

Hybrid Quantum–Classical Computation via Clifford Frame Tracking

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Outline

1 Introduction

2 Theoretical Background

3 Implementation

Introduction

- Quantum hardware is noisy and expensive.
- Physical qubits are prone to decoherence.
- Many quantum gates are classically simulable.

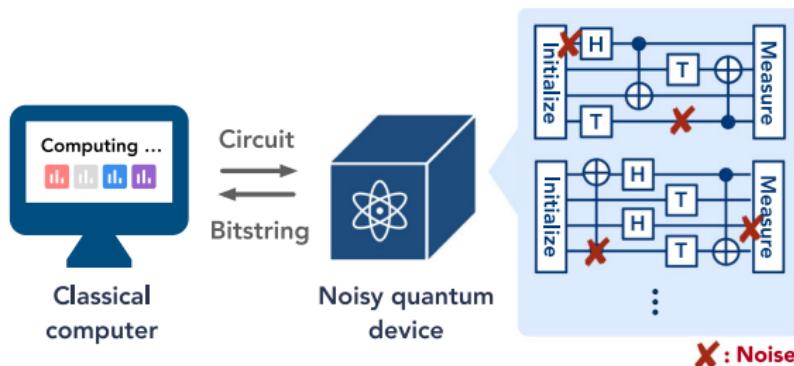


Figure: Schematic of a hybrid algorithm implementation (NISQ) [1]

General quantum gates

Any quantum circuit can be decomposed as

$$U = C_{T+1} N_T C_T \cdots N_1 C_1,$$

where $C_k \in \mathcal{C}_n$ are Clifford unitaries and N_k are non-Clifford unitaries.

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$$\{H, S, CNOT, T\}, \quad T = e^{-i\frac{\pi}{8}Z}.$$

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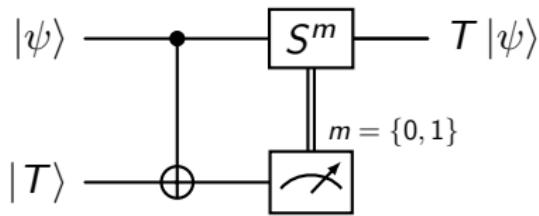
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Physical implementation of T-state injection:

$$|T\rangle = \frac{|0\rangle + e^{i\pi/4}|1\rangle}{\sqrt{2}}$$



Clifford Simulation

From the stabilizer formalism,

$$\text{Stab}(|\psi\rangle) = \langle P_1, \dots, P_n \rangle \iff \text{Stab}(U|\psi\rangle) = \langle UP_1U^\dagger, \dots, UP_nU^\dagger \rangle,$$

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A stabilizer generator $P = P_j$ can be represented by bit-vectors:

$$P = i^{p(x,z)} X^x Z^z, \quad (x|z) = (x_1 \ x_2 \ \cdots \ x_n \mid z_1 \ z_2 \ \cdots \ z_n).$$

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For each new generator UPU^\dagger ,

$$(x|z) \mapsto (x|z) M \in \mathbb{F}_2^{2n}$$

where $M \in \text{Sp}(2n, \mathbb{F}_2)$ is the **symplectic matrix** of the conjugation.

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In the tableau format:

$$(\mathbf{D}, \mathbf{C})^\top = \text{diag}(\mathbf{x}, \mathbf{z}) M \in \mathbb{F}_2^{2n \times 2n},$$

where $\mathbf{x} = (x_1, \dots, x_n)^\top$ and $\mathbf{z} = (z_1, \dots, z_n)^\top$.

Frame Tracking

Define

$$|\phi_k\rangle := F_k |\psi_k\rangle,$$

where $F_k = C_k C_{k-1} \cdots C_1 \in \mathcal{C}_n$ is a sequence of unitaries interleaved by non-Clifford unitaries N .

Time evolution for non-Clifford gate

$$|\psi_{k+1}\rangle = N |\psi_k\rangle, \quad |\phi_{k+1}\rangle = M |\phi_k\rangle,$$

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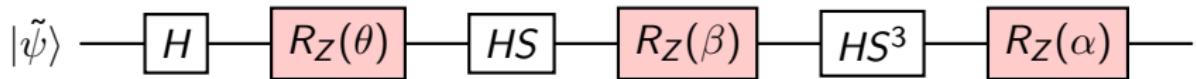
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Pauli rotation as a non-Clifford gate:

$$N = e^{-i\theta/2P}, \quad P \in \mathcal{P}_n.$$

Frame Tracking: Single-qubit circuit

Exact circuit

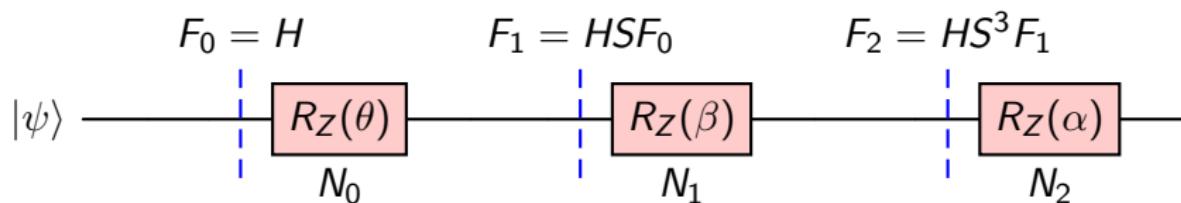


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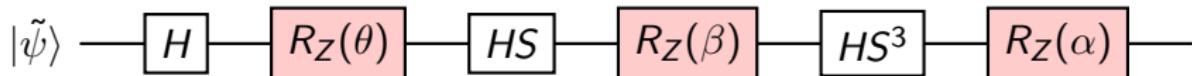


