+1$, weight t , and with $N_e^w \geq 2$ for all $e \in E_w$. We then deduce from (3.1) and (3.2) that

$$\langle \bar{L}_N, x^k \rangle = \sum_{t=1}^{\lfloor \frac{k}{2} \rfloor + 1} \frac{C_{N,t}}{N^{\frac{k}{2}+1}} \sum_{w \in \mathcal{W}_{k,t}} \prod_{e \in E_w^c} \mathbb{E}[Z_{1,2}^{N_e^w}] \prod_{e \in E_w^s} \mathbb{E}[Y_1^{N_e^w}] \quad (3.3)$$

The cardinality of $\mathcal{W}_{k,t}$ is bounded by the number of closed \mathcal{J} -words of length $k+1$ when the cardinality of \mathcal{J} is $t \leq k$

$$|\mathcal{W}_{k,t}| \leq t^k \leq k^k$$

Thus (3.3) and the finiteness of r_k imply

$$\lim_{N \rightarrow \infty} \langle \bar{L}_N, x^k \rangle = \begin{cases} 0 & \text{if } k \text{ odd,} \\ \sum_{w \in \mathcal{W}_{k, \frac{k}{2}+1}} \prod_{e \in E_w^c} \mathbb{E}[Z_{1,2}^{N_e^w}] \prod_{e \in E_w^s} \mathbb{E}[Y_1^{N_e^w}] & \text{if } k \text{ even.} \end{cases}$$

since

$$\lim_{N \rightarrow \infty} \sum_{t=1}^{\lfloor \frac{k}{2} \rfloor + 1} \frac{C_{N,t}}{N^{\frac{k}{2}+1}} \approx \sum_{t=1}^{\lfloor \frac{k}{2} \rfloor + 1} N^{t - \frac{k}{2} - 1}$$

We define a closed word w of length $k + 1 \geq 1$ is a Wigner word if $k = 0$, or k even and w equivalent to an element of $\mathcal{W}_{k, \frac{k}{2}+1}$

Note that if $w \in \mathcal{W}_{k, \frac{k}{2}+1}$, then G_w is a tree, and also $E_w^s = \emptyset$. This implies that

$$\lim_{N \rightarrow \infty} \langle \bar{L}_N, x^k \rangle = |\mathcal{W}_{k, \frac{k}{2}+1}| \quad \text{for } k \text{ even}$$

Now let k be even and choose $\mathcal{W}_{k, \frac{k}{2}+1}$ such that all words $w = v_1, \dots, v_{k+1}$ satisfies for $i = 1, \dots, k + 1$, that $\{v_1, \dots, v_i\}$ is an interval in the integers, beginning with 1. Note each element $w \in \mathcal{W}_{k, \frac{k}{2}+1}$ determines a path $v_1, v_2, \dots, v_k, v_{k+1} = v_1$ length $k + 1$ on a tree G_w called the exploration process for w .

Let $d(v, v')$ denote the distance between v, v' on a tree G_w which is the length of the shortest path between them. Then we set $x_i = d(v_{i+1}, v_1)$ and see that each word $w \in \mathcal{W}_{k, \frac{k}{2}+1}$ defines a Dyck path $D(w) = (x_1, x_2, \dots, x_k)$ of length k .

Conversely, given a Dyck path $x = (x_1, \dots, x_k)$ can construct a word $w = T(x) \in \mathcal{W}_{k, \frac{k}{2}+1}$ by recursively constructing an increasing sequence $w_2, \dots, w_k = w$ of words as follows:

1) $w_2 = (1, 2)$

2) For $i > 2$, if $x_{i-1} = x_{i-2} + 1$, then w_i is obtained by adjoining on the right of w_{i-1} the smallest positive integer not in w_{i-1} . Otherwise w_i is obtained by adjoining on the right of w_{i-1} next to the last letter of w_{i-1}

Note that for all i , G_{w_i} is a tree since G_{w_2} is a tree and at stage i we either add a leaf, or a backtrack is added. Hence the distance in G_{w_i} between the first and last letter of w_i is x_{i-1} , and so $D(w) = (x_1, \dots, x_k)$.

With our choice of representation, $T(D(w)) = w$ as each 'up' in $D(w)$ starting at location $i - 2$ corresponds to the adjointment on the right of w_{i-1} of a new letter, uniquely determined by $\text{supp}(w_{i-1})$, but each 'down' at $i - 2$ corresponds to the joining of penultimate letter in w_{i-1} . Hence there is a bijection between Dyck paths of length k and $\mathcal{W}_{k, \frac{k}{2}+1}$.

