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1  ---------- MODULE Voting ----------
2  EXTENDS Sets
3  ----------------------------------
4  CONSTANT Value, Acceptor, Quorum

6  ASSUME QuorumAssumption ≜
7      ∧  ∀ Q ∈ Quorum : Q ⊆ Acceptor
8      ∧  ∀ Q1, Q2 ∈ Quorum : Q1 ∩ Q2 ≠ {}

10 THEOREM QuorumNonEmpty ≜ ∀ Q ∈ Quorum : Q ≠ {}
11 BY QuorumAssumption

13 Ballot ≜ Nat
14 ----------------------------------
15 VARIABLES votes, maxBal

17 TypeOK ≜ ∧ votes ∈ [Acceptor → SUBSET (Ballot × Value)]
18          ∧ maxBal ∈ [Acceptor → Ballot ∪ {−1}]
19 ----------------------------------
20 VotedFor(a, b, v) ≜ ⟨b, v⟩ ∈ votes[a]

22 DidNotVoteAt(a, b) ≜ ∀ v ∈ Value : ¬ VotedFor(a, b, v)

24 ShowsSafeAt(Q, b, v) ≜
25   ∧ ∀ a ∈ Q : maxBal[a] ≥ b  have promised
26   ∧ ∃ c ∈ −1 .. (b − 1) :
27       ∧ (c ≠ −1) ⇒ ∃ a ∈ Q : VotedFor(a, c, v)
28       ∧ ∀ d ∈ (c + 1) .. (b − 1), a ∈ Q : DidNotVoteAt(a, d)
29 ----------------------------------
30 Init ≜
31     ∧ votes = [a ∈ Acceptor ↦ {}]
32     ∧ maxBal = [a ∈ Acceptor ↦ −1]

34 IncreaseMaxBal(a, b) ≜
35   ∧ b > maxBal[a]
36   ∧ maxBal′ = [maxBal EXCEPT ![a] = b]  make promise
37   ∧ UNCHANGED votes

39 VoteFor(a, b, v) ≜
40     ∧  maxBal[a] ≤ b  keep promise
41     ∧  ∀ vt ∈ votes[a] : vt[1] ≠ b
42     ∧  ∀ c ∈ Acceptor \ {a} :
43          ∀ vt ∈ votes[c] : (vt[1] = b) ⇒ (vt[2] = v)
44     ∧  ∃ Q ∈ Quorum : ShowsSafeAt(Q, b, v)  safe to vote
45     ∧  votes′ = [votes EXCEPT ![a] = votes[a] ∪ {⟨b, v⟩}]  vote
46     ∧  maxBal′ = [maxBal EXCEPT ![a] = b]  make promise
47 ----------------------------------
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$$Next \triangleq$$
$$\quad \exists\, a \in Acceptor,\, b \in Ballot :$$
$$\quad\quad \lor\, IncreaseMaxBal(a,\, b)$$
$$\quad\quad \lor \exists\, v \in Value : VoteFor(a,\, b,\, v)$$

$$Spec \triangleq Init \land \Box[Next]_{\langle votes,\, maxBal\rangle}$$

$$ChosenAt(b,\, v) \triangleq$$
$$\quad \exists\, Q \in Quorum : \forall\, a \in Q : VotedFor(a,\, b,\, v)$$

$$chosen \triangleq \{v \in Value : \exists\, b \in Ballot : ChosenAt(b,\, v)\}$$

$$CannotVoteAt(a,\, b) \triangleq$$
$$\quad \land\, maxBal[a] > b$$
$$\quad \land\, DidNotVoteAt(a,\, b)$$

$$NoneOtherChoosableAt(b,\, v) \triangleq$$
$$\quad \exists\, Q \in Quorum :$$
$$\quad\quad \forall\, a \in Q : VotedFor(a,\, b,\, v) \lor CannotVoteAt(a,\, b)$$

$$SafeAt(b,\, v) \triangleq$$
$$\quad \forall\, c \in 0\,..\,(b-1) : NoneOtherChoosableAt(c,\, v)$$

$$VotesSafe \triangleq$$
$$\quad \forall\, a \in Acceptor,\, b \in Ballot,\, v \in Value :$$
$$\quad\quad VotedFor(a,\, b,\, v) \Rightarrow SafeAt(b,\, v)$$

