```
1  ---------- MODULE PaxosProof ----------
2  EXTENDS TLAPS, PaxosTuple

4  WellFormedMessages ≜ ∀ m ∈ msgs :
5      ∧ m[1] = "1a" ⇒ m[2]     ∈ Ballot
6      ∧ m[1] = "1b" ⇒ ∧ m[2] ∈ Acceptor
7                      ∧ m[3] ∈ Ballot
8                      ∧ m[4] ∈ Ballot ∪ {−1}
9                      ∧ m[5] ∈ Value ∪ {None}
10     ∧ m[1] = "2a" ⇒ m[2] ∈ Ballot ∧ m[3] ∈ Value
11     ∧ m[1] = "2b" ⇒ m[2] ∈ Acceptor ∧ m[3] ∈ Ballot ∧ m[4] ∈ Value
12 ---------------------------------------
13 THEOREM WFmsgs ≜ TypeOK ⇒ WellFormedMessages
14   BY Z3DEFS Ballot, TypeOK, Message, WellFormedMessages

16 THEOREM typing ≜ Spec ⇒ □TypeOK
17 ⟨1⟩.USE DEFS Ballot, TypeOK
18 ⟨1⟩1. Init ⇒ TypeOK
19   BY SMTDEFS Init
20 ⟨1⟩2. TypeOK ∧ [Next]_vars ⇒ TypeOK′
21   PROOF OMITTED
22 ⟨1⟩.HIDE DEFS Ballot, TypeOK
23 ⟨1⟩3. QED
24   BY ⟨1⟩1, ⟨1⟩2, PTL DEF Spec
25 ---------------------------------------
26 StructOK1 ≜ ∀ a ∈ Acceptor : IF maxVBal[a] = −1
27                              THEN maxVal[a] = None
28                              ELSE ⟨maxVBal[a], maxVal[a]⟩ ∈ votes[a]

30 THEOREM Spec ⇒ □StructOK1
31 ⟨1⟩.USE DEFS Ballot, TypeOK, StructOK1
32 ⟨1⟩1. Init ⇒ StructOK1
33   BY Z3DEFS Init
34 ⟨1⟩2. TypeOK ∧ StructOK1 ∧ [Next]_vars ⇒ StructOK1′
35   BY WFmsgs, Z3DEFS Next, Phase1a, Phase2a, Phase1b, Phase2b, Send, votes,
36     WellFormedMessages, vars, Message
37 ⟨1⟩q. QED
38   BY ONLY ⟨1⟩1, ⟨1⟩2, typing, PTL DEF Spec
39 ---------------------------------------
40 StructOK2 ≜ ∀ m ∈ msgs :
41    (m[1] = "1b") ⇒ ∧ maxBal[m[2]] ≥ m[3]
42                    ∧ (m[4] ≥ 0) ⇒ ⟨m[4], m[5]⟩    ∈ votes[m[2]]

44 StructOK3 ≜ ∀ m ∈ msgs : m[1] = "2a" ⇒ ∧ ∃ Q ∈ Quorum : V!ShowsSafeAt(Q, m[2], m[3])
45                                        ∧ ∀ mm ∈ msgs : ∧ mm[1] = "2a"
46                                                        ∧ mm[2] = m[2]
```

$$\Rightarrow mm[3] = m[3]$$

$$StructOK4 \triangleq \forall\, m \in msgs : m[1] = \text{“2b”} \Rightarrow \land \exists\, mo \in msgs : \land\, mo[1] = \text{“2a”}$$
$$\land\, mo[2] = m[3]$$
$$\land\, mo[3] = m[4]$$
$$\land\, maxBal[m[2]] \geq m[3]$$
$$\land\, maxVBal[m[2]] \geq m[3]$$

$$StructOK5 \triangleq \forall\, m \in msgs : m[1] = \text{“1b”} \Rightarrow \forall\, d \in Ballot : m[4] < d \land d < m[3] \Rightarrow$$
$$\forall\, v \in Value : \neg\langle d,\, v\rangle \in votes[m[2]]$$

$$StructOK \triangleq \land\, TypeOK$$
$$\land\, StructOK1$$
$$\land\, StructOK2$$
$$\land\, StructOK3$$
$$\land\, StructOK4$$
$$\land\, StructOK5$$

THEOREM $struct\_lemma \triangleq Spec \Rightarrow \Box StructOK$
$\langle 1\rangle$.USE DEFS $Ballot$, $StructOK$, $TypeOK$, $StructOK1$, $StructOK2$, $StructOK4$, $StructOK5$
$\langle 1\rangle 1$. $Init \Rightarrow StructOK$
BY $Z3$DEFS $Init$
$\langle 1\rangle 2$. $StructOK \land [Next]_{vars} \Rightarrow StructOK'$
PROOF OMITTED
$\langle 1\rangle 3$. QED
BY $\langle 1\rangle 1$, $\langle 1\rangle 2$, PTL DEF $Spec$

THEOREM $Spec \Rightarrow \Box StructOK3$
$\langle 1\rangle$.USE DEFS $Ballot$, $TypeOK$, $StructOK3$
$\langle 1\rangle 1$. $Init \Rightarrow StructOK3$
BY $Z3$DEFS $Init$
$\langle 1\rangle 2$. $TypeOK \land StructOK \land [Next]_{vars} \Rightarrow StructOK3'$
PROOF OMITTED
$\langle 1\rangle 3$. QED
BY ONLY $\langle 1\rangle 1$, $\langle 1\rangle 2$, $struct\_lemma$, PTL DEF $Spec$, $StructOK$

$$Inv \triangleq TypeOK \land StructOK1 \land StructOK2 \land StructOK3 \land StructOK4 \land StructOK5$$