Now we will calculate how many unique Dyck paths of length $2k$ exist for each k . First let B_k be the number of Bernoulli walks S_n length $2k$ such that $S_{2k} = 0$, and let \bar{B}_k be the number of Bernoulli walks S_n length $2k$ such that $S_{2k} = 0$ and $S_t < 0$ for some $t < 2k$. Hence the number of Dyck paths is $B_k - \bar{B}_k$.

Now $B_k = \binom{2k}{k}$ as we must have k 'up' steps and k 'down' steps. Then $\bar{B}_k = \binom{2k}{k-1}$ as must have $k + 1$ 'down' steps and $k - 1$ 'up' steps after the first time going negative. Hence the number of Dyck paths is $\binom{2k}{k} - \binom{2k}{k-1} = C_k$

where C_k is the k^{th} Catalan number.

We finish by showing that $m_{2k} = C_k$ for all integers k . First note that $m_{2k+1} = 0$ by symmetry over the integral. Now,

$$\begin{aligned}
m_{2k} &= \int_{\mathbb{R}} x^{2k} \sigma(x) dx \\
&= \frac{1}{2\pi} \int_{-2}^2 x^{2k} \sqrt{4-x^2} dx \\
&= \frac{1}{2\pi} \int_{-\frac{\pi}{2}}^{\frac{\pi}{2}} (2\sin(\theta))^{2k} \sqrt{4-4\sin^2(\theta)} 2\cos(\theta) d\theta \\
&= \frac{2^{k+1}}{\pi} \int_{-\frac{\pi}{2}}^{\frac{\pi}{2}} \sin^{2k}(\theta) \cos^2(\theta) d\theta \\
&= \frac{2^{k+1}}{\pi} \int_{-\frac{\pi}{2}}^{\frac{\pi}{2}} \sin^{2k}(\theta) (1 - \sin^2(\theta)) d\theta \\
&= \frac{2^{k+1}}{\pi} \int_{-\frac{\pi}{2}}^{\frac{\pi}{2}} \sin^{2k}(\theta) d\theta - (2k+1)m_{2k} \quad \text{from integrating by parts}
\end{aligned}$$

Hence $m_{2k} = \frac{2^{k+1}}{\pi(2k+2)} \int_{-\frac{\pi}{2}}^{\frac{\pi}{2}} \sin^{2k}(\theta) d\theta = \frac{\binom{2k}{k}}{k+1} = C_k$ if fully evaluated.

So to conclude, we have shown

$$\lim_{N \rightarrow \infty} \langle \bar{L}_N, x^k \rangle = |\mathcal{W}_{k, \frac{k}{2}+1}| = C_{\frac{k}{2}} = m_k = \langle \sigma, x^k \rangle$$

and we are done. \square

Lemma 2:

For all $k \in \mathbb{N}$ and $\epsilon > 0$,

$$\lim_{N \rightarrow \infty} \mathbb{P}(|\langle L_N, x^k \rangle - \langle \bar{L}_N, x^k \rangle| > \epsilon) = 0$$

Proof of Lemma 2:

By Chebyshev's Inequality, we can equivalently prove

$$\lim_{N \rightarrow \infty} |\mathbb{E}[\langle L_N, x^k \rangle^2] - \langle \bar{L}_N, x^k \rangle^2| = 0$$

As similar to previous (3.1),

$$\mathbb{E}[\langle L_N, x^k \rangle^2] - \langle \bar{L}_N, x^k \rangle^2 = \frac{1}{N^2} \sum_{\substack{i_1, \dots, i_k=1 \\ i'_1, \dots, i'_k=1}}^N \bar{T}_{\mathbf{i}, \mathbf{i}'}^N \quad (3.4)$$

where $\bar{T}_{\mathbf{i}, \mathbf{i}'}^N = \mathbb{E}[T_{\mathbf{i}}^N T_{\mathbf{i}'}^N] - \mathbb{E}[T_{\mathbf{i}}^N] \mathbb{E}[T_{\mathbf{i}'}^N]$

Given a set \mathcal{J} , a \mathcal{J} -sentence is a finite sequence of \mathcal{J} -words at least one word long. Two sentences are equivalent, $a_1 \sim a_2$, if there is a bijection on \mathcal{J} mapping one to the other. For a sentence $a = (w_1, w_2, \dots, w_n)$, the support $\text{supp}(a) = \bigcup_{i=1}^n \text{supp}(w_i)$, and the weight, $wt(a)$, is the cardinality of $\text{supp}(a)$.