Logical circuit

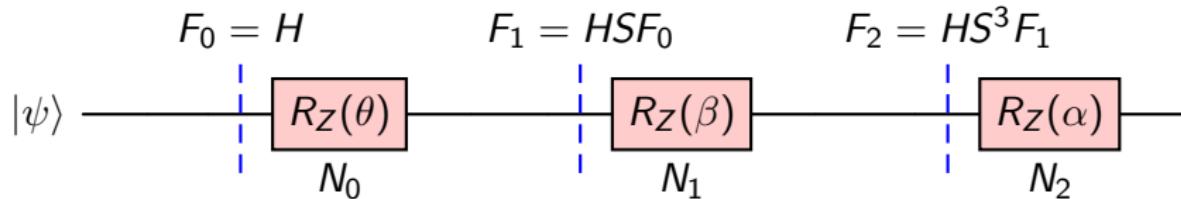


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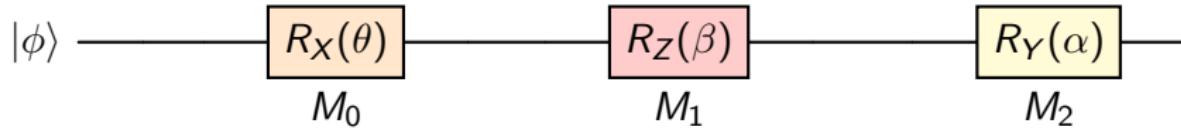


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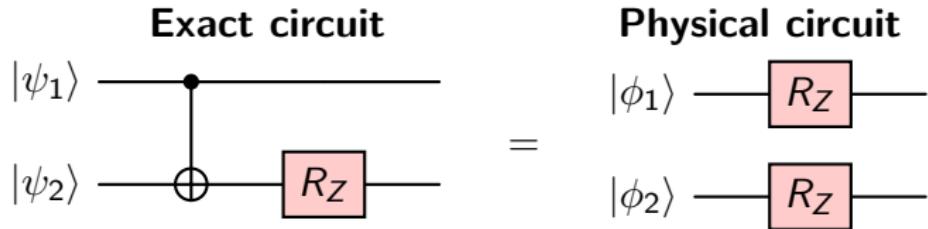
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Physical circuit



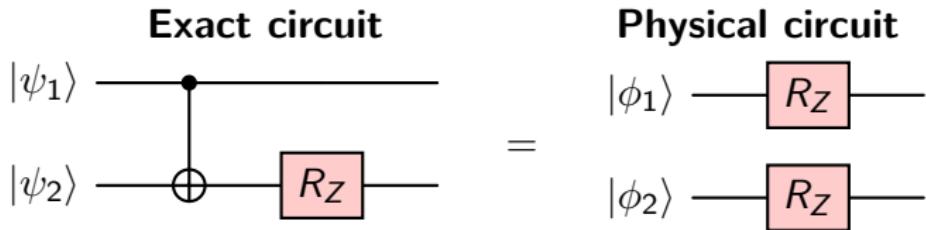
Frame Tracking: Two-qubits circuits

Example 1

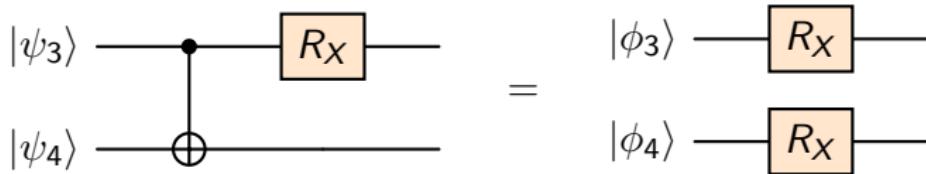


Frame Tracking: Two-qubits circuits

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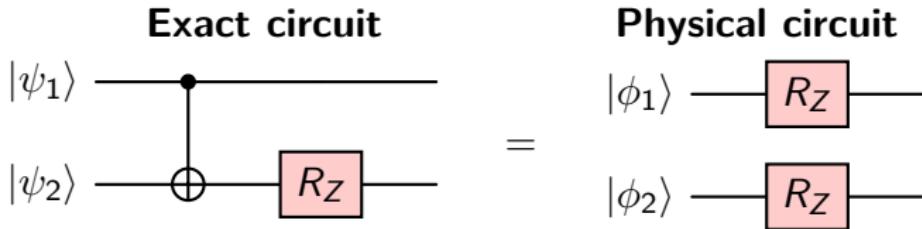