$$OneVote \triangleq$$
$$\quad \forall\, a \in Acceptor,\, b \in Ballot,\, v,\, w \in Value :$$
$$\quad\quad VotedFor(a,\, b,\, v) \land VotedFor(a,\, b,\, w) \Rightarrow (v = w)$$

$$OneValuePerBallot \triangleq$$
$$\quad \forall\, a1,\, a2 \in Acceptor,\, b \in Ballot,\, v1,\, v2 \in Value :$$
$$\quad\quad VotedFor(a1,\, b,\, v1) \land VotedFor(a2,\, b,\, v2) \Rightarrow (v1 = v2)$$

$$Inv \triangleq TypeOK \land VotesSafe \land OneValuePerBallot$$

THEOREM $AllSafeAtZero \triangleq \forall\, v \in Value : SafeAt(0,\, v)$
BY DEF $SafeAt$

THEOREM $ChoosableThm \triangleq$
$$\quad\quad \forall\, b \in Ballot,\, v \in Value :$$
$$\quad\quad\quad ChosenAt(b,\, v) \Rightarrow NoneOtherChoosableAt(b,\, v)$$
BY DEF $ChosenAt$, $NoneOtherChoosableAt$

THEOREM $OneVoteThm \triangleq OneValuePerBallot \Rightarrow OneVote$
BY DEF $OneValuePerBallot$, $OneVote$

```
THEOREM VotesSafeImpliesConsistency ≜
  ASSUME VotesSafe, OneVote, chosen ≠ {}
  PROVE  ∃ v ∈ Value : chosen = {v}
⟨1⟩1. PICK v ∈ Value : v ∈ chosen
  BY DEF chosen
⟨1⟩2. SUFFICES ASSUME NEW w ∈ chosen
               PROVE  w = v
  BY ⟨1⟩1, ⟨1⟩2
⟨1⟩3. ASSUME NEW b1 ∈ Ballot, NEW b2 ∈ Ballot, b1 < b2,
             NEW v1 ∈ Value, NEW v2 ∈ Value,
             ChosenAt(b1, v1) ∧ ChosenAt(b2, v2)
     PROVE  v1 = v2
  ⟨2⟩1. SafeAt(b2, v2)
    BY ⟨1⟩3, QuorumAssumption, SMT DEF ChosenAt, VotesSafe
  ⟨2⟩2. QED
    BY ⟨1⟩3, ⟨2⟩1, QuorumAssumption, Z3
    DEFS CannotVoteAt, DidNotVoteAt, OneVote,
         ChosenAt, NoneOtherChoosableAt, Ballot, SafeAt
⟨1⟩4. QED
  BY QuorumAssumption, ⟨1⟩1, ⟨1⟩2, ⟨1⟩3, Z3
  DEFS Ballot, ChosenAt, OneVote, chosen

THEOREM ShowsSafety ≜
          TypeOK ∧ VotesSafe ∧ OneValuePerBallot ⇒
            ∀ Q ∈ Quorum, b ∈ Ballot, v ∈ Value :
              ShowsSafeAt(Q, b, v) ⇒ SafeAt(b, v)
  BY QuorumAssumption, Z3
  DEFS Ballot, TypeOK, VotesSafe, OneValuePerBallot, SafeAt,
    ShowsSafeAt, CannotVoteAt, NoneOtherChoosableAt, DidNotVoteAt

THEOREM SafeAtStable ≜ Inv ∧ Next ∧ TypeOK′ ⇒
                          ∀ b ∈ Ballot, v ∈ Value :
                            SafeAt(b, v) ⇒ SafeAt(b, v)′
  OMITTED
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THEOREM Invariance ≜ Spec ⇒ □Inv
⟨1⟩ USE DEF Inv
⟨1⟩1. Init ⇒ Inv
  BY DEF Init, TypeOK, VotesSafe, OneValuePerBallot, VotedFor
⟨1⟩2. Inv ∧ [Next]⟨votes, maxBal⟩ ⇒ Inv′
  ⟨2⟩ SUFFICES ASSUME Inv, [Next]⟨votes, maxBal⟩
               PROVE  Inv′
    OBVIOUS
  ⟨2⟩1.CASE Next
    ⟨3⟩ SUFFICES ASSUME NEW a ∈ Acceptor, NEW b ∈ Ballot,
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$\lor$ *IncreaseMaxBal*$(a,\ b)$