THEOREM $OtherMessage \triangleq \forall\, m1,\, m2 \in msgs',\, a,\, b \in \{\text{“1a”},\, \text{“2a”},\, \text{“1b”},\, \text{“2b”}\} :$
$$\land\, m1[1] = a \land m2[1] = b \land a \neq b$$
$$\land\, msgs' = msgs \cup \{m2\}$$
$$\Rightarrow m1 \in msgs$$
BY $Z3$

THEOREM $\forall\, b \in Ballot,\, v \in Value :$
$$Phase2a(b,\, v) \land Inv \Rightarrow \exists\, Q \in Quorum : V\,!\,ShowsSafeAt(Q,\, b,\, v)$$
$\langle 1\rangle 1$. SUFFICES ASSUME NEW $b \in Ballot$,

```
                    NEW v ∈ Value,
                    ¬∃ m ∈ msgs : m[1] = "2a" ∧ m[3] = b,
                    NEW Q ∈ Quorum,
                    LET Q1b ≜ {m ∈ msgs : ∧ m[1] = "1b"
                                          ∧ m[2] ∈ Q
                                          ∧ m[3] = b}
                        Q1bv ≜ {m ∈ Q1b : m[4] ≥ 0}
                    IN  ∧ ∀ a ∈ Q : ∃ m ∈ Q1b : m[2] = a
                        ∧ ∨ Q1bv = {}
                          ∨ ∃ m ∈ Q1bv :
                              ∧ m[5] = v
                              ∧ ∀ mm ∈ Q1bv : m[4] ≥ mm[4],
                    Send(⟨"2a", b, v⟩),
                    UNCHANGED ⟨maxBal, maxVBal, maxVal⟩,
                    Inv
          PROVE  V!ShowsSafeAt(Q, b, v)
  BY SMT DEF Phase2a
⟨1⟩.USE ⟨1⟩1  DEF Ballot, Inv, TypeOK
⟨1⟩2. V!ShowsSafeAt(Q, b, v)!1
  BY SMT DEF Ballot, Send, StructOK2
⟨1⟩.DEFINE Q1b ≜ {m ∈ msgs : ∧ m[1] = "1b"
                             ∧ m[2] ∈ Q
                             ∧ m[3] = b}
⟨1⟩.DEFINE Q1bv ≜ {m ∈ Q1b : m[4] ≥ 0}
⟨1⟩3. V!ShowsSafeAt(Q, b, v)!2
  ⟨2⟩1. SUFFICES ASSUME NEW c ∈ −1 .. b − 1
                 PROVE ∧ c ≠ −1 ⇒ (∃ a ∈ Q : V!VotedFor(a, c, v))
                       ∧ ∀ d ∈ c + 1 .. b − 1, a ∈ Q : V!DidNotVoteAt(a, d)
    BY SMT
  ⟨2⟩2. c ≠ −1 ⇒ (∃ a ∈ Q : V!VotedFor(a, c, v))
  ⟨2⟩3. ∀ d ∈ c + 1 .. b − 1, a ∈ Q : V!DidNotVoteAt(a, d)
  ⟨2⟩q. QED
    BY ⟨2⟩2, ⟨2⟩3, SMT
⟨1⟩q. QED
  BY ⟨1⟩2, ⟨1⟩3, Z3 DEF V!ShowsSafeAt
```

```
THEOREM Next ∧ Inv ⇒ V!Next ∨ UNCHANGED ⟨votes, maxBal⟩
⟨1⟩1. SUFFICES ASSUME Next, InvPROVE V!Next
  BY ⟨1⟩1, SMT
⟨1⟩.USE  DEF Next, V!Next
⟨1⟩2.CASE ∃ b ∈ Ballot : Phase1a(b)
⟨1⟩3.CASE ∃ b ∈ Ballot : ∃ v ∈ Value : Phase2a(b, v)
⟨1⟩4.CASE ∃ a ∈ Acceptor : Phase1b(a)
  BY ⟨1⟩4, Inv, WFmsgs, Z3T(10)
   DEF Phase1b, Inv, WellFormedMessages, Ballot, V!Ballot, V!IncreaseMaxBal, votes, Send
```

139 $\langle 1\rangle 5$.CASE $\exists\, a \in Acceptor : Phase2b(a)$
140 $\langle 2\rangle 1$. PICK $a \in Acceptor$, $m \in msgs$:
141 $\land\ m[1] =$ "2a"
142 $\land\ m[2] \geq maxBal[a]$
143 $\land\ maxBal' = [maxBal$ EXCEPT $![a] = m[2]]$
144 $\land\ maxVBal' = [maxVBal$ EXCEPT $![a] = m[2]]$
145 $\land\ maxVal' = [maxVal$ EXCEPT $![a] = m[3]]$
146 $\land\ Send(\langle$"2b", a, $m[2]$, $m[3]\rangle)$
147 BY $\langle 1\rangle 5$, $Z3$ DEF $Phase2b$
148 $\langle 2\rangle 2$. SUFFICES ASSUME NEW $a\_1 \in Acceptor$, NEW $b \in Nat$
149 PROVE $\lor\ V\,!\,IncreaseMaxBal(a\_1,\ b)$
150 $\lor \exists\, v \in Value : V\,!\,VoteFor(a\_1,\ b,\ v)$
151 BY $Z3$ DEF $V\,!\,Ballot$
152 $\langle 2\rangle$q. QED
153 $\langle 1\rangle$q. QED
154 BY $\langle 1\rangle 1$, $\langle 1\rangle 2$, $\langle 1\rangle 3$, $\langle 1\rangle 4$, $\langle 1\rangle 5$, $Z3$
155