Example 2

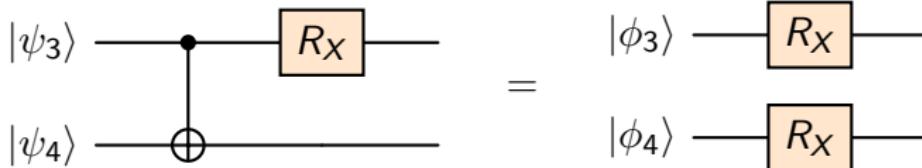


Frame Tracking: Two-qubits circuits

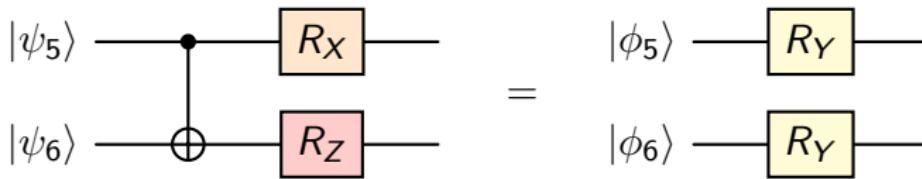
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Example 2



Example 3



CNOT gate

- Clifford propagation through CNOT gates [2].

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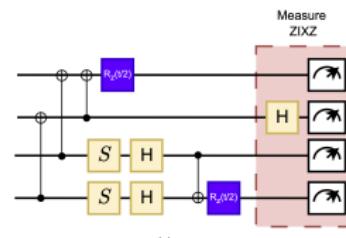
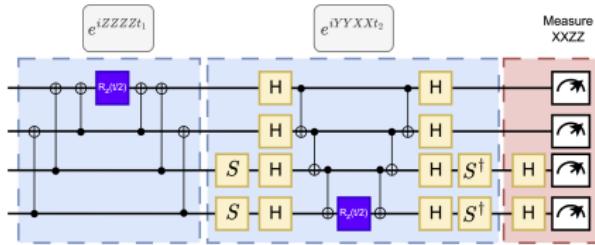
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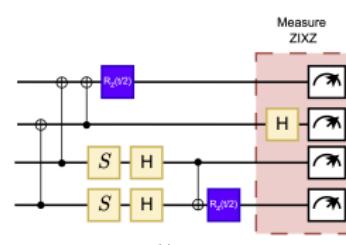
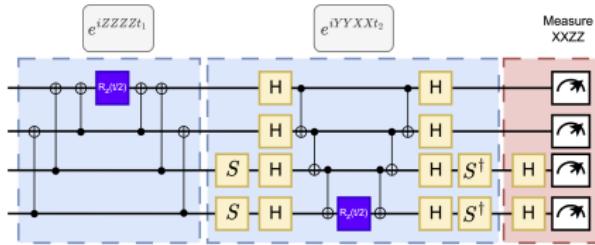


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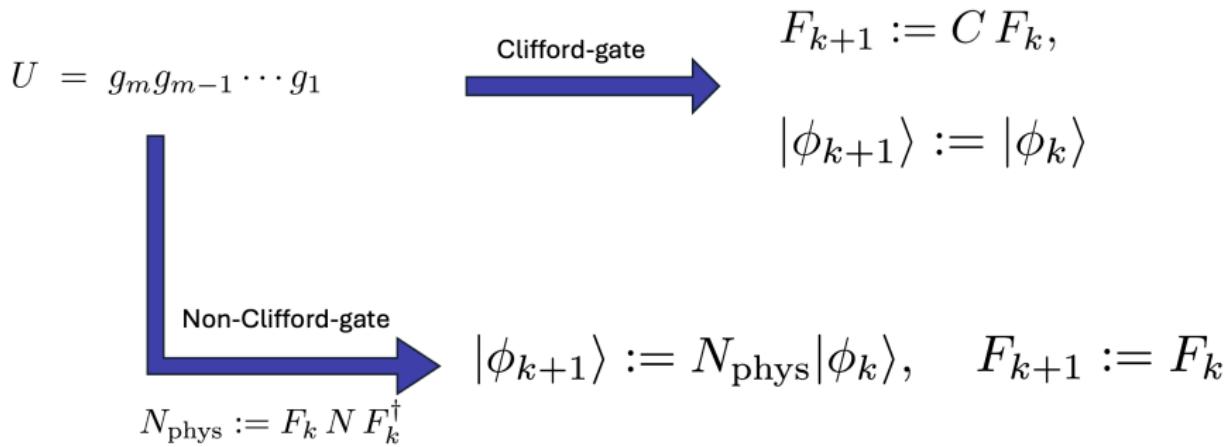
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- Correction with ancilla qubit.

Logic of the implementation



Final state: $|\phi_m\rangle = F_m |\psi_m\rangle = F_m U |\psi_0\rangle$

References I

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References II

4. Liu, J., Gonzales, A., Huang, B., Saleem, Z. H. & Hovland, P. *QuCLEAR: Clifford Extraction and Absorption for Quantum Circuit Optimization.* en. arXiv:2408.13316 [quant-ph]. Feb. 2025. <http://arxiv.org/abs/2408.13316> (2025).