$\lor \exists\, v \in \mathit{Value} : \mathit{VoteFor}(a,\ b,\ v)$
PROVE Inv'
BY $\langle 2\rangle 1$ DEF *Next*
$\langle 3\rangle 1$.CASE *IncreaseMaxBal*$(a,\ b)$
$\langle 4\rangle 1$. TypeOK'
BY $\langle 3\rangle 1$ DEF *TypeOK*, *IncreaseMaxBal*
$\langle 4\rangle 2$. $\mathit{VotesSafe}'$
$\langle 5\rangle$ SUFFICES ASSUME NEW $a\_1 \in \mathit{Acceptor}'$, NEW $b\_1 \in \mathit{Ballot}'$, NEW $v \in \mathit{Value}'$
PROVE $\mathit{VotedFor}(a\_1,\ b\_1,\ v)' \Rightarrow \mathit{SafeAt}(b\_1,\ v)'$
BY DEF *VotesSafe*
$\langle 5\rangle 1$. $\forall\, aa \in \mathit{Acceptor},\ bb \in \mathit{Ballot},\ vv \in \mathit{Value}$:
$\mathit{VotedFor}(aa,\ bb,\ vv) \equiv \mathit{VotedFor}(aa,\ bb,\ vv)'$
BY $\langle 3\rangle 1$ DEF *IncreaseMaxBal*, *VotedFor*
$\langle 5\rangle 2$. $\forall\, aa \in \mathit{Acceptor},\ bb \in \mathit{Ballot}$:
$\mathit{maxBal}[aa] > bb \Rightarrow \mathit{maxBal}'[aa] > bb$
BY $\langle 3\rangle 1$ DEF *IncreaseMaxBal*, *TypeOK*, *Ballot*
$\langle 5\rangle 3$. $\forall\, aa \in \mathit{Acceptor},\ bb \in \mathit{Ballot}$:
$\mathit{DidNotVoteAt}(aa,\ bb) \Rightarrow \mathit{DidNotVoteAt}(aa,\ bb)'$
BY $\langle 3\rangle 1$ DEF *IncreaseMaxBal*, *DidNotVoteAt*, *VotedFor*
$\langle 5\rangle 4$. $\forall\, aa \in \mathit{Acceptor},\ bb \in \mathit{Ballot}$:
$\mathit{CannotVoteAt}(aa,\ bb) \Rightarrow \mathit{CannotVoteAt}(aa,\ bb)'$
BY $\langle 3\rangle 1$, $\langle 5\rangle 2$, $\langle 5\rangle 3$ DEF *IncreaseMaxBal*, *CannotVoteAt*
$\langle 5\rangle 5$. $\forall\, bb \in \mathit{Ballot},\ vv \in \mathit{Value}$:
$\mathit{NoneOtherChoosableAt}(bb,\ vv) \Rightarrow \mathit{NoneOtherChoosableAt}(bb,\ vv)'$
BY $\langle 5\rangle 1$, $\langle 5\rangle 4$, *QuorumAssumption*DEFS *NoneOtherChoosableAt*
$\langle 5\rangle 6$. QED
BY $\langle 5\rangle 1$, $\langle 5\rangle 5$ DEF *TypeOK*, *Ballot*, *VotesSafe*, *SafeAt*
$\langle 4\rangle 3$. $\mathit{OneValuePerBallot}'$
BY $\langle 3\rangle 1$ DEF *IncreaseMaxBal*, *OneValuePerBallot*, *VotedFor*
$\langle 4\rangle 4$. QED
BY $\langle 4\rangle 1$, $\langle 4\rangle 2$, $\langle 4\rangle 3$ DEF *Inv*
$\langle 3\rangle 2$. ASSUME NEW $v \in \mathit{Value}$,
$\mathit{VoteFor}(a,\ b,\ v)$
PROVE Inv'
$\langle 4\rangle$ SUFFICES ASSUME NEW $Q \in \mathit{Quorum}$,
$\mathit{ShowsSafeAt}(Q,\ b,\ v)$
PROVE Inv'
BY $\langle 3\rangle 2$ DEF *VoteFor*
$\langle 4\rangle 1$. TypeOK'
BY $\langle 3\rangle 2$ DEF *TypeOK*, *VoteFor*
$\langle 4\rangle 2$. $\mathit{VotesSafe}'$ Using *OneValuePerBallot* in *SafeAtStable*
$\langle 5\rangle$ SUFFICES ASSUME NEW $aa \in \mathit{Acceptor}'$, NEW $bb \in \mathit{Ballot}'$, NEW $vv \in \mathit{Value}'$,