We can assign graphs to sentences as previous. Given a sentence $a = (w_1, \dots, w_k) w_i = s_1^i, s_2^i, \dots, s_{l(w_i)}^i$ we set $G_a = (V_a, E_a)$ with vertices $V_a = \text{supp}(a)$ and edges $E_a = \{\{s_j^i, s_{j+1}^i\}; j = 1, \dots, l(w_i) - 1, i = 1, \dots, k\}$. We define the set of self edges and connecting edges in the same way as before, but note the graphs may be disconnected this time and sentence a defines k paths G_a . Define N_e^a the number of times the union of the paths traverse edge e . Equivalent sentences generate the same graph and passage counts.

We recall closed words $w_{\mathbf{i}}, w_{\mathbf{i}'}$ of length $k + 1$ and define the two word sentence $a_{\mathbf{i}, \mathbf{i}'} = (w_{\mathbf{i}}, \mathbf{i}')$, then

$$\begin{aligned} \bar{T}_{\mathbf{i}, \mathbf{i}'}^N &= \frac{1}{N^k} \left(\prod_{e \in E_{a_{\mathbf{i}, \mathbf{i}'}}^c} \mathbb{E}[Z_{1,2}^{N_e^{a_{\mathbf{i}, \mathbf{i}'}}}] \prod_{e \in E_{a_{\mathbf{i}, \mathbf{i}'}}^s} \mathbb{E}[Y_1^{N_e^{a_{\mathbf{i}, \mathbf{i}'}}}] \right. \\ &\quad \left. - \prod_{e \in E_{w_{\mathbf{i}}}^c} \mathbb{E}[Z_{1,2}^{N_e^{w_{\mathbf{i}}}}] \prod_{e \in E_{w_{\mathbf{i}}}^s} \mathbb{E}[Y_1^{N_e^{w_{\mathbf{i}}}}] \prod_{e \in E_{w_{\mathbf{i}'}}^c} \mathbb{E}[Z_{1,2}^{N_e^{w_{\mathbf{i}'}}}] \prod_{e \in E_{w_{\mathbf{i}'}}^s} \mathbb{E}[Y_1^{N_e^{w_{\mathbf{i}'}}}] \right) \end{aligned} \quad (3.5)$$

Hence $\bar{T}_{\mathbf{i}, \mathbf{i}'}^N = 0$ unless $N_e^{a_{\mathbf{i}, \mathbf{i}'}} \geq 2$ for all edges, and $E_{w_{\mathbf{i}}} \cap E_{w_{\mathbf{i}'}} \neq \emptyset$

Again if $N \geq t$, there are $C_{N,t}$ N -sentences equivalent to a given N -sentence weight t

Let $\mathcal{W}_{k,t}^{(2)}$ denote the set of representatives for equivalence classes of sentences a with weight t comprised of 2 closed t -words (w_1, w_2) each with length $k + 1$ with $N_e^a \geq 2$ for all $e \in E_a$ and $E_{w_1} \cap E_{w_2} \neq \emptyset$. Then from (3.4) and (3.5),

$$\begin{aligned} \mathbb{E}[\langle L_N, x^k \rangle^2] - \langle \bar{L}_N, x^k \rangle^2 &= \sum_{t=1}^{2k} \frac{C_{N,t}}{N^{k+2}} \sum_{a=(w_1, w_2) \in \mathcal{W}_{k,t}^{(2)}} \left(\prod_{e \in E_{a_{\mathbf{i}, \mathbf{i}'}}^c} \mathbb{E}[Z_{1,2}^{N_e^{a_{\mathbf{i}, \mathbf{i}'}}}] \prod_{e \in E_{a_{\mathbf{i}, \mathbf{i}'}}^s} \mathbb{E}[Y_1^{N_e^{a_{\mathbf{i}, \mathbf{i}'}}}] \right. \\ &\quad \left. - \prod_{e \in E_{w_{\mathbf{i}}}^c} \mathbb{E}[Z_{1,2}^{N_e^{w_{\mathbf{i}}}}] \prod_{e \in E_{w_{\mathbf{i}}}^s} \mathbb{E}[Y_1^{N_e^{w_{\mathbf{i}}}}] \prod_{e \in E_{w_{\mathbf{i}'}}^c} \mathbb{E}[Z_{1,2}^{N_e^{w_{\mathbf{i}'}}}] \prod_{e \in E_{w_{\mathbf{i}'}}^s} \mathbb{E}[Y_1^{N_e^{w_{\mathbf{i}'}}}] \right) \end{aligned} \quad (3.6)$$

In view of (3.6), it now suffices to show that $\mathcal{W}_{k,t}^{(2)}$ is empty for $t \geq k + 2$ since N^{k+2} beats $C_{N,t} \approx N^t$ for $t \geq k + 2$. We will show the stronger $t \geq k + 1$.

