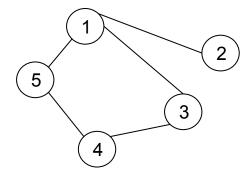
- Another data structure including vertices and edges
- Although trees are quite flexible, they have an inherent limitation in that they can only express hierarchical structures
- Fortunately, we can generalize a tree to form a graph, in which this limitation is removed
- Informally, a graph is a collection of nodes(vertices) and the connections between them(edges)

- Let's consider some definitions:
- A simple graph G = (V, E) consists of a (finite) set denoted by V, and a collection E, of unordered pairs {u, v} of distinct elements from V
- Each element of V is called a vertex or a point or a node, and each element of E is called an edge or a link
- The number of vertices, the cardinality of V, is called the order of graph and devoted by |V|
- The cardinality of E, called the size of graph, is denoted by |E|

- A path between vertices v₁ and v_n is a sequence of edges denoted v₁, v₂, ..., v_{n-1}, v_n
- If $v_1 = v_n$, the path is actually a cycle

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 $1 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 1$ is a cycle

- Two vertices are adjacent or neighbor if there is an edge between them
- The number of edges incident with a vertex v, is the degree of the vertex; if the degree is 0, v is called isolated

• Graphs can be represented in a number of ways, two of which are:

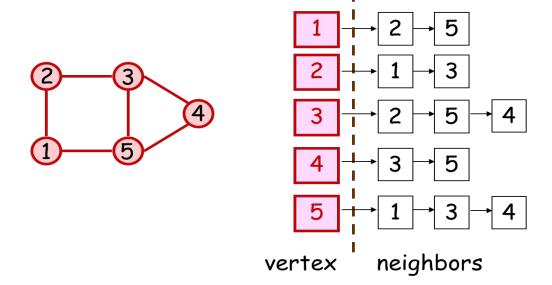
1. **Adjacency List:** One of the simplest is an adjacency list, where each vertex adjacent to a given vertex is listed

2. **Adjacency Matrix** is a $|V| \times |V|$ binary matrix where:

$$a_{ij} = \begin{cases} 1 & \text{if there exists an edge } (v_i v_j) \\ 0 & \text{otherwise} \end{cases}$$

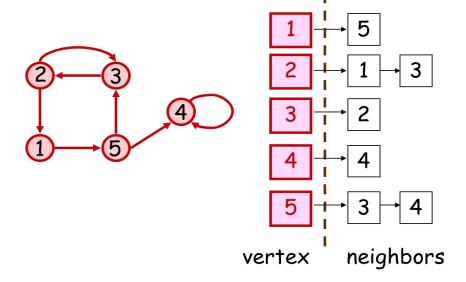
Adjacency List

 For each vertex u, store its neighbors in a linked list



Adjacency List

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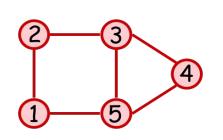


Adjacency List:

- Let G = (V, E) be an input graph
- Using Adjacency List representation :
- Space : O(|V| + |E|)
 - Excellent when |E| is small
- Easy to list all neighbors of a vertex
- Takes O(|V|) time to check if a vertex u is a neighbor of a vertex v

Adjacency Matrix

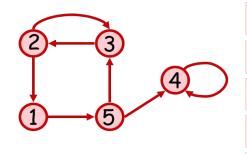
• Use a $|V| \times |V|$ matrix A such that A(u,v) = 1 if (u,v) is an edge A(u,v) = 0 otherwise



	1	2	3	4	5
1	0	1	0	0	1
2	1	0	1	0	0
3	0	1	0	1	1
4	0	0	1	0	1
5	1	0	1	1	0

Adjacency Matrix

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Adjacency Matrix:

- Let G = (V, E) be an input graph
- Using Adjacency Matrix representation :
- Space: O(|V|²)
 - o Bad when |E| is small
- O(1) time to check if a vertex u is a neighbor of a vertex v
- (|V|) time to list all neighbors

How to implement adjacency list using **list** container?

```
#include <list>
class Graph
    int V; // No. of vertices
    // Pointer to an array containing adjacency
    // lists
    list<int> *adj;
```

How to implement adjacency list using **list** container?

```
#include <list>
class Graph
    int V;
    list<int> *adj;
Graph::Graph(int V)
    this->V = V;
    adj = new list<int>[V];
```

How to implement adjacency matrix?

How to implement adjacency matrix? It can be implemented using a 2D array only!