$\mathit{VotedFor}(aa,\ bb,\ vv)'$
PROVE $\mathit{SafeAt}(bb,\ vv)'$

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186           BY DEF VotesSafe
187         ⟨5⟩1.CASE VotedFor(aa, bb, vv)
188           ⟨6⟩1. SafeAt(bb, vv)
189             BY ⟨5⟩1 DEF VotesSafe
190           ⟨6⟩ QED
191             BY ⟨4⟩1, ⟨6⟩1, SafeAtStable DEF Next
192         ⟨5⟩2.CASE ¬VotedFor(aa, bb, vv)
193           ⟨6⟩1. aa = a ∧ bb = b ∧ vv = v ∧ VotedFor(a, b, v)′
194             BY ⟨3⟩2, ⟨4⟩1, ⟨5⟩2 DEF VoteFor, VotedFor, TypeOK
195           ⟨6⟩ QED
196             BY ⟨4⟩1, ⟨6⟩1, ShowsSafety, SafeAtStable DEF VoteFor, Next
197         ⟨5⟩ QED
198           BY ⟨5⟩1, ⟨5⟩2
199       ⟨4⟩3. OneValuePerBallot′
200         BY ⟨3⟩2 DEF VoteFor, OneValuePerBallot, VotedFor, TypeOK
201       ⟨4⟩4. QED
202         BY ⟨3⟩2, ⟨4⟩1, ⟨4⟩2, ⟨4⟩3 DEF Inv
203     ⟨3⟩3. QED
204       BY ⟨2⟩1, ⟨3⟩1, ⟨3⟩2
205   ⟨2⟩2.CASE UNCHANGED ⟨votes, maxBal⟩
206     BY ⟨2⟩2
207     DEFS TypeOK, Next, VotesSafe, OneValuePerBallot,
208          VotedFor, SafeAt, NoneOtherChoosableAt, CannotVoteAt, DidNotVoteAt,
209          IncreaseMaxBal, VoteFor
210   ⟨2⟩3. QED
211     BY ⟨2⟩1, ⟨2⟩2
212 ⟨1⟩3. QED
213   BY ⟨1⟩1, ⟨1⟩2, PTL DEF Spec
214 ├──────────────────────────────────────────────────────────────────────────┤
215 C ≜ INSTANCE Consensus

217 THEOREM Spec ∧ Inv ⇒ C!Spec
218 ⟨1⟩1. Init ⇒ C!Init
219   BY QuorumAssumption, SetExtensionality, IsaM("force")
220    DEF Init, C!Init, chosen, ChosenAt, VotedFor
221 ⟨1⟩2. Next ∧ Inv ⇒ C!Next ∨ UNCHANGED chosen
222   ⟨2⟩1 SUFFICES ASSUME Next, InvPROVE C!Next ∨ UNCHANGED chosen
223     BY ⟨2⟩1
224   ⟨2⟩2. chosen ⊆ chosen′
225     BY ⟨2⟩1, QuorumAssumption, Z3   SMTT(10) fails
226      DEF Next, Inv, TypeOK, IncreaseMaxBal, chosen, ChosenAt, VotedFor, Ballot, VoteFor
227   ⟨2⟩3. chosen′ = {} ∨ ∃ v ∈ Value : chosen′ = {v}
228     ⟨3⟩1. PICK a ∈ Acceptor, b ∈ Ballot :
229               ∨ IncreaseMaxBal(a, b)
230               ∨ ∃ v ∈ Value : VoteFor(a, b, v)
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231         BY ⟨2⟩1  DEF Next
232       ⟨3⟩2.CASE IncreaseMaxBal(a, b)
233       ⟨3⟩3.CASE ∃ v ∈ Value : VoteFor(a, b, v)
234       ⟨3⟩q. QED
235         BY ⟨3⟩1, ⟨3⟩2, ⟨3⟩3, SMT
236     ⟨2⟩q. QED
237       BY ⟨2⟩1, ⟨2⟩2, ⟨2⟩3, OneVoteThm, VotesSafeImpliesConsistency, SetExtensionality, SMT
238        DEF Inv, C!Next
239   ⟨1⟩3. QED
240     PROOF OMITTED
241 ╚═══════════════════════════════════════════════════════════════════════════╝
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