Note that if $a \in \mathcal{W}_{k,t}^{(2)}$, then G_a is a connected graph with t vertices, and at most k edges, which is impossible for $t > k + 1$. Considering when $t = k + 1$, G_a is a tree and each edge must be visited by a exactly twice. But since the path

from w_1 in tree G_a is closed, we must visit each edge an even number of times. Thus the set of edges from w_1 is disjoint to the set of edges from w_2 , creating a contradiction, and we are done. \square

4 Final Proof of Wigner's:

To conclude the proof of Wigner's Semicircular Theorem for real Wigner matrices, we must check that for bounded continuous functions f

$$\lim_{N \rightarrow \infty} \langle L_N, f \rangle = \langle \sigma, f \rangle \quad \text{in probability}$$

By Chebyshev's Inequality,

$$\begin{aligned} \mathbb{P}(\langle L_N, |x|^k \mathbf{1}_{\{|x| > B\}} \rangle > \epsilon) &\leq \frac{1}{\epsilon} \mathbb{E}[\langle L_N, |x|^k \mathbf{1}_{\{|x| > B\}} \rangle] \\ &= \frac{1}{\epsilon} \langle \bar{L}_N, |x|^k \mathbf{1}_{\{|x| > B\}} \rangle \\ &\leq \frac{1}{\epsilon} \langle \bar{L}_N, \frac{x^{2k}}{B^k} \rangle \\ &= \frac{\langle \bar{L}_N, x^{2k} \rangle}{\epsilon B^k} \end{aligned}$$

Hence by Lemma 1 we have,

$$\begin{aligned} \limsup_{N \rightarrow \infty} \mathbb{P}(\langle L_N, |x|^k \mathbf{1}_{\{|x| > B\}} \rangle > \epsilon) &\leq \limsup_{N \rightarrow \infty} \frac{\langle \bar{L}_N, x^{2k} \rangle}{\epsilon B^k} \\ &\leq \frac{\langle \sigma, x^{2k} \rangle}{\epsilon B^k} = \frac{C_k}{\epsilon B^k} \leq \frac{4_k}{\epsilon B^k} \end{aligned}$$

Now we set $B = 5$ and note the LHS is increasing in k

$$0 \leq \limsup_{N \rightarrow \infty} \mathbb{P}(\langle L_N, |x|^k \mathbf{1}_{\{|x| > B\}} \rangle > \epsilon) \leq \frac{1}{\epsilon} \left(\frac{4}{5}\right)^k \xrightarrow{k \rightarrow \infty} 0$$

$$\text{Hence } \limsup_{N \rightarrow \infty} \mathbb{P}(\langle L_N, |x|^k \mathbf{1}_{\{|x| > B\}} \rangle > \epsilon) = 0 \quad (4.1)$$

Fix an f with support $[-5, 5]$ and $\delta > 0$, then by the Weierstrass Approximation Theorem, we can find a polynomial $\mathcal{Q}_\delta(x) = \sum_{i=0}^L c_i x^i$ such that

$$\sup_{x: |x| \leq \delta} |\mathcal{Q}_\delta(x) - f(x)| \leq \frac{\delta}{8}$$

Then,

$$\begin{aligned} \mathbb{P}(|\langle L_N, f \rangle - \langle \sigma, f \rangle| > \delta) &\leq \mathbb{P}(|\langle L_N, \mathcal{Q}_\delta \rangle - \langle \bar{L}_N, \mathcal{Q}_\delta \rangle| > \frac{\delta}{4}) \\ &\quad + \mathbb{P}(|\langle \bar{L}_N, \mathcal{Q}_\delta \rangle - \langle \sigma, \mathcal{Q}_\delta \rangle| > \frac{\delta}{4}) \\ &\quad + \mathbb{P}(|\langle L_N, \mathcal{Q}_\delta \mathbf{1}_{\{|x| > B\}} \rangle| > \frac{\delta}{4}) \end{aligned}$$

The first term tends to zero in N by Lemma 2, the second term tends to zero in N by Lemma 1, and the third term tends to 0 in N by (4.1). Hence this finishes the proof. \square