

# MATH/CS 27700 FALL 2020: HW 4

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**Due: October 30th at 9PM.**

**Note 1:** As noted, the "Take Home Test" Question (in red) must be worked on individually.

**Note 2:** Answers must be fully explained (unless otherwise specified). Follow the Homework policy (see syllabus) and remember to *write up all solutions on your own*.

- (1) Suppose that we have theories in propositional logic [ie. a pair of a language  $L$  and a subset  $\Gamma \subset \text{Form}(L)$ ]  $(L_1, \Gamma_1)$ ,  $(L_2, \Gamma_2)$ , with  $L_1, L_2$  not necessarily disjoint. You may assume  $L_1$  and  $L_2$  are finite, however this actually isn't needed. [ For my original approach I wrote that your approach to HW3 problem 3 may be useful, however for the better approach this is not true.]
  - (a) **Take Home Test Question:** Suppose in the language  $L_1 \cup L_2$  we have that  $\Gamma_1 \models \Gamma_2$ . Show that there exists  $\Gamma_{12} \subset \text{Form}(L_1 \cap L_2)$  such that  $\Gamma_1 \models \Gamma_{12}$  and  $\Gamma_{12} \models \Gamma_2$ .
  - (b) Show that if  $\Gamma_1 = \{\phi\}$ , and  $\Gamma_2 = \{\psi\}$ , we can pick  $\Gamma_{12} = \{\theta\}$  for some  $\theta \in \text{Form}(L_1 \cap L_2)$   
 Bonus: Remove the assumption that  $L_1$  and  $L_2$  are finite in Q1b), Either by using Q2b) [note this doesn't actually prove 1(b) unless we provide an alternative proof of 2(b) which is possible], or otherwise [better].
- (2) In this question assume the result of 1 (b) *without the finiteness assumption*. That is to say assume for  $\phi \in \text{Form}(L_1)$ ,  $\psi \in \text{Form}(L_2)$ , such that  $\phi \models \psi$  (for the language  $L_1 \cup L_2$ ) there is  $\theta \in \text{Form}(L_1 \cap L_2)$  such that  $\phi \models \theta$  and  $\theta \models \psi$ .
  - (a) Suppose that we have consistent theories  $(L_1, \{\phi\})$  and  $(L_2, \{\chi\})$  as before. Suppose that  $(L_1 \cup L_2, \{\phi, \chi\})$  is inconsistent. Show that there exists  $\theta$  in  $\text{Form}(L_1 \cap L_2)$  such that  $\phi \models \theta$ , and  $\chi \models \neg\theta$ .
  - (b) Suppose that we have consistent theories  $(L_1, \Gamma_1)$  and  $(L_2, \Gamma_2)$  as before. Suppose that  $(L_1 \cup L_2, \Gamma_1 \cup \Gamma_2)$  is inconsistent. Show that there exists  $\theta$  in  $\text{Form}(L_1 \cap L_2)$  such that  $\Gamma_1 \models \theta$ , and  $\Gamma_2 \models \neg\theta$ .
- (3) Suppose  $L' \supset L$ . Let  $\phi \in \text{Form}(L')$ . Suppose that for any two truth functions  $t_1, t_2$  (for  $L'$ ) such that  $t_1|_{\text{Form}(L)} = t_2|_{\text{Form}(L)}$ , we have that  $t_1(\phi) = t_2(\phi)$ .  
 Show that<sup>1</sup>  $\models \phi \leftrightarrow \psi$  for some  $\psi \in \text{Form}(L)$ .
- (4) **Take Home Test Question.** In this question do not assume Lindebaum's lemma, or Gödel's completeness Theorem. Show that<sup>2</sup>  $\vdash \alpha \rightarrow ((\neg\alpha) \rightarrow \beta)$ .
- (5) In this question do not assume Lindebaum's lemma, or Gödel's completeness Theorem. Show that if  $\Gamma, \phi \vdash \psi$  then  $\Gamma, \neg\psi \vdash \neg\phi$ .

<sup>1</sup>By  $\phi \leftrightarrow \psi$  I mean  $(\phi \rightarrow \psi) \wedge (\psi \rightarrow \phi)$

<sup>2</sup>Note that this formula is simply  $\alpha \rightarrow (\alpha \vee \beta)$ . Note also that by  $\vdash \phi$  we mean that  $\emptyset \vdash \phi$ .