

Solutions to Re-Exam 2020

Exercise 1

1.1.1 (b)

$$5(n^3)^2 = n^6 \sim \Theta(n^6)$$

$$\sqrt{n^7} = n^{3.5} \sim \Theta(n^{3.5})$$

1.1.2. (d)

$n^2\sqrt{\lg n}$ is $\Theta(n^2\sqrt{\lg n})$, and thus $\Omega(n)$

2.1. (c) This can be solved by the master method, the second case.

2.2. (a) $a = 4$, $b = 4$, and $f(n) = n$, so $f(n) = \Theta(n^{\log 4^4})$.

2.3. (a) $T(n) = \Theta(n^{\log 4^4} \lg n) = \Theta(n \lg n)$.

3.1. (c)

(a) The worst case time complexity of merge sort is $\Theta(n \lg n)$.

(b) Merge sort is not an in place sorting algorithm..

(c) The average case time complexity of Merge sort is $\Theta(n \lg n)$

3.2. (d)

First call on MERGE(,) A[1], A[2] are merged and sorted. Second call on MERGE(,) A[3], A[4] are merged and sorted. Third call, A[1..2] and A[3..4] are merged and sorted.

4. (c)

5. (d)

6.1 (c)

6.2 (a)

Exercise 2

2.1 (10 points) ENQUEUE(List L , int a): when inserting a new element with key being a into a queue, we need to insert the element to the end of the linked list L .

The pseudo code assumes that the list is not empty and has already N elements in the list.

```
ENQUEUE(List  $L$ , int  $a$ )
1  pointer  $x = L.head$ 
2  while  $x.next \neq NULL$  do
3       $x = x.next$ 
4  Create a new element  $y$  such that  $y.next = null$  and  $y.key = a$ .
5   $x.next = y$ ;
```

Asymptotic complexity: $\Theta(N)$

2.2 (10 points) Return the element pointed by the list's head pointer. And move head pointer to its next element.

```
DEQUEUE(List  $L$ )
1  integer  $m = L.head.key$ ;
2   $L.head = L.head.next$ ;
3  return  $m$ ;
```

Asymptotic complexity: $\Theta(1)$

Exercise 3

3.1 (7 point)

```
MINIMUMINTEGER(Array  $A$ , Integer  $p$ , Integer  $r$ )
1  if  $p < r$ 
2       $q = \lfloor \frac{(p+r)}{2} \rfloor$ 
3      Integer  $minleft = MINIMUMINTEGER(A, p, q)$ 
4      Integer  $minright = MINIMUMINTEGER(A, q + 1, r)$ 
5      return  $\min(minleft, minright)$ ;
6  else return  $A[p]$ ;
```

Call MinimumInteger($A, 1, n$) in the beginning.

3.2 (6 point) $T(n) = 2T(n/2) + \Theta(1)$; $T(1) = \Theta(1)$; two subproblems, each subproblem is half size of the original problem. merging is constant time.

3.3 (4 point) Master method, first case. $\Theta(n)$.

3.4 (3 point) Constant space overhead.